

TCP CONGESTION SIGNATURES

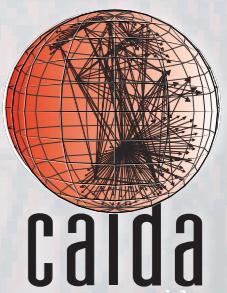
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kc Claffy (CAIDA/UCSD)
Mark Allman (ICSI)



INTERNATIONAL
COMPUTER SCIENCE
INSTITUTE



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UNIVERSITY



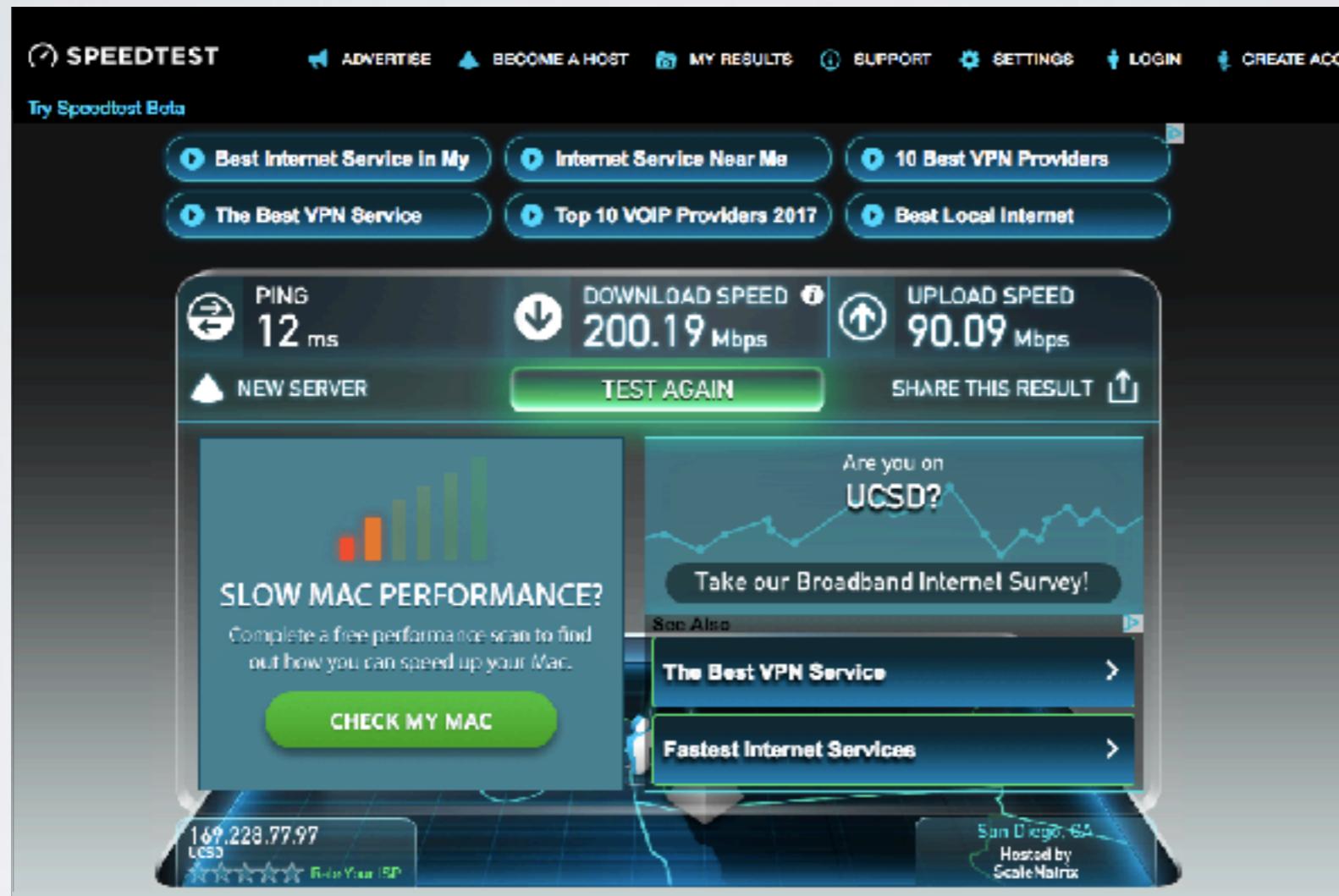
Typical Speed Tests Don't Tell Us Much



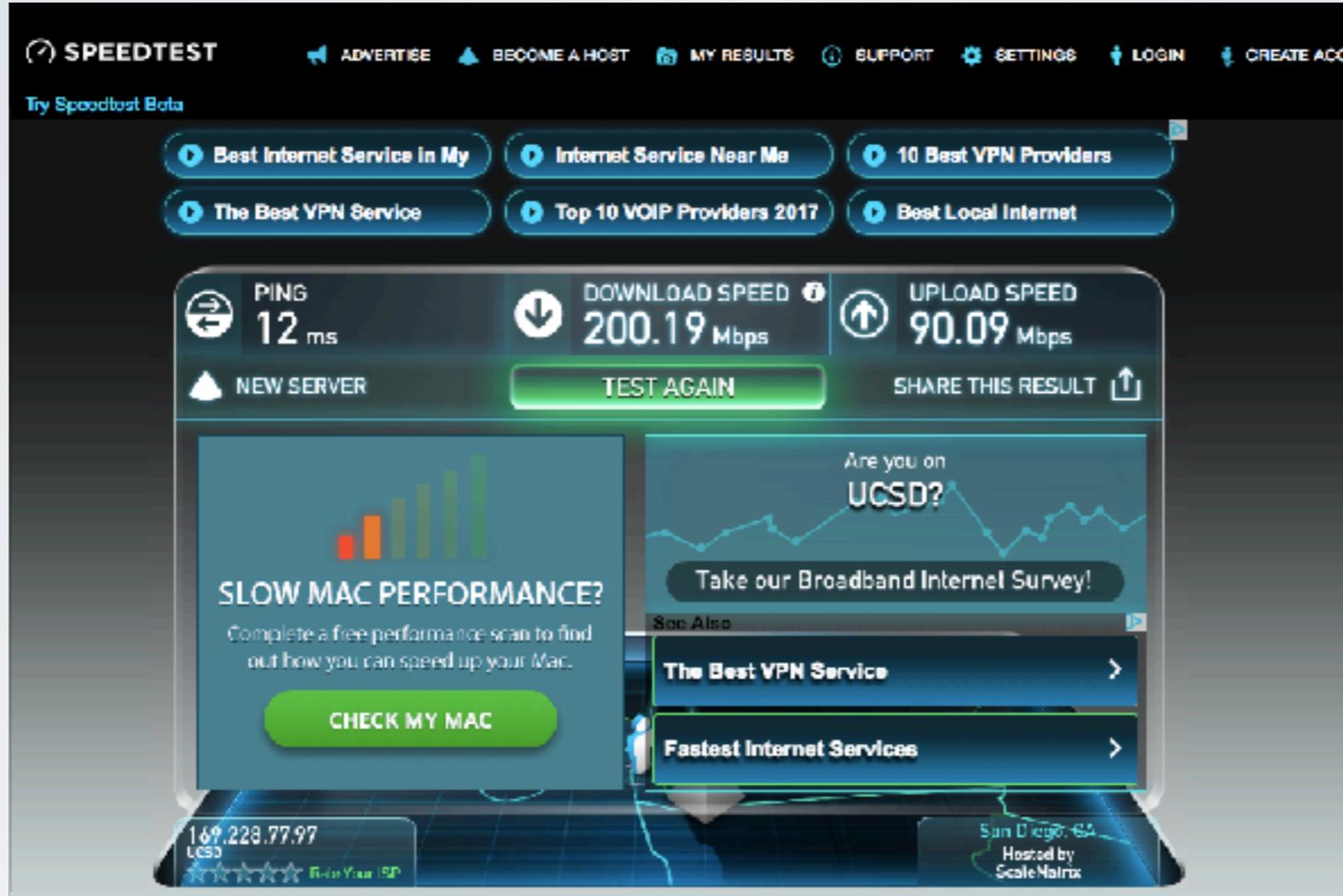
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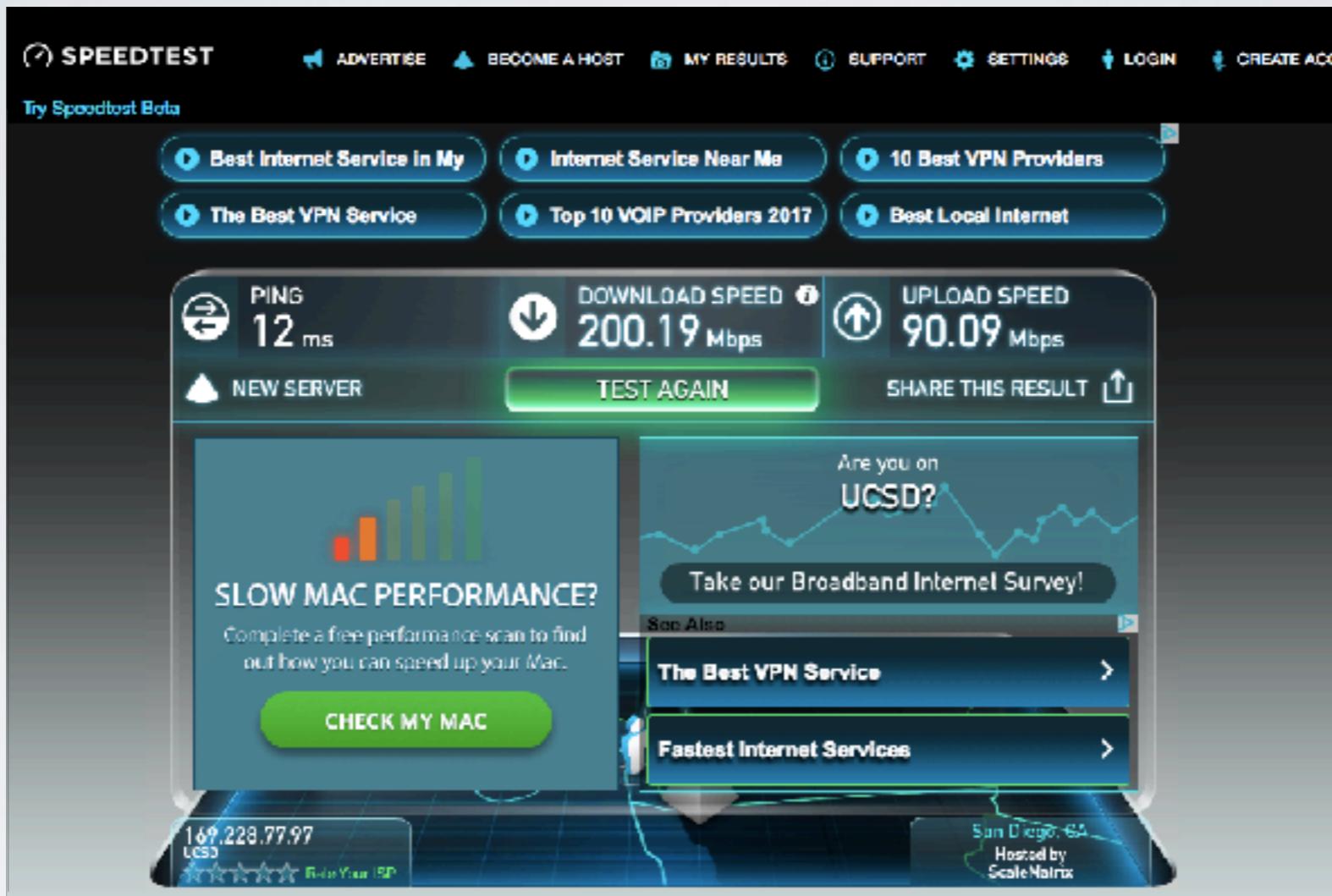


Typical Speed Tests Don't Tell Us Much



- Upload and download throughput measurements: no information beyond that

Typical Speed Tests Don't Tell Us Much



What type of congestion did the TCP flow experience?

Two Potential Sources of Congestion in the End-to-end Path

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- Self-induced congestion
 - Clear path, the flow is able to saturate the bottleneck link
 - eg: last-mile access link

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- External congestion
 - Flow starts on an already congested path
 - eg: congested interconnect

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Distinguishing the two cases has implications for users /
ISPs / regulators

Does Throughput Indicate Type of Congestion?

