

## A Bounding the widths of confidence sets

We present elementary arguments which culminate in a proof of Theorem 3.

**Lemma 4** (Concentration results for  $\sqrt{d_T/n_t(x)}$ ).  
For all finite sets  $\mathcal{X}$  and any  $d_T, \epsilon \geq 0$ :

$$\sum_{t=1}^T \mathbb{1} \left\{ \sqrt{d_T/n_t(x_t)} > h(d_T, \epsilon) \right\} \leq \sum_{t=1}^T \mathbb{1} \left\{ \sqrt{d_T/n_t(x_t)} > \epsilon \right\} + |\mathcal{X}|,$$

Where  $h(d_T, \epsilon) := \sqrt{d_T \epsilon^2 / (d_T + \epsilon^2)}$ .

*Proof.* Let  $(x_{s_1}, \dots, x_{s_K})$  be the largest subsequence of  $x_1^T$  such that  $\sqrt{d_T/n_{s_i}(x_{s_i})} \in (h(d_T, \epsilon), \epsilon] \forall i$ . Now for any  $x \in \mathcal{X}$ , let  $\mathcal{T}_x = \{s_i \mid x_{s_i} = x\}$ . Suppose there exist two distinct elements  $\sigma, \rho \in \mathcal{T}_x$  with  $\sigma < \rho$  so that  $n_\rho(x) \geq n_\sigma(x) + 1$ . We note that for any  $n \in \mathbb{R}_+$ ,  $h(d_T, \sqrt{d_T/n}) = \sqrt{d_T/(n+1)}$  so that:

$$\epsilon \geq \sqrt{d_T/n_\sigma(x)} \implies h(d_T, \epsilon) \geq \sqrt{d_T/(n_\sigma(x) + 1)} \geq \sqrt{d_T/n_\rho(x)}$$

This contradicts our assumption  $\sqrt{d_T/n_\rho(x)} \in (h(d_T, \epsilon), \epsilon]$  and so we must conclude that  $|\mathcal{T}_x| \leq 1$  for all  $x \in \mathcal{X}$ . This means that  $(x_{s_1}, \dots, x_{s_K})$  forms a subsequence of unique elements in  $\mathcal{X}$ , the total length of which must be bounded by  $|\mathcal{X}|$ .  $\square$

We now provide a corollary of this result which allows for episodic delays in updating visit counts  $n_t(x)$ . We imagine that we will only update our counts every  $\tau$  steps.

**Corollary 3** (Concentration results for  $\sqrt{d_T/n_{t_k}(x)}$  in the episodic setting).  
Let us associate times within episodes of length  $\tau$ ,  $t = t_k + i$  for  $i = 1, \dots, \tau$  and  $T = M \times \tau$ . For all finite sets  $\mathcal{X}$  and any  $d_T, \epsilon \geq 0$ :

$$\sum_{k=1}^M \sum_{i=1}^{\tau} \mathbb{1} \left\{ \sqrt{d_T/n_{t_k}(x_{t_k+i})} > h^{(\tau)}(d_T, \epsilon) \right\} \leq \sum_{k=1}^M \sum_{i=1}^{\tau} \mathbb{1} \left\{ \sqrt{d_T/n_{t_k}(x_{t_k+i})} > \epsilon \right\} + 2\tau|\mathcal{X}|,$$

Where  $h^{(\tau)}(d_T, \epsilon)$  is the  $\tau$ -fold composition of  $h(d_T, \cdot)$  acting on  $\epsilon$ .

*Proof.* By an argument of visiting times similar to lemma 4 we can see that the worst case scenario for the episodic case  $\sum_{k=1}^M \sum_{i=1}^{\tau} \mathbb{1} \left\{ \sqrt{d_T/n_{t_k}(x_{t_k+i})} > h^{(\tau)}(d_T, \epsilon) \right\}$  is to visit each  $x$  exactly  $\tau - 1$  times before the start of an episode, and then spend the entirety of the following episode within the state. Here we have upper bounded  $2\tau - 1$  by  $2\tau$  and  $|\mathcal{X}| - 1$  by  $|\mathcal{X}|$  to complete our result.  $\square$

It will be useful to define notion of radius for each confidence set at each  $x \in \mathcal{X}$ ,  $r_{\mathcal{F}_t}(x) := \sup_{f \in \mathcal{F}_t} \|(f - \hat{f}_t)(x)\|$ . By the triangle inequality, we have  $w_{\mathcal{F}_t}(x) \leq 2r_{\mathcal{F}_t}(x)$  for all  $x \in \mathcal{X}$ .

