

# TorFA: Tor Fingerprinting Attack

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## Abstract

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## 1 Introduction

Online anonymous communication methods are becoming more popular and receiving much interest from the research community today. One such method for online anonymity is Tor. Tor is a network that provides a method for creating greater privacy and security in Internet communications. Tor accomplishes this task by using a scheme called onion routing, which uses multiple layers of encryption and route redirection to provide anonymity to a limited set of web applications (e.g. browsing, chat clients). The goal of tor is to disallow an adversary from knowing who is requesting and receiving what information on the Internet.

It has been suggested by the tor developers that Tor could be vulnerable to a website fingerprinting attack. In a general fingerprinting attack, an adversary creates a signature of a website (based upon what the adversary can view on the network) and builds a statistical profile of the website based upon the different characteristics of the traffic he can view. This fingerprint is then stored to a fingerprint database, which will be used later by the attacker to compare against unknown data in order to identify what information a given user is viewing.

This paper introduces a novel method to perform a fingerprinting attack on the tor network. TorFA builds a statistical profile on the total number of packets downloaded from a given website as the fingerprint data. Then, the attacker listens to traffic at the tor entry router, and if the traffic matches the pattern of the fingerprint, the attacker knows what web page the client was viewing.

## 2 Related Work

## 3 Attack Methodology

The focus of our attack was to write a set of programs that performed two primary tasks: create signatures from a set of known data and compare an unknown capture in pcap format to the known set of signatures. The second program outputs the best match from the fingerprint database to the unknown capture file if such a website exists in the database.

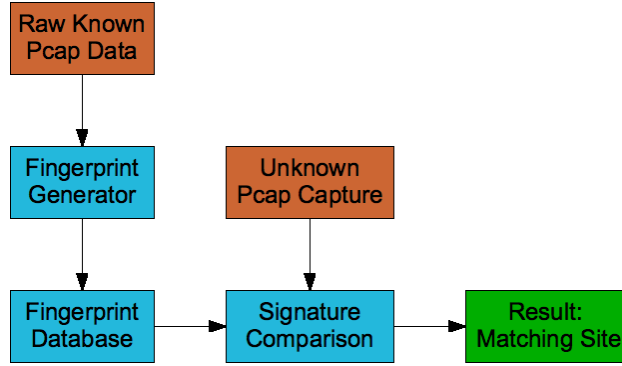


Figure 1: Program data and information flow

### 3.1 Attack Overview

The method by which we generated fingerprints was to intercept traffic between an entry onion router and the client capturing only traffic that was associated with the request of a webpage. Once several request and downloads are captured the data is processed by an application. This application takes as input a directory from which to compute the fingerprint, and several other options to allow for specific control of the application. This data is stored as a file in a fingerprints directory. In this way the fingerprint “database” is a directory where each file represents a webpage fingerprint.

The second application, TorFA, performs the comparison of an unknown webpage to the fingerprints in the database. It takes as input a capture of an unknown webpage (a single stream with only that page’s data), which it parses and subsequently creates a fingerprint. The application then compares that fingerprint to those in the database using K-L divergence. As a result the application outputs the best match from the database, if there exists such a match (i.e., within  $2 \times$  Standard Deviation of the total number of packets in the known fingerprint).

Figure 1 displays the general flow of information in the two applications. Raw known data is sent into the fingerprint creator, which outputs a fingerprint into the fingerprint database. Then an unknown data item is input into the fingerprint comparator, which outputs the fingerprint match.

### 3.2 Fingerprint Generation

The first part of the project included the development of a fingerprinting format that could be used for webpage comparisons. Hjelmvik [1] showed that probability vectors could be created by using frequency analysis of data. TorFA uses this same technique in the creation of probability vectors that reflect the unique data viewable to the attacker. In order to produce such probability vectors we captured and analyzed website traffic data looking for patterns that could be used to distinguish one website download from another.

