

# The C Programming Language

## Lecture 3

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# Programs vs. Processes

- Program

- Source file.
- Compiled source file on disk.

- Process

- “*Instance of a computer program that is being executed*” (Wikipedia).
- Consists of:
  - PID (Process IDentification number).
  - Context (register values).
  - Page table.
  - List of open files, etc.



# Physical Memory (RAM)

- May not be **sufficient** to fulfill our needs.
- Certain **code** portions may be **rarely used**.
- Physical memory should be **abstracted**.
- Usually, processes **should not interfere** with each other.

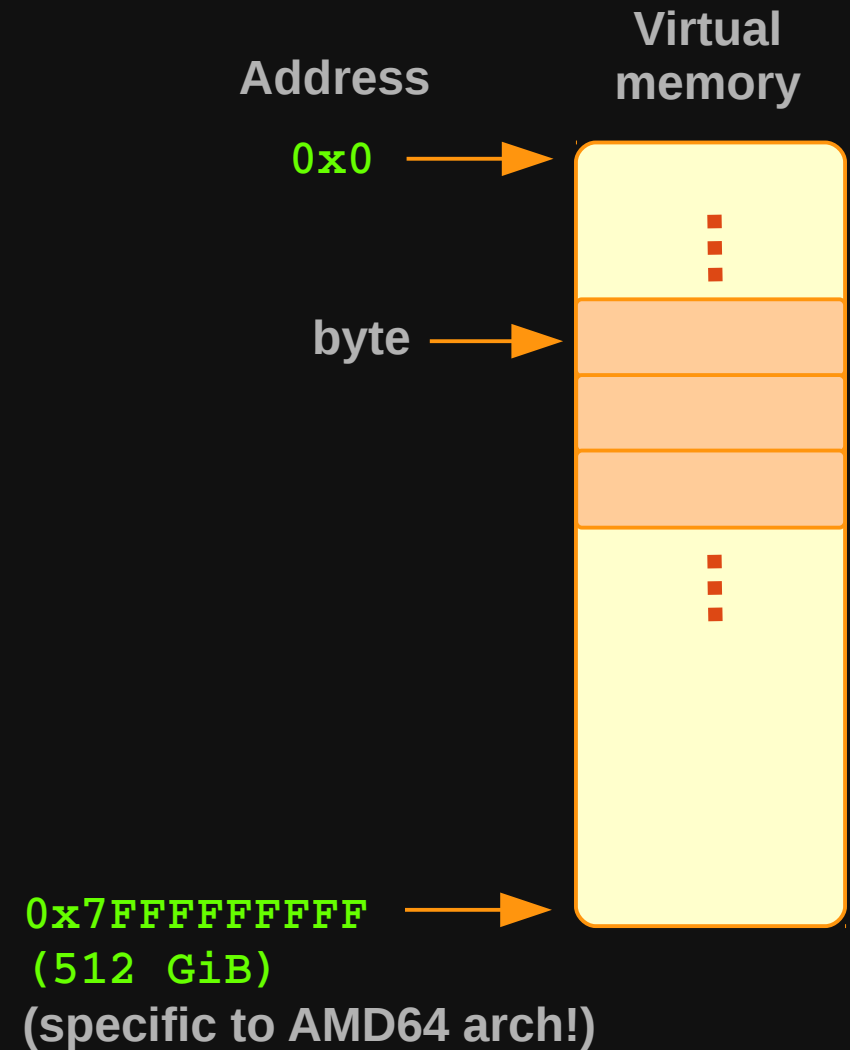
How we could  
handle these problems?

Virtual Memory!