**Lemma 5** (Bounding the number of large radii).

Let us write  $\mathcal{F}_k$  for  $\mathcal{F}_{t_k}$  and associate times within episodes of length  $\tau$ ,  $t = t_k + i$  for  $i = 1, \dots, \tau$  and  $T = M \times \tau$ . For all finite sets  $\mathcal{X}$ , measurable spaces  $(\mathcal{Y}, \Sigma_{\mathcal{Y}})$ , function classes  $\mathcal{F} \subseteq \mathcal{M}_{\mathcal{X}, \mathcal{Y}}$ , non-decreasing sequences  $\{d_t : t \in \mathbb{N}\}$ , any  $T \in \mathbb{N}$  and  $\epsilon > 0$ :

$$\sum_{k=1}^M \sum_{i=1}^{\tau} \mathbb{1} \{r_{\mathcal{F}_k}(x_{t_k+i}) > \epsilon\} < \left( \frac{d_T}{\tau \epsilon^2} + 1 \right) 2\tau|\mathcal{X}|$$

*Proof.* By construction of  $\mathcal{F}_t$  and noting that  $d_t$  is non-decreasing in  $t$ , we can say that  $r_{\mathcal{F}_k}(x_t) \leq \sqrt{d_T/n_{t_k}(x_t)}$  for all  $t = 1, \dots, T$  so that

$$\sum_{k=1}^M \sum_{i=1}^{\tau} \mathbb{1} \{r_{\mathcal{F}_k}(x_{t_k+i}) > \epsilon\} \leq \sum_{k=1}^M \sum_{i=1}^{\tau} \mathbb{1} \left\{ \sqrt{d_T/n_{t_k}(x_{t_k+i})} > \epsilon \right\}.$$

Now let  $g(\epsilon) = \sqrt{d_T \epsilon^2 / (d_T - \tau \epsilon^2)}$  be the  $\epsilon$ -inverse of  $h^{(\tau)}(d_T, \epsilon)$  such that  $g(h^{(\tau)}(d_T, \epsilon)) = \epsilon$ . Applying Corollary 3 to our expression  $n$  times repeatedly we can say:

$$\sum_{k=1}^M \sum_{i=1}^{\tau} \mathbb{1} \left\{ \sqrt{d_T/n_{t_k}(x_{t_k+i})} > \epsilon \right\} \leq \sum_{k=1}^M \sum_{i=1}^{\tau} \mathbb{1} \left\{ \sqrt{d_T/n_{t_k}(x_{t_k+i})} > g^{(n)}(\epsilon) \right\} + 2n\tau|\mathcal{X}|.$$

Where  $g^{(n)}(\epsilon)$  denotes the composition of  $g(\cdot)$   $n$ -times acting on  $\epsilon$ . If we take  $n$  to be the lowest integer such that  $g^{(n)}(\epsilon) > \sqrt{d_T/\tau}$  then,  $\sum_{k=1}^M \sum_{i=1}^\tau \mathbb{1} \left\{ \sqrt{d_T/n_{t_k}(x_{t_k+i})} > g^{(n)}(\epsilon) \right\} \leq 2\tau|\mathcal{X}|$  so that the whole expression is bounded by  $(n+1)2\tau|\mathcal{X}|$ . Note that for all  $N \in \mathbb{R}_+$ ,  $g(\sqrt{d_T/N}) = \sqrt{d_T/(N-\tau)}$ , if we write  $\epsilon = \sqrt{d_T/N_1}$  then  $n \leq N_1/\tau = \frac{d_T}{\tau\epsilon^2}$ , which completes the proof.  $\square$

Using these results we are finally able to complete our proof of Theorem 3. We first note that, via the triangle inequality  $\sum_{k=1}^M \sum_{i=1}^\tau w_{\mathcal{F}_k}(x_{t_k+i}) \leq 2 \sum_{k=1}^M \sum_{i=1}^\tau r_{\mathcal{F}_k}(x_{t_k+i})$ . We streamline our notation by letting  $r_{k,i} = r_{\mathcal{F}_k}(x_{t_k+i})$ . Reordering the sequence  $(r_{1,1}, \dots, r_{M,\tau}) \rightarrow (r_{i_1}, \dots, r_{i_T})$  such  $r_{i_1} \geq \dots \geq r_{i_T}$  we have that:

$$\sum_{k=1}^M \sum_{i=1}^\tau r_{\mathcal{F}_k}(x_{t_k+i}) = \sum_{t=1}^T r_{i_t} \leq 1 + \sum_{i=1}^T r_{i_t} \mathbb{1}\{r_{i_t} \geq T^{-1}\}.$$

