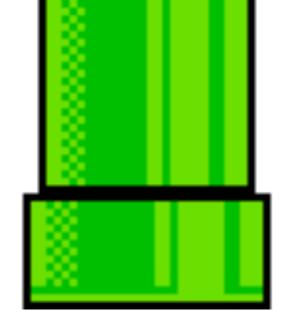
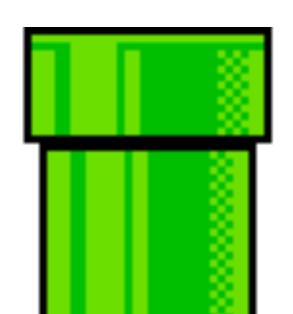
Interpreters





42
Matt Might
University of Utah
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```
function id(x)
{
  return x; // comment
}
```

FUNCTION

IDENT(id)

LPAR

IDENT(x)

RPAR

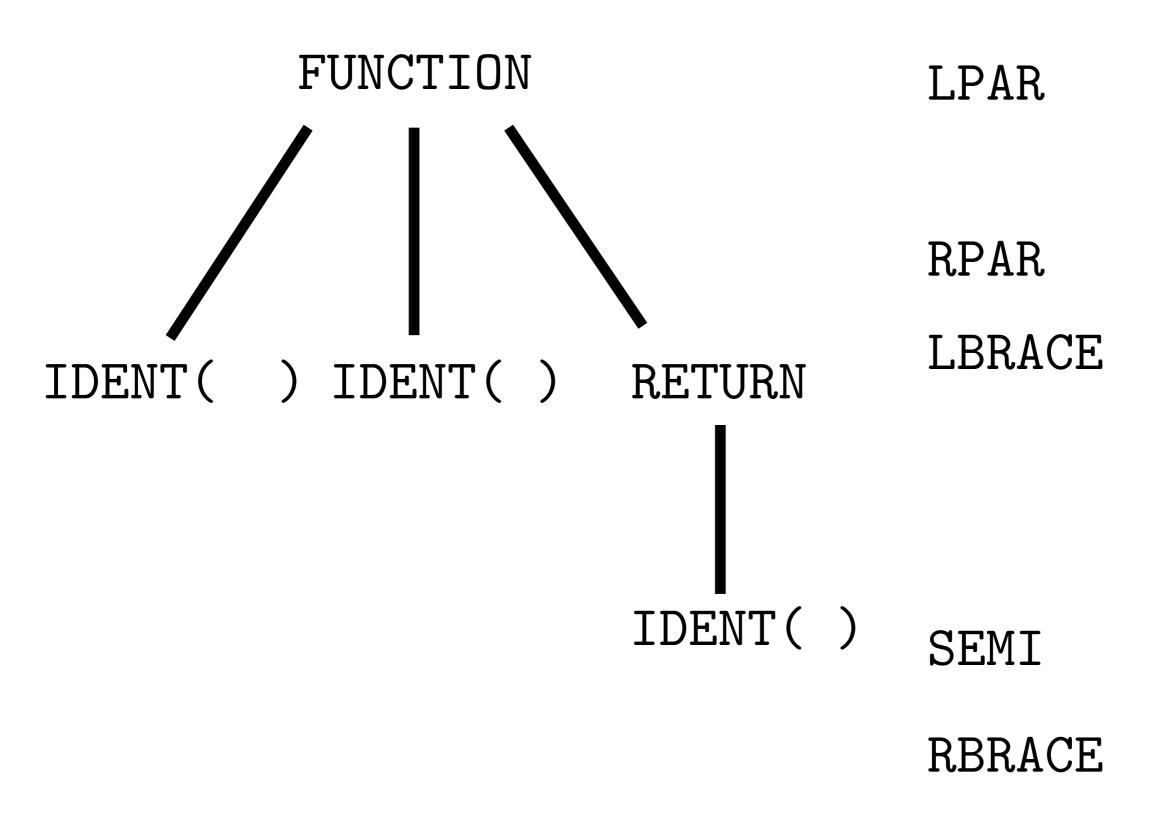
LBRACE

RETURN

IDENT(x)

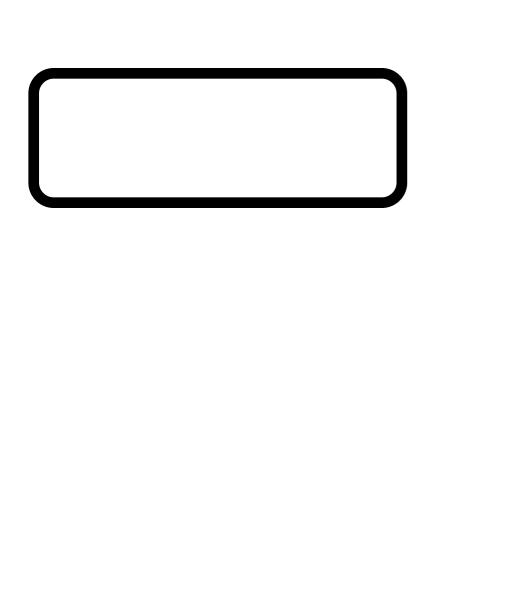
SEMI

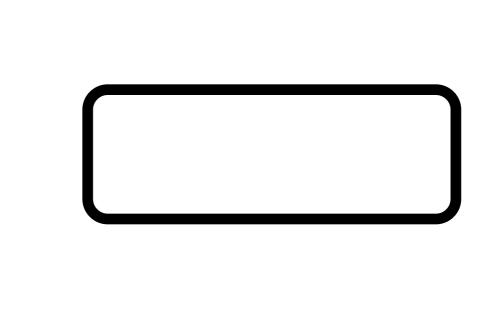
RBRACE

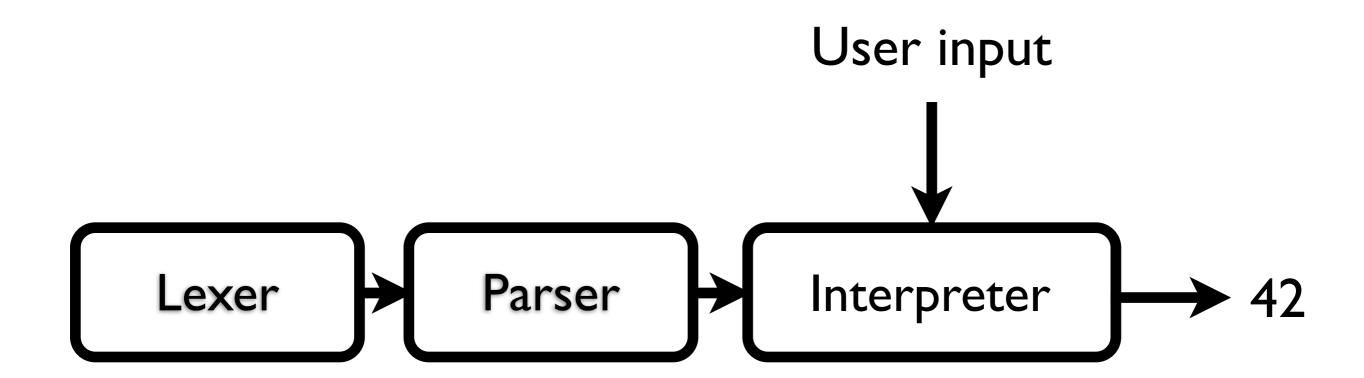


FUNCTION IDENT() IDENT() RETURN IDENT()

