BILKENT UNIVERSITY CS202

HOMEWORK 1

İrem Seven

Section : 1

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Part 1)

a)

We need to find two positive constants c and n0 such that

0 <= 20 n^4 + 20 n^2 + 5 <= n^5

for all n >= n0

dividing both sides with n^5 we can see that to make c constant integer we can put n0 = 1 into n.

Then, with n0 = 1 we get c = 45, for all n >= 1.

b)

Selection Sort:

18	4	47	24	15	24	17	11	31	23	initial
18	4	**23**	24	15	24	17	11	31	**47**	
18	4	23	24	15	24	17	11	**31**	47	
18	4	23	**11**	15	24	17	**24**	31	47	
18	4	23	11	15	**17**	**24**	24	31	47	
18	4	**17**	11	15	**23**	24	24	31	47	
15	4	17	11	**18**	23	24	24	31	47	
15	4	**11**	**17**	18	23	24	24	31	47	
11	4	**15**	17	18	23	24	24	31	47	
11	**4**	15	17	18	23	24	24	31	47	
4	11	15	17	18	23	24	24	31	47	

Bubble Sort:

18	**4**	47	24	15	24	17	11	31	23	initial
4	**18**	**47**	24	15	24	17	11	31	23	
4	18	**47**	**24**	15	24	17	11	31	23	
4	18	24	**47**	**15**	24	17	11	31	23	
4	18	24	15	**47**	**24**	17	11	31	23	
4	18	24	15	24	**47**	**17**	11	31	23	
4	18	24	15	24	17	**47**	**11**	31	23	
4	18	24	15	24	17	11	**47**	**31**	23	
4	18	24	15	24	17	11	31	**47**	**23**	
4	**18**	24	15	24	17	11	31	23	47	
4	**18**	**24**	15	24	17	11	31	23	47	
4	18	**24**	**15**	24	17	11	31	23	47	
4	18	15	**24**	**24**	17	11	31	23	47	
4	18	15	24	**24**	**17**	11	31	23	47	
4	18	15	24	17	**24**	**11**	31	23	47	
4	18	15	24	17	11	**24**	**31**	23	47	
4	18	15	24	17	11	24	**31**	**23**	47	
4	**18**	**15**	24	17	11	24	23	31	47	
4	15	**18**	**24**	17	11	24	23	31	47	
4	15	18	**24**	**17**	11	24	23	31	47	
4	15	18	17	**24**	**11**	24	23	31	47	
4	15	18	17	11	**24**	**24**	23	31	47	
4	15	18	17	11	24	**24**	**23**	31	47	
4	**15**	**18**	17	11	24	23	24	31	47	
4	15	**18**	**17**	11	24	23	24	31	47	
4	15	17	**18**	**11**	24	23	24	31	47	
4	15	17	11	**18**	**24**	23	24	31	47	
4	15	17	11	18	**23**	**24**	24	31	47	

4	15	17	11	18	23	24	24	31	47
4	15	17	11	18	23	24	24	31	47
4	15	17	11	18	23	24	24	31	47
4	15	17	11	18	23	24	24	31	47
4	15	17	11	18	23	24	24	31	47
4	15	17	11	18	23	24	24	31	47
4	15	17	11	18	23	24	24	31	47
4	15	17	11	18	23	24	24	31	47
4	15	17	11	18	23	24	24	31	47
4	15	11	17	18	23	24	24	31	47
4	15	11	17	18	23	24	24	31	47
4	15	11	17	18	23	24	24	31	47
4	15	11	17	18	23	24	24	31	47
4	15	11	17	18	23	24	24	31	47
4	15	11	17	18	23	24	24	31	47
4	15	11	17	18	23	24	24	31	47
4	11	15	17	18	23	24	24	31	47
4	11	15	17	18	23	24	24	31	47

Part 2)