# Virtual Memory

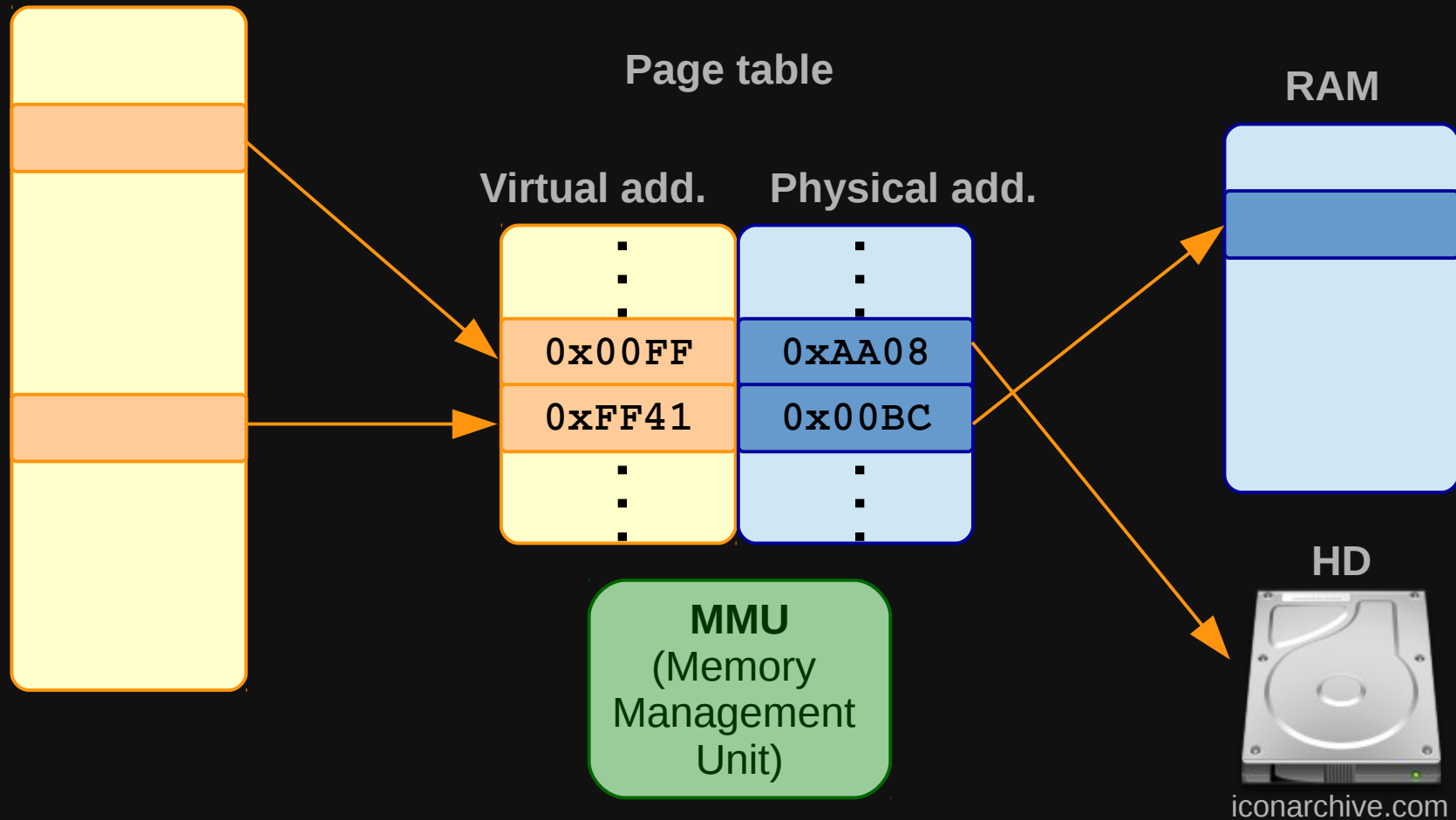
- Main storage appears as a **continuous** address space.
- Maps **virtual** addresses to **physical** addresses.
- The OS **manages** virtual address spaces.
- Address translation is accomplished by the CPU, through dedicated hardware: **Memory Management Unit (MMU)**.
- Even **files** (located on the HD) can be **mapped** into the virtual memory!



# Virtual Memory

Virtual memory

Physical memory

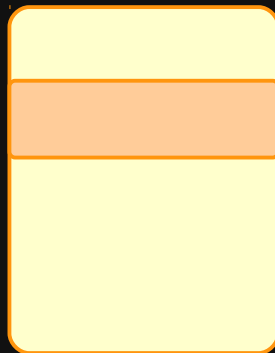


# Virtual Memory

Program A  
(virtual mem.)



Program B  
(virtual mem.)



Page table for prog. A

Virtual add.	Physical add.
⋮	⋮
0x00AA	0xCF11
⋮	⋮

Page table for prog. B

Virtual add.	Physical add.
⋮	⋮
0x00AA	0xEE20
⋮	⋮

Physical memory



How a **program** is  
**organized** within the  
**virtual memory**?



# Physical / Virtual Memory Example

## Checking current physical / virtual memory usage

```
~$ man top
```

```
~$ top
top - 16:14:06 up 2:19, 2 users, load average: 0,57, 0,52, 0,46
Tasks: 232 total, 1 running, 231 sleeping, 0 stopped, 0 zombie
%Cpu(s): 7,8 us, 2,0 sy, 0,0 ni, 90,0 id, 0,0 wa, 0,0 hi, 0,2 si, 0,0 st
KiB Mem: 3949960 total, 2980252 used, 969708 free, 96540 buffers
KiB Swap: 22142972 total, 0 used, 22142972 free. 1168940 cached Mem

  PID USER      PR  NI    VIRT    RES    SHR  S  %CPU  %MEM     TIME+ COMMAND
 1323 root        20   0   250052   84100   46968  S   9,6   2,1    8:10.86 Xorg
 2371 christi+   20   0  1267512  246364  133924  S   9,6   6,2    7:18.72 chrome
 3272 christi+   20   0  1044068  232076   77696  S   7,6   5,9    4:11.55 chrome
 2082 christi+   20   0  1390992  191244   84784  S   6,0   4,8    3:18.66 compiz
```

VIRT : virtual memory size (code, data, shared libraries, etc.).

