Haskell les 4:

Lazyness

Een goede informaticus moet lui zijn

door Pietervdvn

Syntax: Record syntax

```
data Person =
Person Int String Int Bool Bool
```

```
data Person =
Person ID Name Year Gender Bool
```

```
type ID = Int
type Name = String
```

```
data Person =
Person ID Name Year Gender Bool
```

```
type ID = Int
type Name = String
```

```
data Person
  Person {id::ID,
          name::Name,
          birth::Year,
           sex:: Gender,
          member::Bool }
```

Record Syntax: getters

```
:t id
 id :: Person -> Id
:t name
 name ::Person -> Name
:t member
 member :: Person -> Bool
```

Record Syntax: getters

```
person = Person 42 "Pieter" 1993 True
member person
  True
birth person
  1993
```

Record Syntax: setters

```
person = Person 42 "Pieter" 1993 True
person \{id = 43\}
  Person 43 "Pieter" 1993 True
person
  Person 42 "Pieter" 1993 True
```

Theorie

Luiheid

Dingen worden maar uitgerekend wanneer het echt nodig is

```
42 `div` 0
  *** Exception: divide by zero
const 5 (42 `div` 0)
  5
```

```
const 5 (42 `div` 0)
```

```
((x -> (y -> x)) 5) (42 idiv 0)
```

```
(\y -> 5) (42 \div 0)
```

Meeste talen:

```
f = \x -> x * x
f (1 + 1)
```

```
f = \x -> x * x
f(1 + 1)
(\x -> x * x) (1 + 1)
```

```
f = \x -> x * x

f (1 + 1)
(\x -> x * x) (1 + 1)
(1 + 1) * (1 + 1)
```

```
f = \langle x \rightarrow x * x \rangle
f(1 + 1)
(\x -> x * x) (1 + 1)
(1 + 1) * (1 + 1)
2 * (1 + 1)
```

```
f = \langle x \rightarrow x * x \rangle
f(1 + 1)
(\x -> x * x) (1 + 1)
(1 + 1) * (1 + 1)
2 * (1 + 1)
```

```
f = \langle x \rightarrow x * x \rangle
f(1 + 1)
(\x -> x * x) (1 + 1)
(1 + 1) * (1 + 1)
2 * (1 + 1)
```

Dingen worden maar uitgerekend wanneer het echt nodig is:

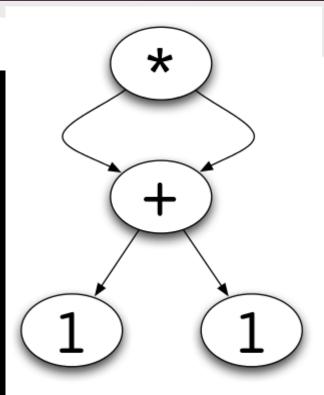
Dingen worden maar uitgerekend wanneer het echt nodig is:

Call by Need = Call by Name + Sharing

```
f = \x -> x * x
f (1 + 1)
(\x -> x * x) (1 + 1)
```

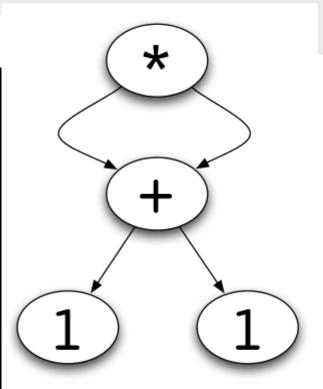
```
f = \x -> x * x
```

```
f (1 + 1)
(\x -> x * x) (1 + 1)
```



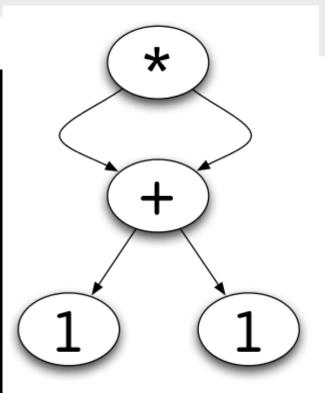
```
f = \x -> x * x
```

```
f (1 + 1)
(\x -> x * x) (1 + 1)
(1 + 1) * (1 + 1)
```

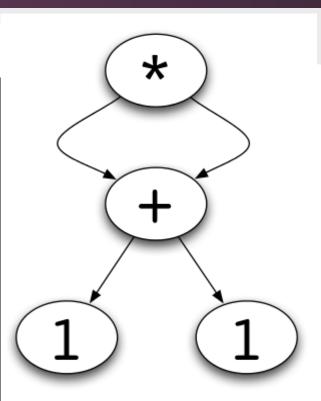


```
f = \x -> x * x
```

```
f (1 + 1)
(\x -> x * x) (1 + 1)
(1 + 1) * (1 + 1)
2 * 2
```



```
f = \x -> x * x
f(1 + 1)
 (x -> x * x) (1 + 1)
       * (1 + 1)
```