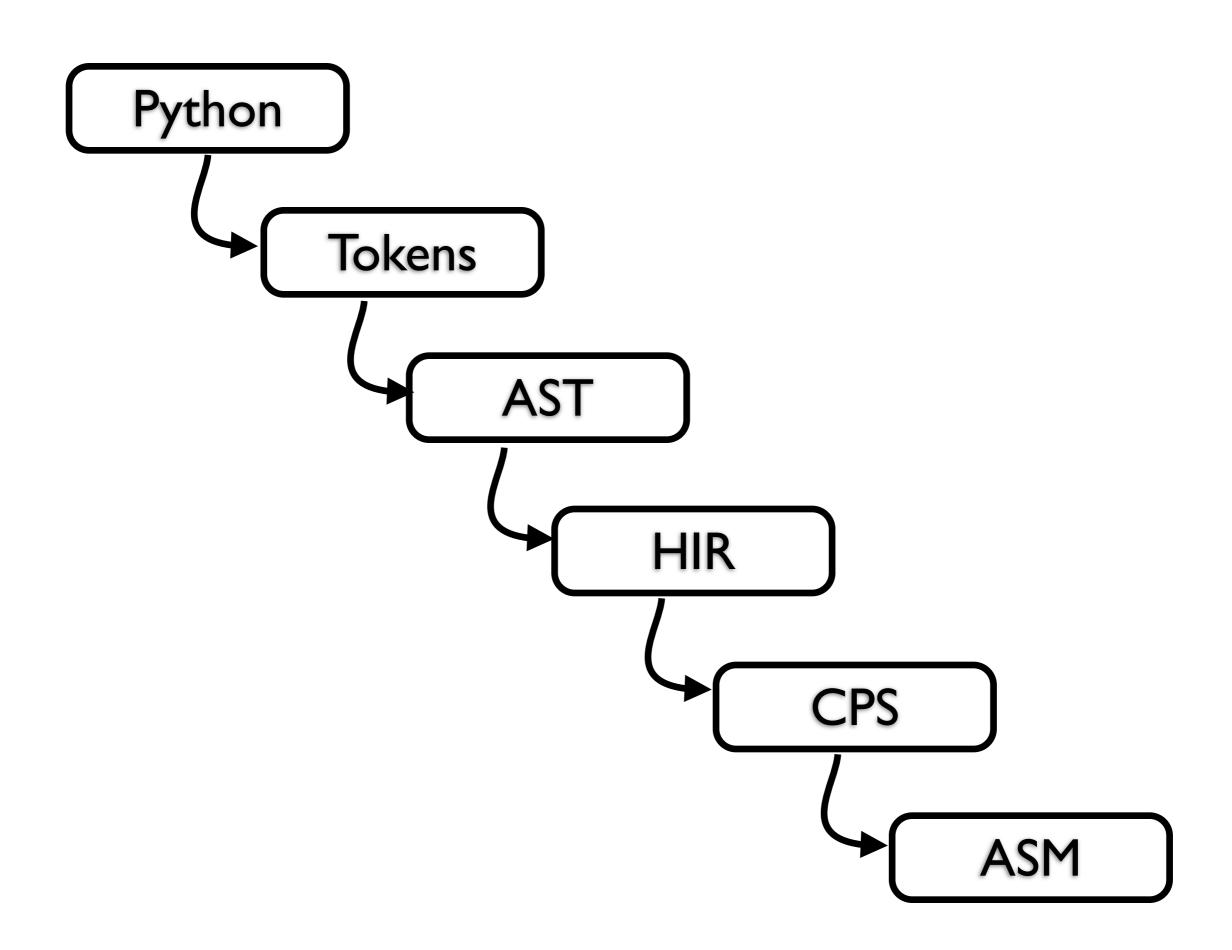
FUNCTION IDENT() IDENT() RETURN IDENT()

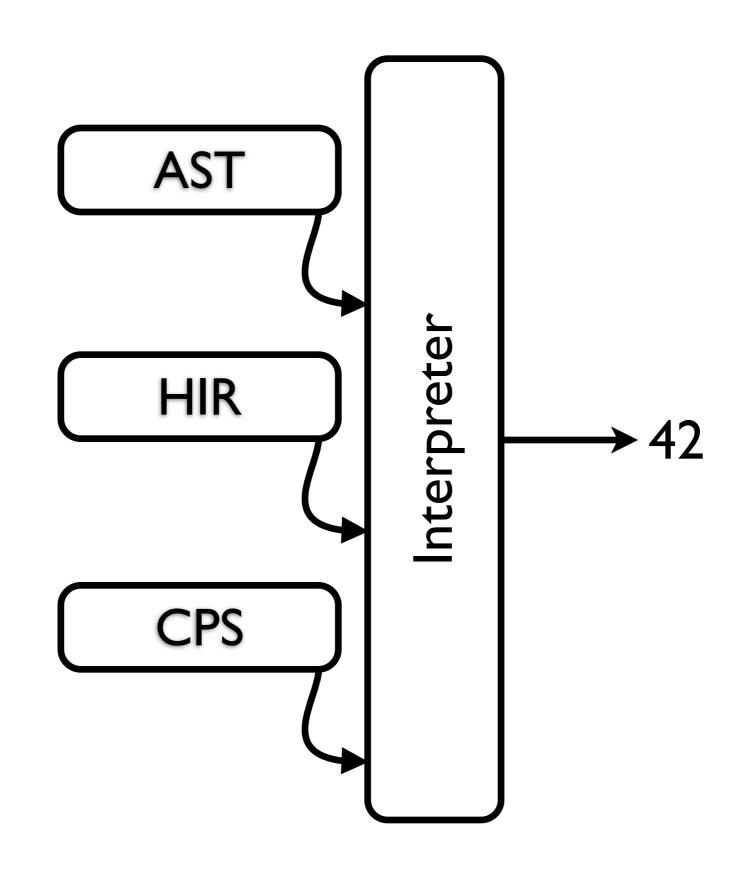






Why study interpreters?





It's a model.

It's easier to write.

It's easier to port.

Compilers > Interpreters?

Interpreters > Compilers?

Neither.

They're equivalent.

How to write interpreters

Syntatic reduction

program

program

 $program \implies program$

 \Longrightarrow program \Longrightarrow program

 $ext{am} \implies ext{program} \implies ext{program}$

$$(1 + 2) * (4 + 5)$$

(3) * (4 + 5)

3 * (4 + 5)

