



# Graphs

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# Why Graphs?

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- ❖ Biological networks

- ❖ Maps

- ❖ Social networks

- ❖ ...

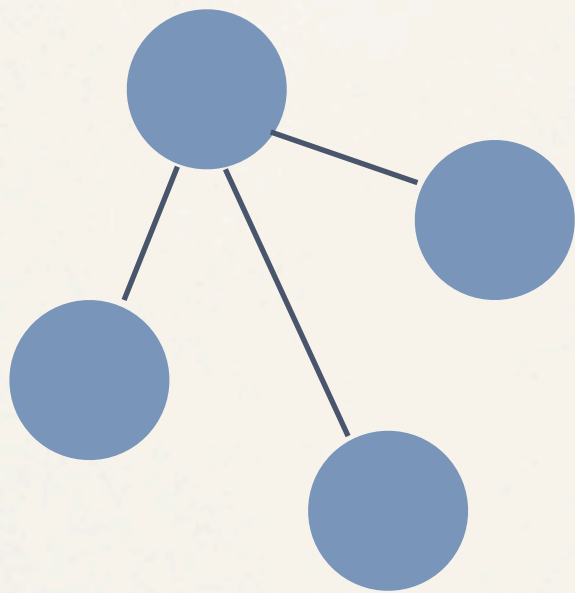
# Graphs: Basic Definitions

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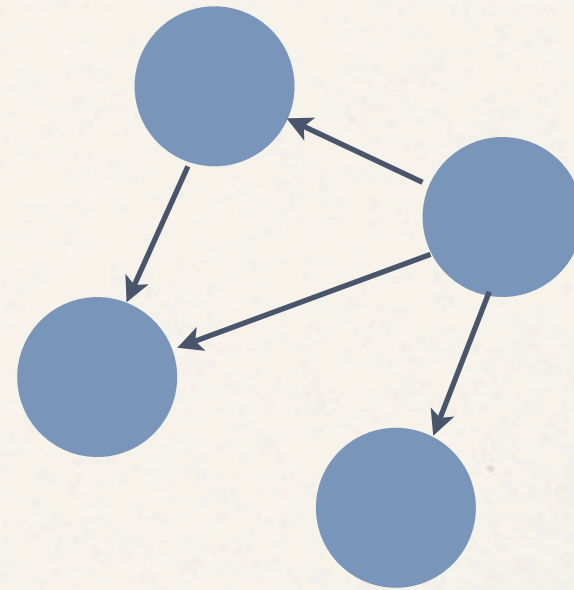


# The Connectivity of Networks

- ❖ Common to all networks we saw is that the connectivity of each of them can be represented via some form of a graph



Undirected graph



Directed graph

1 : node  
(or, vertex)

1 — 2 : edge  
 $\{1,2\}$

1 → 2 : directed edge  
 $(1,2)$

# Graphs

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- ❖ An **undirected graph**, or **graph** for short,  $G$ , is a pair  $(V, E)$ , where
  - ❖  $V = \{0, 1, \dots, n-1\}$  is a **nonempty** set of nodes, and
  - ❖  $E \subseteq \{\{i, j\} : i, j \in V\}$  is a set of unordered pairs, each of which corresponds to an undirected edge in the graph  $G$ .
- ❖ A **directed graph**, or **digraph** for short,  $G$ , is a pair  $(V, E)$ , where
  - ❖  $V = \{0, 1, \dots, n-1\}$  is a **nonempty** set of nodes, and
  - ❖  $E \subseteq (V \times V)$  is a set of ordered pairs, each of which corresponds to a directed edge in the graph  $G$ .

