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Handling failure

P4 out (last Mon)

```
def f():  
    try:  
        return 10  
    finally:  
        print("got here!")  
    return 20
```

```
def f():  
    try:  
        return 10  
    finally:  
        return 20
```

```
def f():  
    try:  
        return 10  
    finally:  
        return 20  
return 30
```

```
def f():  
    try:  
        raise Exception()  
    except:  
        return 10  
    finally:  
        print("got here!")  
    return 20
```

```
def f():  
    try:  
        raise Exception()  
    except:  
        return 10  
    finally:  
        print("got here!")  
        return 20
```

```
while True:
    try:
        continue
    finally:
        print "got here!"
```


return

while

for

try/throw

break

continue

finally

call/ec

diet-call/cc

call/ec =
setjmp +
longjmp

return

```
def fun(args):  
    if easy case  
        return something  
    do something  
    complicated
```

```
(lambda (args)  
  (if easy  
      something  
      do something complicated))
```



```
def fun(args):  
    if easy case:  
        return something  
    if also easy case:  
        return something else  
    do something  
    complicated
```

```
(lambda (args)  
  (if easy  
      something  
      (if also easy  
          something else  
          do something complicated)))
```

```
def contains(e, seq):  
    for x in seq:  
        if x is e:  
            return True  
    return False
```

```
(λ (e seq)
  (for ([x seq])
    (if (eq? x e)
        (return #t))))
#f)
```

```
(λ (e seq)
  (for ([x seq])
    (if (eq? x e)
        (1 #t)))
  #f)
```

What *is* return?

A second-class continuation.

λ/return


```
(λ/return (e seq)
  (for ([x seq])
    (if (eq? x e)
        (return #t))))
#f)
```

