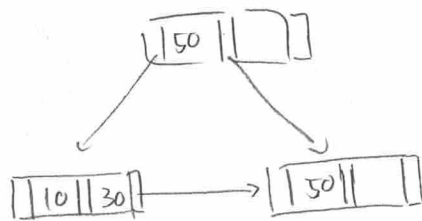
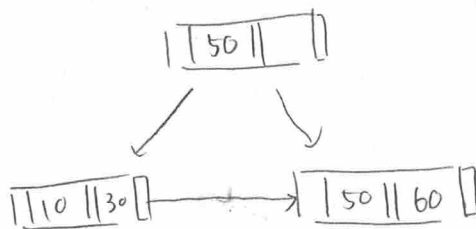


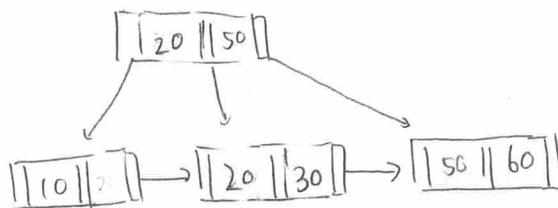
1. (a)



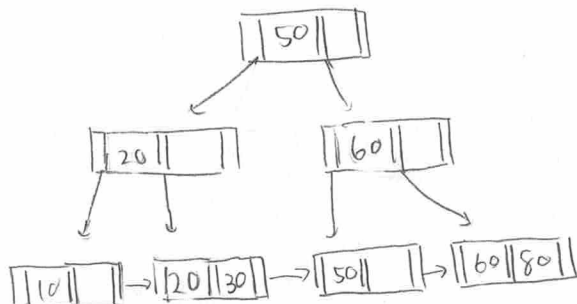
insert 60



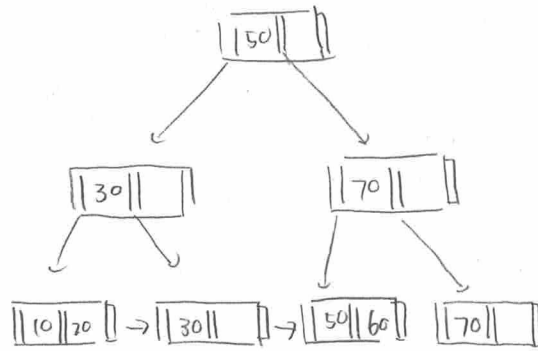
insert 20



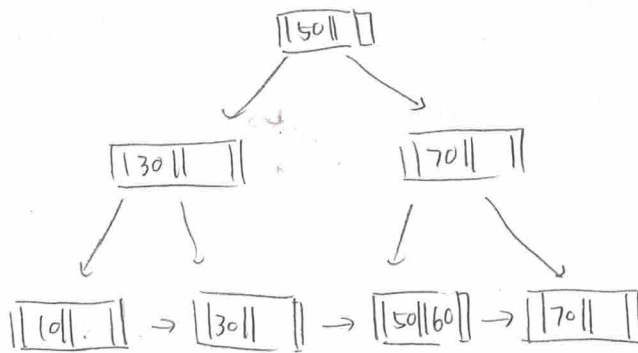
insert 80



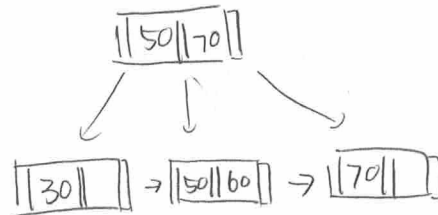
(b)



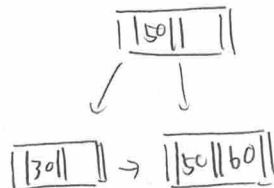
delete 20



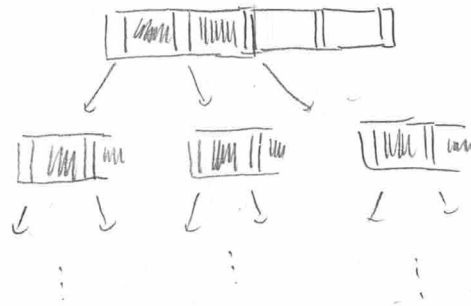
delete 10



delete 70



2.