RES : resident memory size (non-swapped physical memory).

SWAP : swapped size (non-resident portion of the task's address space).

SHR : memory potentially shared with other processes.

top: press 'f' to  
select columns!



# Storage Class and Scope

- Programs are composed of:
  - Executable code.
  - Data.
- Data
  - Pertains to a storage class, which defines its lifespan.
  - Has a scope, which defines its visibility within the program.
  - *“Storage class and scope are assumed from the location of a datum's declaration, and determine its placement within virtual memory.”*





# Storage Class and Scope

- Data declared **outside** any **function** has
  - Global scope.
  - Static duration (it exists for the program lifespan).
- Example

`example_31.c`

```
int x = 7;

void f() {
    x = 19;
}

int main() {

    x = 13;
    f();

    return 0;
}
```

- **x** has **global scope**.
- **x** has **static duration**.



# Storage Class and Scope

- Data declared **inside** a function
  - Local scope.
  - Automatic duration (it exists for the duration of the function call).
- Example

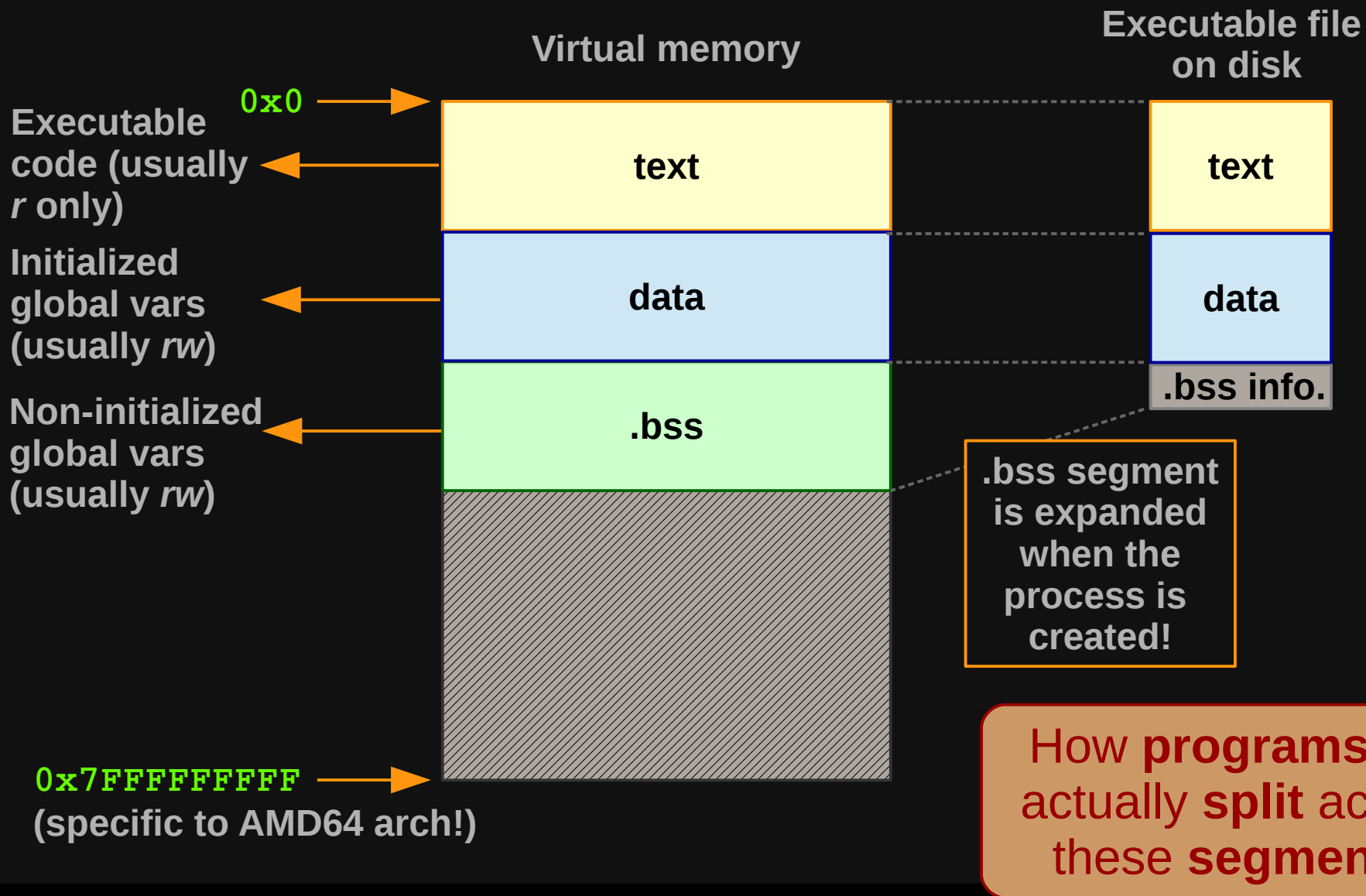
**example\_32.c**

```
void f() {  
    int x = 19;  
    x++;  
}  
  
int main() {  
  
    f();  
    f();  
  
    return 0;  
}
```

- **x** has **local scope**.
- **x** has **automatic duration**.



