



# Solutions to previous round!



11 March 2016



# Counting binary numbers

---

How many binary numbers of length  $N$  have maximum  $K$  consecutive 1s ?

E.g. for  $N = 3$ ,  $K = 1$  the answer is 5

000, 001, 010, 100, 101

# Counting binary numbers solution

---

Let's calculate the following way,  $\text{num}[i][j]$  = how many binary numbers of length  $i$  exist such that they currently have exactly  $j$  1s at the end.

The recurrence is:

- if the last bit is 0,  $\text{num}[i][0] = \text{sum}(\text{num}[i-1][j])$  for all  $j$  from 0 to  $k$
- if the last bit is 1,  $\text{num}[i][j] = \text{num}[i-1][j-1]$  for all  $j$  from 1 to  $k$

# Trap coins

---

You have an array of values  $v$ . Select a subset so that their sum is as big as possible and no 2 consecutive elements are taken.

# Trap coins solution

---

Let's do the following dynamic programming,  $\text{best}[i]$  = how many coins can we get if we only consider the first  $i$  elements.

Well if we want to take the current coin we can have a maximum of  $v[i] + \text{best}[i-2]$  (we can't get  $\text{best}[i-1]$  since it may include  $v[i-1]$  which is forbidden).

If we don't take the current position we simply take  $\text{best}[i-1]$ .

So  $\text{best}[i] = \max(\text{best}[i-1], \text{best}[i-2] + v[i])$

# Tree coins

---

You are given a tree with values in each node, select a subset of nodes so that their sum is as big as possible, and no two selected nodes are neighbours.

# Tree coins solution

---

This is very similar to the previous problem (the previous was a corner case where the tree is a line).

We'll root the tree at 1 to make it easier for us and do the following two dynamic programming:  $take[i]$  -> the maximum value of  $i$ 's subtree if we take  $i$ , and  $notake[i]$  the maximum value of  $i$ 's subtree if we don't take  $i$ . The final solution will be  $\max(take[i], notake[i])$

$take[i]$  will be equal to  $v[i]$  plus the sum of the best notakes of its children.

$notake[i]$  will be equal to the sum of best value of its children (either taken or not taken since both are allowed)

# Tree coins solution

---

So  $\text{take}[i] = v[i] + \sum(\text{notake}[c])$  for all  $c$  children of  $i$

$\text{notake}[i] = \sum(\max(\text{notake}[c], \text{take}[c]))$  for all  $c$  children of  $i$ .



# Chocolate squares

---

You have an  $N \times M$  chocolate, you can split it vertically or horizontally. How many splits will it take to reach only square pieces ?

# Chocolate squares solution

---

Let's do the following dynamic programming  $best[i][j]$ , how many will it take if we have an  $i * j$  matrix.

If  $i$  equals  $j$ , the result is 0.

Otherwise we'll just simulate all the vertical and horizontal splits and pick the one that minimizes the sum of splits needed in the 2 pieces.

$$best[i][j] = 1 + \min(\min(best[i][k] + best[i][j-k]), \min(best[k][j] + best[i-k][j]))$$

# South-East Counting

---

You have a  $N \times M$  matrix with free cells (0) or occupied cells (1). How many ways are there to get from (1,1) to (N,M) ?

# South-East Counting solution

---

Let's count how many ways there are to get from 1,1 to  $i, j$  :  $\text{num}[i][j]$

Well if  $(i, j)$  is occupied, the result is 0.

If it's free then we just add it's north and west neighbour, since those are the only 2 possible paths we can continue:

```
if(free) num[i][j] = num[i-1][j] + num[i][j-1];
```

```
else num[i][j] = 0;
```