$(\lambda/\text{return } (\textit{params}) \textit{ body})$

\Rightarrow

$(\lambda (\textit{params})$

$(\text{call/cc } (\lambda (\text{return})$
 $\textit{body})))$

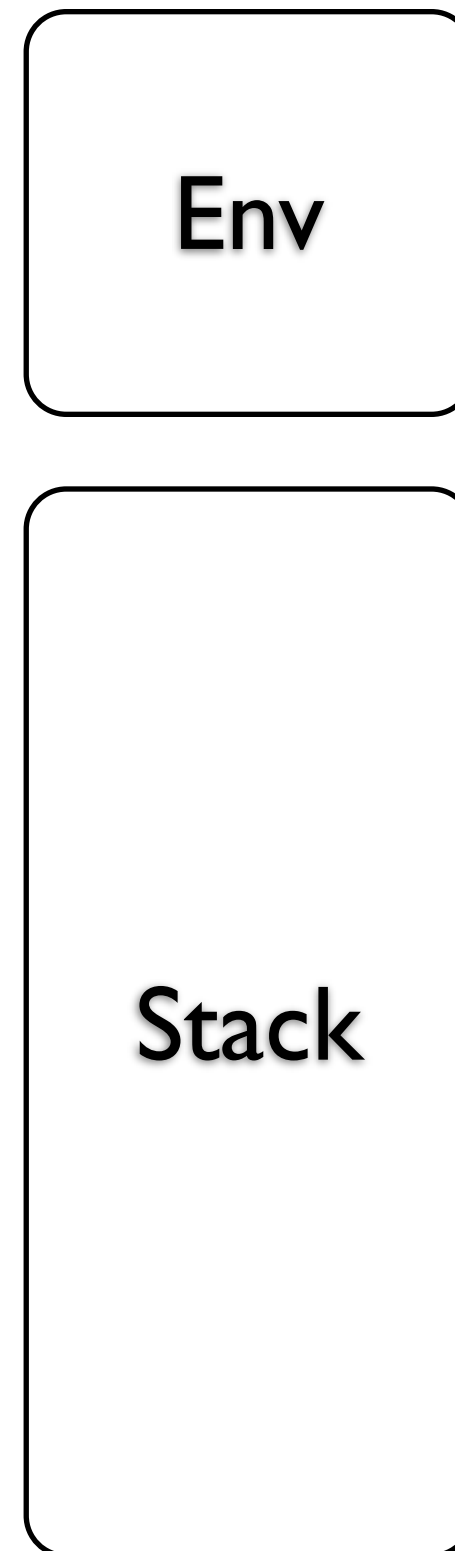
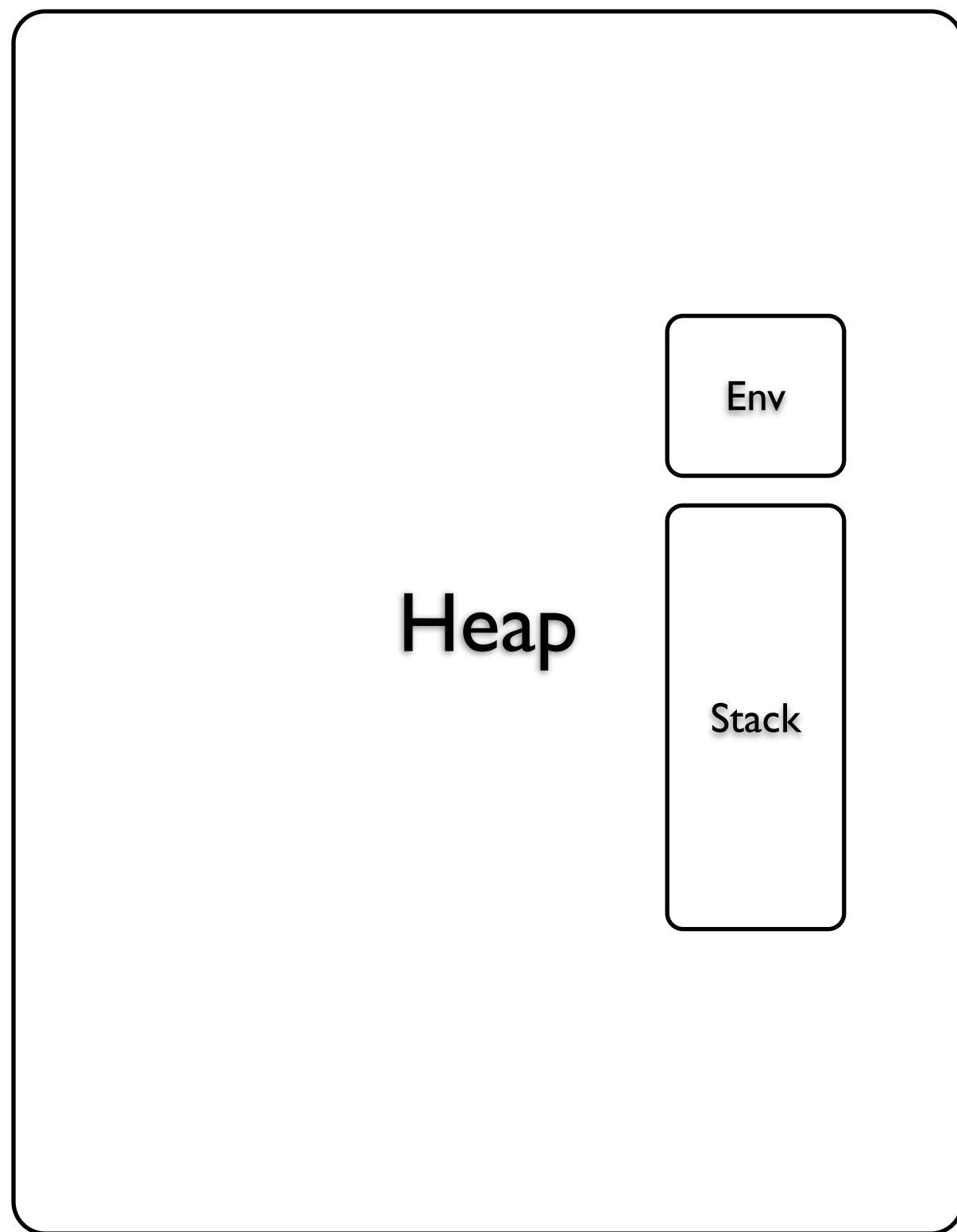
call/cc too heavyweight!