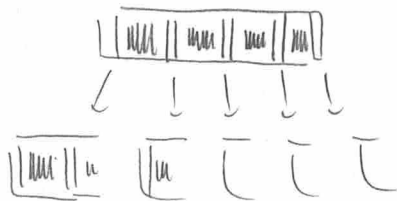
$$2 \cdot 3^0$$

$$2 \cdot 3^1$$

$$2 \cdot 3^4 = 162$$

$$2 \cdot 3^5 = 486$$

max height = 6



$$4 \cdot 5^0$$

$$4 \cdot 5^1$$

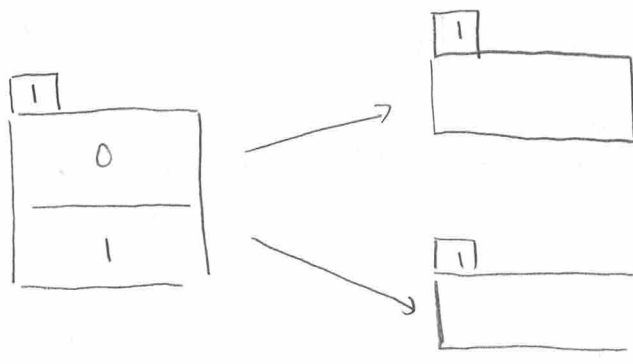
$$4 \cdot 5^2 = 100$$

$$4 \cdot 5^3 = 500$$

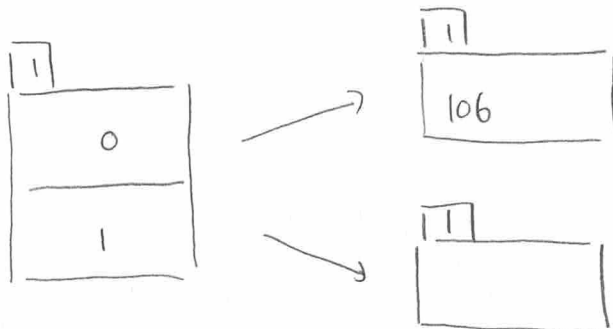
min height = 4

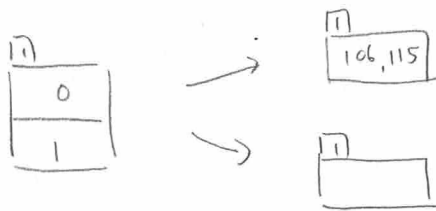
3. ^{hash}

106 → 01101010
 115 → 01110011
 916 → 10010100
 0 → 00000000
 96 → 01100000
 126 → 01111110
 16 → 00010000
 15 → 00001111
 31 → 00011111