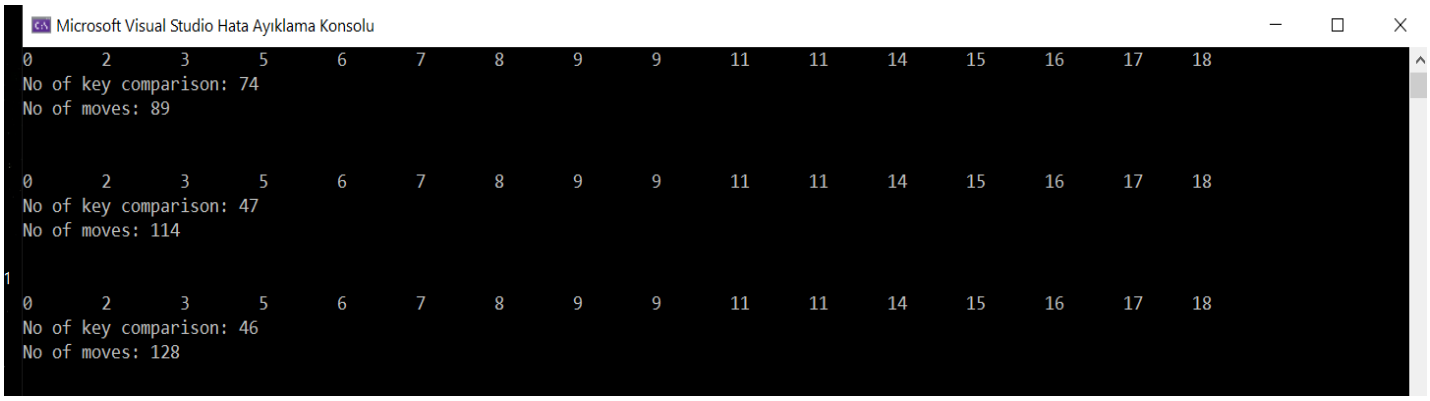


Figure1: Screenshot of Question 2 part a

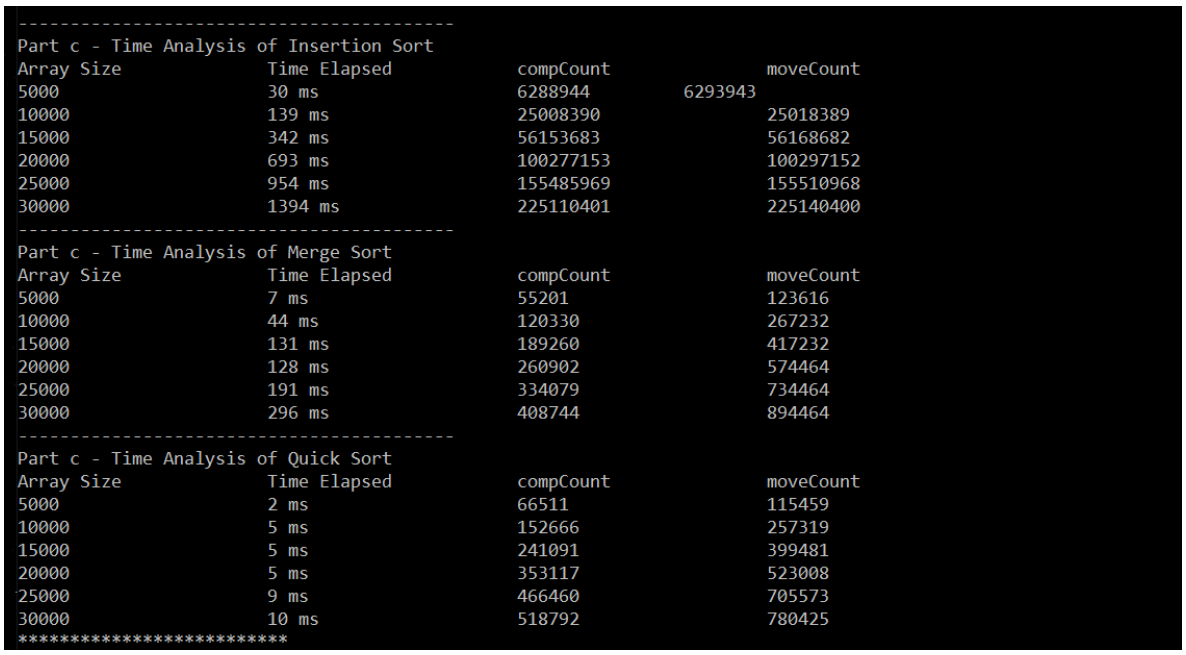


Figure2: Screenshot of Randomly