

# **Stabilizer Formalism**

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# Groups

A group is a (non-empty) set  $G$  together with a function  $\star : G \times G \rightarrow G$  such that

1.  $a \star (b \star c) = (a \star b) \star c \quad \forall a, b, c \in G$
2.  $\exists e \in G$  such that  $e \star a = a \star e = a \quad \forall a \in G$
3.  $\forall a \in G \quad \exists a^{-1} \in G$  such that  $a \star a^{-1} = a^{-1} \star a = e \quad \forall a \in G$

Examples:

1.  $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$  with addition.
2.  $GL_{\mathbb{C}}(n) = \{n \times n \text{ invertible matrices}\}$  with matrix product.
3.  $\mathcal{P}_1 = \{\pm I, \pm iI, \pm X, \pm iX, \pm Y, \pm iY, \pm Z, \pm iZ\}$  with matrix product.
4.  $\mathcal{P}_n = \{P_1 \otimes \cdots \otimes P_n : P_j \in \mathcal{P}_1\}$  with matrix product.

# Groups

Let  $G$  be a group. A non-empty subset  $S \subset G$  is a subgroup if

1.  $x \in S \Rightarrow x^{-1} \in S$ .
2.  $x, y \in S \Rightarrow xy \in S$ .

Let  $G$  be a group and fix  $a \in G$ . The subgroup of  $G$  generated by  $a$  is  $\langle a \rangle = \{a^n : n \in \mathbb{Z}\}$

Let  $S \subset G$ . A word on  $S$  is an element  $w \in G$  of the form

$$w = x_1^{e_1} \cdots x_n^{e_n}$$

such that  $x_i \in S$ ,  $e_i \in \{+1, -1\}$ ,  $n \geq 1$

Theorem: Let  $S \subset G$ . If  $S$  is empty, then  $\langle S \rangle = \{e\}$ . If  $S$  is not empty, then  $\langle S \rangle = \{\text{words on } S\}$

# Groups

A group action  $G \curvearrowright X$  of a group  $G$  on a non-empty set  $X$  is a function  $G \times X \rightarrow X$  denoted by  $(g, x) \mapsto g \cdot x$  such that

1.  $e \cdot x = x \quad \forall x \in X$
2.  $g \cdot (h \cdot x) = (gh) \cdot x \quad \forall x \in X \text{ and } \forall g, h \in G$

Example:

$$\mathcal{P}_n \curvearrowright \mathcal{H}^n$$

Let  $G \curvearrowright X$  and fix  $x \in X$ . The stabilizer of  $x$  is defined as

$$G_x = Stab(x) = \{g \in G : g \cdot x = x\}$$

Example:

Consider the Bell pair  $|\psi\rangle = \frac{|00\rangle + |11\rangle}{\sqrt{2}}$ . We have that  $X_1X_2, Z_1Z_2 \in (\mathcal{P}_2)_{|\psi\rangle}$

# Stabilizers

Let  $S$  be a subgroup of  $\mathcal{P}_n$  with an action  $S \curvearrowright \mathcal{H}^n$ . We denote the sub vector space stabilized by  $S$  by

$$V_S = \{|\psi\rangle \in \mathcal{H}^n : P|\psi\rangle = |\psi\rangle \forall P \in S\}$$

$$Stab(|\psi\rangle) = \{P \in \mathcal{P}_n : P|\psi\rangle = |\psi\rangle\}$$

Note that  $Stab(V_S) = S$  and  $V_S \subset \mathcal{H}^n$  is a sub vector space.

## Example

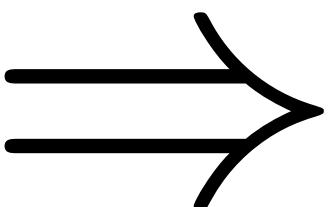
Consider the action  $\mathcal{P}_3 \curvearrowright \mathcal{H}^3$ , and take the subgroup  $S = \langle Z_1Z_2, Z_2Z_3 \rangle = \{I, Z_1Z_2, Z_2Z_3, Z_1Z_3\}$

$$V_{Z_1Z_2} = span\{|\underline{000}\rangle, |\underline{001}\rangle, |\underline{110}\rangle, |\underline{111}\rangle\}$$

