Minimum Spawning Tree
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DPHPC

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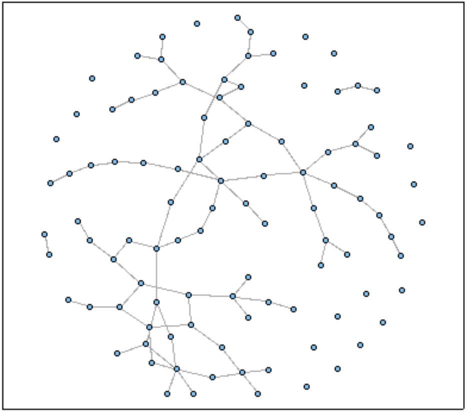
Problem definition

The MST problem

Concepts

Input sets: G(100, 0.02)



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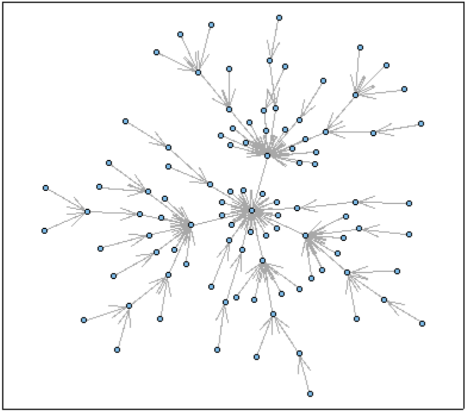
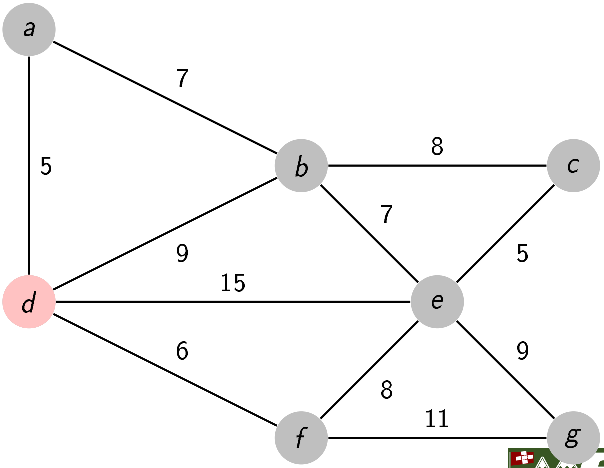


Algorithms

Sequential Prim

- Initialise a tree with a single random vertex.
- Among all the edges connecting the tree to another vertex, find the minimum-weighted one and transfer it to the tree.
- Repeat until all vertices are in the tree.

Sequential Prim



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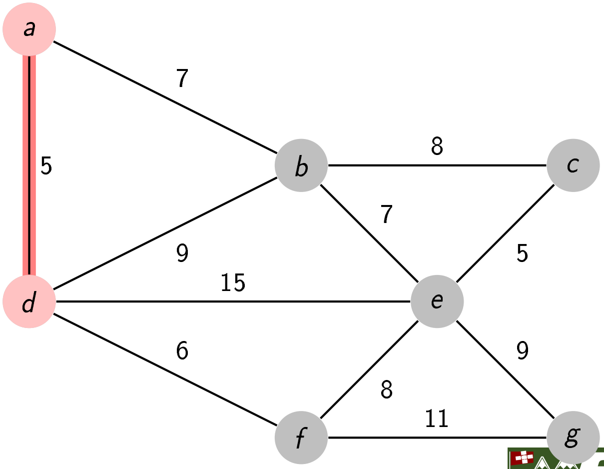

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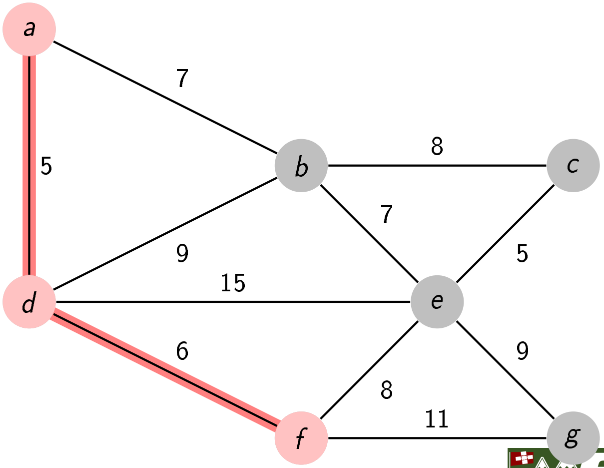