Created Performance Analysis

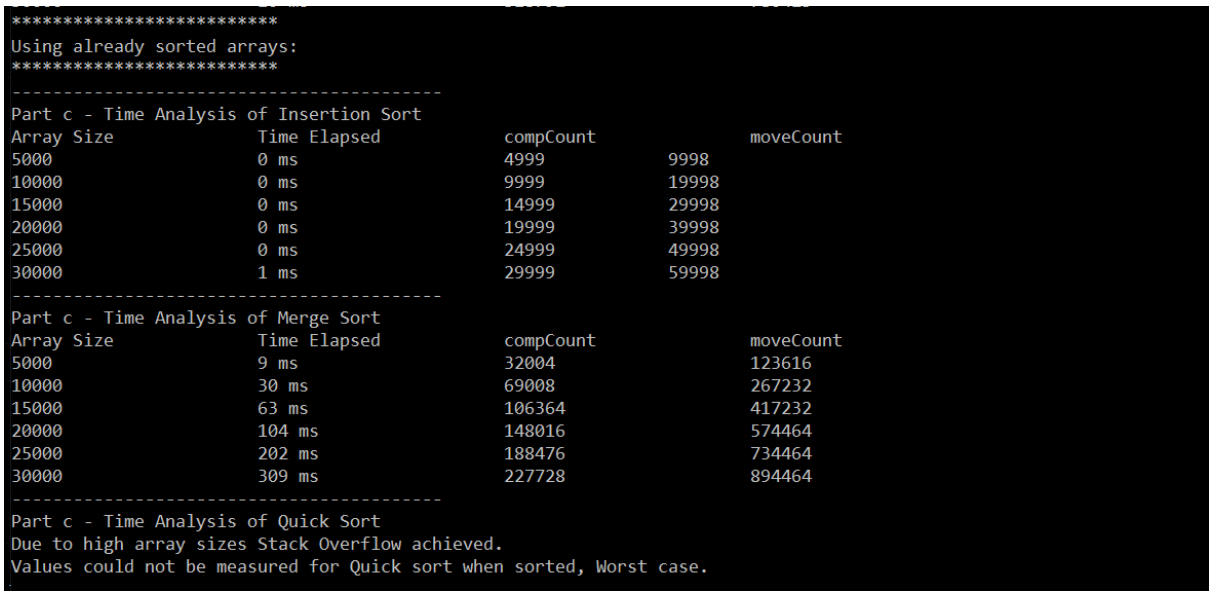
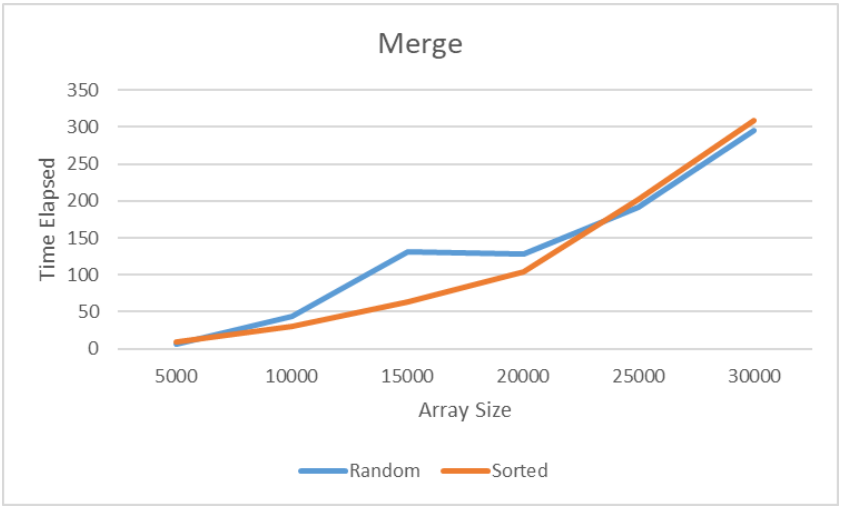
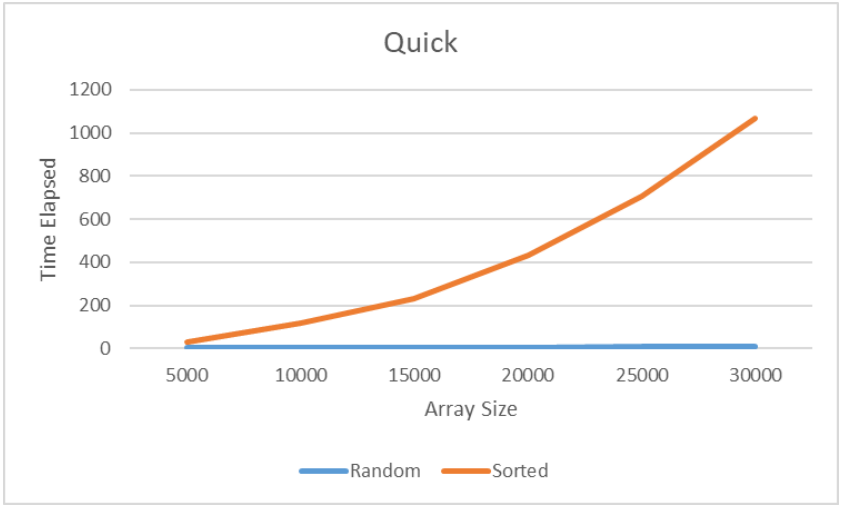