$$V_{Z_2Z_3} = span\{|\underline{000}\rangle, |\underline{011}\rangle, |\underline{100}\rangle, |\underline{111}\rangle\}$$

$$V_{Z_1Z_3} = span\{|\underline{000}\rangle, |\underline{010}\rangle, |\underline{101}\rangle, |\underline{111}\rangle\}$$

$$V_I = \mathcal{H}^3$$



$$V_S = span\{|\underline{000}\rangle, |\underline{111}\rangle\}$$

# Stabilizers

Consider the subgroup  $H = \{\pm I, \pm X\} \subset \mathcal{P}_1$ . We have that

$$-I|\psi\rangle = |\psi\rangle \Rightarrow |\psi\rangle = 0 \Rightarrow V_H = \{0\}$$

Non-physical!!!

We want subgroups  $S \subset \mathcal{P}_n$  such that  $V_S \neq \{0\}$

Proposition:  $-I \notin S \Rightarrow \pm iI \notin S$

Any two elements of the Pauli group  $M, N \in \mathcal{P}_n$  either commute or ~~anticommute~~.

Assume for the moment that  $MN = -NM$  and  $V_S \neq \{0\}$  and take  $|\psi\rangle \in V_S$  such that  $|\psi\rangle \neq 0$

$$|\psi\rangle = MN|\psi\rangle = -NM|\psi\rangle = -|\psi\rangle$$

Therefore, if we want  $V_S \neq \{0\}$  we must have that

1.  $-I \notin S$
2.  $MN = NM \forall M, N \in S$  ( $S$  is commutative/Abelian)

# Stabilizers

Therefore, if we want  $V_S \neq \{0\}$  we must have that

1.  $-I \notin S$
2.  $MN = NM \quad \forall M, N \in S$  (S is commutative/Abelian)

Proposition: Let  $S = \langle g_1, \dots, g_k \rangle$  be a subgroup of  $\mathcal{P}_n$ . Then

$$S \text{ is abelian} \Leftrightarrow g_i g_j = g_j g_i \quad \forall i, j$$

Moreover, if  $-I \notin S$  then  $\underline{g^2 = I \quad \forall g \in S}$

**The whole reason for efficient simulation**

Now, we define a function  $r : \mathcal{P}_n \rightarrow \mathbb{Z}_2^{2n}$  that assigns to an element of the Pauli group its  $ZX$  – factors

Forget coefficient

X Z

$$r(X_1 X_2 X_3) = (1, 1, 1 | 0, 0, 0)$$

$$r(-iY_1 Z_3) = (1, 0, 0 | 1, 0, 1)$$

$$r(-X_1 Y_2) = (1, 1 | 0, 1)$$

$$\langle Z_1 Z_2, Z_2 Z_3 \rangle = \{I, Z_1 Z_2, Z_2 Z_3, Z_1 Z_3\}$$

$$r(I)$$

$$0 0 0 0 0 0$$

$$r(Z_1 Z_2)$$

$$\boxed{0 0 0 1 1 0}$$

$$r(Z_2 Z_3)$$

$$\boxed{0 0 0 0 1 1}$$

$$r(Z_1 Z_3)$$

$$0 0 0 1 0 1 =$$

+

For simulation, use 2 bits to save the coefficient

# Stabilizers

Proposition: Let  $\mathbb{I}_n$  be the  $n \times n$  identity matrix and denote  $\Lambda = \begin{pmatrix} 0_n & \mathbb{I}_n \\ \mathbb{I}_n & 0_n \end{pmatrix}$ . Let  $g, g' \in \mathcal{P}_n$ , then

$$gg' = g'g \Leftrightarrow \boxed{r(g)\Lambda r(g')^T = 0} \quad \text{Symplectic form}$$

Moreover,  $r : \mathcal{P}_n \rightarrow \mathbb{Z}_2^{2n}$  is an epimorphism of groups and  $\ker(r) = \{\pm I^{\otimes n}, \pm iI^{\otimes n}\}$

We say that the generators  $g_1, \dots, g_k$  of a subgroup  $S = \langle g_1, \dots, g_n \rangle$  of  $\mathcal{P}_n$  are independent  
Definition  
if the vectors  $r(g_1), \dots, r(g_k)$  are linearly independent.