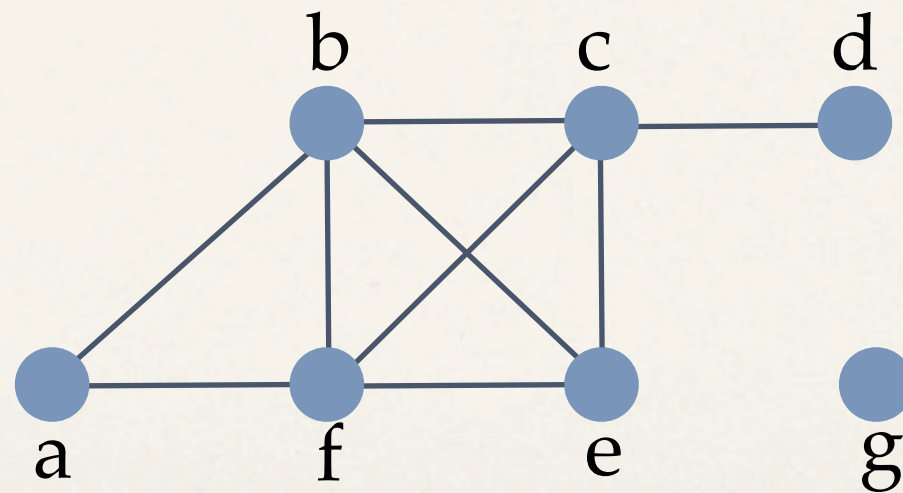


IMPORTANT: In this course, graphs have no self-loops or parallel edges, unless explicitly stated otherwise.

# Basic Terminology

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- ❖ Two nodes  $i$  and  $j$  in a graph  $G=(V,E)$  are called adjacent (or neighbors) in  $G$  if there is an edge between  $i$  and  $j$ ; that is, if  $\{i,j\}\in E$ .
- ❖ The degree of a node in an undirected graph is the number of edges incident with it. The degree of node  $i$  is denoted by  $\text{deg}(i)$ .



What are the degrees of the nodes?



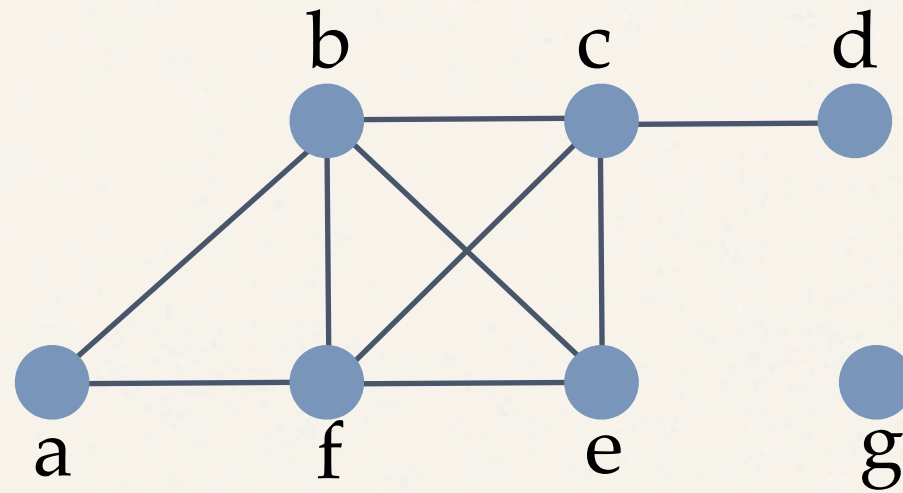
# Degree Distribution

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- ✦ Define  $p_k$  to be the fraction of nodes in the graph that have degree  $k$ .
- ✦ The degree distribution of a graph can be visualized by making a histogram of the  $p_k$  values.