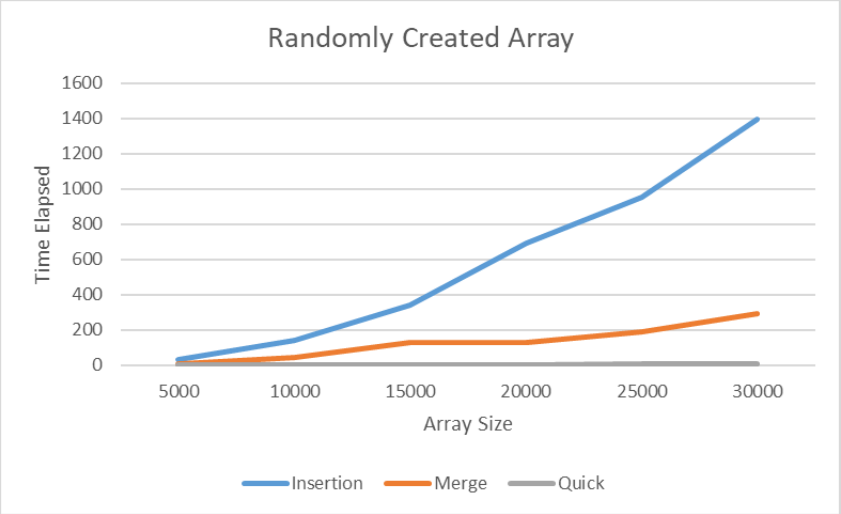
Figure3: Screenshot of Sorted Created Performance Analysis

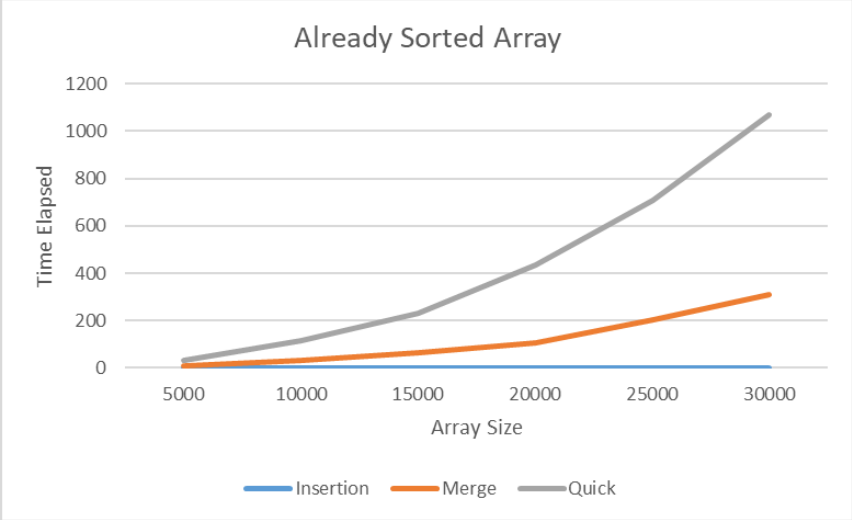
Plots:


Insertion
1600
1400
1200
1000
800
600
400
200
0
Time Elapsed
5000
10000
15000
20000
25000
30000
Array Size
Random
Sorted


Merge
350
300
250
200
150
100
50
0
Time Elapsed
5000
10000
15000
20000
25000
30000
Array Size
Random
Sorted


Quick
1200
1000
800
600
400
200
0
Time Elapsed
5000
10000
15000
20000
25000
30000
Array Size
Random
Sorted





Discussion:

As it can be seen in the plots, when considering randomly created array it can be said that quick sort performs way better than the merge and insertion sort algorithms. Also, merge sort performed way better than the insertion sort algorithm. Theoretically, the results seems logical. Merge sort and Quick sort algorithms have O(nlogn) time complexities which are less than Insertion sort time complexity O(n^2). Although their time complexities are same, quick sort seemed more efficient than the merge sort in the experiment. This can be resulted due to merge sort copying the array inside its merge function.

When already sorted array is used, it can be said that quick sort is the worst. Theoretically, quick sort has its worst case when the array items are already sorted and it gives O(n) time complexity. In experiment, it was not possible to obtain no of comparisons and no of moves since it results in stack over flow. Thus, the values in graphs are obtained separately only for time elapse values apart from the given code. On the other hand, Insertion sort's efficiency is significantly increased when sorted array is used. This is because Insertion sort has its best case when the array elements are sorted, O(n). Regarding merge sort its worst, average and best case have all the same time complexity O(nlogn). Thus, measured time elapse did not change significantly for merge sort in both experiments.

Part 3)

Since the array is nearly sorted choosing merge sort will give the best efficient solution. If we choose key as n/2 it will result time complexity similar to merge sorts best case. Key should be half of the size when entering each recursive function and the target should be as close as possible to key. Thus, k should be as close to make the target as the merge sorts sublist middle. If k is 0 it will result in best.