Proposition: If  $-I \notin S = \langle g_1, \dots, g_k \rangle$  and the generators are independent, then for any  $j$

the subgroup  $\langle g_1, \dots, \hat{g}_j, \dots, g_k \rangle$  is strictly smaller than  $S$ .

# Stabilizers

Theorem: Let  $S = \langle g_1, \dots, g_{n-k} \rangle$  be an abelian subgroup of  $\mathcal{P}_n$  with independent generators and  $-I \notin S$ . Then  $\dim_{\mathbb{C}} V_S = 2^k$

$$X_1 X_2 Z_1 Z_2 = Z_1 Z_2 X_1 X_2$$

Example:

$$|\psi\rangle = \frac{|00\rangle + |11\rangle}{\sqrt{2}}$$

$$\mathcal{P}_2 \curvearrowright \mathcal{H}^2$$

$$S = Stab(|\psi\rangle) = \langle X_1 X_2, Z_1 Z_2 \rangle$$

$k = 0$

$\Rightarrow \dim_{\mathbb{C}}(V_S) = 1$

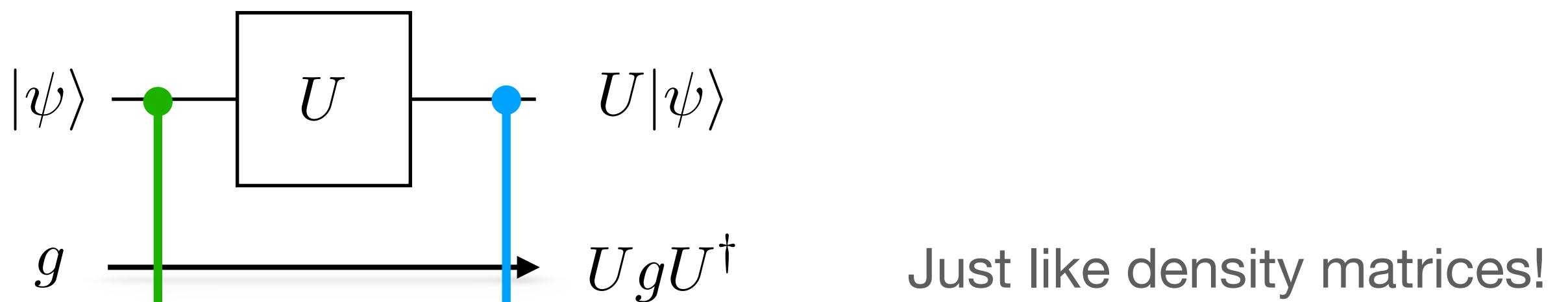
We know the state up to a global phase!

# Evolution

Let  $S = \langle g_1, \dots, g_{n-k} \rangle$  be an abelian subgroup of the Pauli group  $\mathcal{P}_n$  with independent generators such that  $-I \notin S$  and let  $|\psi\rangle \in V_S - 0$ . Let  $U$  be a unitary on  $n$ -qubits. Let  $g \in S$ . Then

$$U|\psi\rangle = Ug|\psi\rangle = UgU^\dagger U|\psi\rangle \quad U|\psi\rangle = (UgU^\dagger)U|\psi\rangle$$

This means that the state  $U|\psi\rangle$  is stabilized by  $UgU^\dagger$

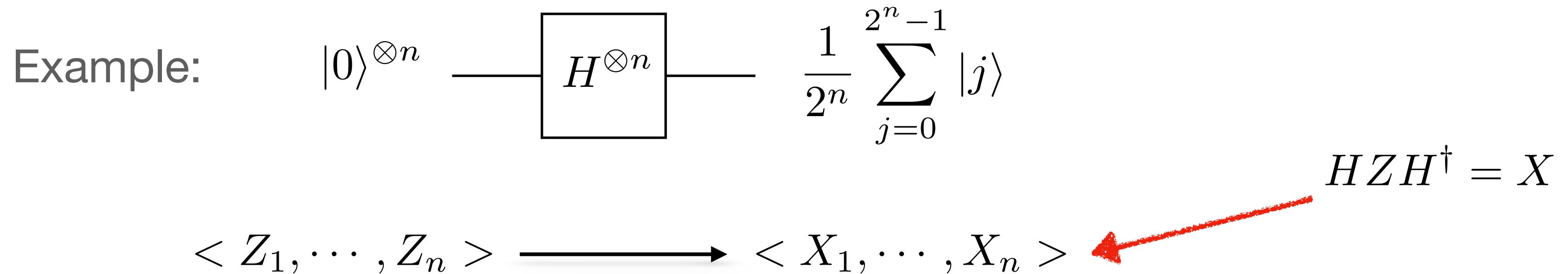


Proposition: The vector space  $UV_S = \{U|\psi\rangle : |\psi\rangle \in V_S\}$  is stabilized by the group  $USU^\dagger = \{UgU^\dagger : g \in S\}$