#### 3.2.1 Data Collection and Analysis

Data collection of traffic in a tor flow turns out to be extremely hard to accomplish. The first step is defining where to capture the data, which for this project was chosen to be the client’s ethernet port. It was determined that this was satisfactory because an attacker could view data anywhere along the route from a client to its associated entry router, and thus without loosing any correctness our approach alleviated the need to implement our own entry router.

Collecting data was an arduous task because the tor implementation does not exactly follow the specification and the encrypted nature of the communications protocol. In experiments it was found that tor allows the entry routers to choose different port numbers for transferring its TLS data. Once this was understood a controlled environment was established to collect download data using Wireshark.

Once we had obtained this data we analyzed it for patterns to uniquely identify a given webpage download. In order to achieve this task a python script was created to parse and analyze the data for packet counting attributes. The script uses the command line wireshark equivalent, tshark, to parse the data for useful information.

It was determined that the primary attribute of the data was the number of packets sent and received from a website of specific size ranges. This is intuitive in the sense that a webpage has an approximately fixed size, which leads to a similar amount of packets for each download of the webpage. The team also considered using packet inter-arrival times as an attribute, but due to the dynamic nature of the tor network an exact time for a webpage download is not reproduced on a consistent basis. The analysis here showed one extremely interesting fact about the tor implementation in that the tor specification states that all traffic is to be sent in fixed 512 KB sized frames, but in the data we collected packets ranged from 512 – 4400 KB size packets. This turns out to benefit the fingerprint attack because it provides more vectors by which the data can be distinguished.

The end result of the analysis was to use a probability vector that included packet counts in the size ranges: 0 – 599, 600 – 799, 800 – 1449, and 1450 – 4500. These ranges were selected because frequency analysis showed that there were localized packet sizes that dominated the frequency counts, and as such the ranges were selected so that each one of these sizes was in a different range bin.

### 3.3 Fingerprint Comparison

The method used in this attack to match an unknown fingerprint to a known fingerprint is Kullback-Leibler Divergence (K-L divergence). K-L divergence is a measure in statistics that quantifies the difference between two probability distributions. It is also referred to as relative entropy. The following equation is used to measure the difference of a probability distribution Q from another probability distribution P:

$$D_{KL}(P||Q) = \sum_i P(i) \log \frac{P(i)}{Q(i)}$$

The greater this number is the less Q is like P. In other words, the closer the K-L divergence measure is to zero the more accurately Q matches P. In the tor fingerprinting attack, P is the probability distribution of the known fingerprints, and the distribution Q is a sample of data captured while listing to client and entry router communications.

One shortcoming of this method is that it does not take into consideration the potential for a given probability distribution to match another distribution and be far off the mean values. E.g., a distribution with packet counts of [200, 200, 200, 200] will match a distribution with packet counts of [1, 1, 1, 1]. Each one of these counts will produce the same probability distribution: [0.25, 0.25, 0.25, 0.25]. In order to account for this method each probability vector includes measures of standard deviation and variance. The standard deviation is used to filter out such occurrences thus reducing the false positive rates of the attack.

## 4 Adaptive Traffic Shaping: Adding and Dropping

One method proposed in literature to combat fingerprint attacks is the use of traffic shaping mechanisms. These methods include adaptive padding (e.g. adding empty packets into a stream) and defensive dropping

(e.g. randomly dropping good packets) of packets. It is unknown how these methods affect the K-L divergence measure of an adversary's attacks in the Tor network. Therefore, we modeled a small portion of the Tor network in Mobius, and simulated traffic shaping methods in order to identify the change in K-L divergence measures compared to the costs incurred in the tor network. The costs we consider are timing delays experienced by a client request, and the extraneous packet increase incurred in the tor network due to the adaptive traffic shaping mechanisms.

The model attempts to simulate a single download of a webpage as captured by Wireshark. It provides a simple implementation of the TCP protocol, namely we have modeled the TCP timeout event and ACK generation. We then model each entity in a minimal tor network: client, server, entry onion router, middle onion router, and exit onion router.