The diagram illustrates a memory layout with three distinct regions. On the left is a large rectangular area labeled 'Heap'. On the right, there are two smaller rectangular areas stacked vertically: the top one is labeled 'Env' and the bottom one is labeled 'Stack'. All three regions are represented by white rounded rectangles with black outlines.

Heap

Env

Stack





The diagram illustrates a memory layout with three distinct regions. On the left is a large rectangular area labeled 'Heap'. On the right, there are two smaller rectangular areas stacked vertically: 'Env' on top and 'Stack' on the bottom. All three regions are represented by white boxes with black outlines and rounded corners.

Heap

Env

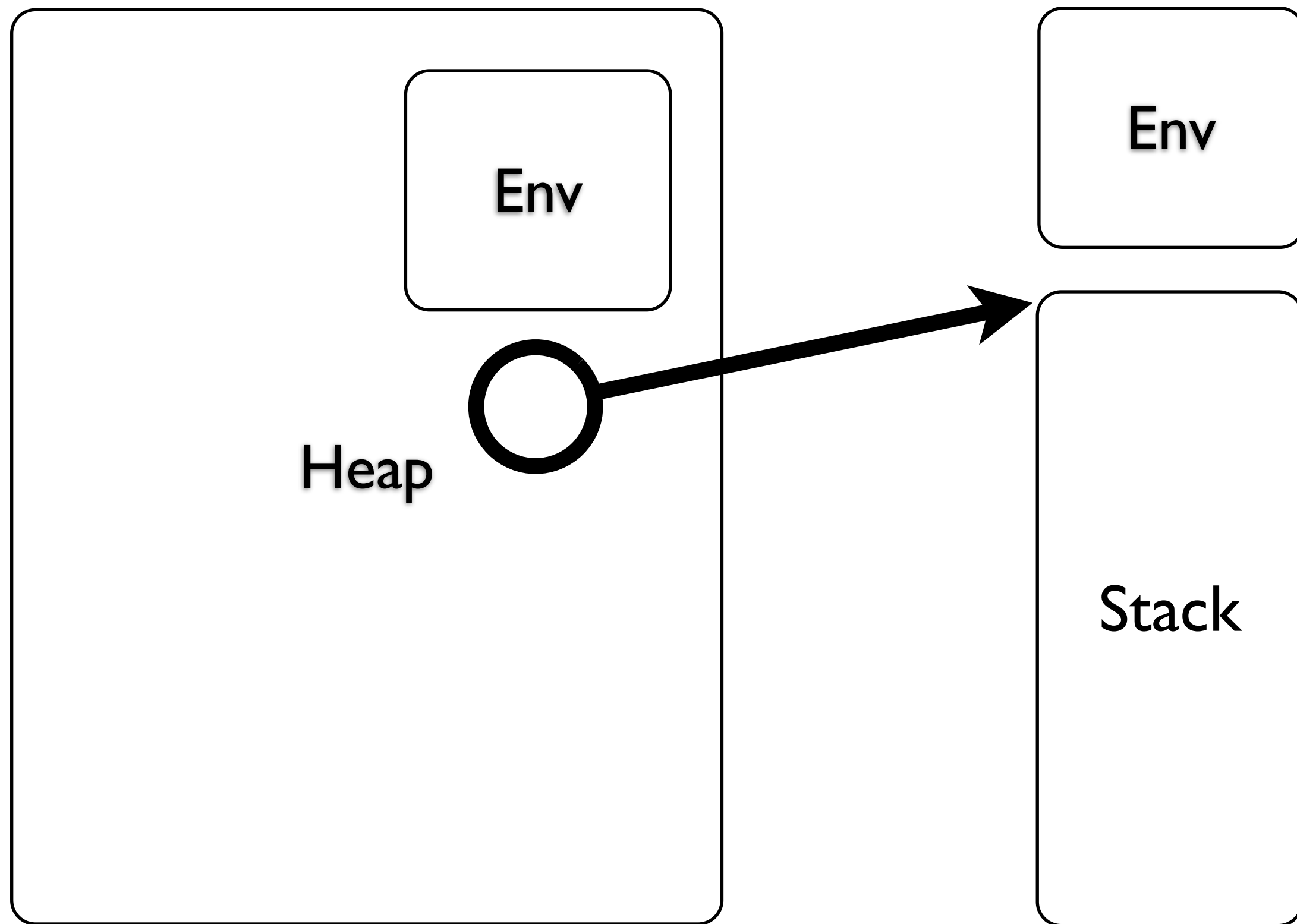
Stack

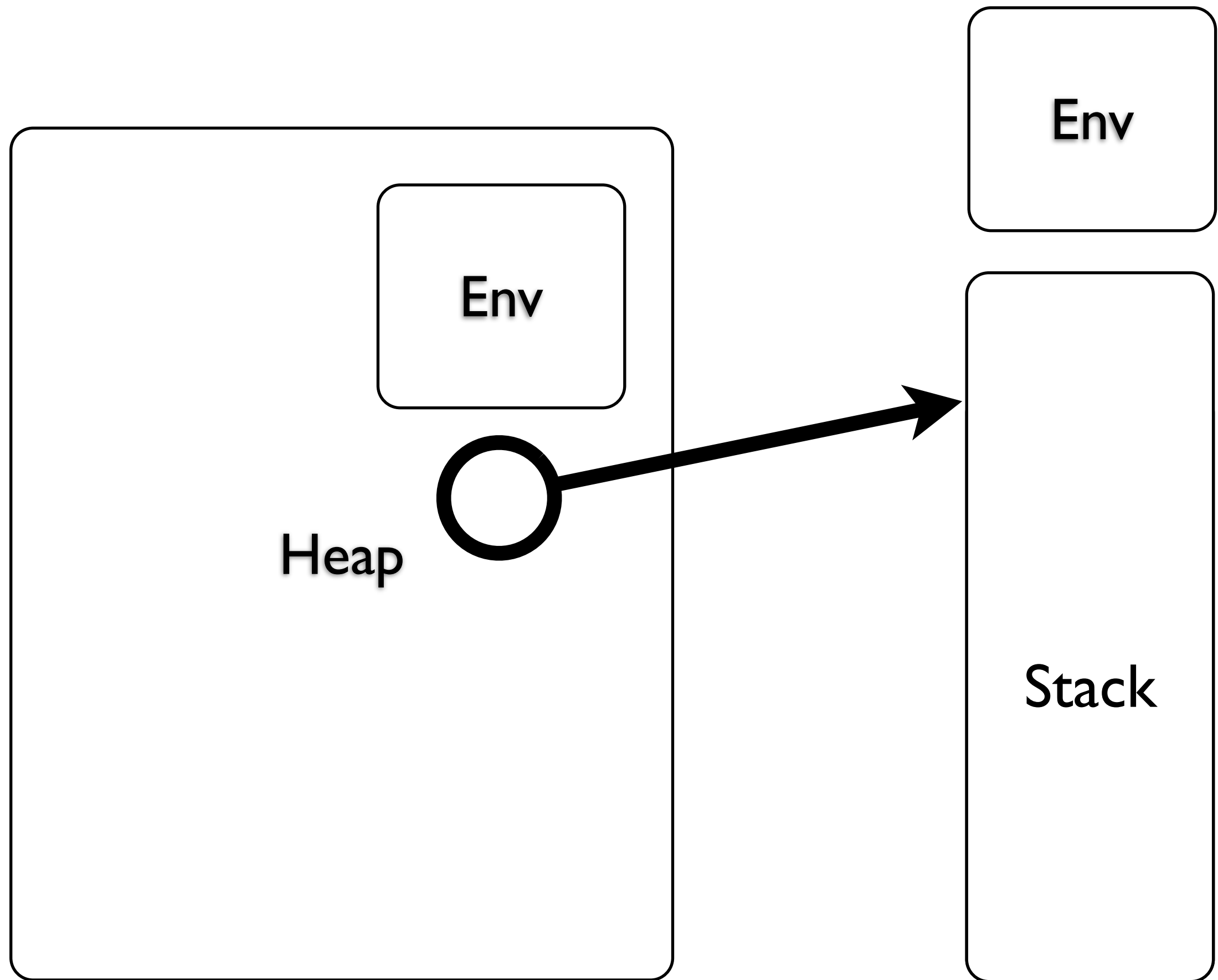
call/ec is just right!

