

Formal Languages week4

Michelle Bergin

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1 Chap 3 (p 98-101): 9, 11, 37, One More

9 Construct a grammar over $\{a, b, c\}$ whose language is $\{a^n b^m c^i \mid 0 \leq n + m \leq i\}$

$$S \rightarrow ASCc|B|\lambda$$

$$A \rightarrow AaBCc|a$$

$$B \rightarrow Bbc|b$$

$$C \rightarrow Cc|c$$

11 Construct a grammar over $\{a, b\}$ whose language is $\{a^m b^i a^n \mid i = m + n\}$

$$S \rightarrow AC$$

$$A \rightarrow aAb|ab$$

$$C \rightarrow bCa|ba$$

37 Construct unambiguous grammars for the languages $L_1 = \{a^n b^n c^m \mid n, m > 0\}$ and

$L_2 = \{a^n b^m c^m \mid n, m > 0\}$. Construct a grammar G that generates $L_1 \cup L_2$. Prove that G is ambiguous. This is an example of an inherently ambiguous language. Explain, intuitively, why

every grammar generating $L_1 \cup L_2$ must be ambiguous.

$L_1 : -$

$$S \rightarrow AC$$

$$A \rightarrow aAb|ab$$

$$C \rightarrow Cc|c$$

$L_2 : -$

$$S \rightarrow AC$$

$$A \rightarrow aA|a$$

$$C \rightarrow bCc|bc$$

$L_1 \cup L_2 : -$

$$S \rightarrow ABC$$

$$A \rightarrow aA|ab$$

$$B \rightarrow bB|b$$

$$C \rightarrow Cc|bc$$

That probably isn't right but the idea is that There are enough rules where you can create the same string with different rule paths.

One More) Find a CFG over a,b that generates the language consisting of strings that have twice as many a's as b's and prove your grammar correct.

$$S \rightarrow ASaab$$

$$A \rightarrow aaAb$$