# Process Memory Layout



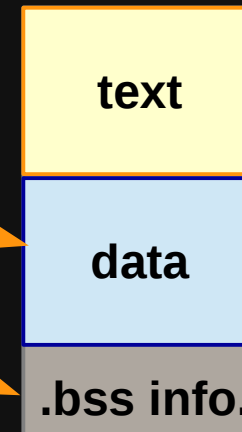
# Program Segmentation Example

- **Executable and Linking Format (ELF)** is the executable file format used by **Linux**.
- **Data and instructions** are aggregated by the compiler in distinct **segments**.

example\_30.c

```
int x = 7;  
int y;  
  
int main() {  
    x = y + 10;  
    return 0;  
}
```

Executable file  
on disk



Executable  
code

Initialized  
global vars

Non-initialized  
global vars



# Process Memory Layout Example

- Example

example\_34.c

```
char x[104857600];  
  
int main (void)  
{  
    return 0;  
}
```

```
~$ gcc example_34.c -o example_34  
~$ ls -l  
-rwxrwxr-x 1 christian christian 8532 Jan 29 20:30 example_34  
-rw-rw-r-- 1 christian christian 55 Jan 29 20:29 example_34.c
```

```
~$ man size
```

```
~$ size example_34  
text      data      bss      dec      hex filename  
1115      552 104857632 104859299 64006a3 example_34
```



# More on Storage Class and Scope!

- Data with **local scope** can present **static duration** if declared with the `static` **qualifier**.
  - It is important when we want to retain the value of a variable between function calls!
- Example

`example_33.c`

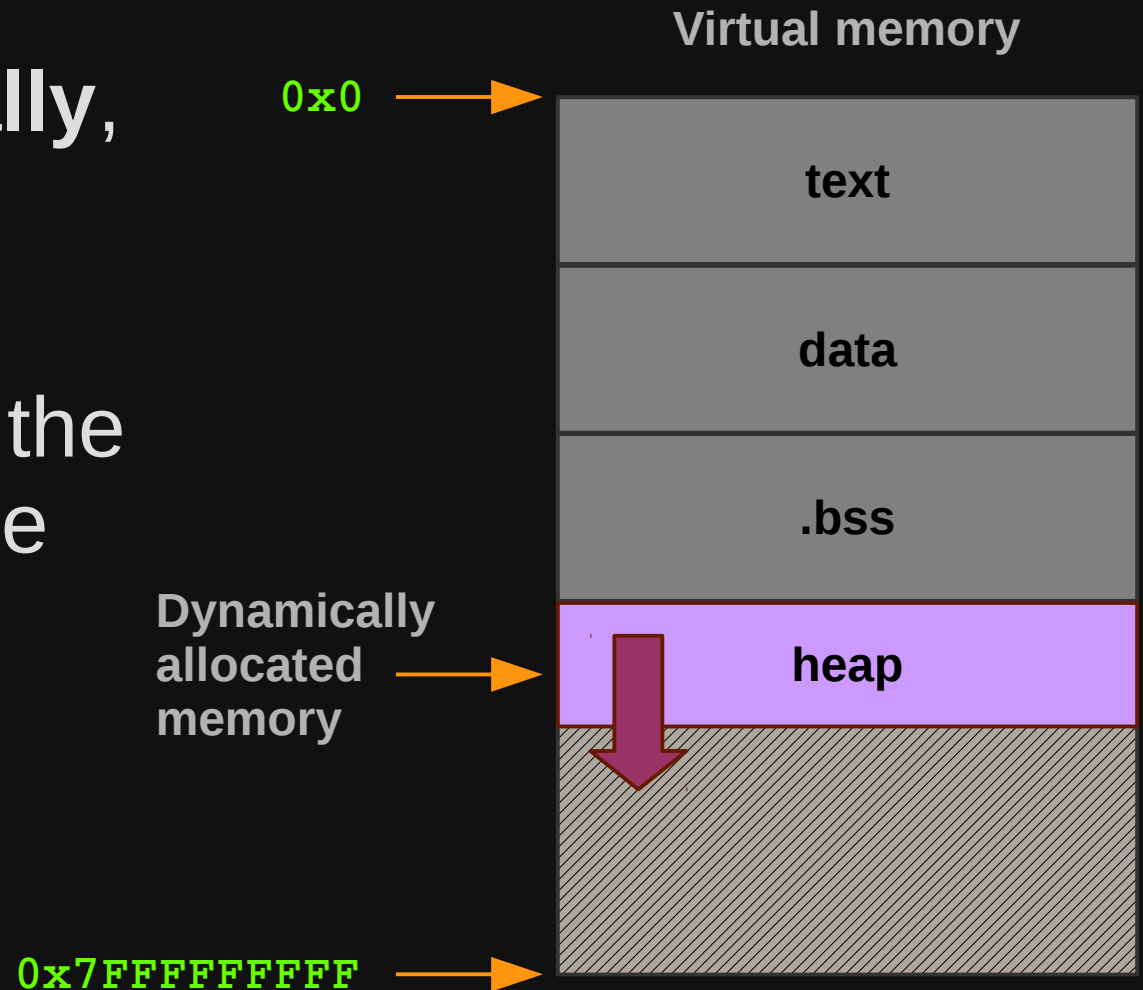
```
void f() {  
    static int x = 19;  
    x++;  
}  
  
int main() {  
  
    f();  
    f();  
  
    return 0;  
}
```