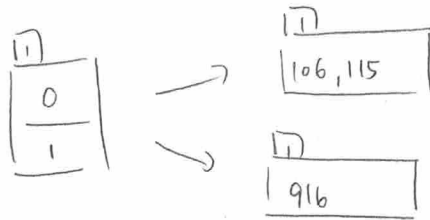


insert 106

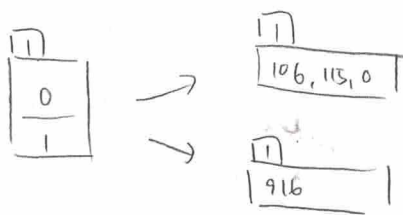




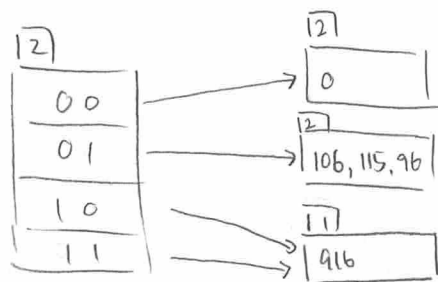
insert 115



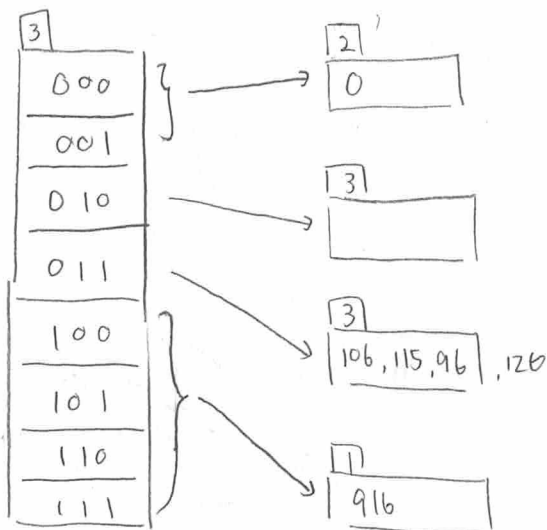
insert 916



insert 0



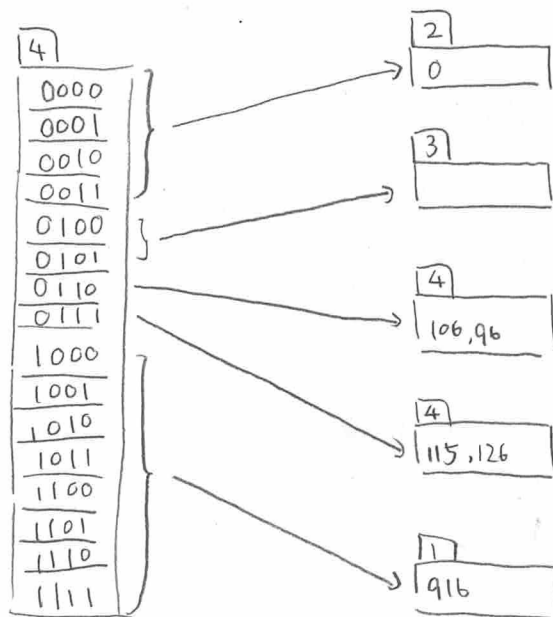
insert 96



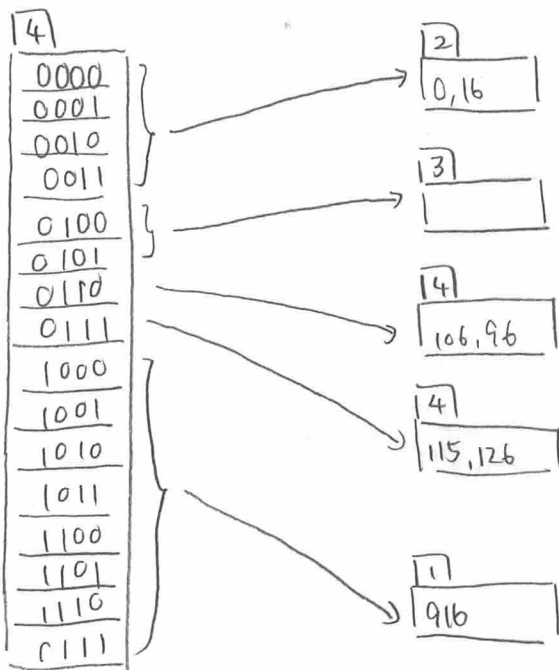
insert 126

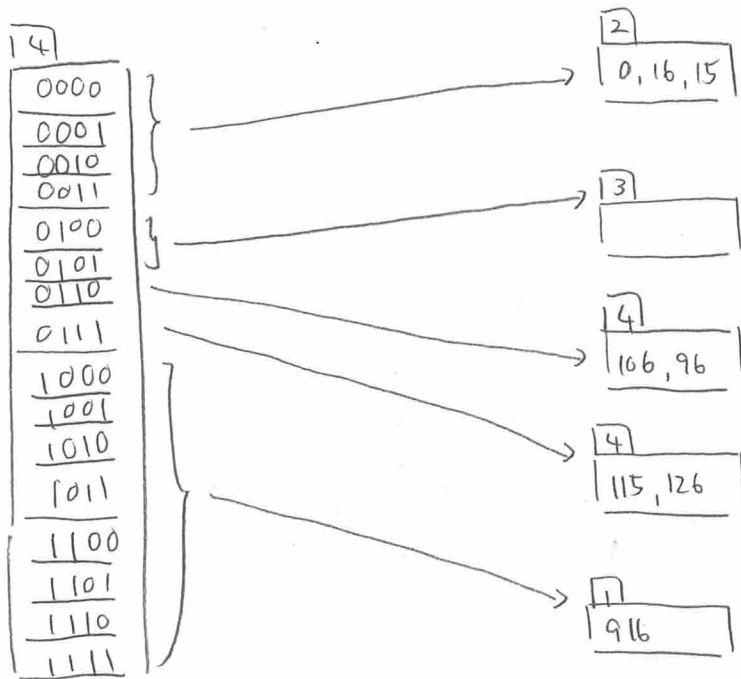
still overflow!