In this case,  $USU^\dagger = \langle Ug_jU^\dagger : j = 1, \dots, n-k \rangle$

$$\langle g_1, \dots, g_{n-k} \rangle \longrightarrow \langle Ug_1U^\dagger, \dots, Ug_{n-k}U^\dagger \rangle$$

# Evolution



EXTREMELY IMPORTANT:  $\dim_{\mathbb{C}}(V_S) = 1$  Same number of generators as qubits

So we know the state up to a global phase from the stabilizers

Yeah, sure.... But what about entanglement?

We do a little bit more algebra!

# Evolution

Let  $U$  be the CNOT gate controlled by qubit 0 and target qubit 1. Then

$$UX_1U^\dagger = X_1X_2$$

$$UX_2U^\dagger = X_2$$

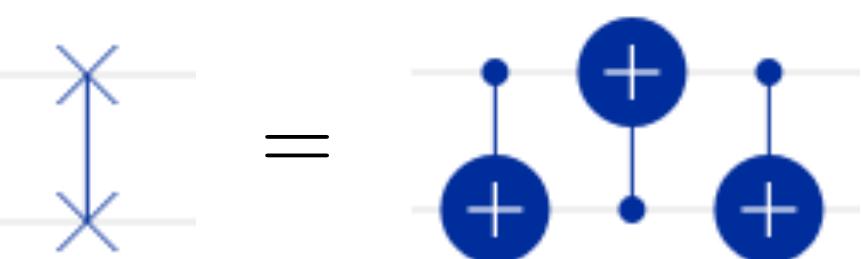
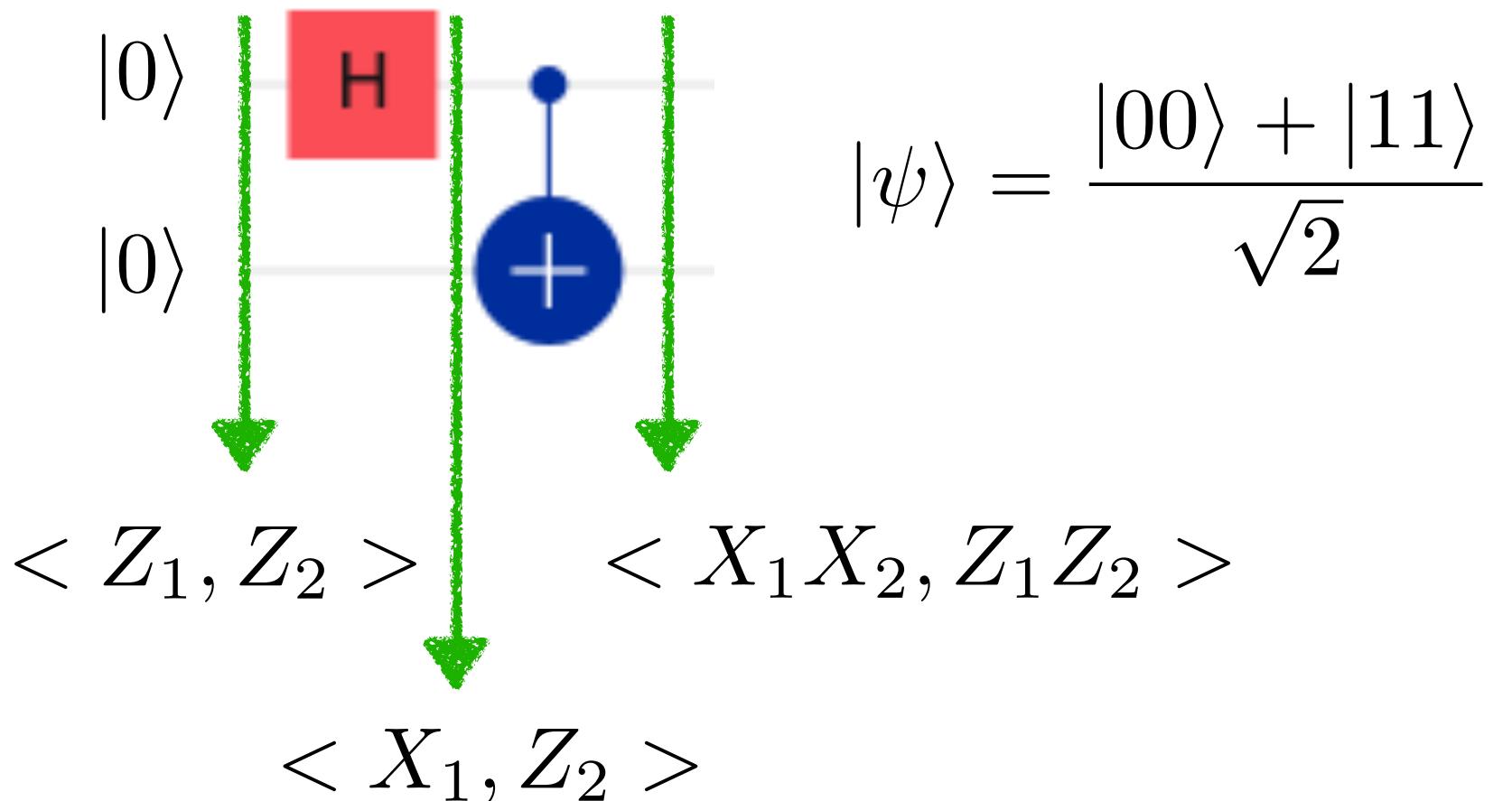
$$UZ_1U^\dagger = Z_1$$

