

# **AMATH 514** Assignment 3

Tyler Chen

**Exercise 2.16 (Stiemke's theorem)**

Prove that there exists a vector  $x > 0$  such that  $Ax = 0$  if and only if for each  $y$  satisfying  $y^T A \geq 0$  one has  $y^T A = 0$ .

The problem is show exactly one of the following is true,

$$\begin{aligned} \exists x > 0 : Ax &= 0 \\ \exists y : y^T A &\geq 0, y^T A \neq 0 \end{aligned}$$

Denote the  $j$ -th column of  $A$  by  $A_j$ . Denote the vector of all ones by  $e$ . Then,

$$y^T A e = \sum_{j=1}^n y^T A_j$$

We first prove a few useful equivalences.

Suppose  $y^T A \geq 0$ , then  $y^j A_k \geq 0$  for all  $j = 1, 2, \dots, n$  so  $y^T A e$  is the sum of non-negative terms. If  $y^T A \neq 0$  at least one term is nonzero (positive). Thus,

$$y^T A \neq 0 \iff y^T (-Ae) = -(y^T A e) < 0$$

Suppose  $\exists x > 0$  with  $Ax = 0$ . We can scale  $x$  so that all entries are at least one. Then  $z = x/(\min_i x_i) - e \geq 0$  and  $Az = -Ae$ .

Suppose  $\exists z \geq 0$  with  $Az = -Ae$ . Then  $A(z + e) = 0$  so  $x = z + e > 0$  solves  $Ax = 0$ .

Thus,

$$\exists x > 0 : Ax = 0 \iff \exists z \geq 0 : Az = -Ae$$

We can now apply Farkas Theorem. Indeed, start with,

$$\exists y : y^T A \geq 0, y^T A \neq 0$$

As explained above this is equivalent to,

$$\exists y : y^T A \geq 0, y^T (-Ae) < 0$$

Applying Farkas Theorem, this is equivalent to,

$$\nexists z \geq 0 : Az = -Ae$$

Again, as explained above this is equivalent to,

$$\nexists x > 0 : Ax = 0$$

This is the desired result. □

**Exercise 2.26**

Give an example of a matrix  $A$  and vectors  $b$  and  $c$  for which both  $\{x \mid Ax \leq b\}$  and  $\{y \mid y \geq 0; y^T A = c^T\}$  are empty.

Trivially we can pick  $A = [0]$ ,  $b = c = [-1]$ . Then  $Ax = 0 \not\leq -1$  and  $y^T A = 0 \neq -1$ .

We can easily characterize all matrices  $A \in \mathbb{R}^{2 \times 2}$ ,  $b \in \mathbb{R}^{2 \times 1}$ ,  $c \in \mathbb{R}^{1 \times 2}$  such that these sets are empty.

Visually,  $\{x \mid Ax \leq b\}$  corresponds to the intersection of two half planes in  $\mathbb{R}^2$ .

Suppose  $a_1x_1 + a_2x_2 \leq b_1$  is one of the half planes. Then we require the other half plane to have the form  $a_1x_1 + a_2x_2 \geq b_2$ , where  $b_2 > b_1$  so that their intersection is empty.

That is,  $\{x \mid Ax = b\}$  will be empty if and only if,

$$A = \begin{bmatrix} a_1 & a_2 \\ -a_1 & -a_2 \end{bmatrix}, \quad x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}, \quad b = \begin{bmatrix} b_1 \\ -b_2 \end{bmatrix}, \quad b_2 < -b_1$$

Now observe,

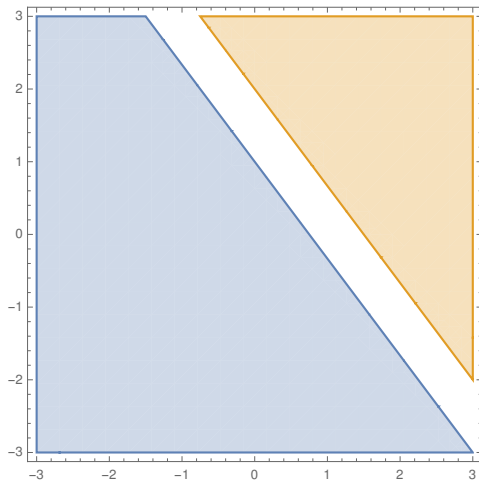
$$y^T A = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}^T \begin{bmatrix} a_1 & a_2 \\ -a_1 & -a_2 \end{bmatrix} = \begin{bmatrix} a_1y_1 - a_1y_2 \\ a_2y_1 - a_2y_2 \end{bmatrix}^T = \begin{bmatrix} a_1(y_1 - y_2) \\ a_2(y_1 - y_2) \end{bmatrix}^T$$

