

Functional and Logic Programming

Exercise Set 4

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Parsing

1. Write a function

$$\textit{satisfy} :: (a \rightarrow \textit{Bool}) \rightarrow \textit{Parser } a \rightarrow \textit{Parser } a$$

which creates a new parser that only succeeds if the parsed value satisfies the predicate. Implement *lower* and *char* in terms of *satisfy*.

2. Implement a parser

$$\textit{token} :: \textit{String} \rightarrow \textit{Parser String}$$

that recognizes a complete string, e.g. a keyword.

3. Write a function

$$\textit{mapParser} :: (a \rightarrow b) \rightarrow \textit{Parser } a \rightarrow \textit{Parser } b$$

that given a function and a parser produces a new parser for the same language with a transformed result value. Implement *digit* in terms of *mapParser*.

4. Implement

$$\textit{parens} :: \textit{Parser } a \rightarrow \textit{Parser } a$$

that recognizes an expression that is enclosed in parentheses. For example

```
Prelude> runP (parens number) "(12)"  
Just 12
```

5. Implement

$$\textit{tuple} :: \textit{Parser } a \rightarrow \textit{Parser } b \rightarrow \textit{Parser } (a, b)$$

that recognizes tuple expressions. For example

```
Prelude> runP (tuple number number) "(12,42)"  
Just (12,42)
```

6.

(a) Implement the combinator

$$many :: Parser\ a \rightarrow Parser\ [a]$$

that recognizes zero or more elements of the given parser and the combinator

$$some :: Parser\ a \rightarrow Parser\ [a]$$

that recognizes one or more elements.

(b) Write a parser

$$list :: Parser\ a \rightarrow Parser\ [a]$$

that recognizes list expressions. For example

```
Prelude> runP (list number) "[1,2,3]"
Just [1,2,3]
Prelude> runP (list number) "[]"
Just []
```

Note: You don't need to recognize list expressions of the form $1 : 2 : 3 : []$.

7. Write a parser for the grammar

$$E ::= N \mid E + E$$

where the non-terminal N denotes number literals and produce a value of the following datatype:

```
data Exp = N Int | Exp :+ Exp
deriving Show
```

8. Write a parser for the grammar

$$E ::= L \mid E + E \mid E * E \mid (E)$$

where the non-terminal L denotes number literals and produce a value of the following datatype:

```
data Expr = L Int | Expr :+: Expr | Expr :*: Expr
deriving Show
```

Make sure that the precedence of multiplication over addition is handled correctly, e.g. `"1+2*3"` yields the same parse result as `"1+(2*3)"`.

Lazy evaluation

9. In the course you have seen the `unfoldr` function that builds (possibly infinite) lists.

(a) Implement a similar function

$$unfoldStream :: (s \rightarrow (a, s)) \rightarrow s \rightarrow [a]$$

that always produces infinite lists.

(b) Implement the following functions in terms of *unfoldStream*:

- *repeat* $a = a : a : a : \dots$
- *nats* $= 0 : 1 : 2 : \dots$
- *iterate* $f\ a = a : f\ a : f\ (f\ a) : \dots$
- *primes* $= 2 : 3 : 5 : \dots$
Use the sieve of Eratosthenes.

10.

(a) Implement the function

$$\text{partialSums} :: \text{Num } a \Rightarrow [a] \rightarrow [a]$$

that creates for a sequence of numbers the sequence of non-empty partial sums. For example

```
Prelude> partialSums [1..]  
[1, 3, 6, 10, 15, 21, 28, 36, 45, 55, 66, 78, 91, 105, 120, ...]
```

Use the same style that you have seen for the list of Fibonacci numbers:

```
fibs :: [Int]  
fibs = 0 : 1 : zipWith (+) fibs (tail fibs)
```

(b) Dropping every second element from the positive natural numbers gives us the list of uneven numbers

1, 3, 5, 7, 9, 11, 13, 15, ...

It is well known, that the partial sums of the uneven numbers are the squared numbers

1, 4, 9, 16, 25, 36, 49, 64, ...

In a similar fashion we can obtain the cubes. Start with the positive natural numbers and drop every **third** element:

1, 2, 4, 5, 7, 8, 10, 11, 13, 14, 16, ...

then take partial sums:

1, 3, 7, 12, 19, 27, 37, 48, 61, 75, 91, ...

drop every **second** element:

1, 7, 19, 37, 61, 91, ...

and form the partial sums again:

1, 8, 27, 64, 125, 216, ...

The above procedure of alternating between dropping elements and forming the partial sums can be used to obtain the streams of powers of n for any $n \geq 1$. This is known as Moessner's theorem. Implement this procedure in a function

$$\text{moessner} :: \text{Int} \rightarrow [\text{Int}]$$

that calculates the stream of powers of n .