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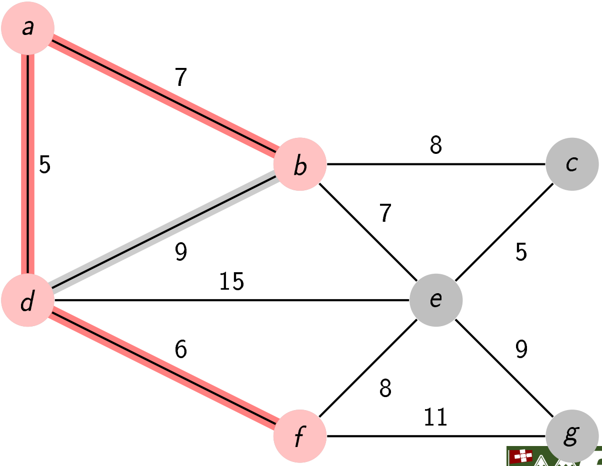

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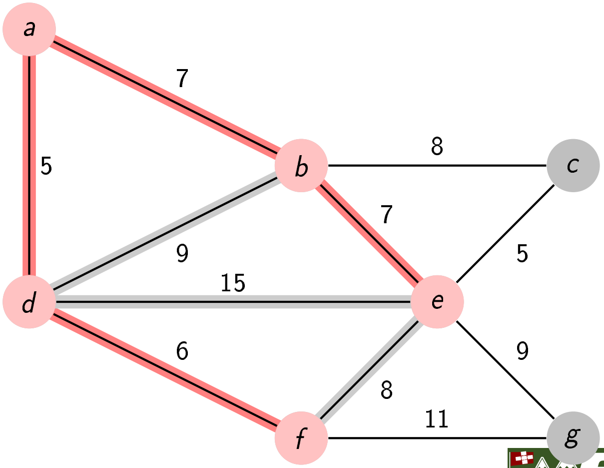

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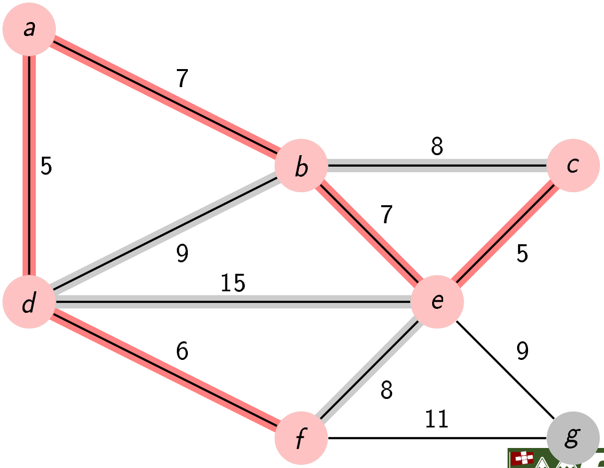

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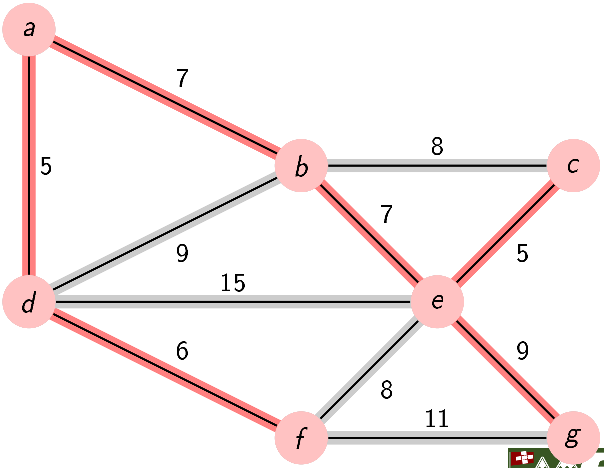

Sequential and parallel Prim

Sequential complexity:

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We can parallelize the search of the edge of minimum weight by dividing the vertices and edges between processors to compute local minima.

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- Sort all edges by growing weight.
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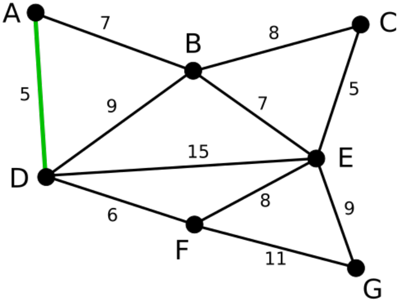

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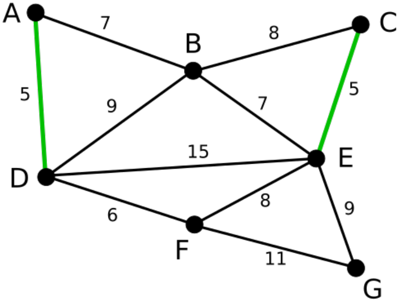

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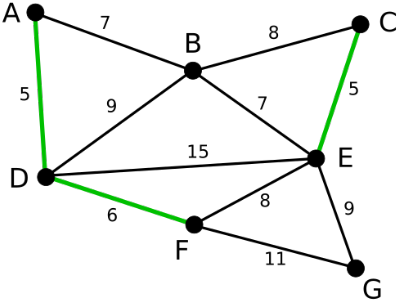