- Cannot distinguish using just throughput numbers
 - Access plan rates vary widely, and are typically not available to content / speed test providers
 - eg: Speed test reports 5 Mbps – is that the access link rate (DSL), or a congested path?

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We can use the dynamics of TCP's startup phase, i.e.,
Congestion Signatures

TCP's RTT Congestion Signatures

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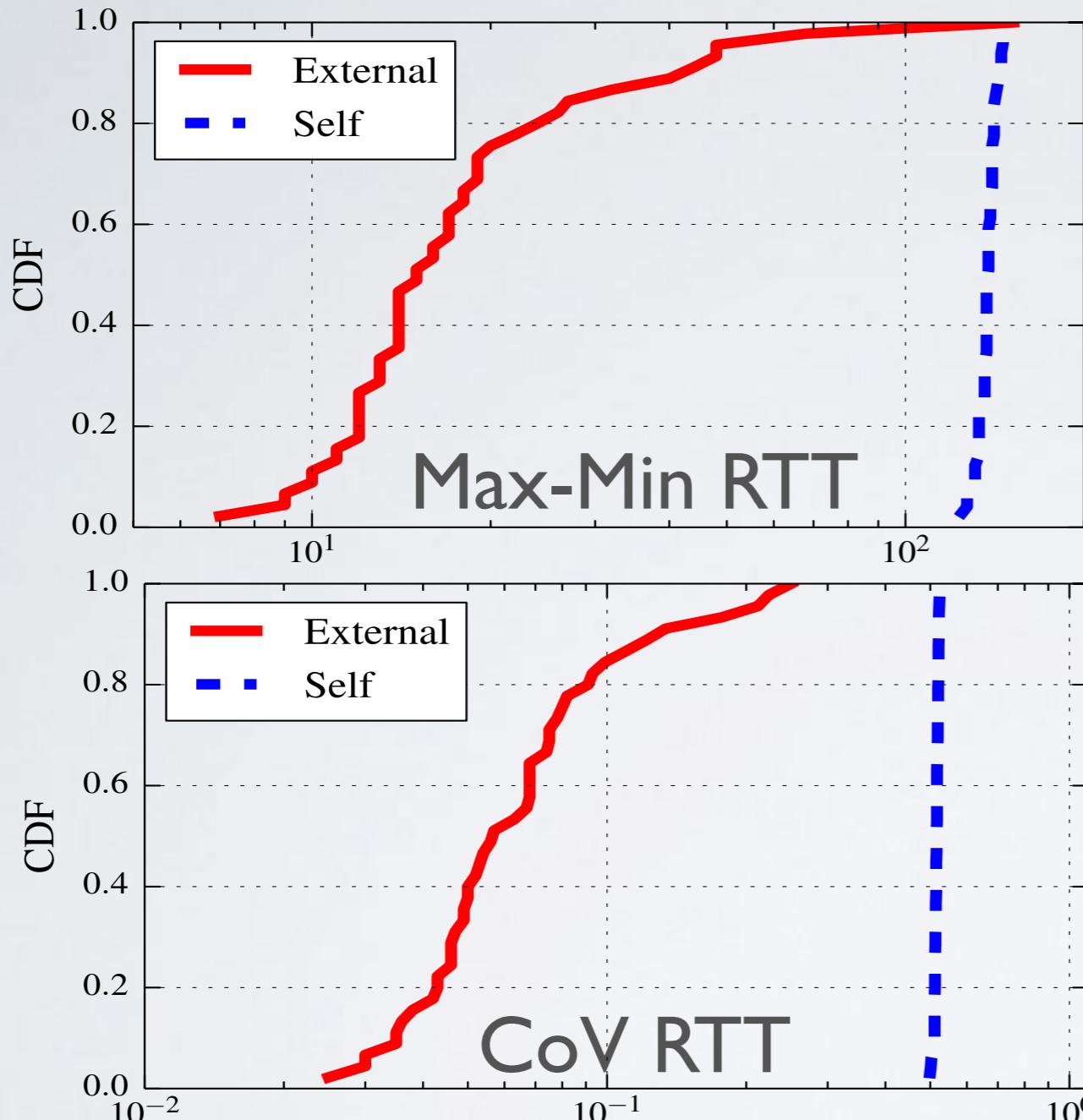
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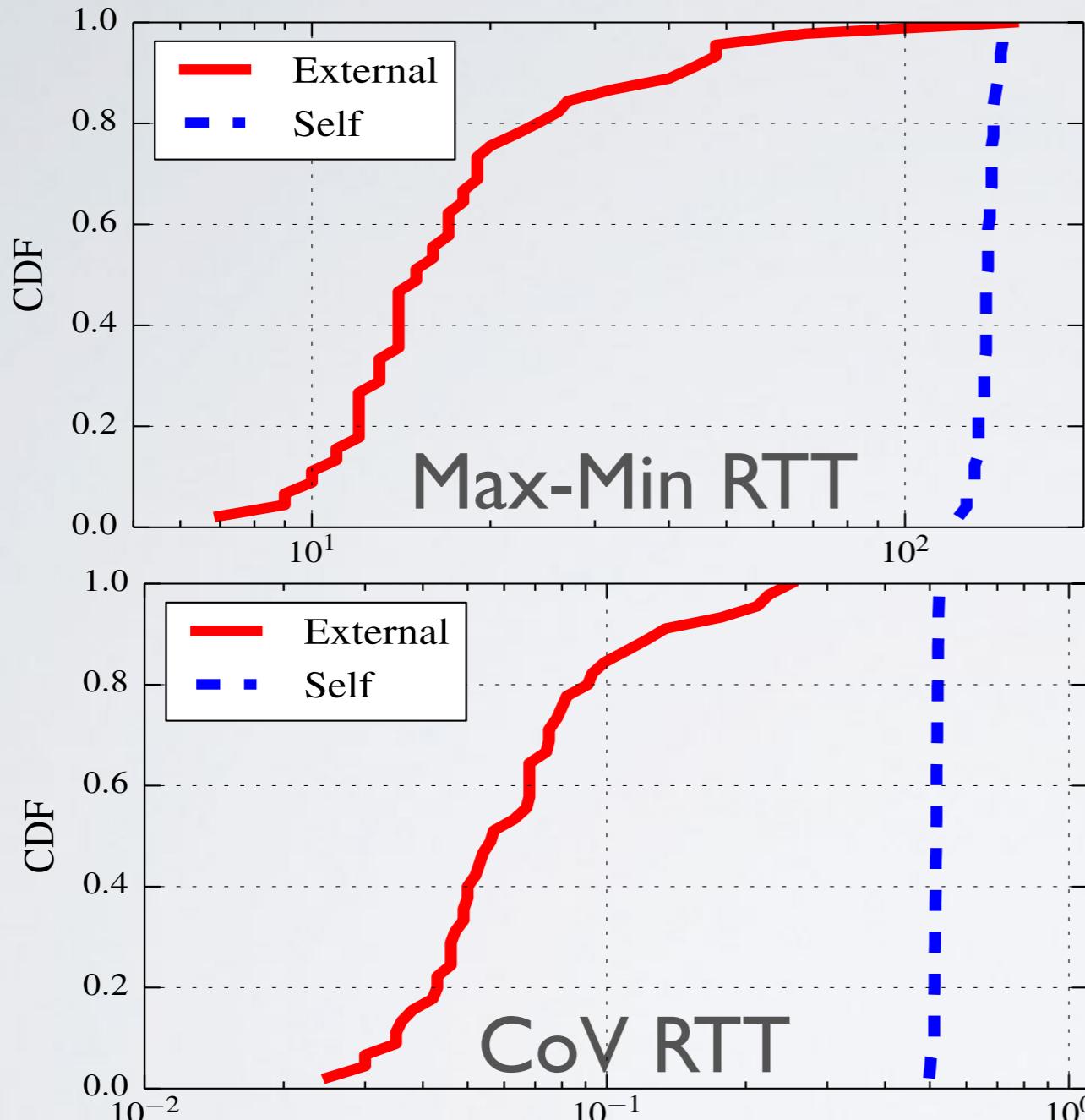
We can quantify this using Max-Min and CoV of RTT

Example Controlled Experiment



- 20 Mbps “access” link with 100 ms buffer
- 1 Gbps “interconnect” link with 50 ms buffer
- Self-induced congestion flows have higher values for both metrics and are clearly distinguishable

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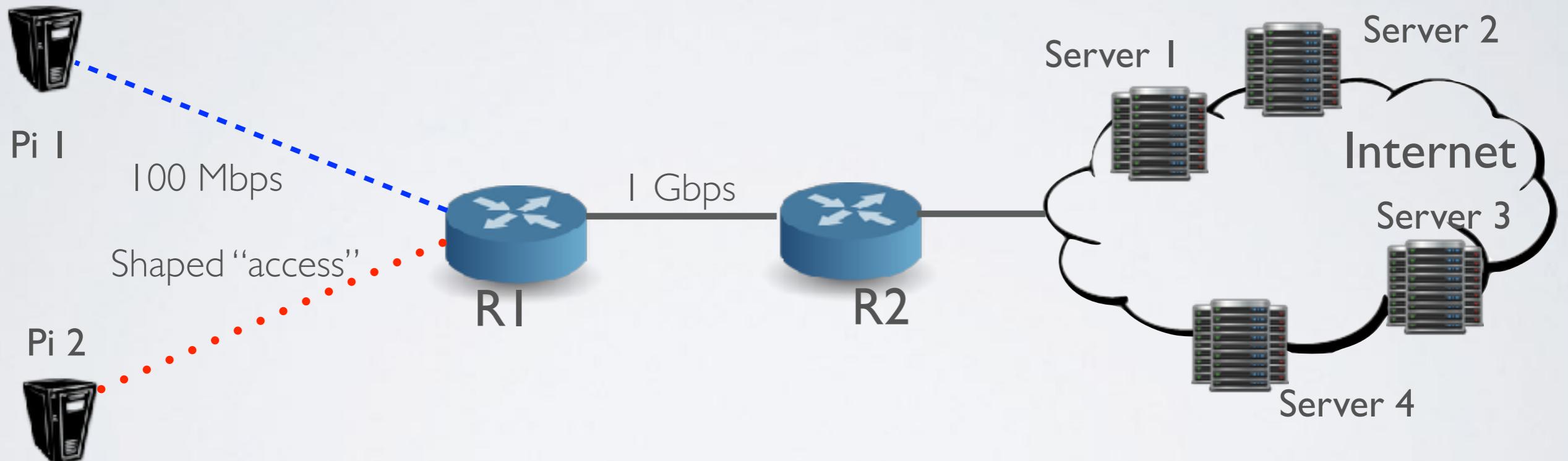
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The two types of congestion exhibit widely contrasting behaviors

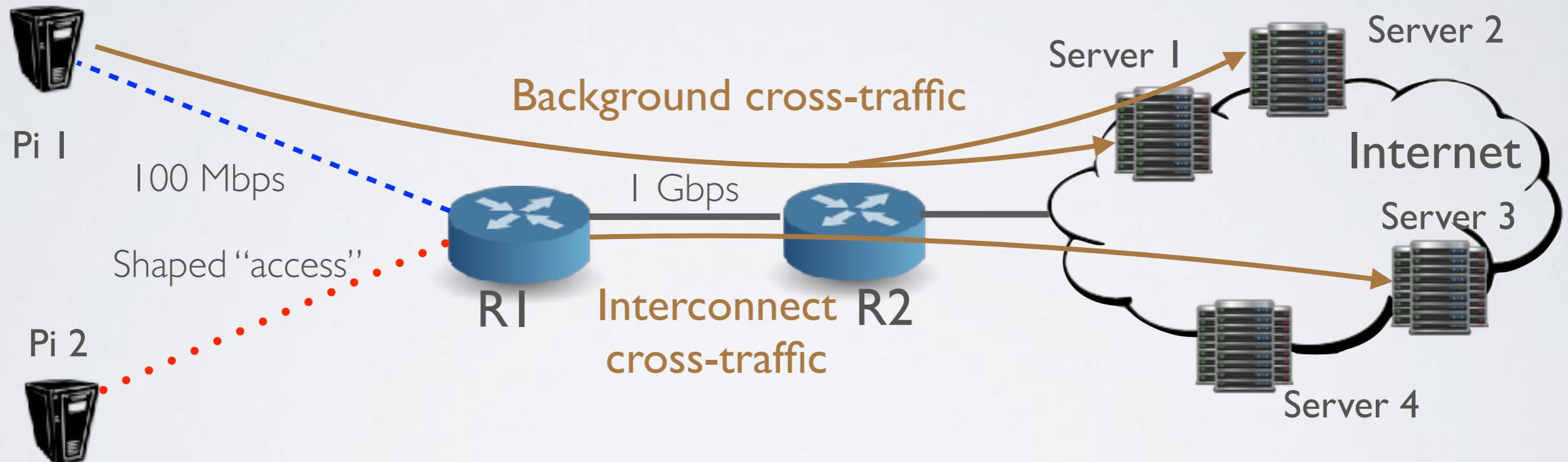
Model

- Max-min and CoV of RTT derived from RTT samples during slow start
- We feed the two metrics into a simple Decision Tree
 - We control the depth of the tree to a low value to minimize complexity
- We build the decision tree classifier using controlled experiments and apply it to real-world data

