

I: Feature Extraction and Similarity Computing:

1. The image content is encoded through the utilization of a feature-extraction process.

2. Let D be an image descriptor defined as a tuple (ϵ, δ) where:

⊛ $\epsilon : O_i \rightarrow \mathbb{R}^d$ is a function which extracts a feature vector \underline{v}_i from the image object O_i .

⊛ $\delta : \mathbb{R}^d \times \mathbb{R}^d \rightarrow \mathbb{R}^+$ is a function which computes the distance between two image objects O_i and O_j .

⊛ Thus, we may write that: $\begin{cases} \epsilon(O_i) = \underline{v}_i \in \mathbb{R}^d \\ \epsilon(O_j) = \underline{v}_j \in \mathbb{R}^d \end{cases}$ and

$$\delta(\epsilon(O_i), \epsilon(O_j)) = \delta(\underline{v}_i, \underline{v}_j) = \|\underline{v}_i - \underline{v}_j\|$$

⊛ The similarity measure $\delta(\cdot, \cdot)$ can be used for ranking tasks, by forming a fundamental component of a similarity measure $p(O_i, O_j)$. $\left[p(O_i, O_j) \propto \frac{1}{\|\underline{v}_i - \underline{v}_j\|} \right]$

IS: Image Retrieval and Rank Model:

1. Let $C = \{O_1, O_2, \dots, O_n\}$ be a collection of images.

2. The target of image retrieval refers to the process of determining image objects from C based on their content.

3. That is, given a query object $O_q \in C$, obtain a ranked list $Z_q = (O_1, O_2, \dots, O_n)$ as a response based on the similarity measure $p(\cdot, \cdot)$.

4.

The ranked list z_q is a permutation of the image set G . #2

Thus, z_q may be defined as a bijection from the set G to the set $[n] = \{1, 2, \dots, n\}$:

$$z_q : G \rightarrow [n]$$

In this setting, $z_q(i)$ indicates that object o_i is the i -th most similar object according to the ranked list induced by object o_q .

Moreover, if $z_q(i) < z_q(j) \rightarrow p(o_q, o_i) \geq p(o_q, o_j)$

5. z_q can be very expensive to compute, especially when n is very high. Therefore, the computed ranked lists may be constrained to consider only a subset $G_L \subset G$ from the complete image set. ($|G_L| = L$)

In this context,

$$z_q : G_L \rightarrow [L] \text{ where } L \ll n.$$

6. Every object $o_i \in G$ can be taken as a query object o_q in order to obtain a set of ranked lists:

$$\mathcal{T} = \{z_1, z_2, \dots, z_n\}$$

7. Based on the ranked list z_q of each query object, we can define its k -nearest neighbours as:

$$N_k(q) = \{z_q(1), z_q(2), \dots, z_q(k)\}$$

III: Manifold Ranking:

#3

⊛ The main objective of your assignment is to exploit the similarity information encoded in the set of ranked lists \mathcal{T} , capturing the manifold structure of the dataset.

⊛ Thus, obtain a more effective set of ranked lists \mathcal{T}_r

$$\mathcal{T}_r = f(\mathcal{T})$$

in an unsupervised learning setting.

IV: Hypergraph Manifold Ranking:

- A. Rank Normalization
- B. Hypergraph Construction
- C. Hyperedge Similarities
- D. Cartesian Product of Hyperedge Elements
- E. Hypergraph-Based Similarity.

- 1 We may obtain a symmetric (normalized) similarity measure $P_n(\cdot, \cdot)$ between a pair of images O_i and O_j , as:

$$P_n(O_i, O_j) = 2L - (Z_i(j) + Z_j(i))$$

- 2 $P_n(\cdot, \cdot)$ is a symmetric similarity measure since:

$$P_n(O_i, O_j) = P_n(O_j, O_i)$$

⊕ We could define a symmetric ranking measure of the form $\hat{Z}_i(j) = \hat{Z}_j(i) = \frac{Z_i(j) + Z_j(i)}{2}$

B: Hypergraph Construction

- 1: • A hypergraph is a generalization of a traditional graph, in which the edges are non-empty subsets of the vertex set and, therefore, can connect any number of vertices.
- 2: • Let $V = \{v_1, v_2, \dots, v_n\}$ be a finite set of vertices and
- E is a family of subsets of V such that

$$\bigcup_{e \in E} e = V$$

which denotes the set of hyperedges.

- In this context, the hypergraph is defined as the triplet:

$$G = (V, E, w)$$

- Each hyperedge $e_i \subset V$ is associated with a positive weight $w(e_i) > 0$, which denotes the confidence of the relationships established by the hyperedge e_i .

3: Each vertex $v_i \in V$ is associated with an object $o_i \in \mathcal{O}$.

4: A hyperedge $e_i \in E$ is said to be incident with a vertex $v_j \in V$ when $v_j \in e_i$.

• In this way, a hypergraph may be represented as an H_0 $|E| \times |V|$ incident matrix as:

$$h_b(e_i, v_j) = \begin{cases} 1, & v_j \in e_i; \\ 0, & \text{otherwise.} \end{cases}$$

5: A hyperedge e_i may be defined as a subset of vertices $e_i = \{v_1, v_2, \dots, v_m\}$. Matrix H_b allows only for a binary assignment of a vertex to a hyperedge.

6: A probabilistic hypergraph model may be defined by considering a $|V| \times |V|$ affinity matrix \underline{W} over V such that $W(i, j) \in [0, 1]$.