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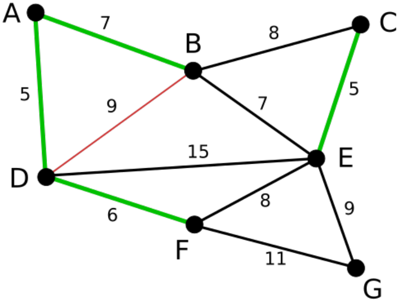

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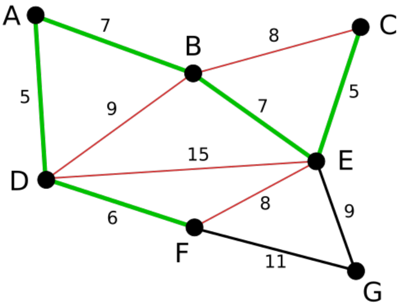

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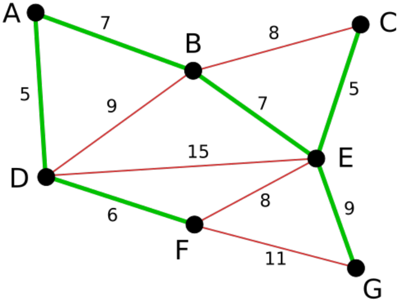

Data structure: Union-find

Represents the connected components of the graph given our MST.
Implementation with an array *parent* and 3 operations:

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  2. Return $parent[x]$.
- Union(x, y): Unite two components.
  1. $parent[find(x)] = find(y)$.

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We can parallelize the sort on $O(\log E)$ processors: $O(E)$ (in practice).

A better parallel approach: Filter-Kruskal

Similar to a quick sort[1]:

1. If $E < \textit{threshold}$, solve using classical Kruskal
2. Choose a pivot (edge)
3. Partition in two sets $E_{\leq}, E_{>}$ (weight)
4. Recursive call to solve problem with $E_{\leq}$
5. Filter out the edges of $E_{>}$ that connect two vertices of the same component
6. Recursive call to solve problem with $E_{>}$

[1] Osipov, Sanders, and Singler 2009.

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Better for parallelization since we can distribute the edges for filtering and partitioning.

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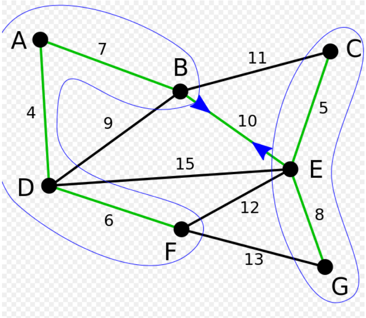
Borůvka (Sollin)

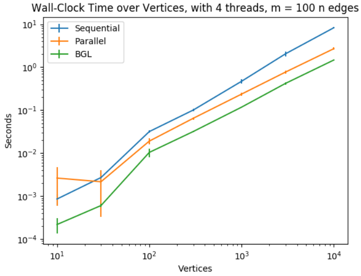


1) Init V as independant sets.
2) Initialize MST as empty.
3) While #sets > 1, do:
   a) Find closest E from this set to another.
   b) Add this E to MST if not already added.
4) Return MST.

- Prim (Parallel & Seq)
- Kruskal (Parallel & Seq), Kruskal filter
- Sollin (Parallel & Seq)

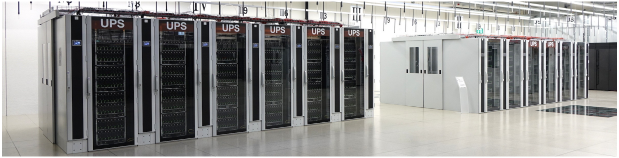
- Randomization

- Correctness



Environment

Architecture



EULER Cluster

Xeon E$x, x \in \{3, 5, 7\}$
x86_64 architecture

Source : https://scicomp.ethz.ch/wiki/Euler

Tools

- CMake
  v3.3+
- C++11
  GCC v4.9.2+
- OpenMPI (shared memory)
  v1.6.5+

Benchmarking

Tools

- **Measures :** LibSciBench library
- **Interpretation :**
  - LibSciBench's R scripts
  - (Custom python scripts)

Ref : https://spcl.inf.ethz.ch/Research/Performance/LibLSB/

Baseline

https://en.wikipedia.org/wiki/Otakar_Bor%C5%AFvka

Borůvka's serial algorithm

$O(E \cdot log(V))$

Osipov, V., P. Sanders, and J. Singler (2009). "The Filter-Kruskal Minimum Spanning Tree Algorithm". In: *Proceedings of the Meeting on Algorithm Engineering & Expermiments*. New York, New York: Society for Industrial and Applied Mathematics, pp. 52–61.