MATH676 — Tropical Geometry

Based on the lectures by Renzo Cavalieri

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Please note that these notes were not provided or endorsed by the lecturer and have been significantly altered after the class. They may not accurately reflect the content covered in class and any errors are solely my responsibility.

This is a topics course on this stuff

Requirements

Knowledge on stuff

Contents

Chapter 1

Combinatorial Shadow of Algebraic Geometry

1.1 Day 1 | 20230821

Think of an algorithm where the input is an algebraic variety and the output is a combinatorial object, a piecewise linear object.

Example 1.1.1. Consider as an input a line in the plane. Say $V(x + y - 1)$, then an output would be a tropical line. If we remain in the plane and consider a higher degree polynomial, say an elliptic curve, as an output we obtain a tropical cubic.

Leaving the plane behind and thinking of abstract nodal curves, we can think of a sphere attached to a torus which is attached to a genus 2 torus, then the corresponding object is what we call the dual graph.

Right now we do not know the specific algorithm, but we can observe that the outputs are *more simple* than the inputs. So the important question is:

> *What algebraic information does the simplified object remember? How do we extract the information the object remembers? And once we know how to work with this objects, can we return to algebraic geometry from any kind of these objects?*

Observe that the number of ends which go to infinity corresponds with the degree.

1.2 Interim

There are many ways to obtain the previous combinatorial pictures:

Algebraic Geometry on $\mathbb{T}$

We call $(\mathbb{T}, \oplus, \odot)$ the tropical semifield. A polynomial $p(\underline{x})$ gives rise to its *tropicalization*, a piecewise linear function

$$\mathrm{Trop}(p) : \mathbb{R}^n \to \mathbb{R}.$$

The tropical hypersurface $V(\mathrm{Trop}(p))$ is the codimension 1 locus in $\mathbb{R}^n$ where the function is non-linear (corner locus).

Valued fields

Definition 1.2.1. The field of Puiseux series over $\mathbb{C}$ is $\mathbb{C}\{\{t\}\}$ where the elements are of the form

$$\alpha(t) = \sum_{i=k_0}^{\infty} a_i t^{i/n}.$$

The lower bound k_0 could be negative and the exponents, as in the case of Laurent series, are rational with bounded denominators.

The valuation of $\alpha(t)$ is $\mathrm{val}(\alpha(t)) = \frac{k_0}{n}$, it is the smallest exponent for the natural order of the rational numbers and the coefficient a_{k_0} is called the valuation coefficient.

If K is the field of Puiseux series, then for a polynomial $p \in K[\underline{x}]$, and here

$$\begin{array}{ccc}
\text{Val: } K^n & \longrightarrow & \mathbb{R}^n \\
\cup\vert & & \cup\vert \\
V(p) & \longmapsto & \overline{\mathrm{Val}(V(p))} \\
\Vert & & \Vert \\
\{\mathbf{x} : p(\mathbf{x}) = 0\} & & \mathrm{Trop}(V(p))
\end{array}$$

$\mathrm{Trop}(V(p))$ is the tropical hypersurface for p.

Amoebas

Suppose now $p \in \mathbb{C}[\underline{x}]$ and $V(p) \subseteq \mathbb{C}^n$. We may consider the function

$$\log_t : \mathbb{C}^n \to \mathbb{R}^n, \quad (z_1, \ldots, z_n) \mapsto (\log_t |z_1|, \ldots, \log_t |z_n|).$$
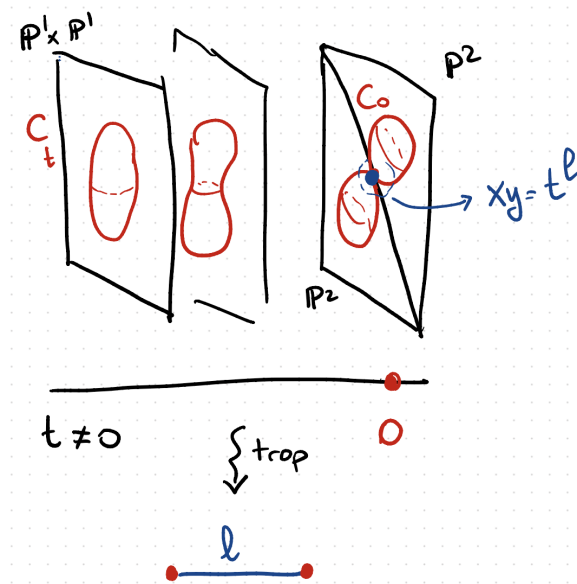
The amoeba is $\log_t(V(p)) \subseteq \mathbb{R}^n$.

Example 1.2.2. When $p = x + y - 1$ then we have a figure which I can't draw.

Observe that this is the tropical hypersurface associated to p. In other words we have that the tropical hypersurface is $\lim_{t \to \infty} \log_t(V(p))$.

Degenerations

Consider a family of of what, what is this family of?! Stuff? Curve in P1xP1 which eventually becomes P2?

It is too early to understand this point of view. We will set everything up to get to it.

In general, the big idea will be to explore and understand these perspectives in the case of plane curves. We want to show how they are equivalent and then recover classical algebraic geometry results in terms of tropical geometry.

