

# Time & Space Complexity and Big O Notation | Lecture 25.

① Time Complexity: No. of operations as a function of Input.

Asymptotic Analysis. (also known as)

\* Worst case ~~time~~ time complexity  $\rightarrow$  Big O,  $O(n)$  (ingenual)

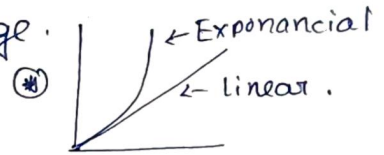
\* Best case time complexity  $\rightarrow$  Big omega,  $\Omega(1)$

\* ~~Average~~ Average case time complexity  $\rightarrow$  Big theta  $\Theta(n)$

\* ~~In 1 second~~ In 1 second  $\rightarrow$   $10^8 - 10^9$  operations can be performed on an average.

①  $O(1) \rightarrow$  constant time complexity.

\* TLE means time limit Exceeded.



\* Time complexity for traversing an array of length  $N$  is  $\rightarrow$  Big  $O(n)$ .

\* Time complexity when traversing 2 individual array of length  $m$  &  $N$  respectively  $\rightarrow O(m+n)$ .

\* Nested loops time complexity  $\rightarrow O(n^2)$ . [where  $i < n$  &  $j < n$ ]  
 $O(n \times n)$  [where  $i < n$  &  $j < \sqrt{n}$ ]

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\* Time complexity for traversing the Array and multiply the increment pointer by 2.  $\Rightarrow O(\log n)$

$$\begin{aligned} * \log_a x^n &= n \log_a x \\ * \log_b b &= 1 \\ * \log_a b &= \frac{\log_c b}{\log_c a} \end{aligned}$$

$\hookrightarrow$  `int val = 0;`  
`for (int i = 1; i <= N; i *= k) {`  
`val++;`  
`}`

$\otimes$  multiply the increment  
pointer by  $k$ .

Time complexity =  $O(\log_k N)$

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$\otimes$  Space Complexity: extra memory / space used by an algorithm with respect to the input size.  $\propto$  input size.

$\hookrightarrow$  Input & Output not considered as space. (clarify that)

$\hookrightarrow$  Asymptotic Analysis.

$\hookrightarrow$  Space complexity for an array of length  $N$ .  $\rightarrow O(n)$

$\hookrightarrow$  Space complexity for a 2-d matrix of  $N$  rows and  $M$  columns  $\rightarrow O(m \times n)$

$\hookrightarrow$