

091M4041H - Assignment 1

Algorithm Design and Analysis

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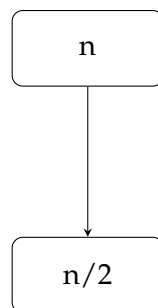
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1 determine the median of $2n$ values(1)

1.1 Algorithm Description

The problem is to find the median of two databases with n elements. We can separately obtain the median of database a and the median of database b . If the median of database a is less than the median of database b , then the overall median may only appear in the right interval of a and the left interval of b . If the median of a is greater than or equal to the median of b , then the overall median may only appear in the left interval of a and the right interval of b . So that the range of seeking will continue to shrink.

1.2 subproblem reduction graph



```

1: function FINDMEDIANUM(databaseA, beginA, databaseB, beginB, n)
2:   if n == 1 then
3:     return min(A[beginA], B[beginB])
4:   end if
5:   mediumA = A[beginA + n/2 - 1]
6:   mediumB = B[beginB + n/2 - 1]
7:   if mediumA < mediumB then
8:     FINDMEDIANUM(A, beginA + n/2, B, beginB, n/2)
9:   else
10:    FINDMEDIANUM(A, beginA, B, beginB + n/2, n/2)
11:   end if
12: end function

```

1.3 Prove the correctness

Set up two databases are A, B, each containing n elements. If we remove k numbers smaller than the median and k numbers bigger than the median, the median of the resulting sub-array is same as the median of the origin array. Then

$$Medium(A, B) = f(n)_{(A_{0-n}, B_{0-n})} = \begin{cases} f(n/2)_{(A_{n/2-n}, B_{0-n/2})} & A_{n/2} < B_{n/2} \\ f(n/2)_{(A_{n/2-n}, B_{0-n/2})} & A_{n/2} \geq B_{n/2} \end{cases}$$

1.4 time complexity

$$T(n) = T(n/2) + c = O(\log n)$$

2 Find the k^{th} largest element in an unsorted array(2)

2.1 Algorithm Description

Randomly choose element A_i , Comparing the other elements with A_i , if bigger than A_i put into set S_+ , if smaller put into set S_- . If $|S_-| = k - 1$, the result is A_i . If $|S_-| > k - 1$, the problem can be simplified as find the k^{th} largest element in S_- . If $|S_-| < k - 1$, the problem can be simplified as find the $(k - |S_-| + 1)^{th}$ largest element in S_+ .

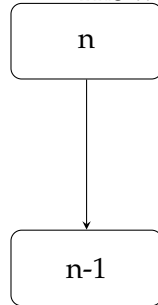
2.2 pseudo-code

```
1: function FINDKTHELEMENT( $A, k$ )
2:   Randomly choose  $A_i$ 
3:   for  $j = 0 \rightarrow n - 1$  do
4:     put  $A[j]$  in  $S_-$  if  $A[j] < A[i]$ 
5:     put  $A[j]$  in  $S_+$  if  $A[j] > A[i]$ 
6:   end for
7:   if  $|S_-| = k - 1$  then
8:     return  $A_i$ 
9:   else
10:    if  $|S_-| > k - 1$  then
11:      FINDKTHELEMENT( $S_-, k$ )
12:    end if
13:    FINDKTHELEMENT( $S_+, k - |S_-| + 1$ )
14:  end if
15: end function
```

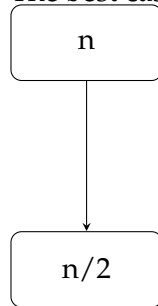
2 FIND THE K^{TH} LARGEST ELEMENT IN AN UNSORTED ARRAY(2)4

2.3 subproblem reduction graph

The worst case:



The best case:



2.4 Prove the correctness

$$f(k)_A = \begin{cases} A_i & |S - | + 1 = k \\ f(k)_{S-} & |S - | + 1 > k \\ f(k - |S - | + 1)_{S+} & |S - | + 1 < k \end{cases}$$

2.5 time complexity

The worst case: A_i is the biggest/smallest element of the array, $T(n) = T(n-1) + cn = O(n^2)$

The best case: A_i is the median element of the array, $T(n) = T(n/2) + cn = O(n)$

3 Divide convex polygon into traingles(5)

3.1 Algorithm Description

Numbering the vertices of the convex polygon from 1 to n , choose the vertex 1, n and any one of vertices 2 to $n-1$. The 3 vertices choosed can divide the polygon into a polygon of i edges, a traingle and a polygon of $n-i+1$ edges. The total ways to divide the origin polygon equal to the product of the ways to divide the polygon of i edges and the ways to divide the polygon of $n-i+1$ edges. To avoid repeat caculation, we can store the ways to divide a polygon of k edges into the k^{th} element of array `Walked[]`.