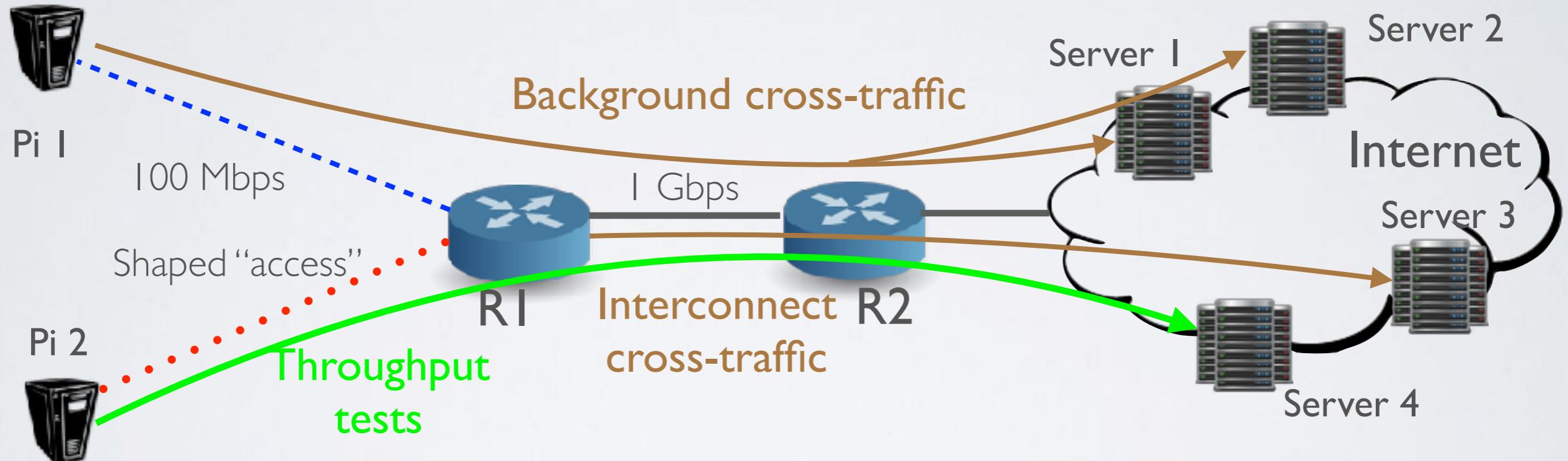
Validating the Method: Step I - Controlled Experiments



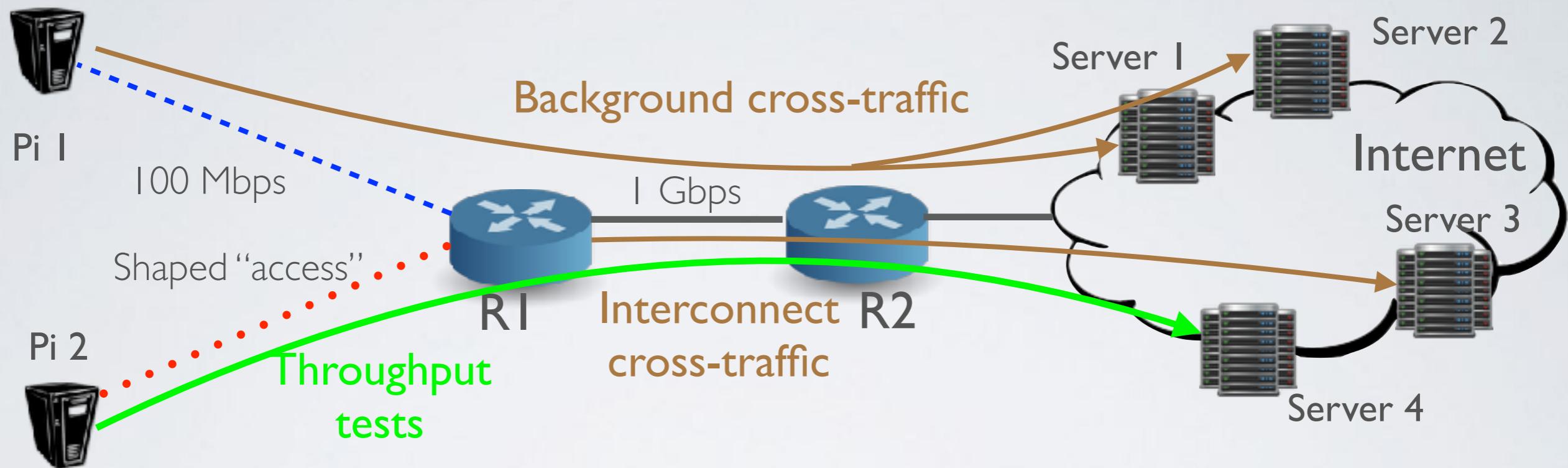
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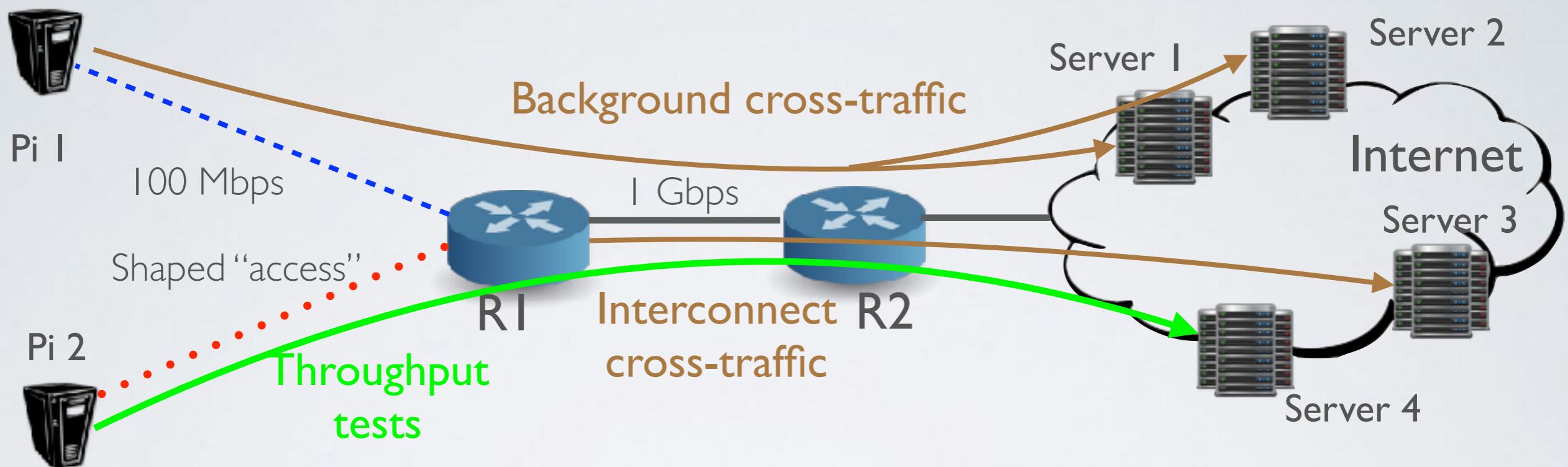


Validating the Method: Step I - Controlled Experiments



- Emulated “access” link + “core” link
 - Wide range of access link throughputs, buffer sizes, loss rates, cross-traffic (background and congestion-inducing)
 - Can accurately label flows in training data as “self” or “externally” congested