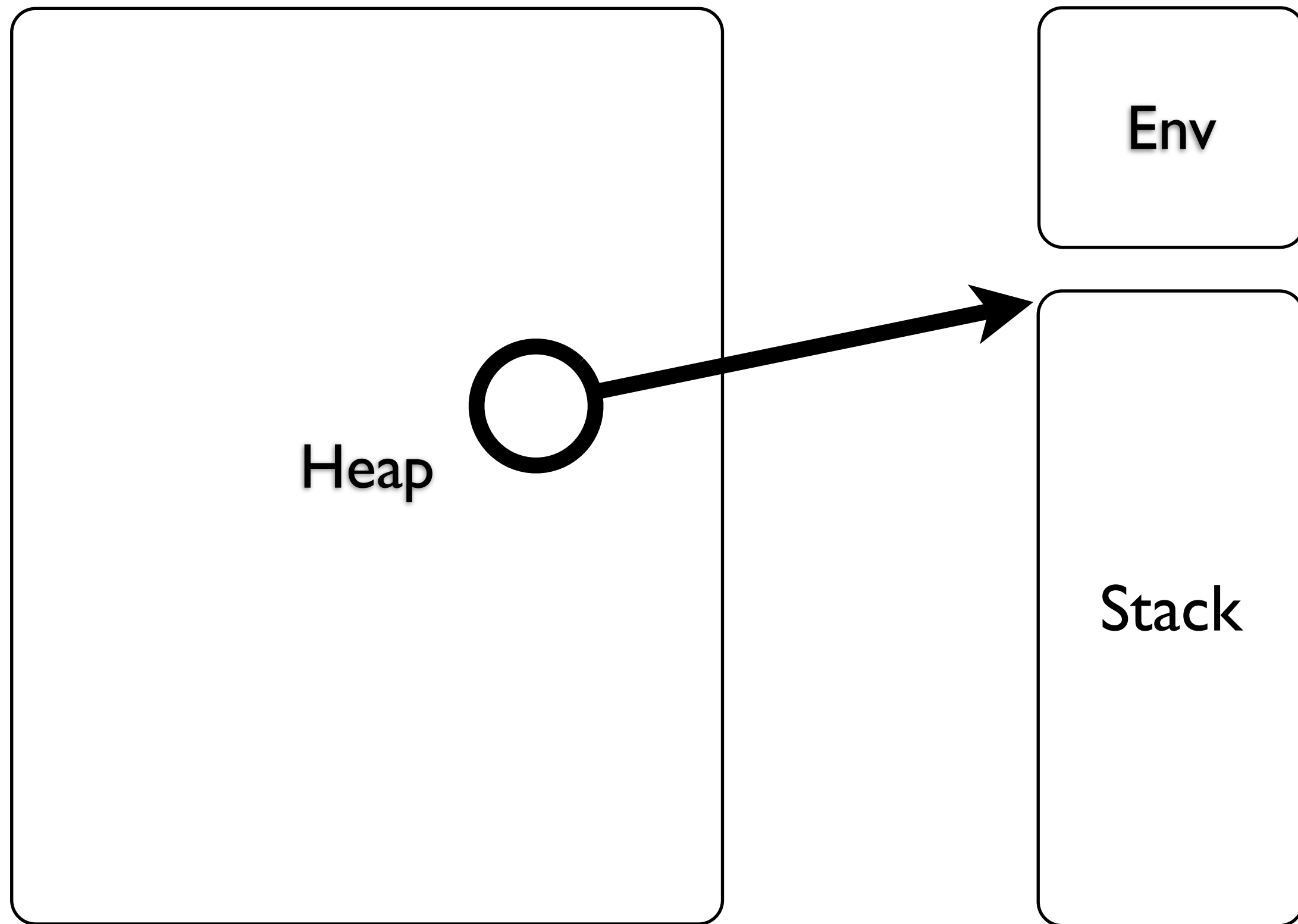
Heap

Env

Stack







$(\lambda / \text{return } (\textit{params}) \textit{ body})$

\Rightarrow

$(\lambda (\textit{params})$

$(\text{call/ec } (\lambda (\text{return})$
 $\textit{body})))$

while

```
while cond:  
    do-this
```

```
(letrec ([loop (λ ()  
  (when cond  
    do-this  
    (loop)))])  
  loop))
```

`while cond:`

do-this

`continue`

do-that


```
(letrec ([loop (λ ()  
  (when cond  
    do-this  
    (continue)  
    do-that  
    (loop)))])  
  (loop))
```

```
(letrec ([loop (λ ()  
  (when cond  
    (call/ec (λ (continue)  
      do-this  
      (continue)  
      do-that))  
    (loop))))])  
  (loop))
```

`while cond:`

do-this

`break`

do-that

```
(letrec ([loop (λ ()  
  (when cond  
    (call/ec (λ (continue)  
      do-this  
      (break)  
      do-that))  
    (loop)))))]  
  (loop))
```

```
(call/ec (λ (break)
  (letrec ([loop (λ ()
    (when cond
      (call/ec (λ (continue)
        do-this
        (break)
        do-that))
      (loop)))
    (loop))))
```

`while cond:`

`do-this`

`else:`

`do-that`

```
(call/ec (λ (break)
  (letrec ([loop (λ ()
    (when cond
      (call/ec (λ (continue)
        do-this))
      (loop)))
    (loop)
    do-that))
```

for


```
for i in seq:  
    do-this
```

(for-each *seq* (λ (*i*) *do-this*))

```
(call/ec (λ (break)
  (for-each seq (λ (i)
    (call/ec (λ (continue)
      do-this))))))
```

for *i* in *seq*:

do-this

else:

do-that

```
(call/ec (λ (break)
  (for-each seq (λ (i)
    (call/ec (λ (continue)