We can see that  $r_{i_t} > \epsilon \geq T^{-1} \iff \sum_{i=1}^T \mathbb{1}\{r_{i_t} \geq \epsilon\} \geq t$ . From Lemma 5 this means that  $t \leq \left(\frac{d_T}{\tau\epsilon^2} + 1\right) 2\tau|\mathcal{X}|$ , so that  $\epsilon \leq \sqrt{\frac{2|\mathcal{X}|d_T}{t-2\tau|\mathcal{X}|}}$ . This means that  $r_{i_t} \leq \min\{C_{\mathcal{F}}, \sqrt{\frac{2|\mathcal{X}|d_T}{t-2\tau|\mathcal{X}|}}\}$ . Therefore,

$$\begin{aligned} \sum_{i=1}^T r_{i_t} \mathbb{1}\{r_{i_t} \geq T^{-1}\} &\leq 2\tau C_{\mathcal{F}}|\mathcal{X}| + \sum_{t=2\tau|\mathcal{X}|+1}^T \sqrt{\frac{2d_T|\mathcal{X}|}{t-2\tau|\mathcal{X}|}} \\ &\leq 2\tau C_{\mathcal{F}}|\mathcal{X}| + \int_0^T \sqrt{\frac{2d_T|\mathcal{X}|}{t}} dt \\ &\leq 2\tau C_{\mathcal{F}}|\mathcal{X}| + 2\sqrt{2d_T|\mathcal{X}|T} \end{aligned}$$

Which completes the proof of Theorem 3.

## B Clean bounds for the symmetric problem

We now provide concrete clean upper bounds for Theorems 1 and 2 in the simple symmetric case  $l+1 = m$ ,  $C = \sigma = 1$ ,  $|\mathcal{S}_i| = |\mathcal{X}_i| = K$  and  $|Z_i^R| = |Z_i^P| = \zeta$  for all suitable  $i$  and write  $J = K^\zeta$ . For a non-trivial problem setting we assume that  $K \geq 2$ ,  $m \geq 2$ ,  $\tau \geq 2$ .

From Section 7.3 we have that

$$\begin{aligned} \mathbb{E} [\text{Regret}(T, \pi_\tau^{\text{PS}}, M^*)] &\leq 4 + 2\sqrt{T} + m \left\{ 4(\tau J + 1) + 4\sqrt{8 \log(4mJT^2/\tau)JT} \right\} \\ &\quad + \mathbb{E}[\Psi] \left( 1 + \frac{4}{T-4} \right) m \left\{ 4(\tau J + 1) + 4\sqrt{8K \log(4mJT^2/\tau)JT} \right\} \end{aligned}$$

Through looking at the constant term we know that the bounds are trivially satisfied for all  $T \leq 56$ , from here we can certainly upper bound  $4/(T-4) \leq 1/13$ . From here we can say that:

$$\begin{aligned} \mathbb{E} [\text{Regret}(T, \pi_\tau^{\text{PS}}, M^*)] &\leq \left\{ 4 + 4m \left( 1 + \frac{14}{13} \mathbb{E}[\Psi] \right) (\tau J + 1) \right\} \\ &\quad + \sqrt{T} \left\{ 2 + 4\sqrt{8J \log(4mJT^2/\tau)} + 4\sqrt{8JK \log(4mJT^2/\tau)} \frac{14}{13} \mathbb{E}[\Psi] \right\} \\ &\leq 5(1 + \mathbb{E}[\Psi]) m \tau J + \sqrt{T} \left\{ 12\sqrt{J \log(2mJT)} + 12\mathbb{E}[\Psi] \sqrt{JK \log(2mJT)} \right\} \\ &\leq 5(1 + \mathbb{E}[\Psi]) m \tau J + 12m \left( 1 + \mathbb{E}[\Psi] \sqrt{K} \right) \sqrt{JT \log(2mJT)} \\ &\leq \min(5m\tau^2 J, T) + 12m\tau \sqrt{JKT \log(2mJT)} \\ &\leq 15m\tau \sqrt{JKT \log(2mJT)} \end{aligned}$$

Where in the last steps we have used that  $\Psi \leq \tau$  and  $\min(a, b) \leq \sqrt{ab}$ . We now repeat a similar procedure of upper bounds for UCRL-Factored, immediately replicating  $D$  by  $\tau$  in our analysis to

say that with probability  $\geq 1 - 3\delta$ :

$$\begin{aligned}
\text{Regret}(T, \pi_{\tau}^{\text{UC}}, M^*) &\leq \tau \sqrt{2T \log(2/\delta)} + 2\sqrt{T} + m \left\{ 4(\tau J + 1) + 4\sqrt{8 \log(4mJT/\delta)JT} \right\} \\
&\quad + \tau m \left\{ 4(\tau J + 1) + 4\sqrt{8K \log(4mJT/\delta)JT} \right\} \\
&\leq (1 + \tau)m4(\tau J + 1) + \\
&\quad \sqrt{T} \left\{ \tau \sqrt{2 \log(2/\delta)} + 2 + m4\sqrt{8 \log(4mJT/\delta)J} + \tau m4\sqrt{8 \log(4mJT/\delta)JK} \right\} \\
&\leq 5(1 + \tau)m\tau J + 12m(1 + \tau\sqrt{K})\sqrt{JT \log(4mJT/\delta)} \\
&\leq 15m\tau\sqrt{JKT \log(4mJT/\delta)}
\end{aligned}$$

Where in the last step we used a similar argument