Validating the Method: Step I - Controlled Experiments



High accuracy: precision and recall > 90%
in most settings

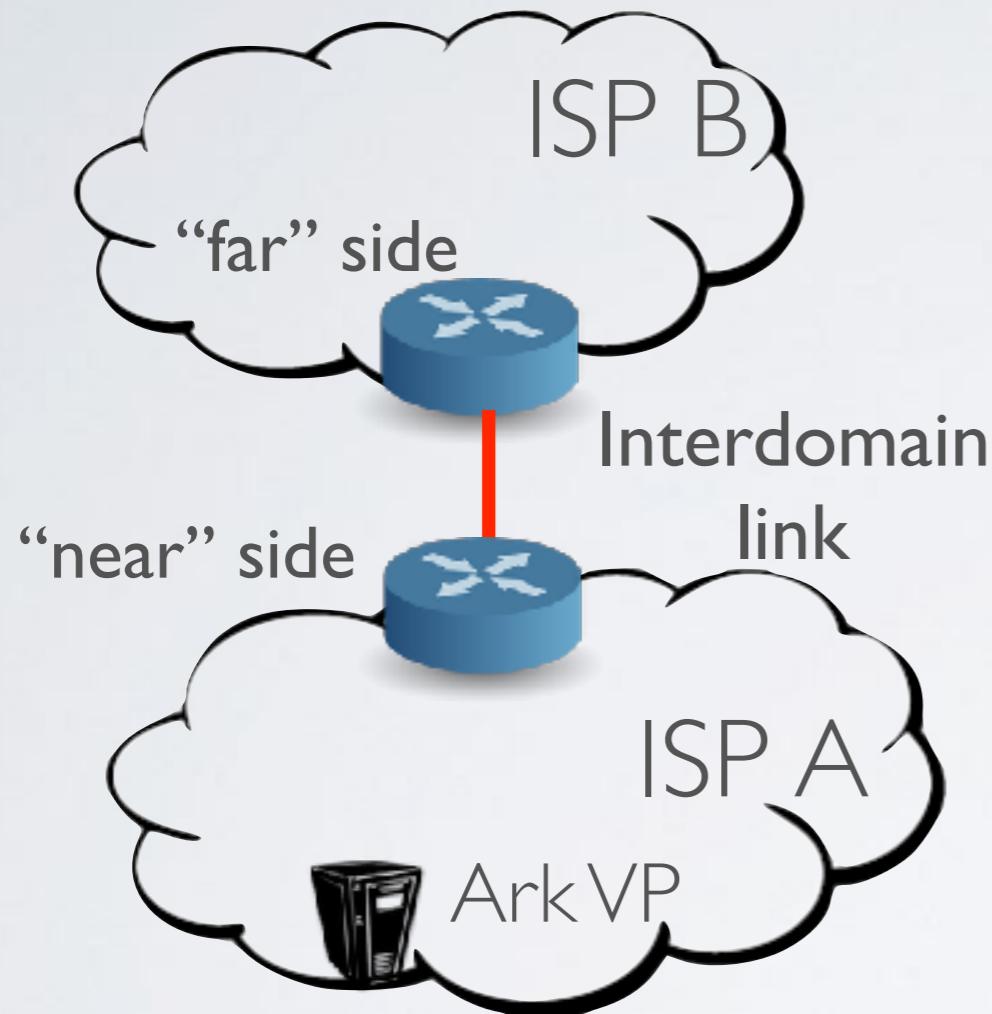
Validating the Method: Step 2



- From Ark VP in ISP A identified congested link with ISP B using TSLP*

*Dhamdhere et al. "Inferring Persistent Interdomain Congestion", SIGCOMM 2018

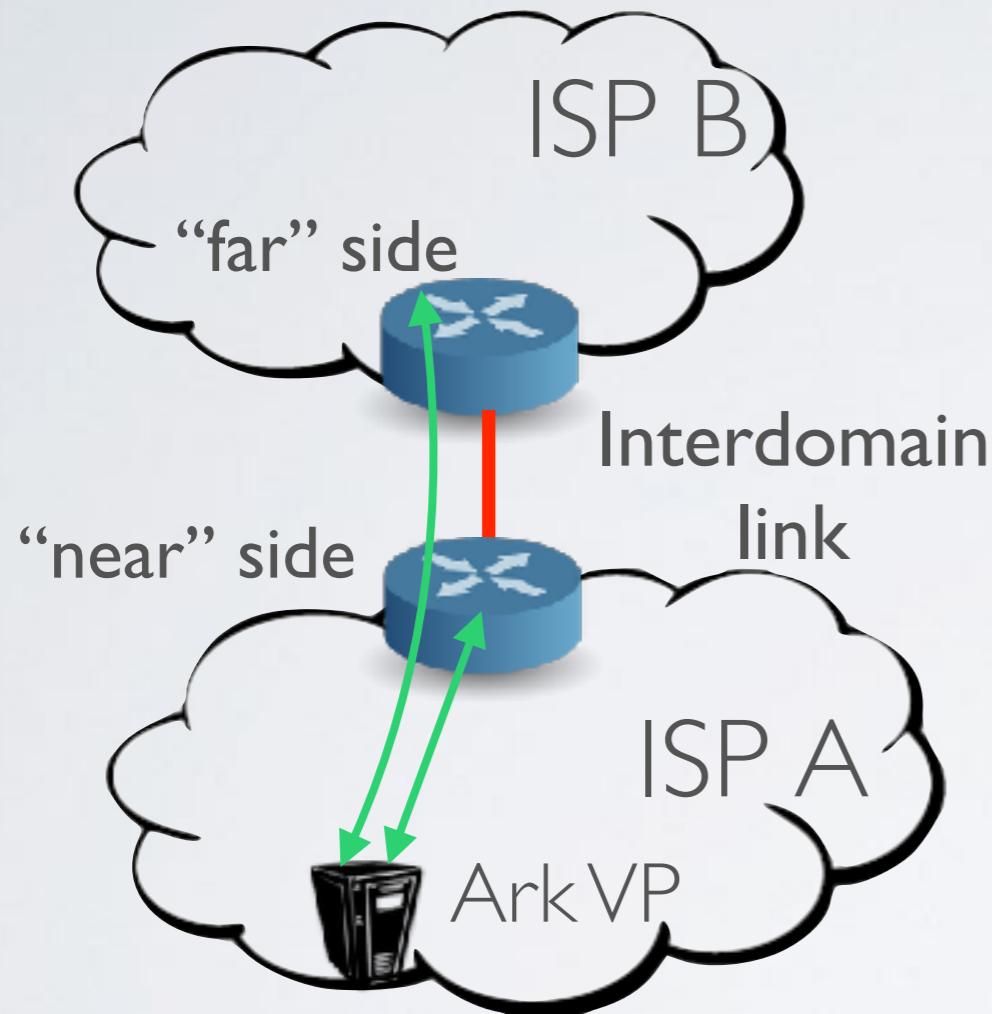
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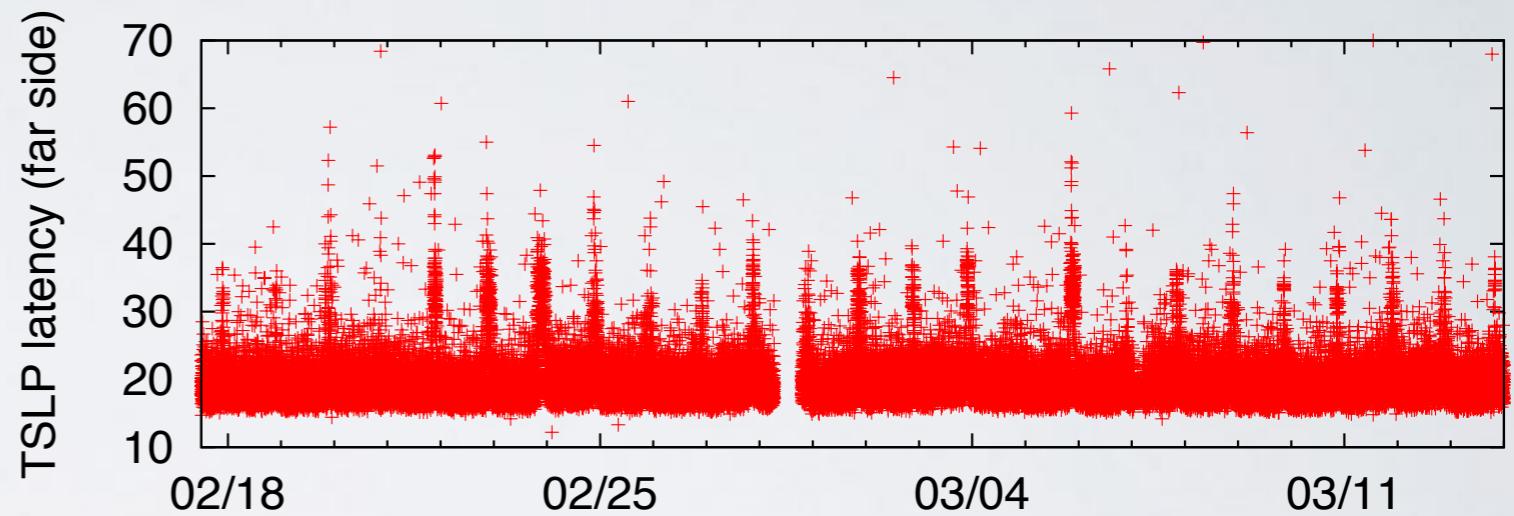
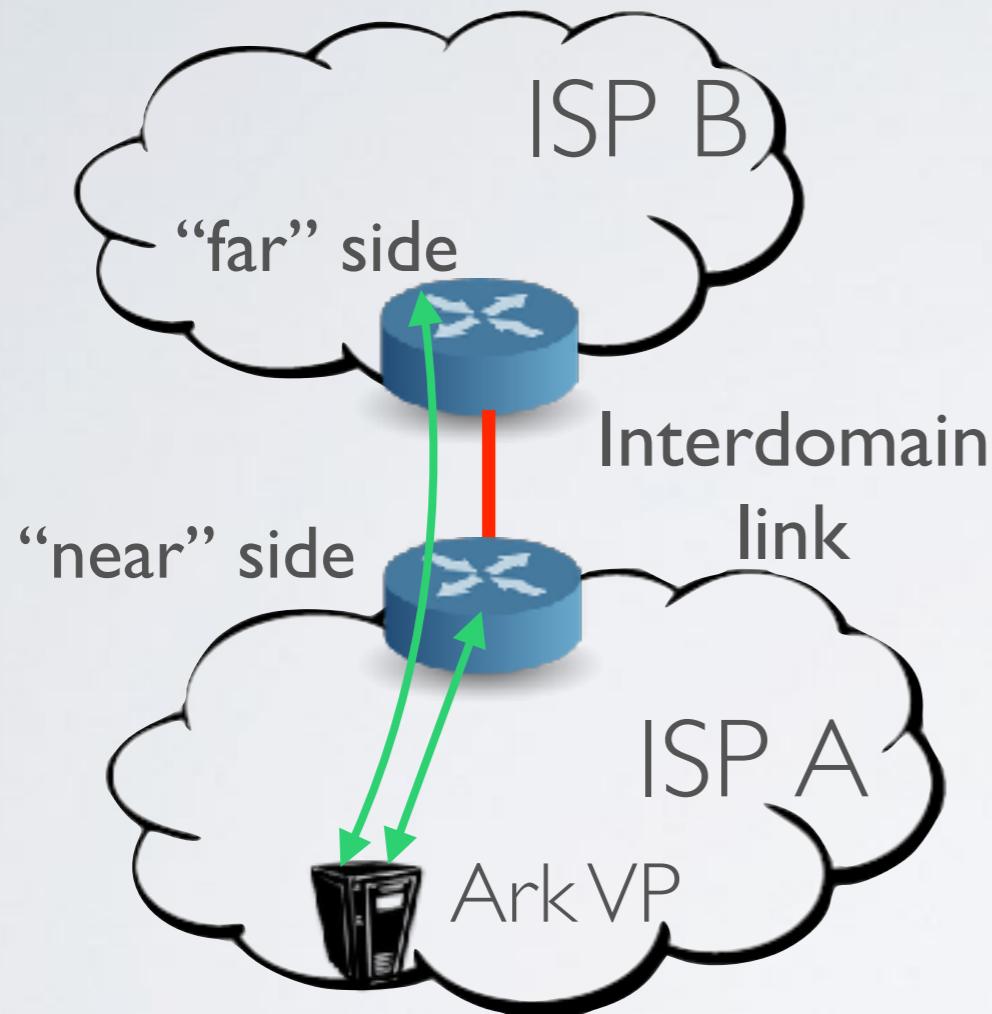