- `x` has **local scope**.
- `x` has **static duration**.



# Dynamic Memory Allocation

- Memory can be allocated **dynamically**, during **runtime**.
- Dynamic memory is allocated right **after** the **data segment**, in the **heap segment**.



# Dynamic Memory Allocation

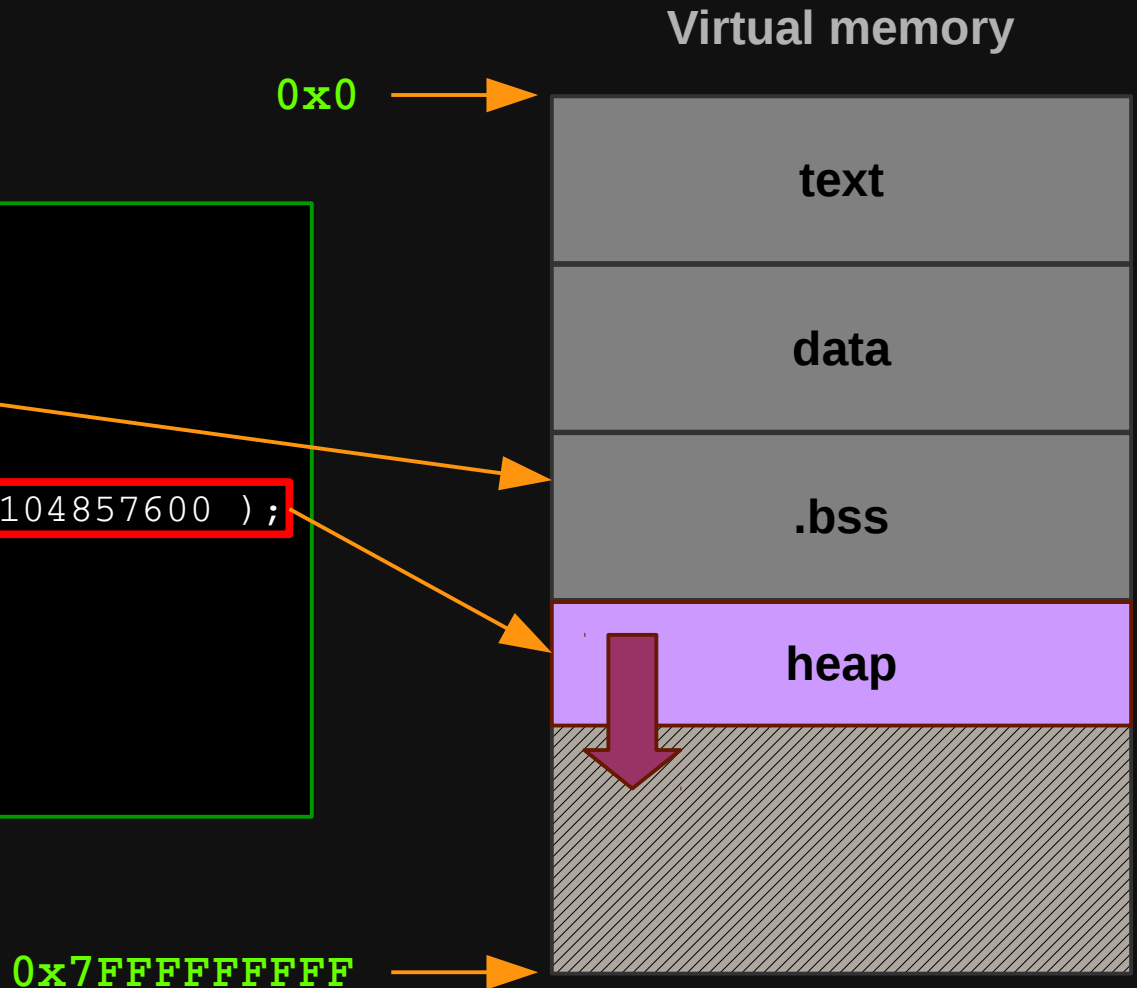
- Example

example\_35.c

```
#include<stdlib.h>

char * pc;

int main (void)
{
    pc = malloc( sizeof( char ) * 104857600 );
    free( pc );
    return 0;
}
```