Lazyness: if

```
ifthenelse :: Bool -> a -> a
ifthenelse True a0 _ = a0
ifthenelse False _ a1 = a1
```

```
ifthenelse True 42 (1 * 2 + fac 43 * fib
```

```
ifthenelse True 42 (1 * 2 + fac 43 * fib 42
```

```
ifthenelse True 42 (5 `div` 0)
42
```

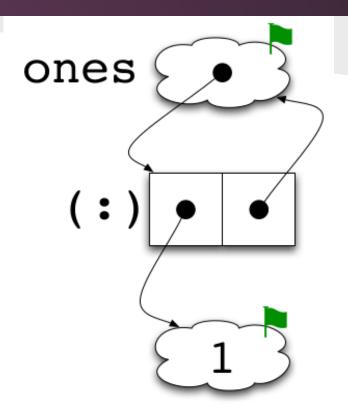
```
ifthenelse False 42 (5 `div` 0)
***Exception: divide by zero
```

Lazyness: oneindige lijst

ones = 1:ones

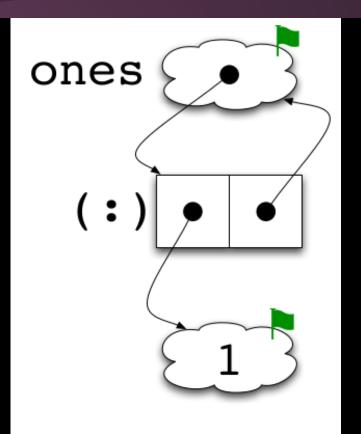
```
ones
```

ones = 1:ones



ones

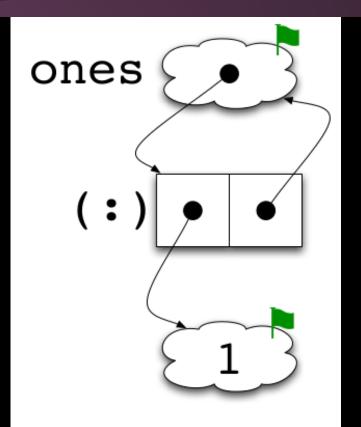
1:ones



ones

1:ones

1:1:ones

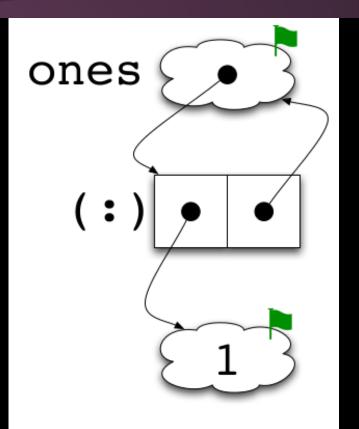


ones

1:ones

1:1:ones

1:1:1:ones



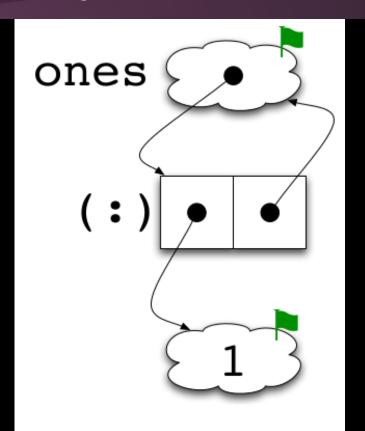
ones

1:ones

1:1:ones

1:1:1:ones

1:1:1:1:ones



ones

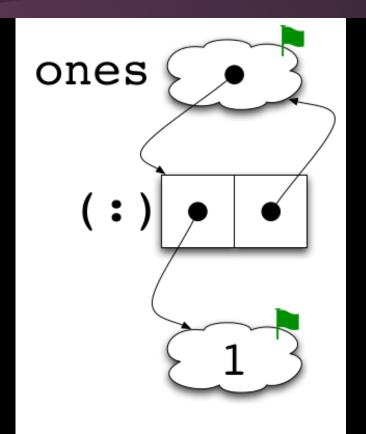
1:ones

1:1:ones

1:1:1:ones

1:1:1:1:ones

... ad infinitum



```
take 10 ones
[1,1,1,1,1,1,1,1]
```

```
nats = 0:map(+1) nats
```

```
nats
0:map (+1) nats
```

```
nats
0:map (+1) (0:map (+1) nats)
```

```
nats
0:1:map (+1) (drop 1 nats)
```

```
nats
0:1:map (+1) (drop 1 (0:1: ... ):
    map (+1) drop 2 nats
```

```
nats
0:1:map (+1) (1: ...):
    map (+1) drop 2 nats
```

```
nats
```

```
0:1:2:map (+1) drop 2 nats
```

```
nats
```

```
0:1:2:3:4:5:...
```

```
nats
0:1:2:3:4:5:...
-- syntactische suiker:
[0..]
```