Latency measurements
to “near” and “far” side of
interdomain link over
time

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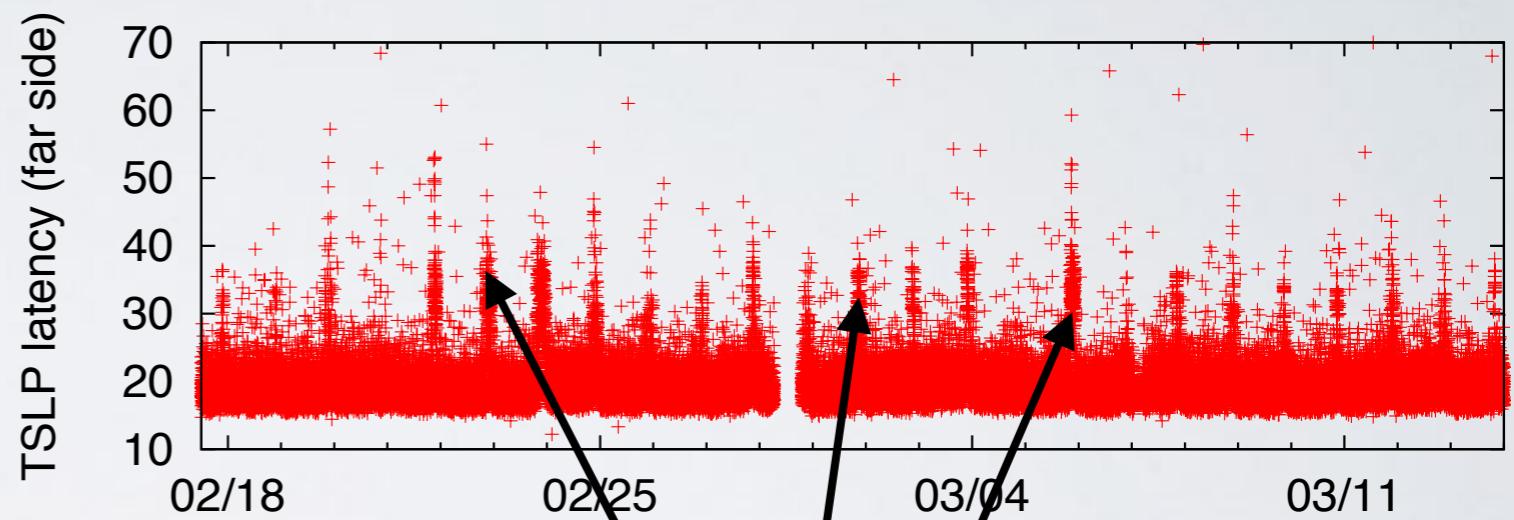
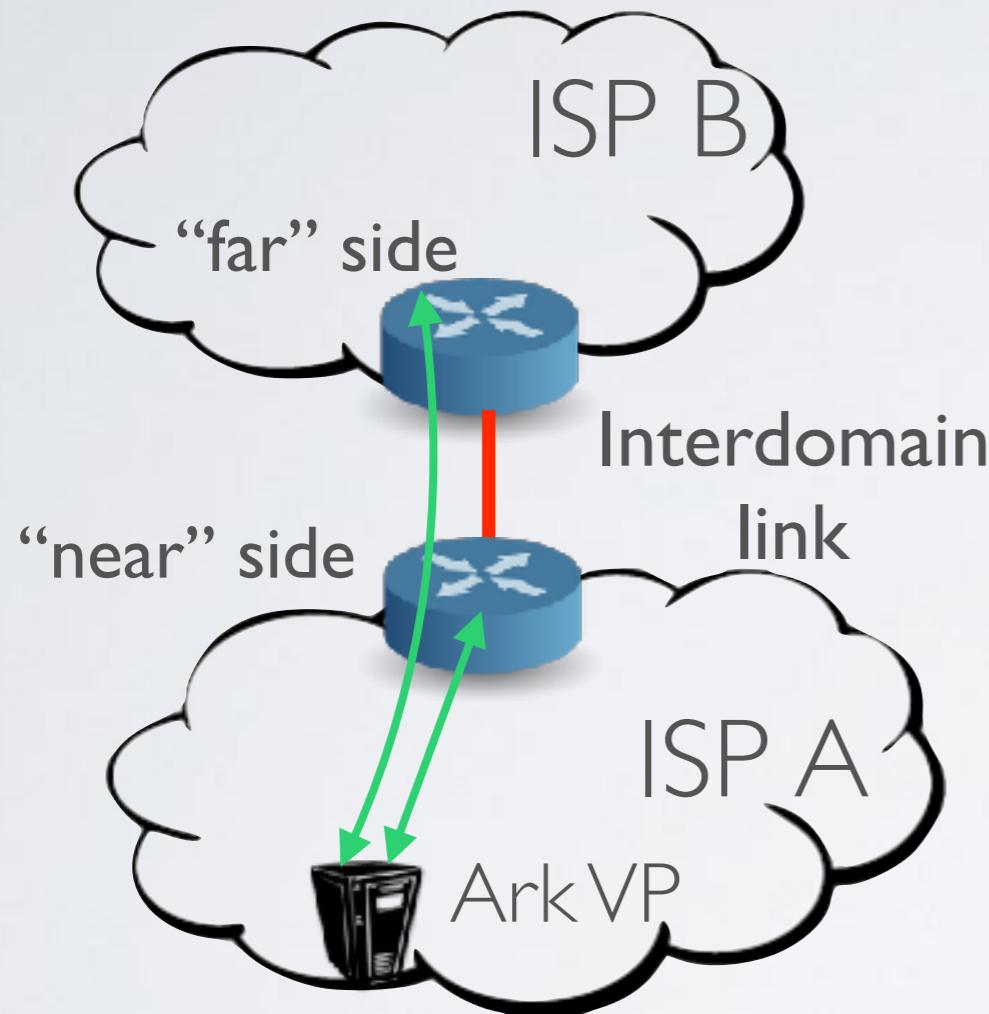
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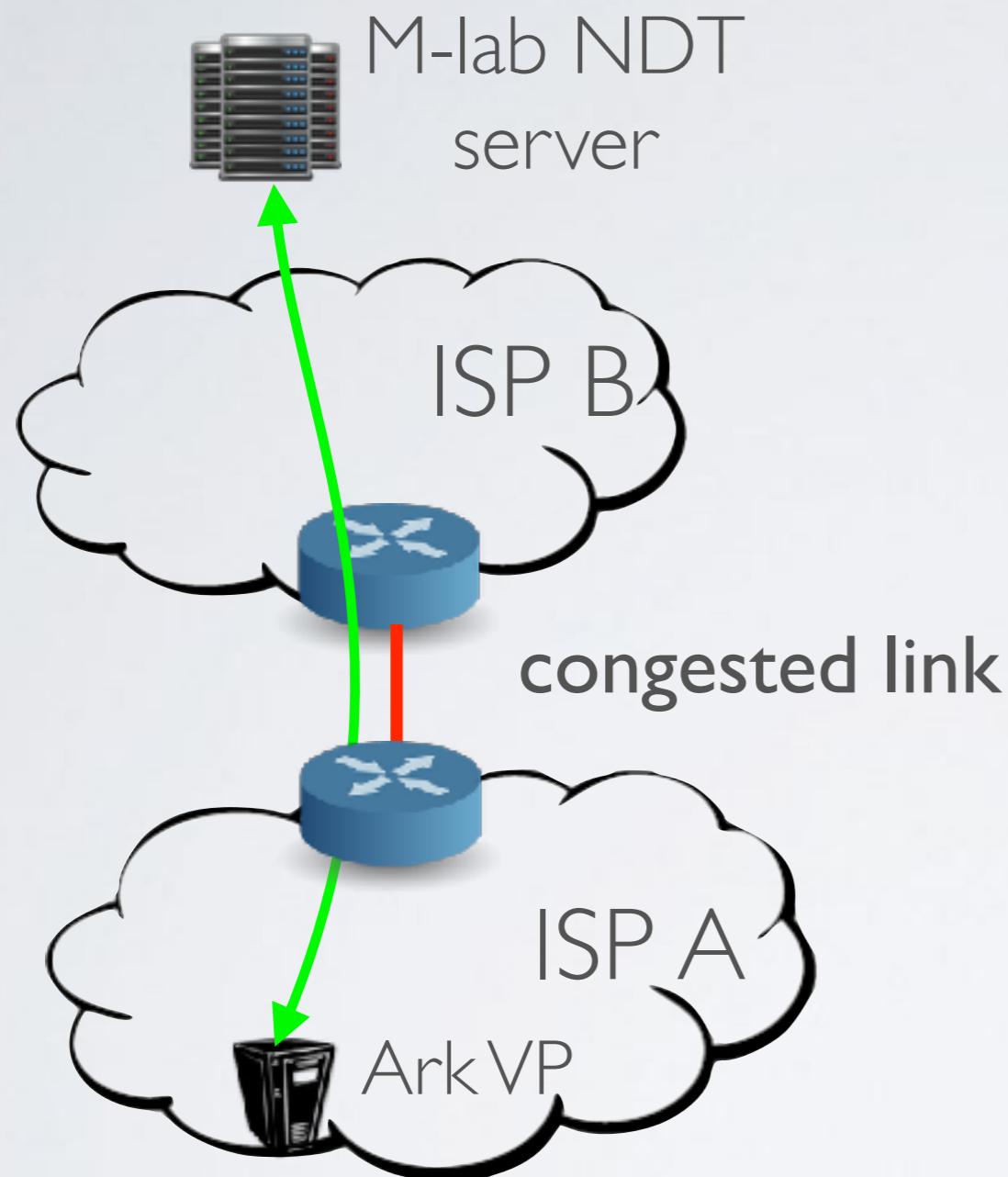


Diurnal latency elevation
indicates congestion

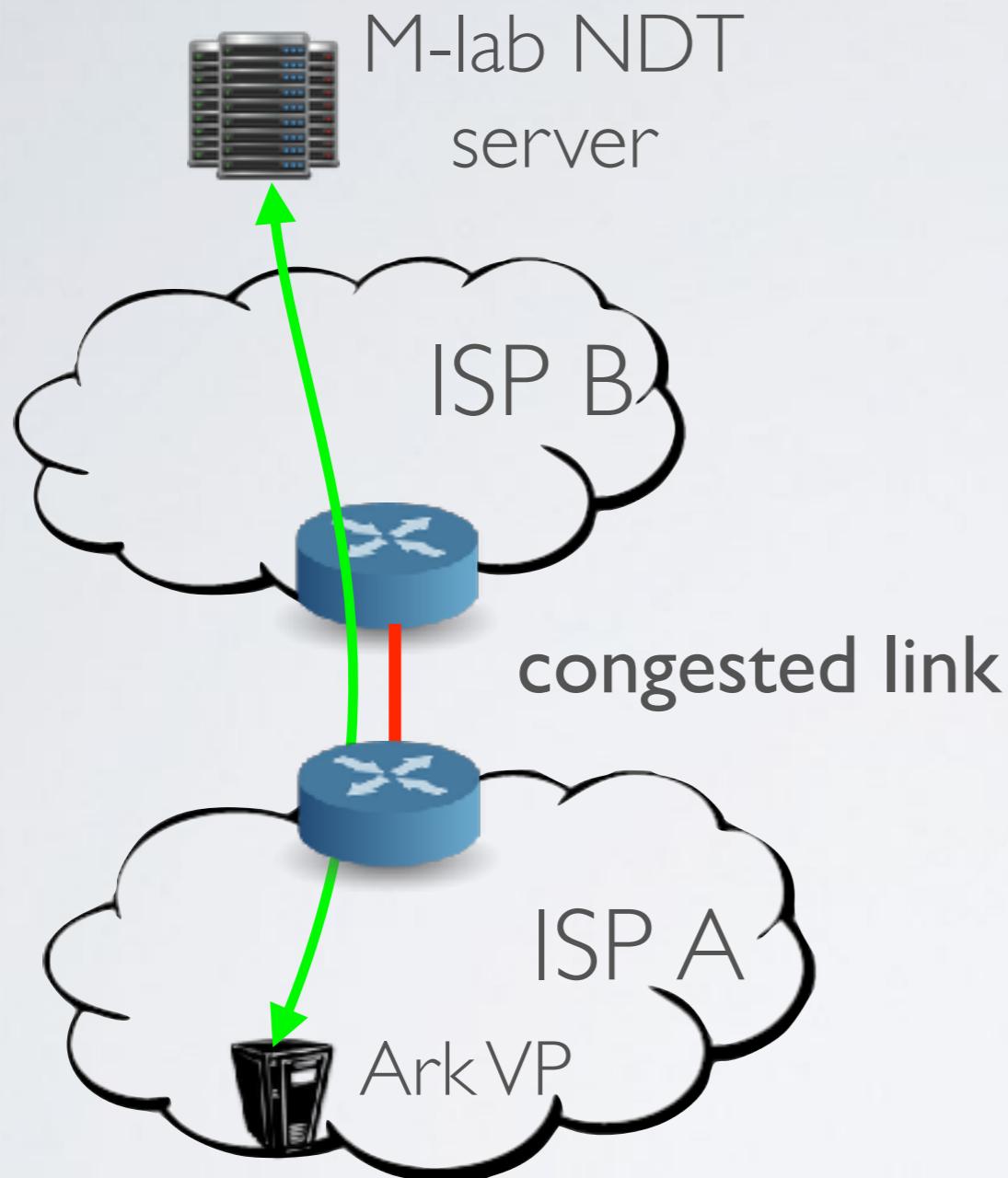
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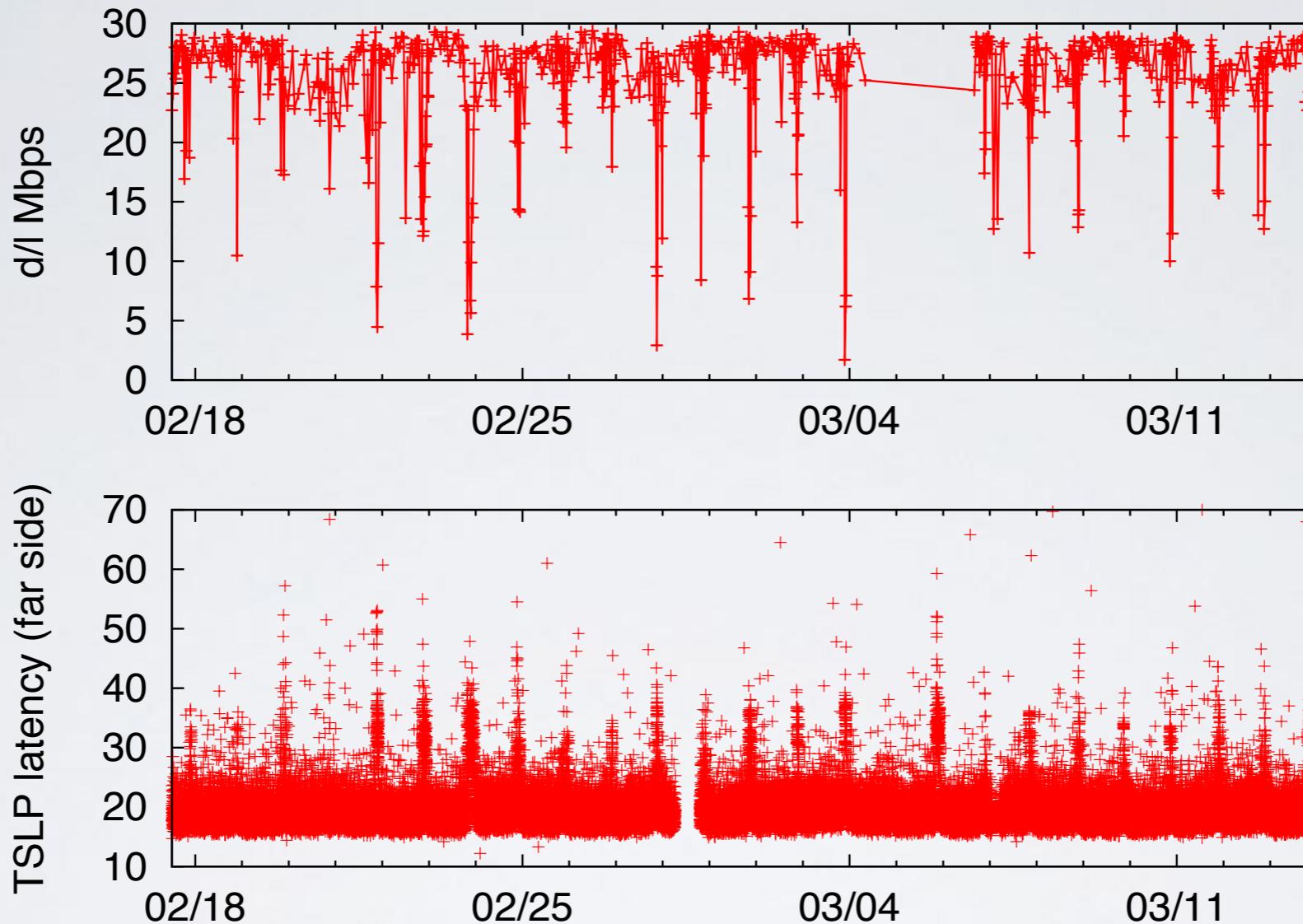