Finally, pick  $c = [c_1 \ c_2]$  such that  $a_1/a_2 \neq c_1/c_2$  (for instance, pick  $c_1 = a_1, c_2 \neq a_2$ ).

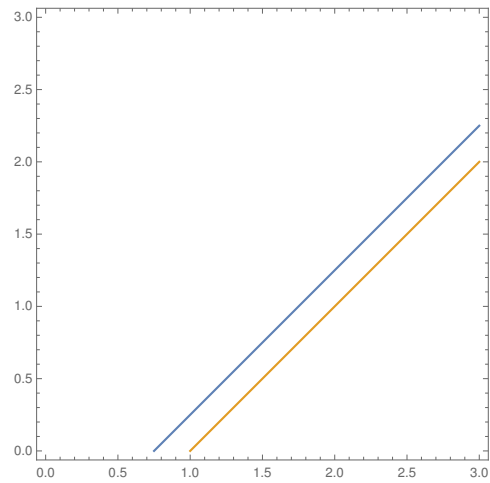
As an example,

$$A = \begin{bmatrix} 4 & 3 \\ -4 & -3 \end{bmatrix}, \quad b = \begin{bmatrix} 3 \\ -6 \end{bmatrix}, \quad c^T = \begin{bmatrix} 3 \\ 3 \end{bmatrix}$$

The intersection of the regions in Figures 1a and ?? show  $\{x \mid Ax \leq b\}$  and  $\{y \mid y \geq 0, y^T A = c^T\}$ . As we showed above, these intersections are both empty.



(a)  $Ax \leq b$



(b)  $y^T A = c^T$

**Exercise 2.27**

Let  $\tilde{x}$  be a feasible solution of  $\max\{c^T x \mid Ax \leq b\}$  and let  $\tilde{y}$  be a solution of  $\min\{y^T b \mid y \geq 0; y^T A = c^T\}$ . Prove that  $\tilde{x}$  and  $\tilde{y}$  are the optimum solutions of the minimum and maximum, respectively if and only if for each  $i = 1, 2, \dots, m$  one has:  $\tilde{y}_i = 0$  or  $a_i \tilde{x} = b_i$ .

Denote the  $i$ -th row of  $A$  by  $a_i$ .

First, note that if  $y$  is feasible, we have  $y^T A = c^T$  so that  $y^T Ax = c^T x$ .

Second, note also that for any  $i = 1, 2, \dots, m$ , if  $x$  is feasible, we have  $a_i x \leq b$  so that  $a_i x - b \geq 0$  and if  $y$  is feasible we have  $y_i \geq 0$ . Thus,  $y_i(a_i x - b) \geq 0$ .

We prove both directions at once. Indeed, suppose  $\tilde{x}$  and  $\tilde{y}$  are feasible.

By duality,  $\tilde{x}$  and  $\tilde{y}$  are the optimum solutions of the maximum and minimum respectively if and only if  $\tilde{y}^T b = c^T \tilde{x}$  which by the first note above is equivalent to,

$$\tilde{y}^T Ax = c^T x = \tilde{y}^T b$$

We can rearrange to find  $\tilde{y}^T (A\tilde{x} - b) = 0$ . Written in sum notation using the definition of matrix/vector multiplication we have,

$$\sum_{i=1}^m \tilde{y}_i (a_i \tilde{x} - b) = 0$$

Every term in this sum is non-negative by the second note above. Thus, the sum is zero if and only if each term is zero. That is, if and only if,

$$y_i(a_i \tilde{x} - b) = 0 \quad \forall i = 1, 2, \dots, m$$

Equivalently, for each  $i = 1, 2, \dots, m$  one has:  $\tilde{y}_i = 0$  or  $a_i \tilde{x} = b_i$ . □