↓

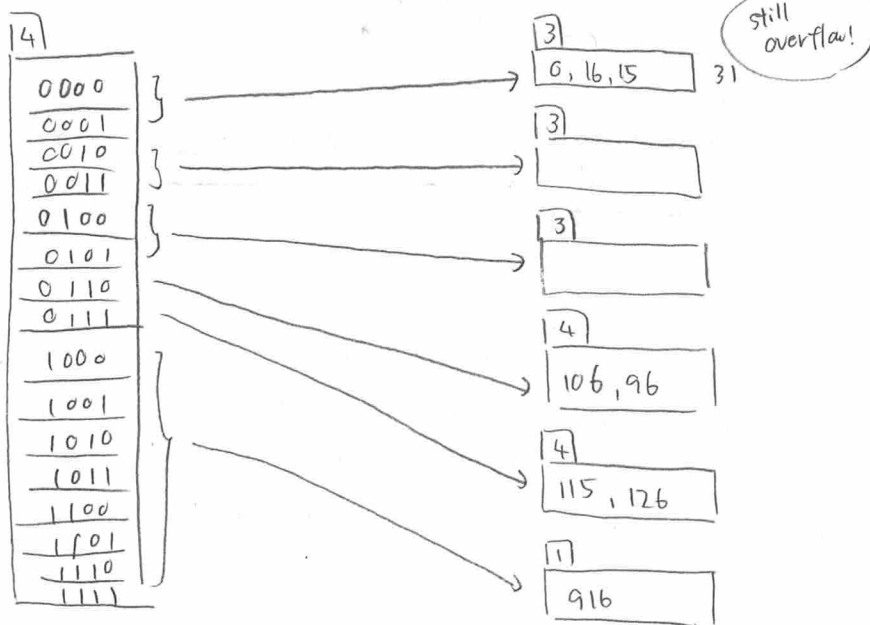


insert 16

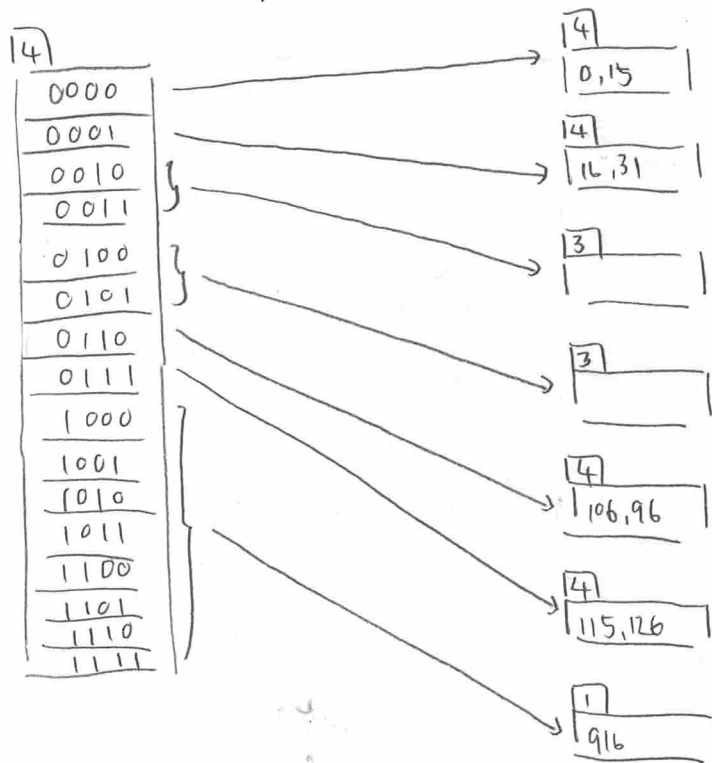




insert 15



insert 31



4. 1.

$R(A, B) \rightarrow 5000 \text{ tuples and } 20 \text{ bytes/tuple}$

$S(B, C) \rightarrow 500 \text{ tuples and } 190 \text{ bytes/tuple}$

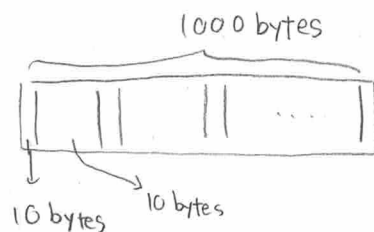
$R: 20 \text{ bytes/tuple} = 1000 \text{ bytes} / 50 \text{ tuples} = 100 \text{ disk blocks} / 5000 \text{ tuples}$

$S: 190 \text{ bytes/tuple} = 950 \text{ bytes} / 5 \text{ tuples} \Rightarrow 5 \text{ tuples / disk block}$
 $\Rightarrow 500 \text{ tuples} / 100 \text{ disk blocks}$

$$\text{So } 100 + 100 = \boxed{200 \text{ [disk blocks]}}$$

4.2.

1 node



n pointers and $n-1$ keys

$$10n + 10(n-1) \leq 1000$$