# Degree Distribution

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$$\begin{aligned} p_0 &= 1/7 & p_1 &= 1/7 & p_2 &= 1/7 \\ p_3 &= 1/7 & p_4 &= 3/7 \end{aligned}$$

Notice that if  $m$  is the highest node degree, then:

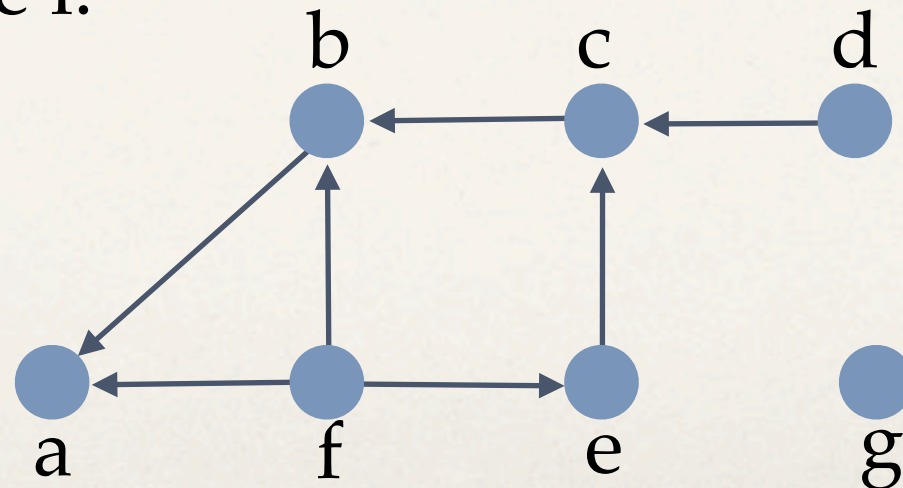
$$\sum_{k=0}^m p_k = 1$$



# Basic Terminology

$$i \rightarrow j$$

- ❖ If  $e=(i,j)$  is a directed edge from node  $i$  to node  $j$ , we say that  $i$  is the tail (or, initial node) of  $e$  and  $j$  is the head (or, terminal node) of  $e$ .
- ❖ The in-degree of a node  $i$  in a directed graph is the number of edges whose head is the node  $i$ . The in-degree of node  $i$  is denoted by  $\text{indeg}(i)$ .
- ❖ The out-degree of  $i$ , denoted by  $\text{outdeg}(i)$ , is the number of edges whose tail is the node  $i$ .



What are the in- and out-degrees of the nodes?



# In- and Out-Degree Distributions

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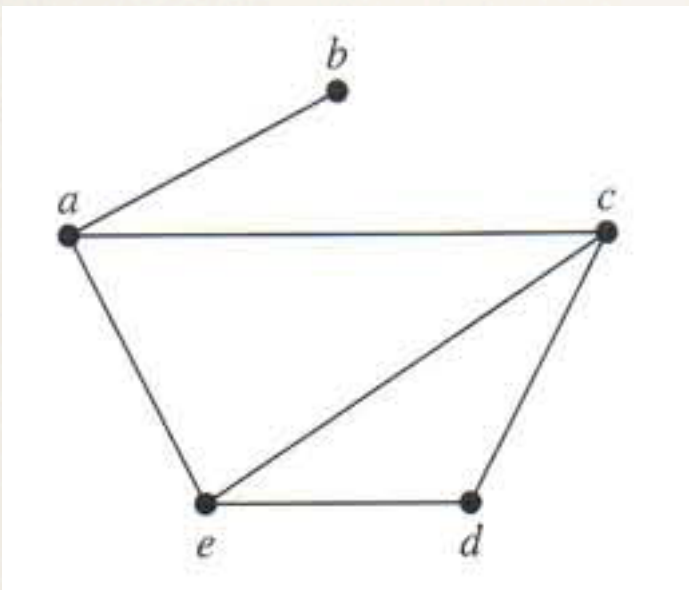
- ❖ For in-degree distribution: Define  $p_k$  to be the fraction of nodes in the graph that have in-degree  $k$ .
- ❖ For out-degree distribution: Define  $q_k$  to be the fraction of nodes in the graph that have out-degree  $k$ .
- ❖ The in- and out-degree distributions of a graph can be visualized by making a histogram of the  $p_k$  and  $q_k$  values, respectively.