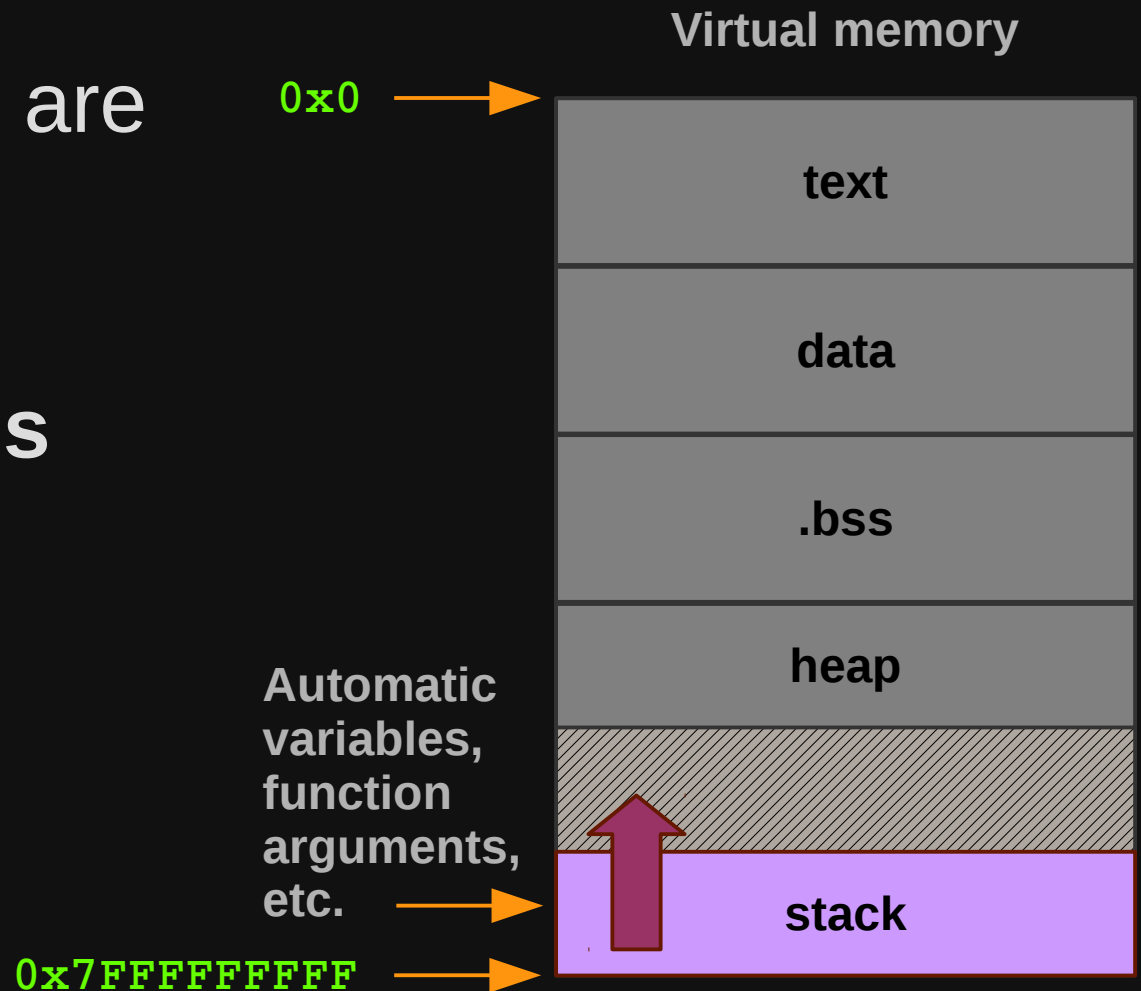




# Automatic Variables

- **Local variables and function arguments are allocated automatically.**
- **Automatic variables are pushed onto the stack.**

More on **parameter passing** later!