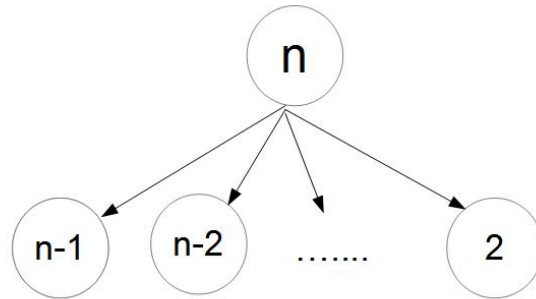
3.2 pseudo-code

```

1: int Walked[N] = 0
2: function DIVIDETRAINGLE( $n$ )
3:   if  $n \leq 3$  then
4:     return 1
5:   end if
6:   if  $Walked[n] \neq 0$  then
7:     return  $Walked[n]$ 
8:   end if
9:    $int\ count = 0$ 
10:  for  $i = 2 \rightarrow n - 1$  do
11:     $count += DivideTraingle(i) * DivideTraingle(n - i + 1)$ 
12:  end for
13:   $Walked[n] = count$ 
14: end function

```

3.3 subproblem reduction graph



3.4 Prove the correctness

$$f(n) = \begin{cases} 1 & n \leq 3 \\ \sum_{i=1}^N f(i) * f(n+1-i) & n > 3 \end{cases}$$

3.5 time complexity

$$T(n) = T(n-1) + T(n-2) + T(n-3) + \dots + T(2) = O(n-2) + O(n-3) + \dots + O(1) = O(n^2)$$

4 Sort-and-Count algorithm(8)

4.1 8-1 Source code

```

1 #include <iostream>
2 #include <fstream>
3 #include <string>
4 #include <time.h>

```

```
5
6 using namespace std;
7 const int N = 100001;
8
9 int num[N]; //store the elements need to be sorted
10 int tmp[N]; //temp array after merge
11 long long ans; //number of inversions
12
13 void merge_sort(int l, int m, int r)
14 {
15     int i = l;
16     int j = m + 1;
17     int k = l;
18     while(i <= m && j <= r)
19     {
20         //i < j, a[i] > a[j] -> inversion
21         if(num[i] > num[j])
22         {
23             tmp[k++] = num[j++];
24             //each element after a[i] -> inversions
25             ans += m - i + 1;
26         }
27         else
28         {
29             tmp[k++] = num[i++];
30         }
31     }
32
33     while(i <= m) tmp[k++] = num[i++];
34     while(j <= r) tmp[k++] = num[j++];
```

```
35
36     for(int i=1;i<=r;i++)
37         num[i] = tmp[i];
38     }
39
40 void sort_and_count(int l,int r)
41 {
42     if(l < r)
43     {
44         int m = (l + r) / 2;
45         sort_and_count(l,m);
46         sort_and_count(m+1,r);
47         merge_sort(l,m,r);
48     }
49 }
50
51 int main()
52 {
53     clock_t start,end;
54     //Read the number from the file line by line
55     ifstream infile;
56     infile.open("8.txt");
57     if(!infile) cout<<"error"<<endl;
58     int t1;
59
60     int len = 0;
61     cout<<"store in the array"<<endl;
62     while(infile>>t1)
63     {
64         num[len] = t1;
```



```

65         len++;
66     }
67
68     /*print the array*/
69     int i;
70     for (i=0;i<len;i++)
71         cout<<num[i]<<endl;
72     cout<<endl;
73
74     /*sort the array and count inversions*/
75     ans = 0;
76     start = clock();
77     sort_and_count(0,len-1);
78     end = clock();
79
80     cout << "The number of inversions:" << ans;
81     cout << "time to merge sort and count:" \
82     << (double)(end-start)/CLK_TCK <<endl;
83     return 0;
84 }