# Graph Representation: Adjacency Lists

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- ❖ Adjacency lists specify the nodes that are adjacent to each node of the graph.



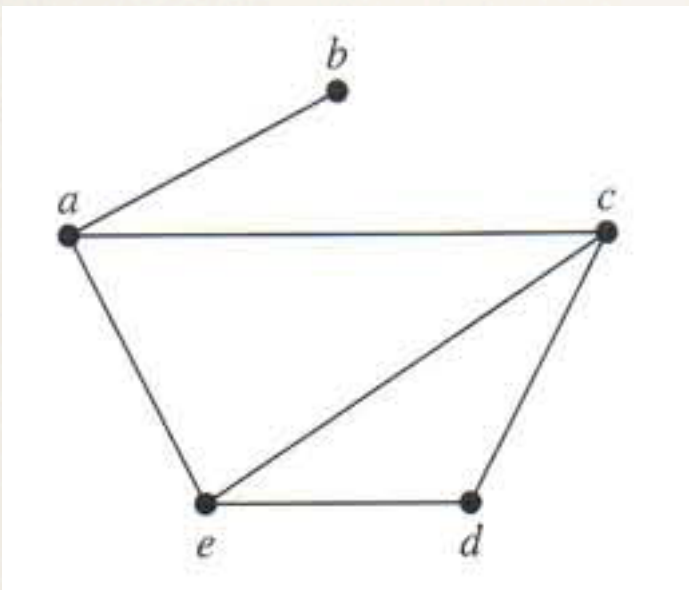
Node	Adjacent nodes
<i>a</i>	<i>b, c, e</i>
<i>b</i>	<i>a</i>
<i>c</i>	<i>a, d, e</i>
<i>d</i>	<i>c, e</i>
<i>e</i>	<i>a, c, d</i>

While “adjacency list” is a historical name, the adjacent nodes of a given node form a **set**; so, you can think of this representation as an “adjacency set.”



# Graph Representation: Adjacency Lists

- Adjacency lists specify the nodes that are adjacent to each node of the graph.



Node	Adjacent nodes
a	b, c, e
b	a
c	a, d, e
d	c, e
e	a, c, d

Notice the **redundancy!**  
(ensures symmetry)

In the case of digraphs,  
there is no redundancy.

While “adjacency list” is a historical name, the adjacent nodes of a given node form a set; so, you can think of this representation as an “adjacency set.”



# Adjacency Lists and Their Dictionary Representation

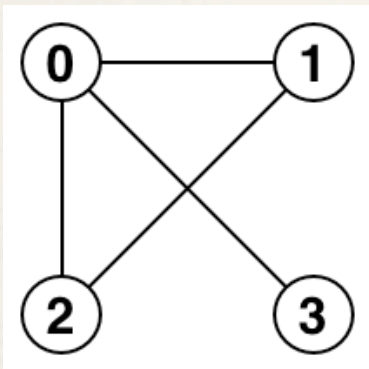
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Node	Adjacent nodes
<i>a</i>	<i>b, c, e</i>
<i>b</i>	<i>a</i>
<i>c</i>	<i>a, d, e</i>
<i>d</i>	<i>c, e</i>
<i>e</i>	<i>a, c, d</i>

{*a*: set([*b,c,e*]),  
*b*: set([*a*]),  
*c*: set([*a,d,e*]),  
*d*: set([*c,e*]),  
*e*: set([*a,c,d*])}

# Graph Representation: Adjacency Matrices

- ✧ Let  $G=(V,E)$  be a graph with  $V=\{0,1,\dots,n-1\}$
- ✧ The adjacency matrix of  $G$ , denoted by  $A_G$ , is the  $n \times n$  0-1 matrix with 1 as its  $(i,j)^{\text{th}}$  entry when  $i$  and  $j$  are adjacent, and 0 as its  $(i,j)^{\text{th}}$  entry when  $i$  and  $j$  are not adjacent.