      do-this))))
  do-that)))
```

```
for i in seq:
```

```
    do-this
```

```
else:
```

```
    do-that
```

```
print(i)
```

```
(call/ec (λ (break)
  (for-each seq (λ ($v)
    (set! i $v)
    (call/ec (λ (continue)
      do-this)))
    do-that))
```

throw


```
(define (throw ex)
  ($handler ex))
```

try

Two approaches

- Designated global handler
- Second continuation in CPS

try:

do-this

except:

do-that

(try
 do-this
 (λ (ex)
 do-that))

(try
do-this
handler)

```
(let ([sold $handler])
  (call/ec (lambda (ec)
    (set! $handler
      (λ (ex)
        (set! $handler sold)
        (ec (handler ex))))))
  (let ([rv do-this])
    (set! $handler sold)
    rv))))
```

`while` *cond*:

`try`:

`break`

`except`:

do-that


```
(let* ([sold $handler]
       [break (λ ()
                 (set! $handler $old)
                 (break))])
  (call/ec (lambda (ec)
             (set! $handler
                   (λ (ex)
                     (set! $handler $old)
                     (ec (handler ex))))))
  (let ([rv do-this])
    (set! $handler $old)
    rv))))
```

finally

(try
 do-this
 handler
 final-code)

```
(let ([$fin (λ () (void))])
  (call/ec (λ (finally)
    (let* ([$old $handler]
      [break (λ ()
        (set! $handler $old)
        (set! $fin break)
        (finally))])
      (call/ec (lambda (ec)
        (set! $handler
          (λ (ex)
            (set! $handler $old)
            (ec (handler ex))))))
        (let ([rv do-this])
          (set! $handler $old)
          rv))))))
final-code
($fin))
```

Questions?