$$UZ_2U^\dagger = Z_1Z_2$$

And this is enough because  $\mathcal{P}_n = \langle X_1, \dots, X_n, Z_1, \dots, Z_n \rangle$

$$UY_2U^\dagger = iUX_2Z_2U^\dagger = iUX_2U^\dagger UZ_2U^\dagger = iX_2Z_1Z_2 = Z_1Y_2$$

Example:



# Limitations

It is all about group theory

Let  $G$  be a group and  $H$  a subgroup of  $G$ . The normalizer of  $H$  on  $G$  is

$$N_G(H) = \{U \in G : UHU^\dagger = H\}$$

Clifford group

Examples:

$$N_{U(2^n)}(\mathcal{P}_n) = \{U \in U(2^n) : U\mathcal{P}_nU^\dagger = \mathcal{P}_n\} = \langle H_i, CX(i, j), S_i : 1 \leq i \neq j \leq n \rangle$$

Let  $U$  denote the Toffoli gate controlled by qubits 0 and 1 and target qubit 2. Let  $T = \begin{pmatrix} 1 & 0 \\ 0 & e^{\frac{i\pi}{4}} \end{pmatrix}$

$$UZ_1U^\dagger = Z_1$$

$$UX_2U^\dagger = X_2 \otimes \frac{I + Z_2 + X_3 - Z_2X_3}{2}$$

$$UX_3U^\dagger = X_3$$

$$U \notin N_{U(2^n)}(\mathcal{P}_n)$$

$$TZT^\dagger = Z$$

$$TXT^\dagger = \frac{X + Y}{\sqrt{2}}$$

$$T \notin N_{U(2^n)}(\mathcal{P}_n)$$

# Measurements

Assume, without loss of generality, that we measure only in the computational basis.

Let  $g \in \mathcal{P}_n$  and assume for the moment that it does not have factors  $-1, \pm i$ .

Suppose that the system is in state  $|\psi\rangle$  with stabilizer group  $S = \langle g_1, \dots, g_n \rangle$ .

We have two possibilities.