```

4.2 8-2 Souce code

It's possible to use the Quick-Sort idea.

source code:

```

1  #include <iostream>
2  #include <fstream>
3  #include <string>
4  #include <time.h>
5

```

```
6 using namespace std;
7
8 const int N = 100001;
9 long long ans; //number of inversions
10
11 void quick_sort_and_count(int *a, int l, int r)
12 {
13     int left_num = 0;
14     int right_num = 0;
15
16     /*S- set, store the elements
17     smaller than pivot*/
18     int *left = (int *)calloc(r-l, sizeof(int));
19     /*S+ set, store the elements
20     bigger than pivot*/
21     int *right = (int *)calloc(r-l, sizeof(int));
22
23     int pivot;
24     /*put smaller element into s-,
25     add the number of s+ now to
26     the inversion number.
27     after a round, add the number
28     of s- to the inversion number*/
29     if(l < r)
30     {
31         pivot = a[l];
32         int i;
33         for(i = l + 1; i <= r; i++)
34         {
35             if(a[i] < pivot)
```

```
36         {
37             left[left_num++] = a[i];
38             ans += right_num;
39         }
40         else
41         {
42             right[right_num++] = a[i];
43         }
44
45     }
46     ans += left_num;
47
48     quick_sort_and_count(left, 0, left_num - 1);
49     quick_sort_and_count(right, 0, right_num - 1);
50 }
51 }
52
53
54 int main()
55 {
56     clock_t start, end;
57     //Read the number from the file line by line
58     ifstream infile;
59     infile.open("8.txt");
60     if (!infile) cout << "error" << endl;
61     int t1;
62
63     int len = 0;
64     cout << "store in the array" << endl;
65     while (infile >> t1)
```

```
66     {
67         num[len] = t1;
68         len++;
69     }
70
71     /*print the array*/
72     int i;
73     for(i=0;i<len;i++)
74         cout<<num[i]<<endl;
75     cout<<endl;
76
77     /*sort the array and count inversions*/
78     ans = 0;
79     start = clock();
80     quick_sort_and_count(0, len - 1);
81     end = clock();
82
83     cout << "The number of inversions:" << ans;
84     cout << "time to quick sort and count:" \
85     << (double)(end-start)/CLK_TCK << endl;
86     return 0;
87 }
```

4.3 8-2 comparing

result for merge sort and count:

C:\Users\w\Desktop\算法作业1\8-1\bin\Debug\8-1.exe

```
存入数组
文件中数字逆序对数为: 2500572073
time to merge sort and count : 0.047

Process returned 0 (0x0)   execution time : 2.582 s
Press any key to continue.
```

result for quick sort and count:

C:\Users\w\Desktop\8_2\bin\Debug\8_2.exe

```
存入数组
文件中数字逆序对数为: 2500572073
time to quick sort and count : 2.265

Process returned 0 (0x0)   execution time : 2.828 s
Press any key to continue.
```

5 the closest pair problem(9)

5.1 Source code

```
1 #include <iostream>
2 #include <math.h>
3 #include <algorithm>
4
5 using namespace std;
6 typedef struct point//point in the plane
7 {
8     double x;
```

```

9  double y;
10
11 //sort the points by x coordinate
12 friend bool operator < \
13 (const point &a, const point &b){
14     if(a.x == b.x)
15         return a.y < b.y;
16     return a.x < b.x;
17 }
18 }point;
19
20 const int N = 100;
21 point points[N];
22 point a,b;//the closest point pair
23
24 //use quick sort to sort the points by y
25 void y_QuickSort(point y_sort_points[],\
26     int l,int r)
27 {
28     if (l < r)
29     {
30         int i = l, j = r;
31         point x = y_sort_points[l];
32         while (i < j)
33         {
34             while(i < j &&\
35                 y_sort_points[j].y>= x.y)
36                 j--;
37             if(i < j)
38                 y_sort_points[i++] = y_sort_points[j];