The primary purpose of the model was to quantify the efficacy of adaptive traffic techniques in terms of the K-L divergence gain verses the cost incurred in the Tor network. As such, two models and types of experiments were developed to obtain the previously mentioned reward variables. The first set of experiments were developed to collect the necessary information for constant add and drop rate traffic model. The simulation is set up to vary the add and drop packet probabilities from 0.0 to 0.3 at .025 length intervals (produces 169 data points). In this way, we could obtain divergence and cost measurements for multiple combinations of add and drop rates. We created a similar study for the variabilistic add and drop method which randomly changed this add/drop rate for each new packet as opposed to the constant rate for every packet.

## 5 Experimentation

It is important to note a few aspects of the way that we developed the experiments in TorFA.

1. We created the known fingerprints from approximately twenty to thirty samples each. This means that the sample packet counts the produce the probability vector could be bolstered by including several more samples per fingerprint.
2. We collected data on two different local host machines. This causes great variation in the fingerprints because they differ from one machine to the next. To some degree this choice creates a more robust system because the victim will be downloading webpages from different machines and host operating systems.

## 6 Results and Analysis

### 6.1 TorFA Results

Table 1 displays data from the experiments. It show the test case data and the solution provided by TorFA, and also displays the divergence measures. As can be seen, most of the websites were detected correctly, but this is an extremely small set of sample points.

It is important to recognize that the numbers of K-L divergence in the table by themselves are not a complete indicator of how well one sample matched against another sample. The estimate cannot give cross correlation between the different trials. When looking at multiple data output for a given run it can be seen that relatively the K-L divergence shifts. Meaning that on the smaller K-L divergence measurements it appears as though most of the measurements for that data run are smaller as well.

Test	Website	Match	KL-Divergence	Match
1	Facebook_1	facebook	0.02912	yes
2	Facebook_2	facebook	0.00708	yes
3	Facebook_3	facebook	0.01343	yes
4	ask.com_1	ask.com	0.00256	yes
5	ask.com_2	ask.com	0.00800	yes
6	ask.com_3	ask.com	0.00363	yes
7	Meebo_	meebo	0.00593	yes
8	Tsquare_1	tsquare	0.10467	yes
9	Tsquare_2	yahoo	0.02075	no
10	wikipedia main page	NONE	----	yes
11	fox	fox	0.00688	yes
12	apple	apple	0.09161	yes
13	amazon	msn	0.00397	no
14	gmail	gmail	0.01717	yes
15	Twitter	Twitter	0.00179	yes

Table 1: TorFA Results

Another result that is important to recognize is the Wikipedia main website download. The result for this data point came back as no match because this webpage is not in the fingerprint database. This proves to be a successful trial, but even more than that it shows that even though a webpage might be similar to another, for example the Tiananmen Square wikipedia webpage, there is still enough difference that it will not automatically match.

## 6.2 Adaptive Traffic Shaping

### 6.2.1 Results

Figures 2, 3, and 4 display the change in divergence at each constant add and drop packet rate. Figures 5, 6, and 7 show the percent increase of the labeled data. Figures 8 and 9 show the percent increase of the divergence, timing costs, and total transmitted packets in a bar chart so that the respective differences can be easily viewed.

### 6.2.2 Analysis

This data is full of interesting modeling and simulation analysis as well as pure fingerprint attack analysis. Some portions of the analysis will be left out for brevity's sake.

One important analysis is that the percent increase in total number of packets displays only the increase in packets for the single download of data as seen by the attacker. Involved in this data that is missing from the results are the added packets that do not make it through the communications channel of the host to entry router. This is significant because the increase would be much the same as the time's increase, but even beyond that the tor network would become even more congested than these results suggest because the percent increases here would need to be multiplied by the total number of active downloads in the network to correctly model the congestion.

Another component of these results is that they do not compare the attacker success rate verses the resultant costs. The divergence gain is two hundred percent in some cases, but what is not evident is how this affects TorFA.