1.  $[g, g_j] = 0 \forall j$

$$\dim_{\mathbb{C}}(V_S) = 1$$

We get that  $g \in S$  because  $g_j g |\psi\rangle = gg_j |\psi\rangle = g |\psi\rangle$ , which gives that  $g |\psi\rangle \in V_S$  is a multiple of  $|\psi\rangle$

We also know that  $\underline{g^2 = I}$  so  $\underline{g |\psi\rangle = \pm |\psi\rangle}$ . If  $\boxed{g \in S}$ , the measurement of  $g$  gives +1 with probability 1

and the measurement does not change the state. Trick:  $g = +\frac{I+g}{2} - \frac{I-g}{2}$

If  $\boxed{-g \in S}$ , then  $-g |\psi\rangle = |\psi\rangle$  and the measurement of  $-g$  gives +1 with probability 1

Therefore, measuring  $g$  does NOT change the state and it leaves the same stabilizers.

# Measurements

2.  $\{g, g_j\} = 0$  for at least one  $j$ .

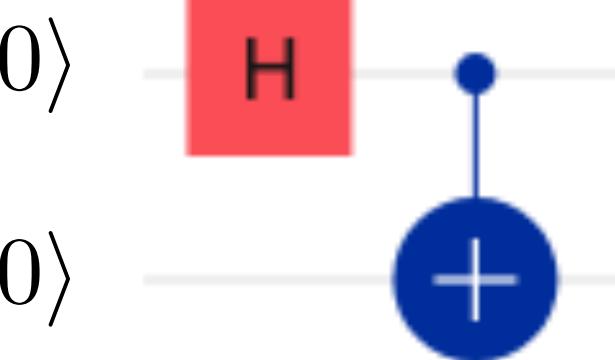
$$\begin{aligned}
 \mathbb{P}_{|\psi\rangle}(\pm 1) &= Tr \left( \frac{I \pm g}{2} |\psi\rangle\langle\psi| \right) \\
 &= Tr \left( \frac{I \pm g}{2} g_j |\psi\rangle\langle\psi| \right) \\
 &= Tr \left( g_j \frac{I \mp g}{2} |\psi\rangle\langle\psi| \right) \\
 &= Tr \left( \frac{I \mp g}{2} |\psi\rangle\langle\psi| g_j \right) \\
 &= Tr \left( \frac{I \mp g}{2} |\psi\rangle\langle\psi| \right) \\
 &= \frac{1}{2}
 \end{aligned}$$

- $g_j \in S$
- $\{g, g_j\} = 0$
- Trace is cyclic
- $g_j^\dagger = g_j \in S$

If we measure +1, the state changes to  $|\psi^+\rangle \propto P^+|\psi\rangle$   
whose stabilizer group is  $\langle g_1, \dots, g_{j-1}, g, g_{j+1}, \dots, g_n \rangle$

If we measure -1, the state changes to  $|\psi^-\rangle \propto P^-|\psi\rangle$   
whose stabilizer group is  $\langle g_1, \dots, g_{j-1}, -g, g_{j+1}, \dots, g_n \rangle$

# Measurements

Example:  $|0\rangle$    $|\psi\rangle = \frac{|00\rangle + |11\rangle}{\sqrt{2}}$

Measure  $Z_1 Z_2$ :

Gives +1 with probability 1 and the state and stabilizers remain the same.

Measure  $Z_1$ :

Gives  $\pm 1$  with probability 1/2 each. If the measurement gives +1, the state changes to  $|\psi^+\rangle$  and the stabilizer group changes to  $\langle Z_1, Z_1 Z_2 \rangle = \langle Z_1, Z_2 \rangle$

If the measurement gives -1, the state changes to  $|\psi^-\rangle$  and the stabilizer group changes to  $\langle -Z_1, Z_1 Z_2 \rangle = \langle -Z_1, -Z_2 \rangle$

# Gottesman-Knill Theorem

Theorem: Assume a quantum computation is performed that uses only

- 1.- State preparation in the computational basis.
- 2.-  $H_j, S_j, CX(i, j), P \in \mathcal{P}_n \quad \forall 0 \leq i \neq j \leq n - 1$
- 3.- Measurement of any  $P \in \mathcal{P}_n$ .
- 4.- Classical control conditioned on the results of these measurements.

Such a quantum computation is efficiently simulable on a classical computer.

Proof: Follow the stabilizers.

## Implementations

[https://github.com/MarcoArmenta/info\\_quantique\\_theorique](https://github.com/MarcoArmenta/info_quantique_theorique)