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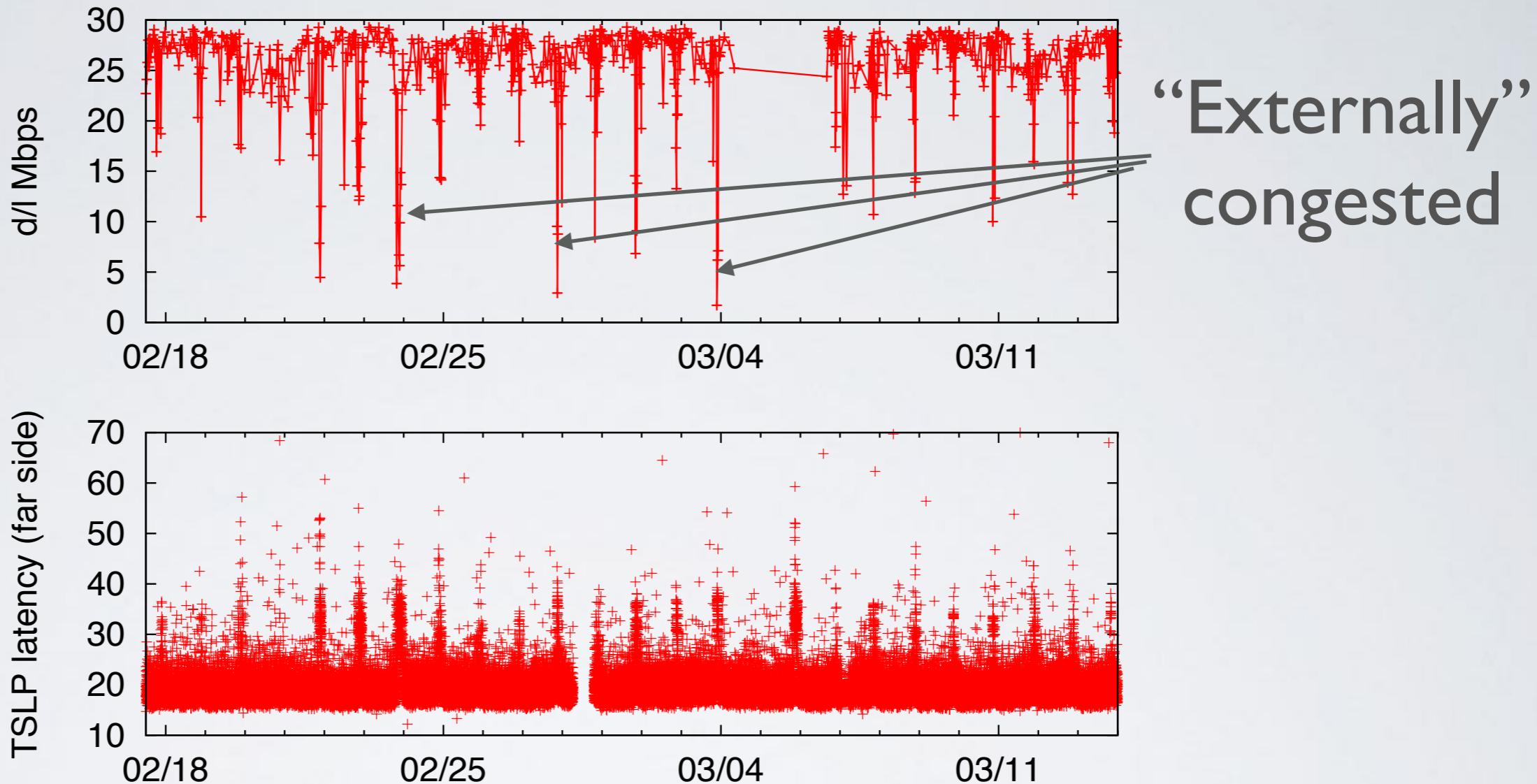
Throughput
measurements from Ark
VP to M-lab NDT server
traversing congested
interdomain link

Validation of the Method: Step 2



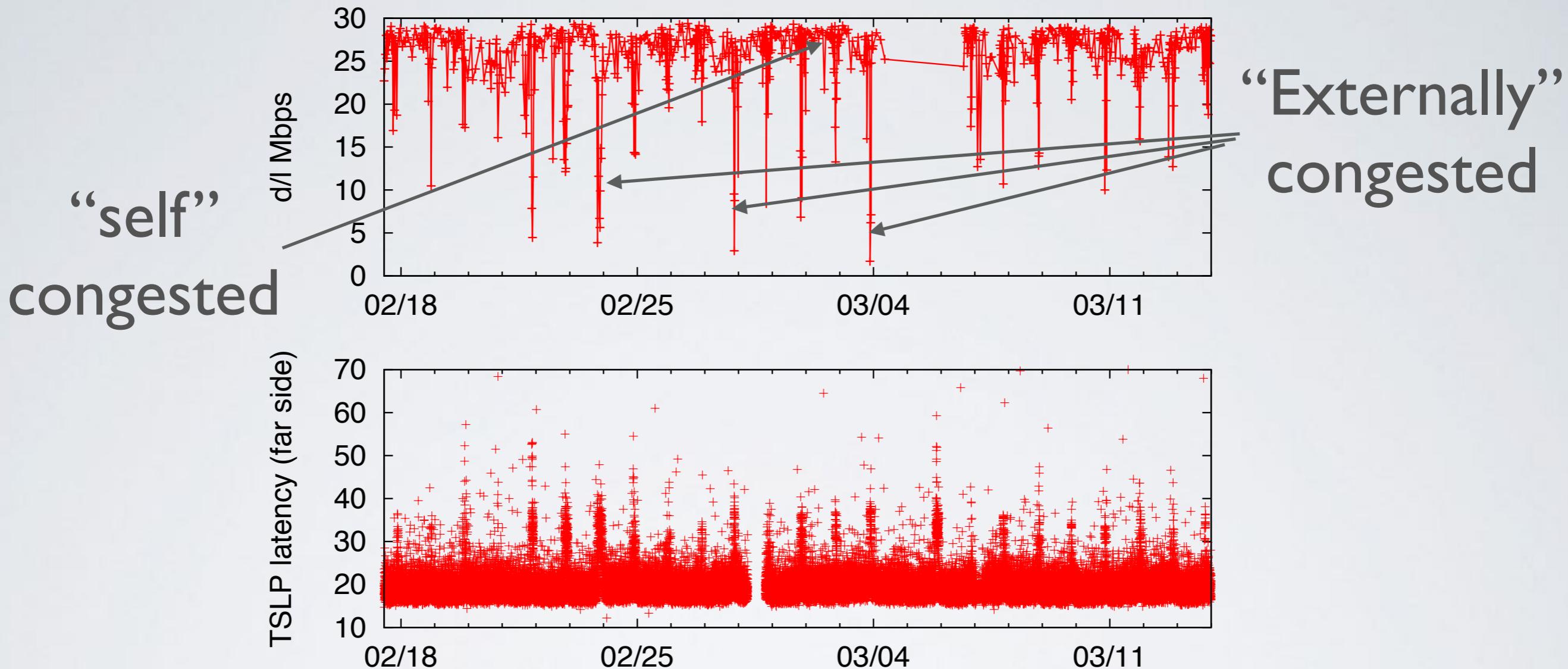
Strong correlation between throughput and TSLP latency: flows during elevated TSLP latency labeled as “externally” congested

