

# Formal Languages week3

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## 1 Chap 2 (pp58-60): 1,5,6,10,12,13,14,23

1 Give a recursive definition of the length of a string over  $\Sigma$ . Use the operation from the definition of a string.

i) Basis:  $\lambda \in \Sigma$

ii) Recursive step:  $u$  and  $v$  are any  $w \in \Sigma^*$  or concatenation of any number of letters so  $w^n$  is the length..?

I have no idea what I am doing...

5 Let  $L$  be the set of strings over  $\{a, b\}$  generated by the recursive definition

i) Basis:  $b \in L$

ii) Recursive step: if  $u$  is in  $L$  then  $ub \in L$ ,  $uab \in L$ , and  $uba \in L$ , and  $bua \in L$ .

(a) List the elements in the sets  $L_0$ ,  $L_1$ , and  $L_2$ .

$$L_0 = \{b\}$$

$$L_1 = \{ub, uab, uba, bua, b\}$$

$$L_2 = \{bb, ubb, bab, uabab, ubaba, bbuaa, uabb, ubab, buab, ubaab, buaab, buaba, b\}$$

(b) Is the string  $bbaaba$  in  $L$ ? If so, trace how it is produced. If not explain why not.

$uba$  or  $bua$

$buaba$  or  $buaba$

$bbuaaba$

No I can not create that string because I can not turn  $u$  into  $\lambda$ . There will always be one more  $b$  than  $a$ 's

(c) Is the string  $baaaaabb$  in  $L$ ? If so, trace how it is produced. If not explain why not.

$ub$

$uabb$

*buabb*

The same problem exists. There will always be ONE more *b* than *a*.

- 6 I am not writing the whole problem out but I am confused. It says that the rules must always be that there is at least ONE *b* and an even number of *a*'s before the first *b*. So why is *bab* used as an acceptable answer? I think this must be a typo and they meant *aab*.

i) Basis:  $b \in L$

ii) Recursive step: if  $u$  is in  $L$  then  $aa * u \in L$ ,  $aa * ub * \in L$ ,  $aa * b * u \in L$ .

iii) Closure: a string  $v$  is in  $L$  only if it can be obtained from the basis by a finite number of iterations of the recursive step.

10

12

- 13 Let  $L_1 = \{aaa\}^*$ ,  $L_2 = \{a, b\}\{a, b\}\{a, b\}\{a, b\}$ , and  $L_3 = L_2^*$ . Describe the strings that are in the languages  $L_2$ ,  $L_3$ , and  $L_1 \cap L_3$

I don't know what it means to describe but I hope this is close:

$L_2$  is the language of only the set of sets of *a* and *b* only.

$L_3$  is the same as above but each iteration will be the list of  $L_2 \times L_2$

$L_1 \cap L_3$  Will be all elements of  $L_1$  and  $L_3$  as since I don't see an overlap... :/

- 14 The set of strings over  $\{a, b, c\}$  in which all the *a*'s precede the *b*'s, which in turn precedes the *c*'s. It is possible that there are no *a*'s, *b*'s, or *c*'s.

$$\{a^n b^m c^r \mid n = m = r \geq 0\}$$

- 23 The set of strings over  $\{a, b, c\}$  that begin with *a*, contain exactly two *b*'s, and end with *cc*.

Basis:  $\{abbcc\}$

$$\{a\}\{a, b, c\} * \{bb\}\{a, b, c\}^* \{cc\}$$

Something like that....