# Automatic Variables

- Example

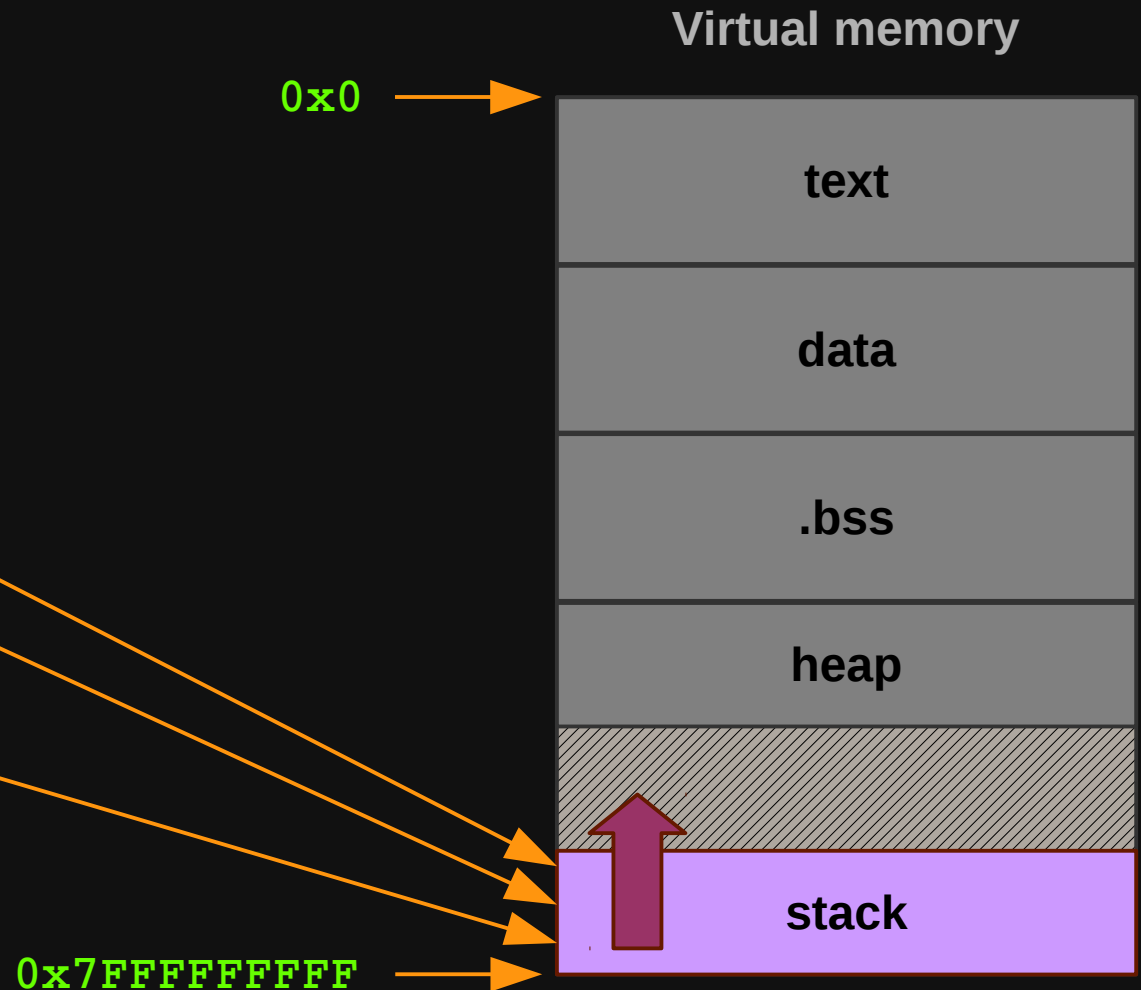
example\_37.c

```
#include <stdio.h>

void f( int i ) {
    int y = i + 20;
    printf( "value: %i\n", y );
}

int main() {
    int x = 10;
    f( x );

    return 0;
}
```



# Stack Facts

- Memory can be allocated **dynamically** on the current **stack frame** with `alloca()`.
  - There is no need to explicitly free `alloca()`'ed memory. It is automatically released at the end of the current function call.
- Stacks usually have a **maximum** prescribed **size**.
  - Thus, take care when using `alloca()`!
- Stacks allow for **recursion**!



More on “*the terrible story of `alloca()` and the `inline` function*” later!



# Parameter Passing and Stack Frames

- C supports **by-value** parameter passing!
  - Everything passed to functions are a copies, including pointers!