Validation of the Method: Step 2



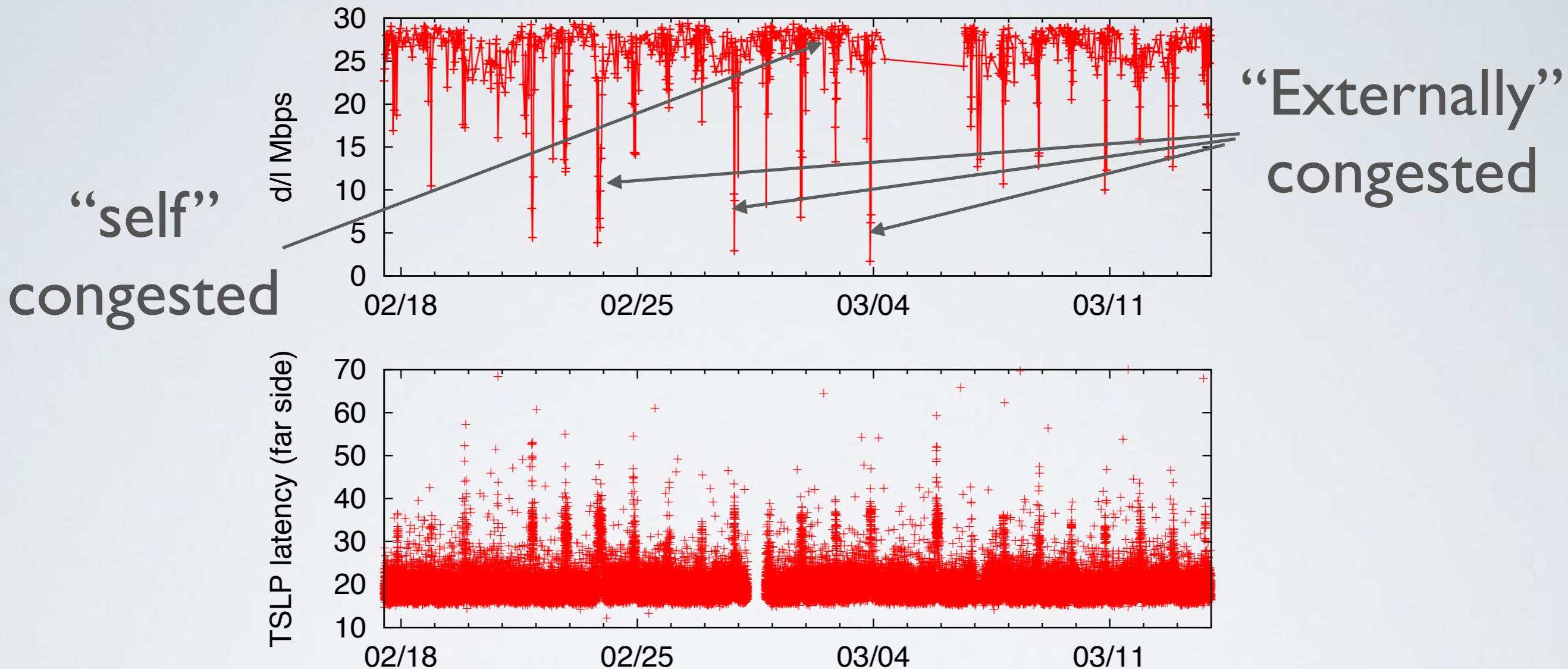
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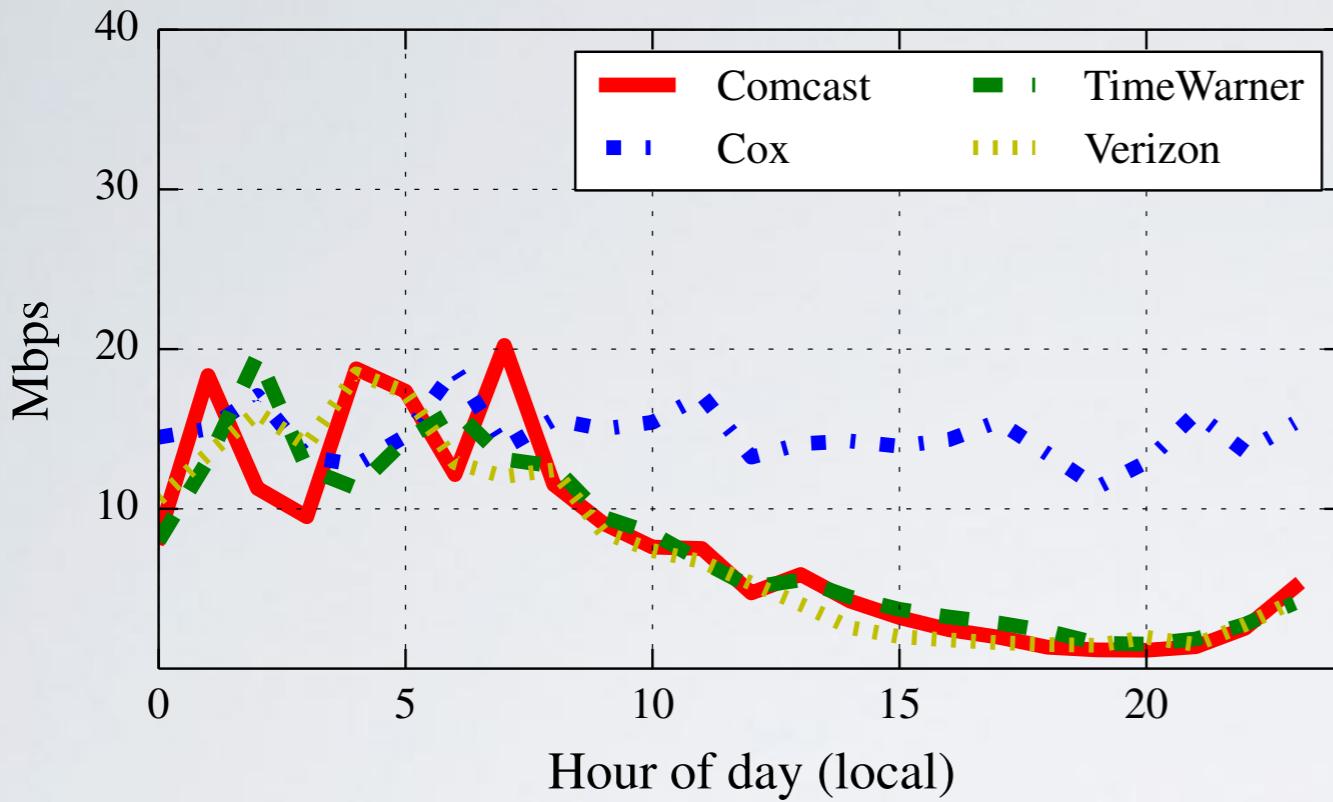


75%+ accuracy in detecting external congestion,
100% accuracy for self-induced congestion

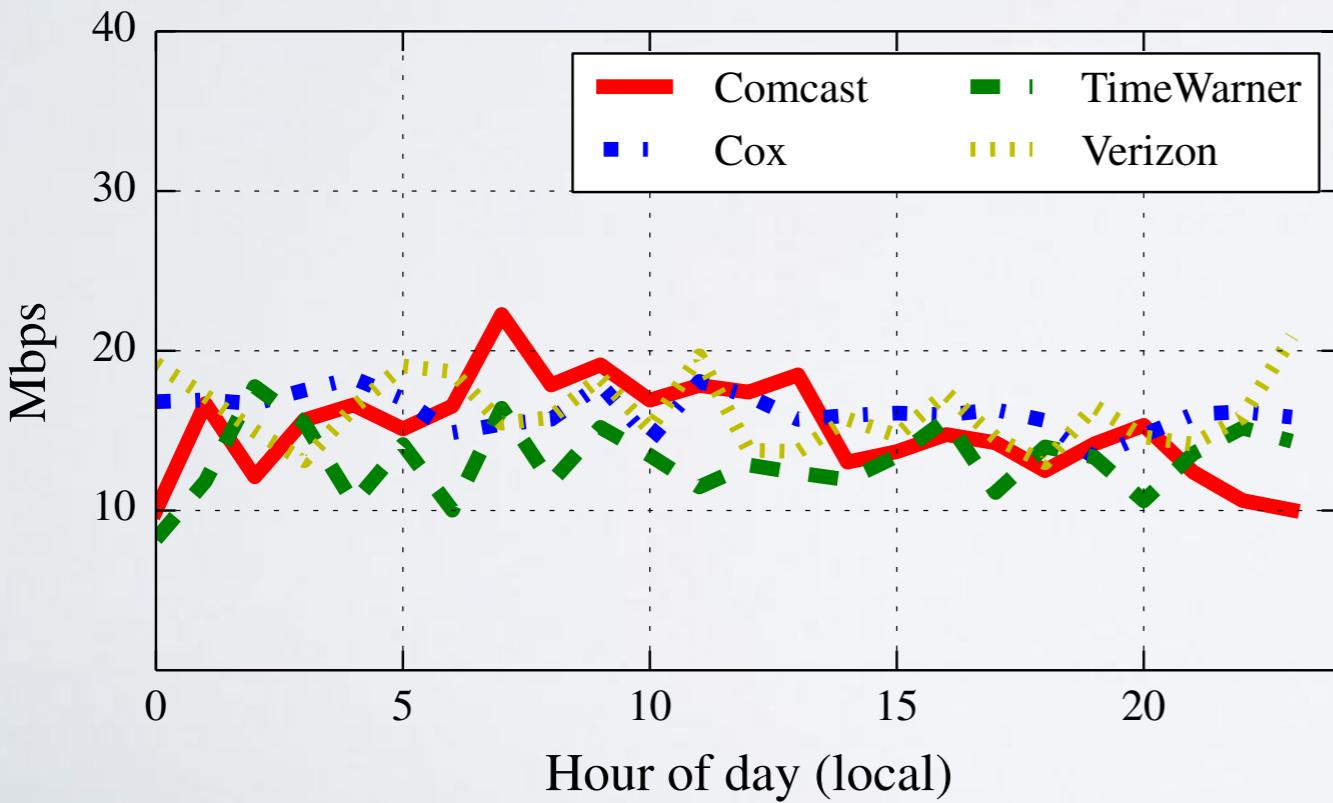
Validation of the Method: Step 3

- M-lab's NDT test data for real-world validation
- Cogent interconnect issue in late 2013/early 2014
 - NDT tests to Cogent M-lab servers from several major U.S. ISPs saw significantly lower throughput during peak hours: Comcast, TWC, Verizon
 - Cox was notably not affected
 - Underlying cause was congested interconnects