Compacte notatie

```
sum $ take 10 [1..]
55
take 10 $ map fac [1..]
[1,2,6,24,120,720,5040,...,3628800]
```

```
repeat 1
 zip [1..5] $ repeat 1
 [(1,1),(2,1),(3,1),(4,1),(5,1)]
zipWith (+1) [1..5] $ repeat 1
 [2,3,4,5,6]
```

```
fibs = 1 : 1 : zipWith (+)fibs (tail fibs)
```

One pass

```
data Tree a = Leaf a
| Node a (Tree a) (Tree a)
```

Is het mogelijk om een boom te overlopen in 1 pass, en overal het minimum in te stellen?

```
-- Vervangt elke waarde door het minimum a
repMin :: a -> Tree a -> (Tree a, a)
repMin a (Leaf b) = (Leaf a, b)
repMin a (Node b left right)
     = let (left', minLeft) = repMin a left
           (right', minRight) = repMin a right
     (Node a left' right',
          min3 b minLeft minRight)
```

```
(newTree, min) = repMin min tree
```

Lazyness: issues

Strictness

Lazyness is traag voor sommige algoritmen (vooral numerieke)

Gebruik stricte code

Strictness

```
e.g
Node (count + 1) ...
Node (0 +1 +1 +1 +1 ...)
```

Strictness

```
let count = oldCount + 1 in
    seq count $ foo count
-- evalueert 'count' voordat deze als parameter wordt
doorgegven
```

Higher order programming

State

Toestand

Voor sommige algoritmes is toestand handig (bv. Dijkstra, queue)

Toestand

Toestand

```
data State s a = State (s -> (s,a))
runstate:: State s a -> s -> (s,a)
runstate (State f) = f
-- door eta-reductie
runstate (State f ) s = f s
```

```
put :: s -> State s ()
put s = State (\_ -> (s,()) )

get :: State s s
get = State (\s -> (s,s))
```

runstate (put 5) 0

```
put s = State (\_ -> (s,()))
```

```
runstate (State $ \_ -> (5,()) 0
```

```
put s = State (\_ -> (s,()))
```

```
(\_ -> (5,()) ) 0
```

```
put s = State (\_ -> (s,()))
```

```
(5,())
```

```
put s = State (\_ -> (s,()) )
```

runstate get 42

```
get = State (\s -> (s,s))
```

```
runstate (State $ \s -> (s,s)) 42
```

```
get = State (\s -> (s,s))
```

```
(\s -> (s,s)) 42
get = State (\s -> (s,s))
```

```
(42,42)
```

```
get = State (\s -> (s,s))
```

runstate (put 5 `andThen` const get) 42

```
runstate (put 5 `andThen` const get) 42
andThen :: State s a -> (a -> State s b) -> State s b
andThen action a2actionb
= State $ \beginState ->
  let (midState, a) = runstate (put 5) beginState
      (endState, b) = runstate ((\ -> get) a) midState in
      (endState, b)
```

```
runstate (put 5 `andThen` const get) 42
andThen :: State s a -> (a -> State s b) -> State s b
andThen action a2actionb
 = State $ \beginState ->
  let (5, ()) = runstate (put 5) beginState
      (5, 5) = runstate ((\setminus -> get) ()) 5 in
      (5, 5)
```

```
runstate (put 5 `andThen` const get) 42
andThen :: State s a -> (a -> State s b) -> State s b
andThen action a2actionb
= State $ \beginState -> (5,5)
```

```
runstate (State $ \ -> (5,5,)) 42
```

(5,5)

```
foo :: State ([Node], Map Node Node,
Graph) ()
foo = get (\ (toVisit:queue, paths, graph) -
        let cn = getClosest toVisit graph in
        put (queue, insert toVisit cn, graph)
```

```
data Context = Ctx {queue :: [Node],
                      paths:: Map Node Node,
                      graph:: Graph}
foo :: State Context ()
foo = get (\ ctx ->
        let node = head $ queue ctx in
        let cn = getClosest node (graph ctx) in
        put $ ctx {queue = tail queue,
                   paths = insert node cn}
```

Tooling

quickcheck

Unit-test voor haskell

Unit-test voor haskell

Unit-test = quickcheck voor Java

Unit-test voor haskell

Stel een aantal eigenschappen op

$$\sl > S == S$$

Unit-test voor haskell

Stel een aantal eigenschappen op

$$\sl > S == S$$

Test deze

QuickCheck

```
import Test.QuickCheck
quickCheck ((\str -> str == str) :: [Char] -> Bool)
+++ OK, passed 100 tests.
```

Copyright

Met dank aan Prof. Tom Schrijvers, van wie ik de afbeeldingen (en oefeningen) overgenomen heb

Oefeningen

github.com/pietervdvn/Haskell/Les4