7: Given that $v_i \in V \leftrightarrow o_i \in \mathcal{O}$, we may define each hyperedge e_i as:

$$e_i = \{v_j \in V : o_j \in N_k(o_i)\}$$



$$W(i, j) = w(e_i, v_j)$$

8: Given that $w(i, j) = 1 - \log_{x+1}[Z_i(j)]$

a non-binary version of the original incident matrix can be defined as:

$$h(e_i, v_j) = \begin{cases} w(e_i, v_j), & v_j \in e_i; \\ 0, & \text{otherwise.} \end{cases}$$

q: Each hyperedge e_i is formed by the k -th nearest neighbours of o_i including o_i . Thus, each hyperedge contains exactly k vertices.

⊛ Since each e_i is defined relative to each v_i , it is easy to deduce that $|E| = |V|$

10: The weight assigned to hyperedge e_i may be computed as

$$w(e_i) = \sum_{v_j \in e_i} w(e_i, v_j)$$

C: Hyperedges Similarities

1: Compute similarity between any pair of hyperedges e_i and e_j as:

$$\begin{array}{ccc} \underline{S_H} & = & \underline{H} \underline{H}^T \\ \downarrow & & \downarrow \downarrow \\ |E| \times |E| & & (|E| \times |V|) \times (|V| \times |E|) \end{array}$$

2: Compute similarity between any pair of vertices v_i and v_j as:

$$\begin{array}{ccc} \underline{S_V} & = & \underline{H}^T \underline{H} \\ \downarrow & & \downarrow \downarrow \\ |V| \times |V| & & (|V| \times |E|) \times (|E| \times |V|) \end{array}$$

3: Compute combined hyperedge and vertices similarity:

$$\underline{S} = \underline{S_H} \circ \underline{S_V}$$

Hadamard product

↓
Element-wise Multiplication

D: Cartesian Product of Hyperedge Elements

1: Given two hyperedges e_q and e_r , the Cartesian product between them can be defined as:

$$e_q \times e_r = \{ (v_x, v_y) : v_x \in e_q \wedge v_y \in e_r \}$$

2: For each pair of vertices $(v_i, v_j) \in e_q^2$ a pairwise similarity relationship $p: E \times V \times V \rightarrow \mathbb{R}^+$ as

$$p(e_q, v_i, v_j) = w(e_q) * w(e_q, v_i) * w(e_q, v_j)$$

3: A similarity measure can be defined based on the Cartesian product for any pair of vertices irrespective of the hyperedge they belong to:

$$C(v_i, v_j) = \sum_{e_q \in E \wedge (v_i, v_j) \in e_q^2} p(e_q, v_i, v_j)$$

$$\underline{C} \in M_{|V| \times |V|}$$

4: Define a new overall weight matrix $\underline{\hat{W}} \in M_{|V| \times |V|}$ for each pair of vertices, as:

$$\begin{array}{ccc} \hat{W} & = & C \circ S \\ \downarrow & & \downarrow \\ |V| \times |V| & & |V| \times |V| \end{array}$$

S: Based on the new affinity measure which quantifies similarity between any pair of images, a new set of ranked lists \mathcal{J} can be obtained. Therefore, the following procedure may be followed:

$$\begin{array}{ccccccc}
 \mathcal{J}^{(0)} & \longrightarrow & \mathcal{J}^{(1)} & \longrightarrow & \dots & \longrightarrow & \mathcal{J}^{(t)} & \longrightarrow & \dots & \longrightarrow & \mathcal{J}^{(T)} \\
 \underline{\underline{W}}^{(0)} & \longrightarrow & \underline{\underline{W}}^{(1)} & \longrightarrow & & \longrightarrow & \underline{\underline{W}}^{(t)} & \longrightarrow & & \longrightarrow & \underline{\underline{W}}^{(T)}
 \end{array}$$

EUCLIDEAN DISTANCE
BASED

After a given number of iterations we obtain the final set of ranked lists \mathcal{J}_r .

Let $C = \{o_1, o_2, o_3, o_4, o_5, o_6\}$ the set of available images.

Each image o_j is associated with a vertex v_j of the hypergraph

$G = (V, E, w)$ where $V = \{v_1, v_2, v_3, v_4, v_5, v_6\}$.

Let's assume that the ^{initial} ranked lists $z_1, z_2, z_3, z_4, z_5, z_6$ for each image according to the pairwise ^{euclidean} distances of their feature vectors are given below:

$$\begin{aligned} \leftarrow k=3 \rightarrow \\ z_1 &= (o_1, o_2, o_3, o_4, o_5, o_6) \\ z_2 &= (o_2, o_3, o_1, o_5, o_6, o_4) \\ z_3 &= (o_3, o_1, o_4, o_2, o_5, o_6) \\ z_4 &= (o_4, o_3, o_5, o_1, o_2, o_6) \\ z_5 &= (o_5, o_6, o_4, o_3, o_2, o_1) \\ z_6 &= (o_6, o_5, o_4, o_1, o_2, o_3) \end{aligned}$$

Assume that the number of nearest neighbours has been set

to $k=3$, such that:

$$w(o_i, o_j) = 1 - \log_2 \hat{z}_j(o_i)$$

$$N_3(o_1) = (o_1, o_2, o_3)$$

$$N_3(o_2) = (o_2, o_3, o_1)$$

$$N_3(o_3) = (o_3, o_1, o_4)$$

$$N_3(o_4) = (o_4, o_3, o_5)$$

$$N_3(o_5) = (o_5, o_6, o_4)$$

$$N_3(o_6) = (o_6, o_5, o_4)$$

Thus, the graph could have the following structure:

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