$G$

$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

$A_G$

Notice the redundancy!

In the case of digraphs,  $A_G$  is not necessarily symmetric.

Draw a graph whose  
adjacency matrix is

$$\begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}$$



# Trade-offs Between Adjacency Lists and Adjacency Matrices

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- ❖ When the graph is sparse, i.e., contains relatively few edges (relative to what?), it is usually preferable to use adjacency lists (Why?)
- ❖ When the graph is dense, i.e., contains relatively many edges (again, relative to what?), it is usually preferable to use adjacency matrices (Why?)



# Graph Connectivity: Paths

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- ❖ Let  $k$  be a nonnegative integer and  $G$  a graph.
  - ❖ A path of length  $k$  from node  $v_0$  to node  $v_k$  in  $G$  is a sequence of  $k$  edges  $e_1, e_2, \dots, e_k$  of  $G$  such that  $e_1 = \{v_0, v_1\}$ ,  $e_2 = \{v_1, v_2\}$ , ...,  $e_k = \{v_{k-1}, v_k\}$ , where  $v_0, \dots, v_k$  are all nodes in  $V$ , and  $e_1, \dots, e_k$  are all edges in  $E$ .
- ❖ We usually denote such a path by its **node sequence**  $(v_0, v_1, \dots, v_k)$ .
- ❖ A path is simple if it does **not** contain the same node more than once.
- ❖ A **cycle** is a simple path that begins and ends at the same node.
- ❖ A path (**not necessarily simple**) that begins and ends at the same node is called a **circuit**.

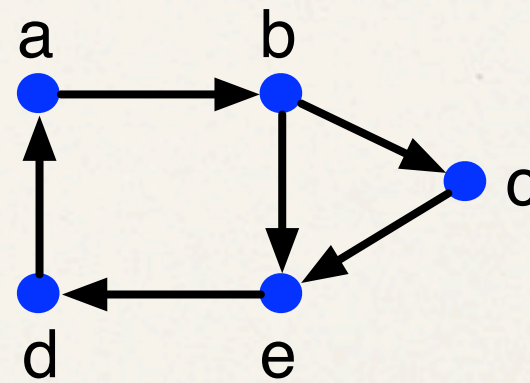
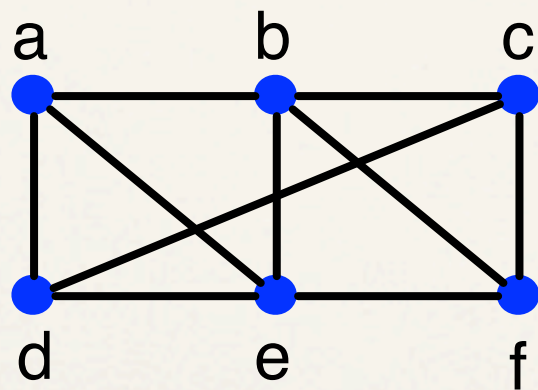