3 * (9)

3 * 9

Code

(if cond then else)

Reduce first

```
(if cond then else)
```

then

else

(let var exp body)

 $\{exp/var\}body$

 $\{e/v\}n$

n

 $\{e/v\}v$

e

 $\{e/v\}v'$



$$\{e/v\}(+ a b)$$

 $(+ \{e/v\}a \{e/v\}b)$

 $\{e/v\}$ (if a c o)

```
(if \{e/v\}a
\{e/v\}c
\{e/v\}o)
```

 $\{e/v\}(\text{let }v a b)$

(let $v \{e/v\}a b$)

 $\{e/v\}$ (let v' a b)

 $(let v' \{e/v\}a \{e/v\}b)$

Code

 $(\lambda v.b)$

 $((\lambda v.b)a)$

 $\{a/v\}b$









Environment-based

What's an environment?

variable => value

hash-table, balanced tree

eval & apply

 $eval: Exp \times Env \rightarrow Value$

eval & apply

 $\mathtt{apply}: \mathit{Value} \times \mathit{Value} \to \mathit{Value}$

apply: Value imes Value o Value

Expects procedure

 $Value = Clo + \mathbb{Z}$

 $Clo = Lambda \times Env$

 $Env = Var \rightarrow Value$

Code

(set! var exp body)

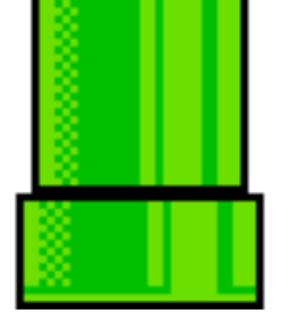
(call/cc exp)

MARIO HORLD TIME 058900 X37 1-4 220

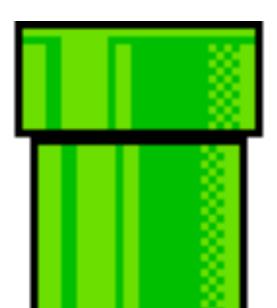
THANK YOU MARIO!

BUT OUR PRINCESS IS IN ANOTHER CASTLE!





Questions?



PHP