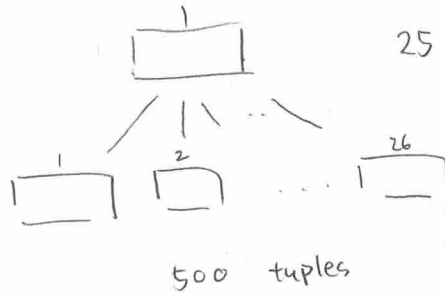
$$20n \leq 1010$$

$$n \leq 50.5$$

$$\boxed{50}$$

4.3.

(a) $n = 50$



$$25 (26)^0 = 25 \text{ keys}$$

$$25 \cdot (26)^1 = 650 \text{ keys} > 500$$

We need depth = 2

Note $25 \cdot 20 = 500$

So we need 20 nodes in level 2

Thus $1 (\text{root}) + 20 = \boxed{21 \text{ nodes}}$

(b) Read one block at a time.

Reading the tuples from S: 100

Writing the constructed B+ tree : 21

Thus total of $\boxed{121 \text{ I/Os}}$.

4.

$$\text{Cost} = \lceil b_r / (M-2) \rceil \times b_s + b_r$$

$$\lceil 100 / 28 \rceil \times 100 + 100 = \boxed{500 \text{ (disk I/Os)}}$$

It's worth while because it would have taken

only $121 + 21 + 100 = 242$ disk I/Os

for optimized algorithm. even considering the

index construction overhead.