# Graph Connectivity: Paths

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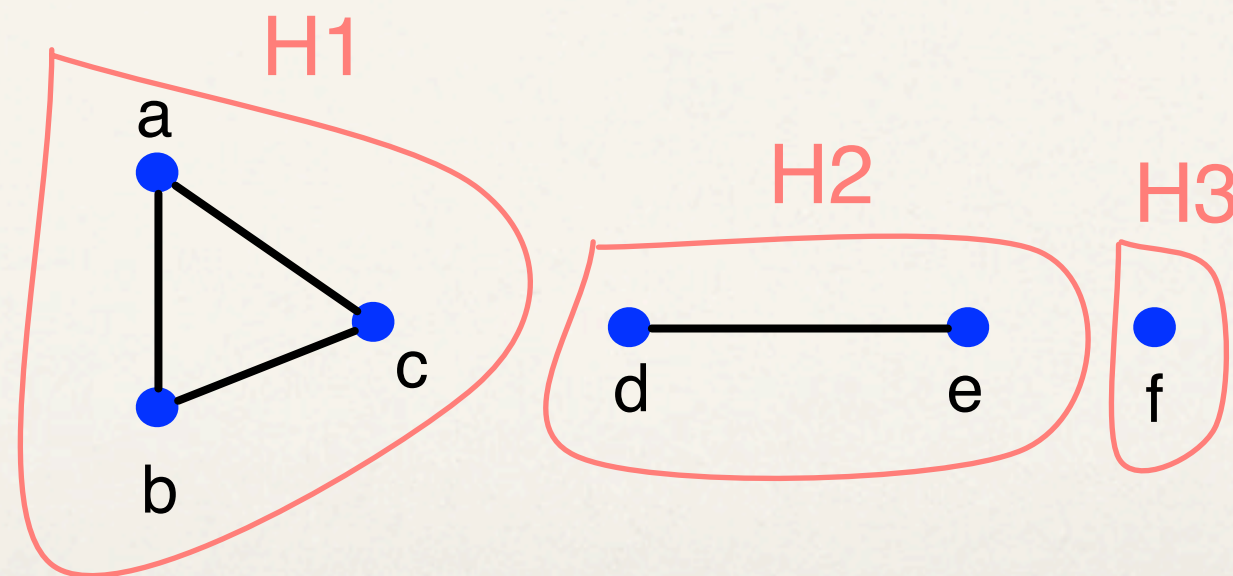
- ❖ If  $G$  is a directed graph, a path must traverse edges in their respective directions.

Show a simple path from node  $e$  to node  $d$  in each of the following two graphs:



# Graph Connectivity

- ❖ An **undirected** graph is called connected if there is a path between **every pair** of distinct nodes of the graph.
- ❖ A connected component (CC) of a graph  $G$  is a connected subgraph of  $G$  that is not a proper subgraph of another connected subgraph of  $G$ .

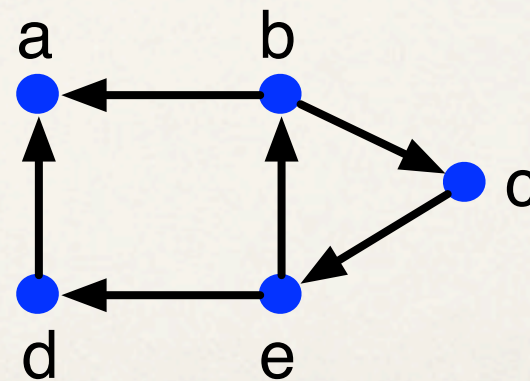
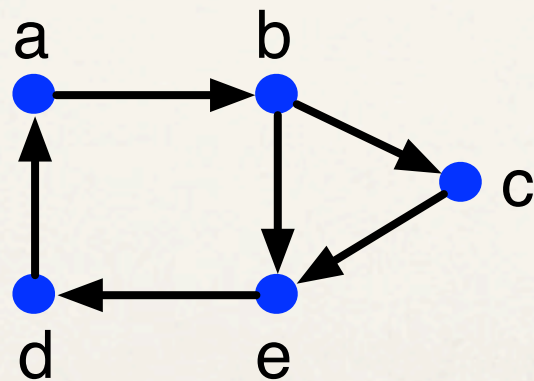


A graph with 3 CCs



# Graph Connectivity

- ❖ A digraph is **strongly connected** if there is a path from  $i$  to  $j$  for every pair of nodes  $i$  and  $j$  of the digraph (note that there must be a path from  $i$  to  $j$  and another from  $j$  to  $i$ ).
- ❖ A digraph is **weakly connected** if there is a path between every two nodes in the underlying undirected graph.
- ❖ The subgraphs of a directed graph  $G$  that are strongly connected but not contained in larger strongly connected subgraphs are called the strongly connected components (SCC) of  $G$ .



Strongly connected? Weakly connected? What are the SCCs?