Tropical Arithmetics

Minimizing Tolls

Consider a set of cities connected by a network of toll-ways: If we only care about minimizing toll expenses, what would be the cheapest way to go from one given city to another? Let us record the information as an incidence matrix.

$$M = \begin{pmatrix} 0 & \infty & 2 \\ x & 0 & y \\ \infty & 1 & 0 \end{pmatrix}$$

In this matrix, the rows determine the outbound city, while the columns are the destination. Each entry records the cost of a toll and tolls are considered to be infinite
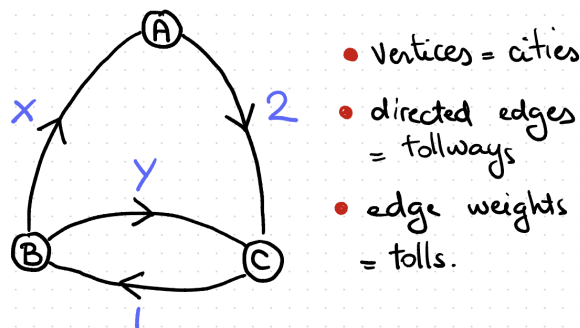
when the road does not exist. We can also think of M as recording the cheapest toll to go from one city to another with at most one move.

But if we wanted to find the cheapest way from one city to another in **two** moves, we could use M^2 with standard matrix multiplication. However we must redefine our basic operations as follows:

$$+ = \min, \quad \cdot = +$$

$$
\begin{pmatrix} 0 & \infty & 2 \\ x & 0 & y \\ \infty & 1 & 0 \end{pmatrix}^2 =
\begin{pmatrix}
\min \begin{pmatrix} 0+0 \\ \infty+x \\ 2+\infty \end{pmatrix} & \min \begin{pmatrix} 0+\infty \\ \infty+0 \\ 2+1 \end{pmatrix} & \min \begin{pmatrix} 0+2 \\ \infty+y \\ 2+0 \end{pmatrix} \\
\min \begin{pmatrix} x+0 \\ 0+x \\ y+\infty \end{pmatrix} & \min \begin{pmatrix} x+\infty \\ 0+0 \\ y+1 \end{pmatrix} & \min \begin{pmatrix} x+2 \\ 0+y \\ y+0 \end{pmatrix} \\
\min \begin{pmatrix} \infty+0 \\ 1+x \\ 0+\infty \end{pmatrix} & \min \begin{pmatrix} \infty+\infty \\ 1+0 \\ 0+1 \end{pmatrix} & \min \begin{pmatrix} \infty+2 \\ 1+y \\ 0+0 \end{pmatrix}
\end{pmatrix}
$$

$$
= \begin{pmatrix}
0 & 3 & 2 \\
x & \min(0, y+1) & \min(x+2, y) \\
1+x & 1 & \min(0, 1+y)
\end{pmatrix}.
$$

Observe that $1 + y$ can be the minimum in the diagonal when we allow *negative tolls*.

Remark 1.2.3. If we disallow negative tolls, the products M^n eventually stabilize to a matrix whose entries record the cheapest way to get from one city to another in n steps.

This gives us an intuition that minimization problems correspond to linear algebra problems over $(\mathbb{T}, +, \cdot)$ which is precisely $(\mathbb{R} \cup \{\infty\}, \min, +)$.

Forgetting phases

Recall that any complex number can be written as $z = re^{i\theta}$ where $r \geqslant 0$. Consider the map $T_t : \mathbb{C} \to \{-\infty\} \cup \mathbb{R}, \quad z \mapsto \log_t(r)$. This map is surjective, and this we can see by
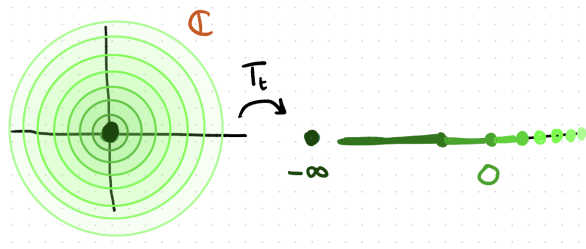
checking it is right-invertible. Observe that:

$$\begin{cases} T_t^{-1}(x) = \{t^x e^{i\theta}\} \subseteq \mathbb{C}, & \text{for} \quad x \in \mathbb{R}, \\ T_t^{-1}(-\infty) = 0. \end{cases}$$

With this in hand, we wish to define an exotic addition and multiplication on $\{-\infty\} \cup \mathbb{R}$ using T_t. We will dequantize!

We begin with **hyper-addition**, the output will be a subset of $\{-\infty\}$ so it's not a binary operation by itself.

$$x \diamondsuit_t y := T(T_t^{-1}(x) + T_t^{-1}(y)) = \left[\log_t(|t^x - t^y|), \log_t(t^x + t^y)\right].$$

This is an interval in $\{-\infty\} \cup \mathbb{R}$, in order to make $\diamondsuit_t$ into an operation we take a limit:

$$
\begin{array}{ccc}
x \diamondsuit_t y & \xrightarrow{\lim_{t\to\infty}} & x \diamondsuit y = \lim_{t\to\infty} x \diamondsuit_t y \\
\Big\downarrow{\scriptstyle \max} & & \Big\downarrow{\scriptstyle \max} \\
x +_t y & \xrightarrow{\lim_{t\to\infty}} & x + y = \max(x, y)
\end{array}
$$

1.3   Day 2 | 20230823

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Bibliography

[1] D. Maclagan and B. Sturmfels. *Introduction to Tropical Geometry*. Graduate Studies in Mathematics. American Mathematical Society, 2021.