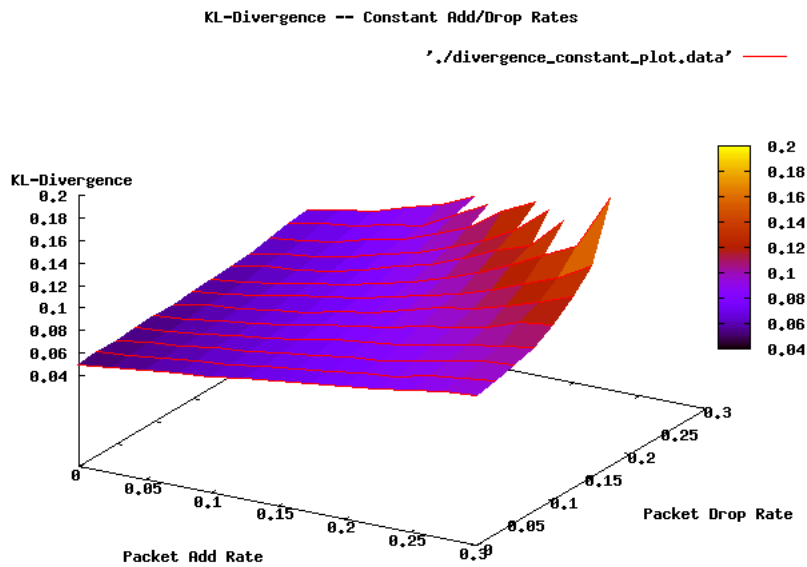


Figure 2: K-L Divergence Verses Constant Add/Drop Rates

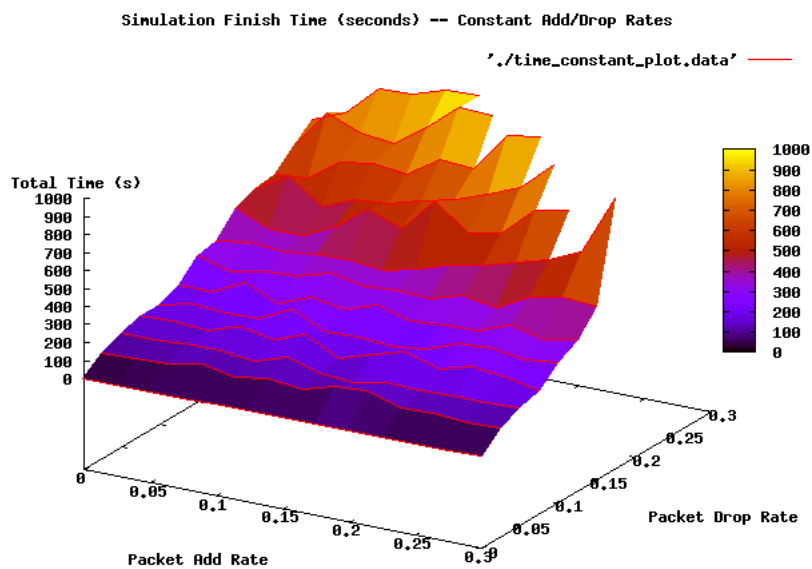


Figure 3: Time Verses Constant Add/Drop Rates

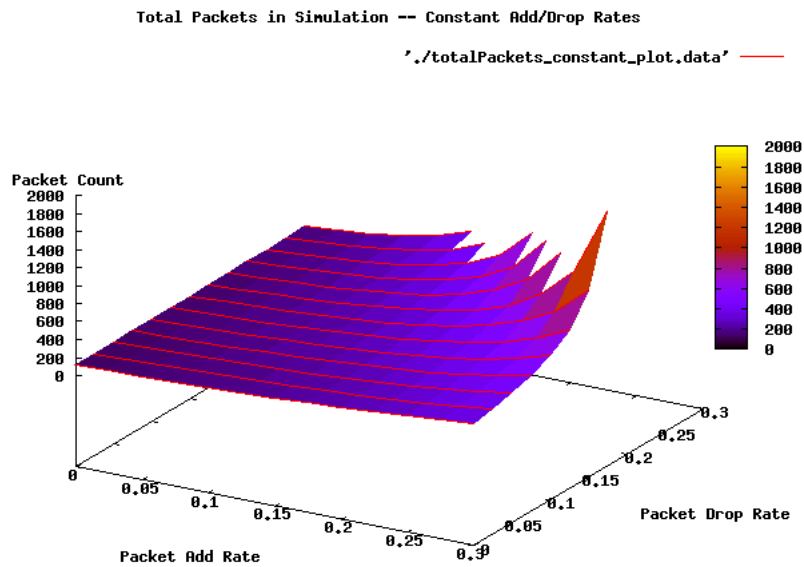


Figure 4: Number of Packets Verses Constant Add/Drop Rates

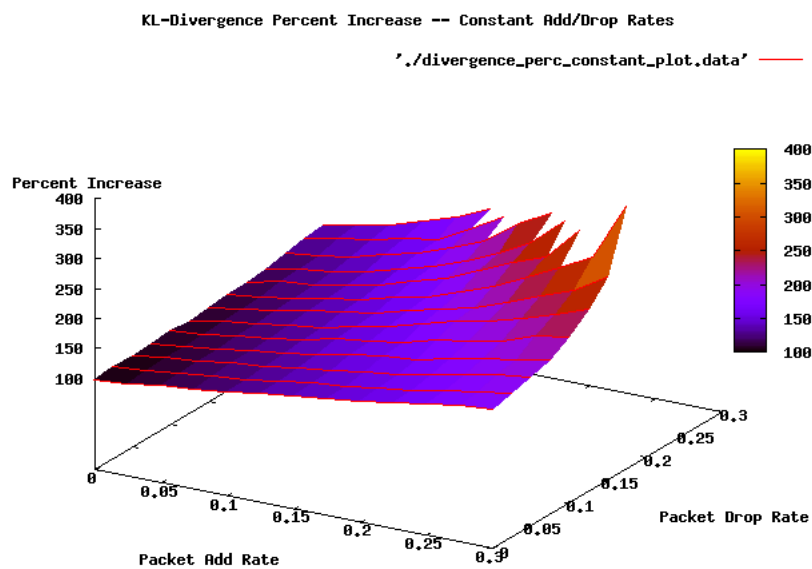


Figure 5: Percent Increase in the Divergence for Download Simulation

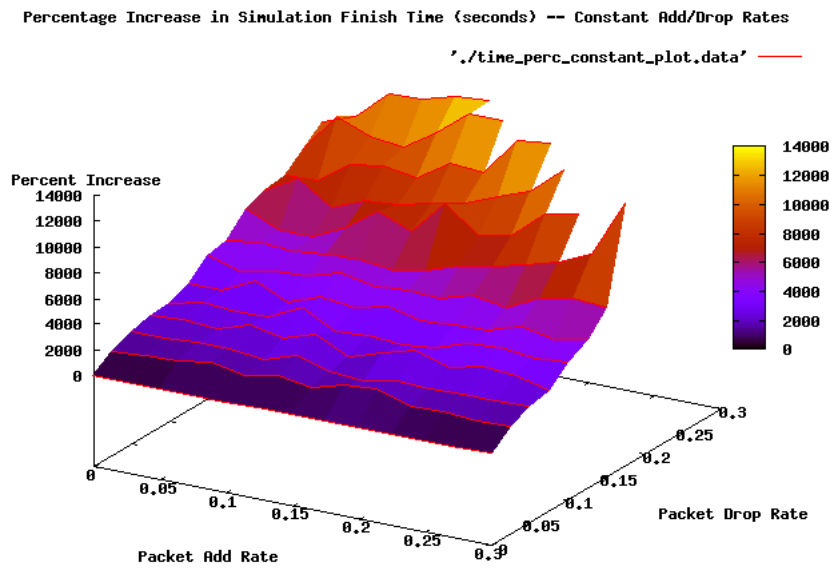


Figure 6: Percent Increase in Total Time of Download

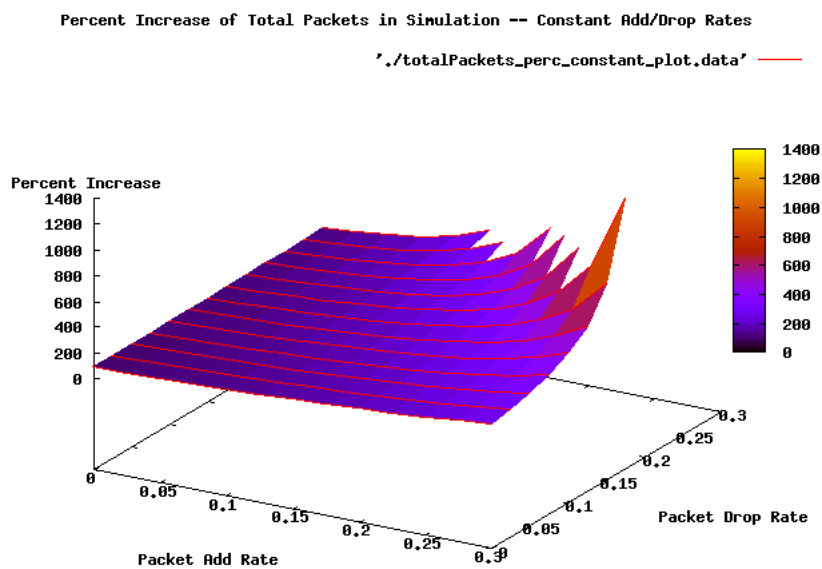


Figure 7: Percent Increase in Total Number of Packets



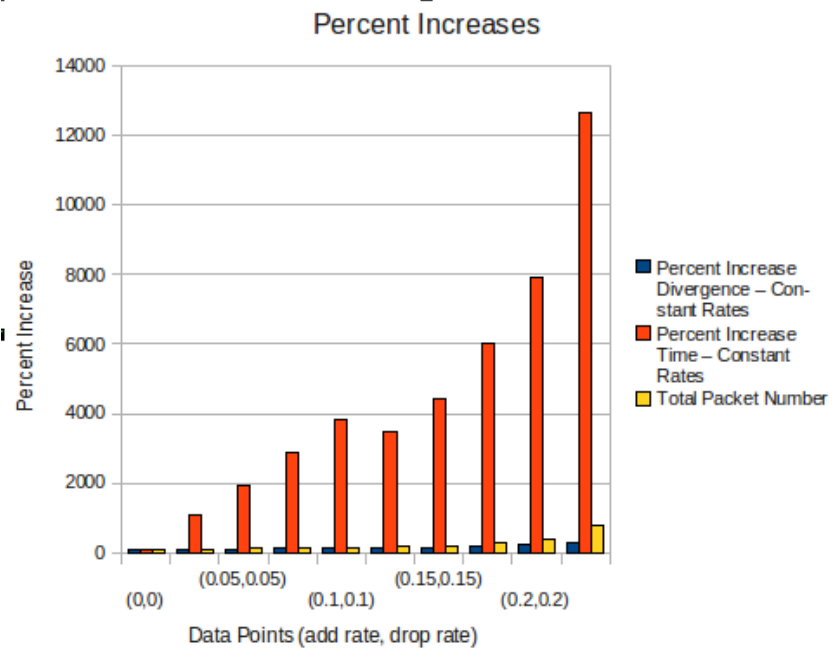


Figure 8: Two Dimensional Percent Increase: Divergence, Time, and Total Packets

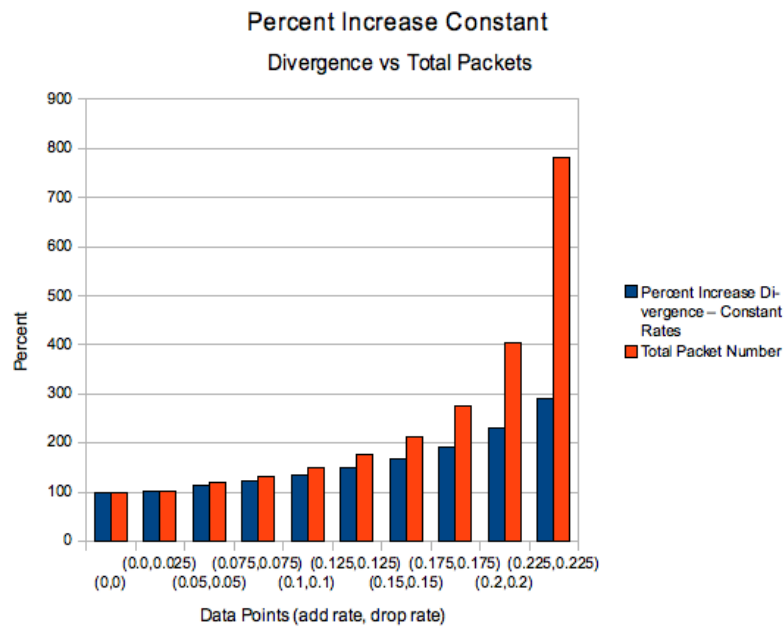


Figure 9: Two Dimensional Percent Increase: Divergence and Total Packets

## 7 Conclusion and Future Work

The TorFA attack has been shown to be successful at identifying a small set of webpage downloads. These results are skewed by the fact that we do not have a significant amount of data. The next