

# PROGRAMMING LANGUAGES CONCEPTS AND IMPLEMENTATION

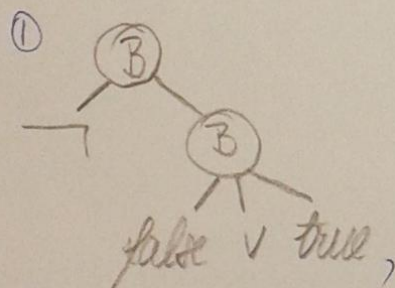
(Mandatory Exercise, Week 6, Student: Irina Alina Gabriela Luca)

## THEORETICAL EXERCISE

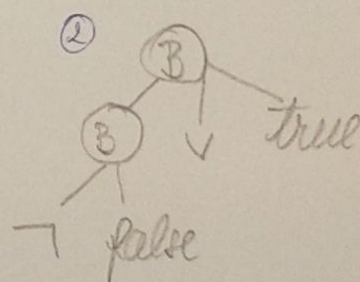
Our ambiguous grammar is  $\Rightarrow$

- $B = \text{true}$
- $\mid \text{false}$
- $\mid B \vee B$
- $\mid B \wedge B$
- $\mid \neg B$
- $\mid (B)$

@ The grammar is ambiguous because for an expression like  $\neg \text{false} \vee \text{true}$ , for instance, we can create two parse trees, as it follows:



interpreted as  $\neg (\text{false} \vee \text{true})$



interpreted as  $(\neg \text{false}) \vee \text{true}$

@ In order to construct an unambiguous grammar from the above presented one, we have to take into account two aspects  $\Rightarrow$

- $\Rightarrow$  { - operators' precedence (see step 1)
- encode associativity (see step 2)

STEP 1  $\Rightarrow$

- $B = B \vee B \mid \text{true}$
- $S = S \wedge S \mid T$
- $T = \neg T \mid U$
- $U = (B) \mid \text{true} \mid \text{false}$

made sure operators' hierarchy is applied

STEP 2  $\Rightarrow$

- $B = S \vee B \mid S$
- $S = S \wedge T \mid T$
- $T = \neg T \mid U$
- $U = (B) \mid \text{true} \mid \text{false}$

encoded associativity, therefore obtaining an unambiguous form

© Considering the unambiguous grammar from ©, we will construct the parse tree for the following expression:

$\text{false} \vee \neg \text{true} \wedge (\text{false} \vee \text{true}) \vee \text{false}$

