

# Hrishikesh Mishra

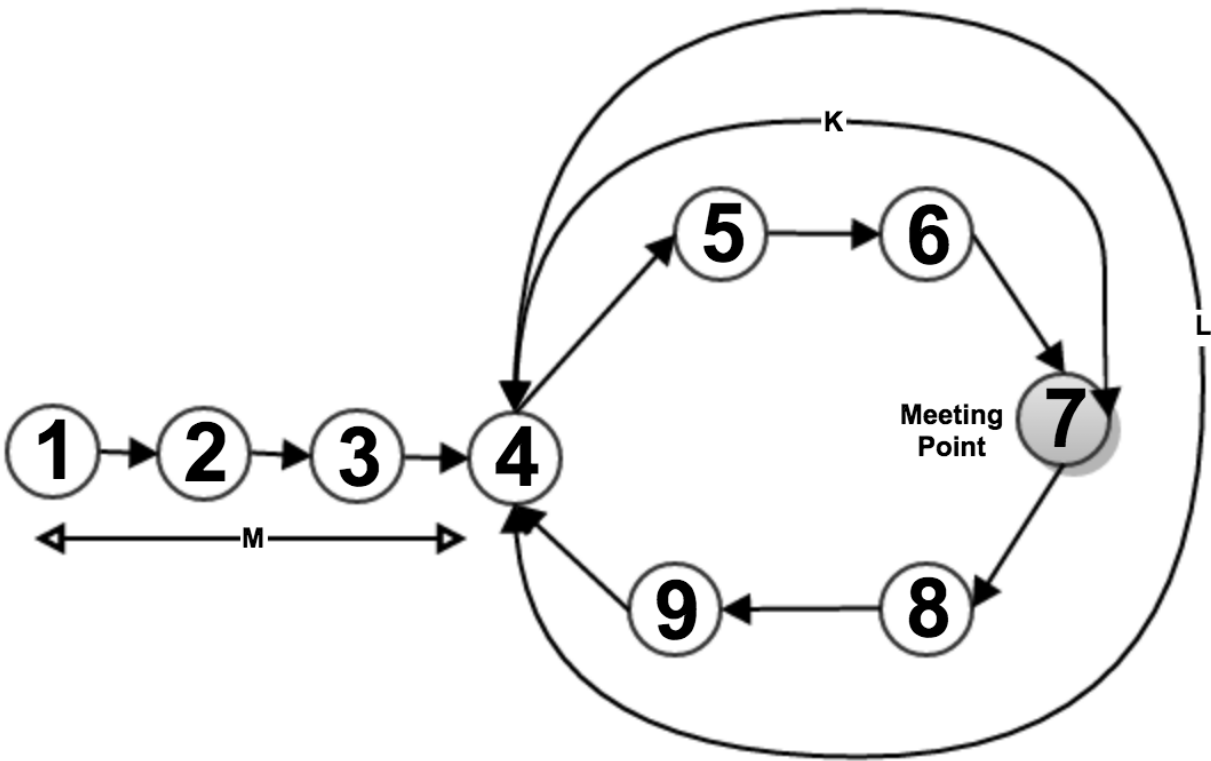
A Code Artist

## How to find the starting point of a loop in a linked list.

### How to find the starting point of a loop in a linked list.

It is based on [Floyd's algorithm for cycle detection](#).  
First we try to find out, is there any loop in list or not. If loop exists then we try to find out starting point of loop. For this we use two pointers namely *slowPtr* and *fastPtr*. In first detection (checking for loop), *fastPtr* moves two steps at once but *slowPtr* moves by one step ahead at once.

Example:



<i>slowPtr</i>	1	2	3	4	5	6	7
<i>fastPtr</i>	1	3	5	7	9	5	7

It is clear that if there is any loop in list then they will meet at point (Point 7 in above image), because *fastPtr* pointer is running twice faster than other one.

Now, we come to second problem of finding starting point of loop.

Suppose, they meet at *Point 7* (as mentioned in above image). Then, *slowPtr* comes out of loop and stands at beginning of list means at *Point 1* but *fastPtr* still at meeting point (*Point 7*). Now we compare both pointers value, if they same then it is starting point of loop otherwise we move one step at ahead (here *fastPtr* is also moving by one step each time) and compare again till we find same point.

<b><i>slowPtr</i></b>	1	2	3	4
<b><i>fastPtr</i></b>	7	8	9	4

Now one question comes in mind, how is it possible. So there is good mathematical proof.

Suppose:

```
m => length from starting of list to starting of loop (i.e 1-2-3-4)
l => length of loop (i.e. 4-5-6-7-8-9)
k => length between starting of loop to meeting point (i.e. 4-5-6-7)
```

```
Total distance traveled by slowPtr = m + p(l) + k
where p => number of repetition of circle covered by slowPtr
```

```
Total distance traveled by fastPtr = m + q(l) + k
where q => number of repetition of circle covered by fastPtr
```

Since,  
*fastPtr* running twice faster than *slowPtr*

Hence,  
Total distance traveled by *fastPtr* = 2 X Total distance traveled by *slowPtr*

i.e

$$m + q(l) + k = 2 * ( m + p(l) + k )$$

$$\text{or, } m + k = q(l) - p(l)$$

$$\text{or, } m + k = (q-p) l$$

$$\text{or, } m = (q-p) l - k$$

So,

If *slowPtr* starts from beginning of list and travels "m" length then, it will reach to *Point 4* (i.e.

and

*fastPtr* start from *Point 7* and travels " (q-p) l - k " length then, it will reach to *Point 4* (i.e. because "(q-p) l" is a complete circle length with " (q-p) " times.

Source code:

Github: [LoopDetector.java](#)

### Output:

```
[1=>2, 2=>3, 3=>4, 4=>5, 5=>6, 6=>7, 7=>8, 8=>9, 9=>4, ]  
Loop Starting point : ListNode(4)
```

This entry was posted in Data Structure and Algorithm, Linked List and tagged algorithm, data structure, linked list on July 3, 2016 [<http://hrishikeshmishra.com/how-to-find-the-starting-point-of-a-loop-in-a-linked-list/>] .

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