phase of this work will be to develop automatic methods of performing data capture and processing. In this way a large enough sample space can be evaluated, thus allowing us the surety of statistical tests/methods that the attack does in fact work. Results such as average success rate, entropy levels of a given packet size, false-positive/false-negative probabilities, and etc. should be obtained so that a high level of confidence can be placed in these results.

The use of adaptive traffic shaping methods increases the KL-divergence measurements, but consequently incurs a cost in the form of client request timing and extra number of packets in the simulation. The first cost, client timing, is negligible because a client could have the extreme need for privacy and be willing to wait five to ten minutes to get information if it meant that they had a higher probability of being undetected. Of course, as this increases very quickly, it may not be useful in a practical setting.

On the other hand the second cost is an extremely important factor because of the effect that it has on the rest of the tor network. It is important to note that the tor relay routers are provided by users of the system offering up their bandwidth and hardware for free. As such, any extra cost to them would be a great burden and potentially reduce participation in the network, thus reducing the overall privacy of the users in the network, which depends upon having a high number of participants to guarantee privacy. This problem is further complicated by the fact that our data only includes the extra packets that are sent along the fingerprinting attack vector (i.e., the client and entry node communication) and not including the internal tor network. These added costs would rise similarly to the timing costs and thus be too great of a burden upon the tor network.

Additionally, it can be seen that a KL-divergence increase may not be that useful if the attacker's view includes this increase across all of the fingerprinted websites. More data collection and analysis are needed in this direction in order to quantify the true effects of a KL-divergence increase. This would come in the form of a relative measurement and require the explicit knowledge of the adversary's attack methodology to quantify accurately. This relative measurement is essential for further understanding of the fingerprint attack and defenses against it.

A few other future work items include:

- Develop a pcap parser to detect single streams inside a multistreamed data set. The TorFA code currently expects a pcap file that is already filtered for one website of traffic. The next step in this direction is to research the use of a Fourier transform technique in order to gain some information about individual data streams. This is important because tor multiplexes multiple TCP streams inside of one tor circuit.
- Implement the attack as an entry router in the Tor network in order to develop more accurate results.
- As described earlier, more data collection and testing is necessary.

## References

- [1] Erik Hjelmvik, *The SPID Algorithm: Statistical Protocol IDentification*. Swedish Internet Infrastructure Foundation, 2008.