- Example

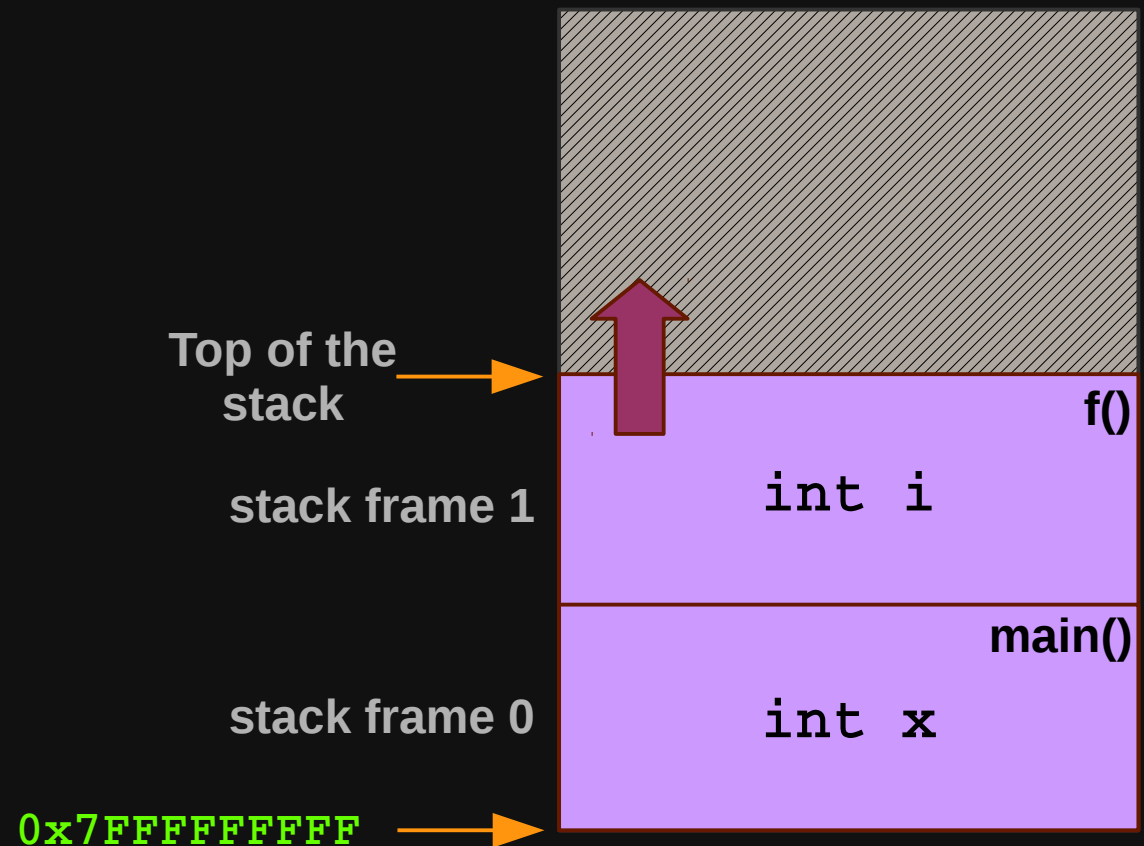
`byvalue.c`

```
#include <stdio.h>

void f( int i ) {
    i= i + 1;
    printf( "value: %i\n", i );
}

int main() {
    int x = 10;
    f( x );
    printf( "value: %i\n", x );

    return 0;
}
```



# Parameter Passing and Stack Frames

- And how about **by-reference** parameter passing?
  - Actually, it is by-value!!
- Example

**byreference.c**

```
#include <stdio.h>

void f( int *i ) {
    *i= *i + 1;
    printf( "value: %i\n", *i );
}

int main() {
    int x = 10;
    f( &x );
    printf( "value: %i\n", x );

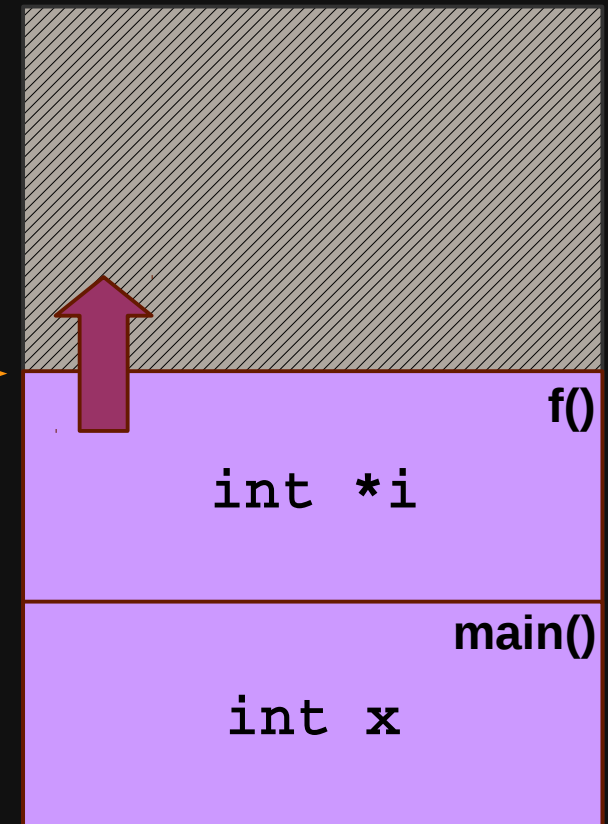
    return 0;
}
```

Top of the  
stack

stack frame 1

stack frame 0

0x7FFFFFFF



# References

- **Understanding Memory.** University of Alberta, 2008.
- **Virtual Memory.** Wikipedia.
  - [https://en.wikipedia.org/wiki/Virtual\\_memory](https://en.wikipedia.org/wiki/Virtual_memory)
- **Memory.** Florent Bruneau. Intersec TechTalk. 2013.
  - <https://techtalk.intersec.com/2013/07/memory-part-1-memory-types>
- **Linux Memory Management System Source Code**
  - <http://git.kernel.org/cgit/linux/kernel/git/torvalds/linux.git/tree/mm>



# Further Reading

- **What Every Programmer Should Know About Memory.** Ulrich Drepper. Red Hat, Inc. 2007.