```

```
39         while(i < j && \
40             y_sort_points[i].y < x.y)
41             i++;
42         if(i < j)
43             y_sort_points[j--] = y_sort_points[i];
44     }
45     y_sort_points[i] = x;
46     y_QuickSort(y_sort_points, l, i - 1);
47     y_QuickSort(y_sort_points, i + 1, r);
48 }
49 }
50
51 double Distance(point a, point b)
52 {
53     return sqrt((a.x - b.x)*(a.x - b.x) \
54         + (a.y - b.y) * (a.y - b.y));
55 }
56
57 double closest(int l, int r)
58 {
59     if(l == r)
60     {
61         a.x = points[l].x;
62         a.y = points[l].y;
63
64         b.x = points[r].x;
65         b.y = points[r].y;
66         return 0;
67     }
68     else if(l == r - 1)
```

```
69     {
70         a.x = points[l].x;
71         a.y = points[l].y;
72
73         b.x = points[r].x;
74         b.y = points[r].y;
75
76         return Distance(points[l], points[r]);
77     }
78
79     int m = (l + r)/2;
80     /*caculate the closest pair
81     in left and right*/
82     double res = min(closest(l,m), closest(m,r));
83     double tmp;
84     int i, j;
85     int node_inside_num = 0;
86     point node_inside[N];
87
88     /*store the points in 2d width strip
89     in temp array*/
90     for(i = m - 1; i >= l && \
91         points[m].x - points[i].x < res; i--){
92         node_inside[node_inside_num++] = points[i];
93     }
94
95     for(i = m + 1; i <= r && \
96         points[m].x - points[i].x < res; i++){
97         node_inside[node_inside_num++] = points[i];
98     }
```



```
99
100 //sort the points in 2d width strip by y
101 y_QuickSort(node_inside,0,node_inside_num-1);
102
103 /*for each points in 2d width
104 strip sorted by y,caculate the
105 diatance with 7 points after it ,
106 if closer ,update the result*/
107 for(i = 0;i < node_inside_num;i++)
108 {
109     int end_inside;
110     if(node_inside_num - i - 1 < 8)
111         end_inside = node_inside_num;
112     else end_inside = i + 8;
113     for(j = i + 1;j < end_inside;j++)
114     {
115         tmp = Distance(node_inside[i], \
116 node_inside[j]);
117         if(tmp < res)
118         {
119             res = tmp;
120             a.x = node_inside[i].x;
121             a.y = node_inside[i].y;
122             b.x = node_inside[j].x;
123             b.y = node_inside[j].y;
124         }
125     }
126 }
127 return res;
128 }
```

```
129 int main()
130 {
131     int n;
132     //read the count of points in the plain
133     cin>>n;
134     int i;
135     for(i=0; i<n ; i++){
136         //read the coordinate of the n points
137         cin>>points[i].x>>points[i].y;
138     }
139
140     //sort the points by x coordinate
141     sort(points , points+n);
142     double result = closest(0,n - 1);
143     cout<<"The closest point pair is : (" ;
144     cout<<a.x<<" , "<<a.y<<" ) _ ("<<b.x<<" , "<<b.y ;
145     cout<<" ) _ the closest distance _ is : "<<result ;
146
147     return 0;
148 }
```

5.2 result

test case for 5 points:

```
5
1 5
2 3
4 9
10 6
18 3
node_inside[0]:18,3
node_inside[1]:4,9
node_inside[0]:1,5
node_inside[1]:4,9
node_inside[0]:2,3
node_inside[1]:18,3
node_inside[2]:10,6
最近点对为: (1,5) (2,3) 距离为: 2.23607
```

test case for 20 points:

```
20
5 8
4 6
9 2
15 19
7 16
3 52
9 6
15 23
14 9
18 7
22 6
25 3
27 14
12 5
2 5
7 4
3 8
15 6
24 37
14 14
```

```
node_inside[2]:3,8
node_inside[3]:5,8
node_inside[0]:9,2
node_inside[1]:7,4
node_inside[2]:12,5
node_inside[3]:4,6
node_inside[4]:9,6
node_inside[5]:7,16
node_inside[0]:25,3
node_inside[1]:22,6
node_inside[2]:15,6
node_inside[3]:18,7
node_inside[4]:14,9
node_inside[5]:27,14
node_inside[6]:14,14
node_inside[7]:15,19
node_inside[8]:15,23
node_inside[9]:24,87
最近点对为: (3,8) (5,8) 距离为: 2
```