Using the M-lab Data



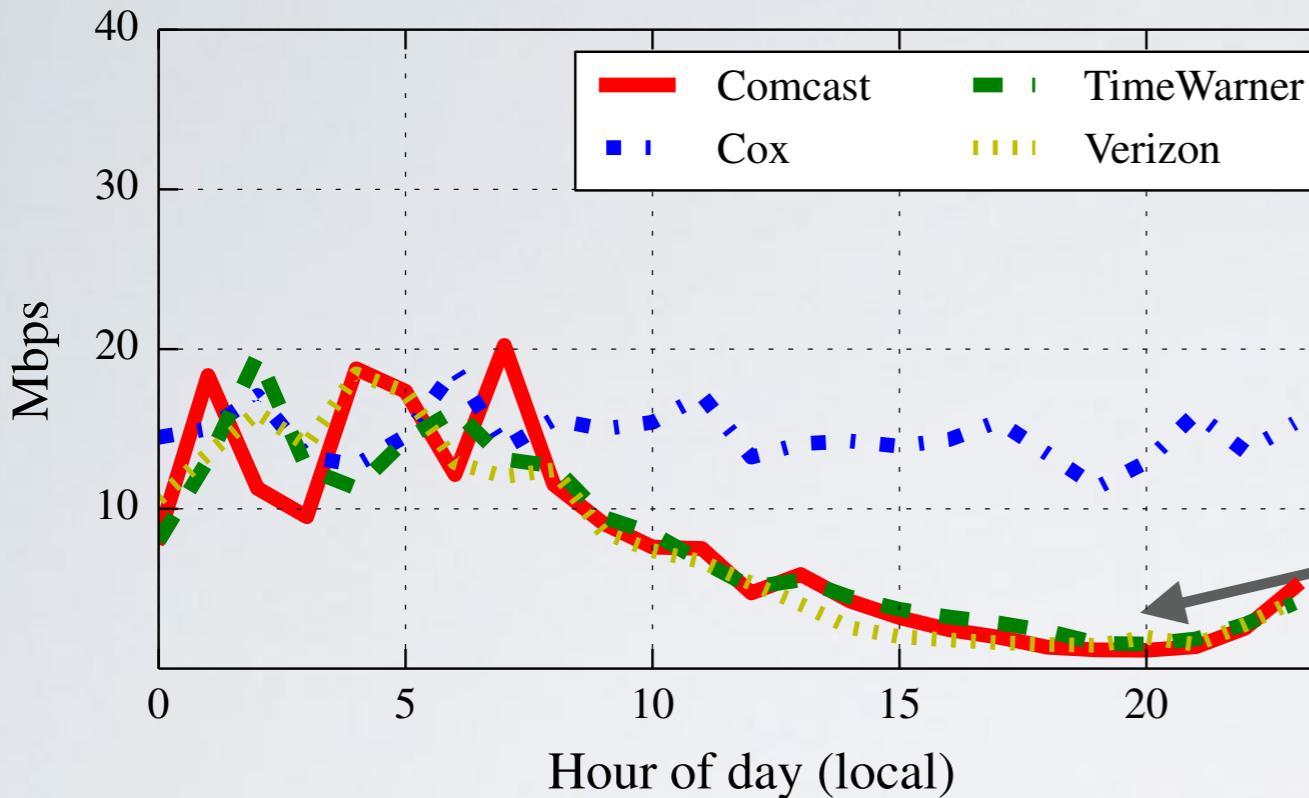
January 2014



April 2014

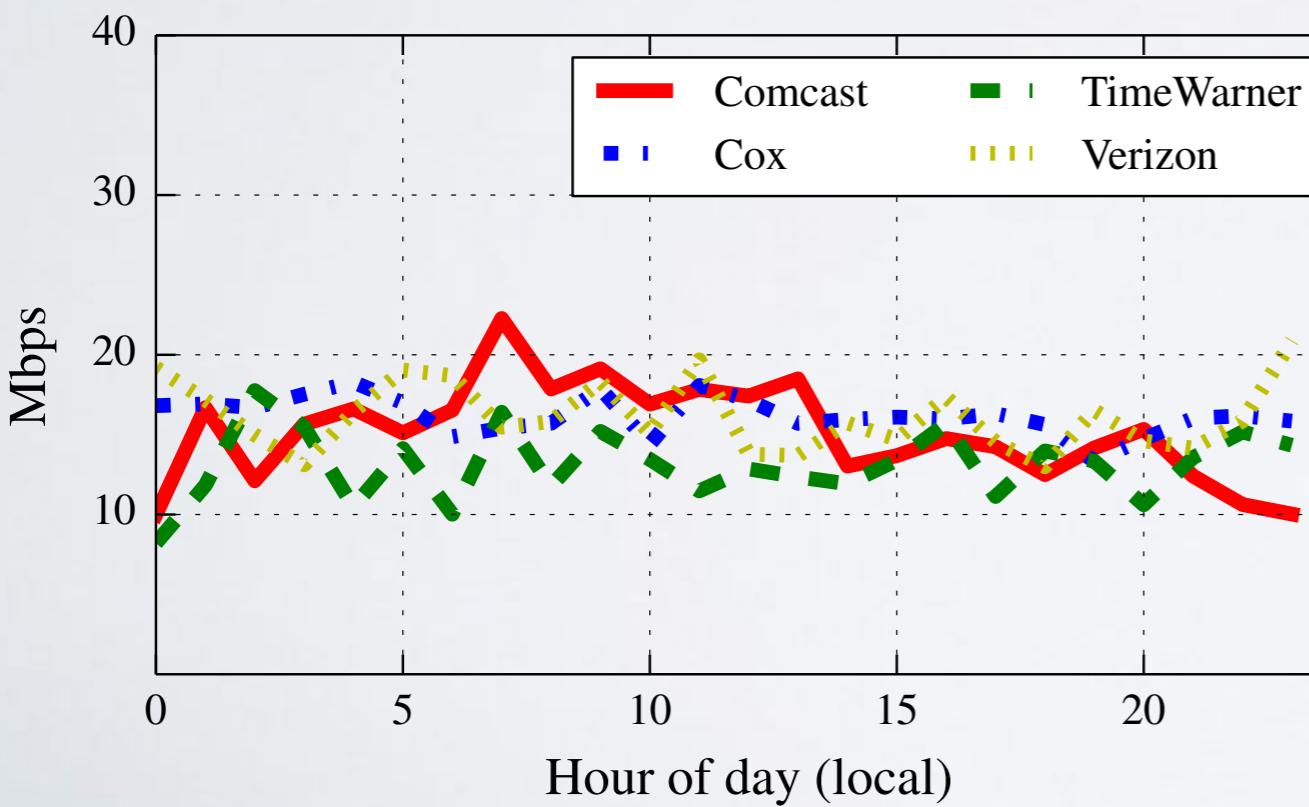


Using the M-lab Data



January 2014

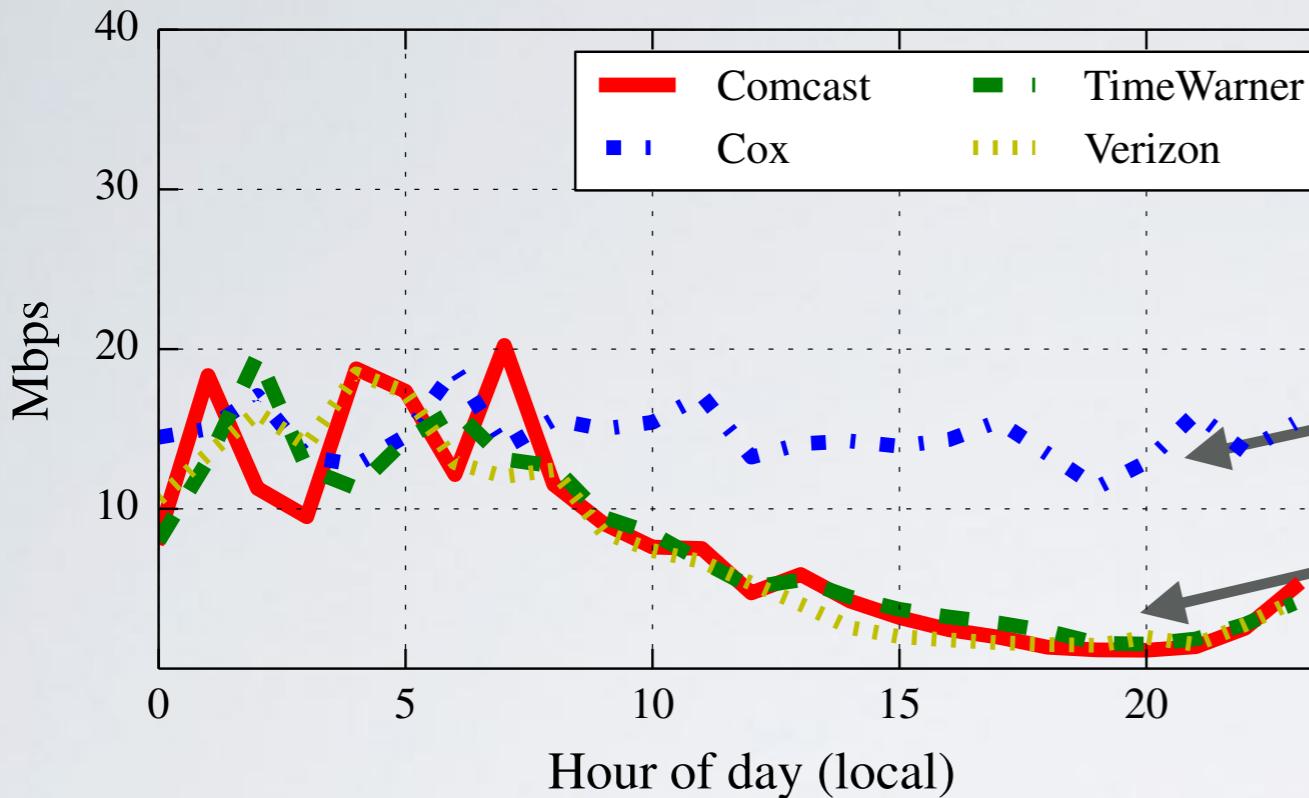
Drop in peak-hour throughput for
for Comcast, TWC, Verizon



April 2014

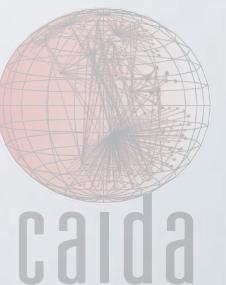
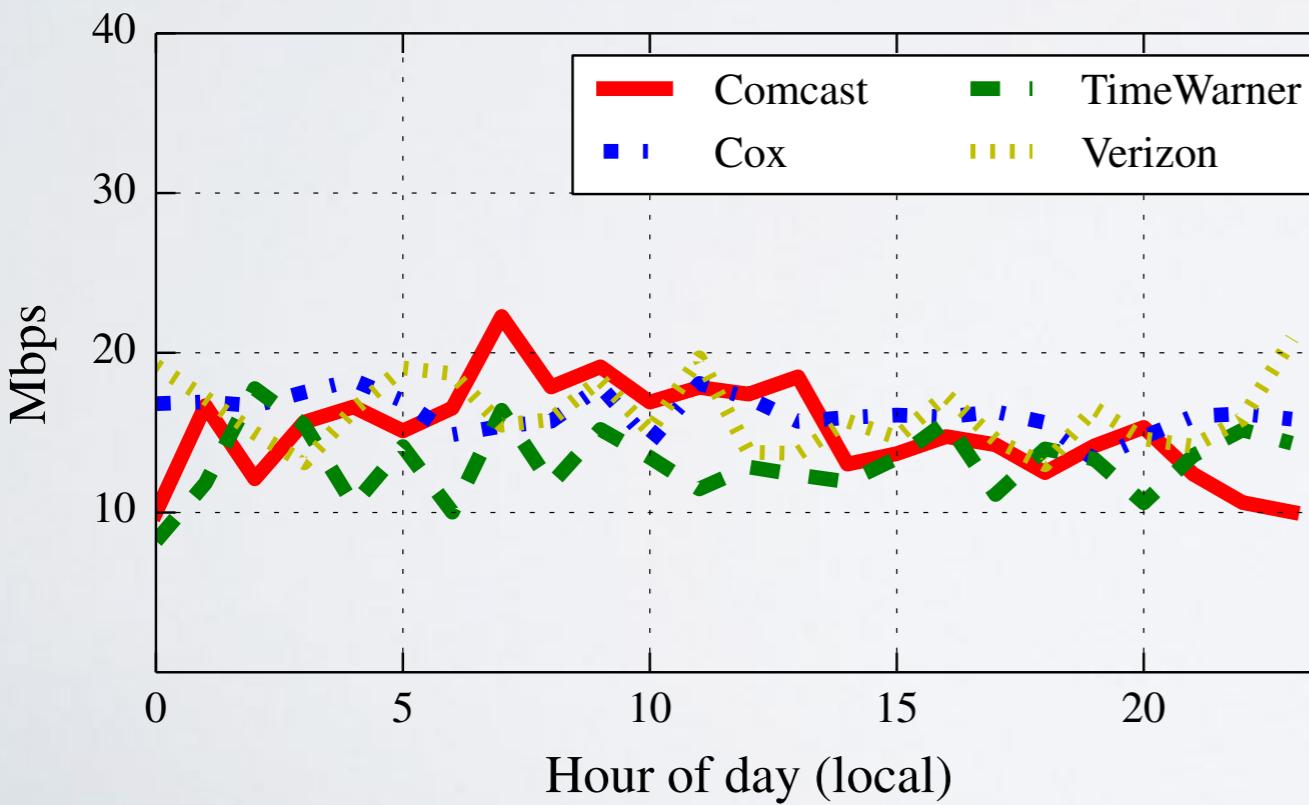


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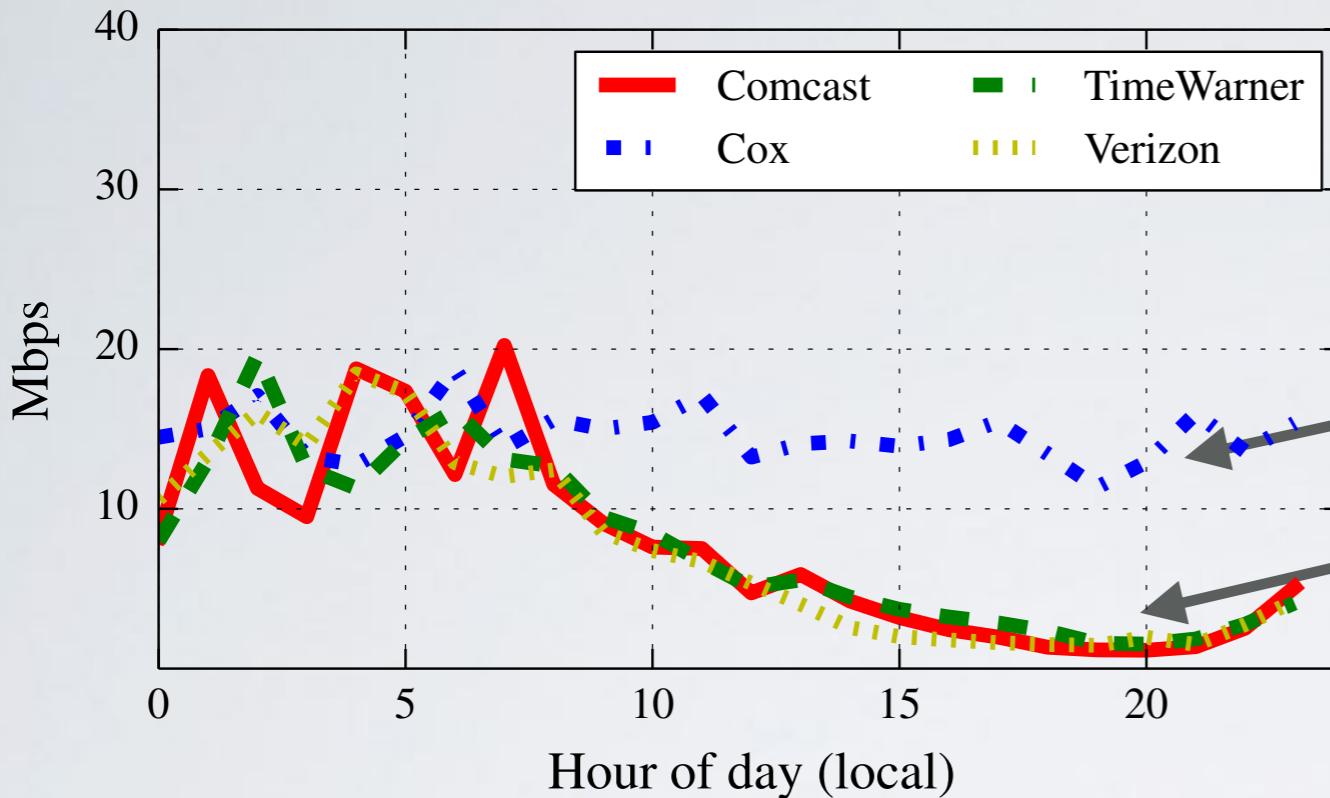


Cox not affected

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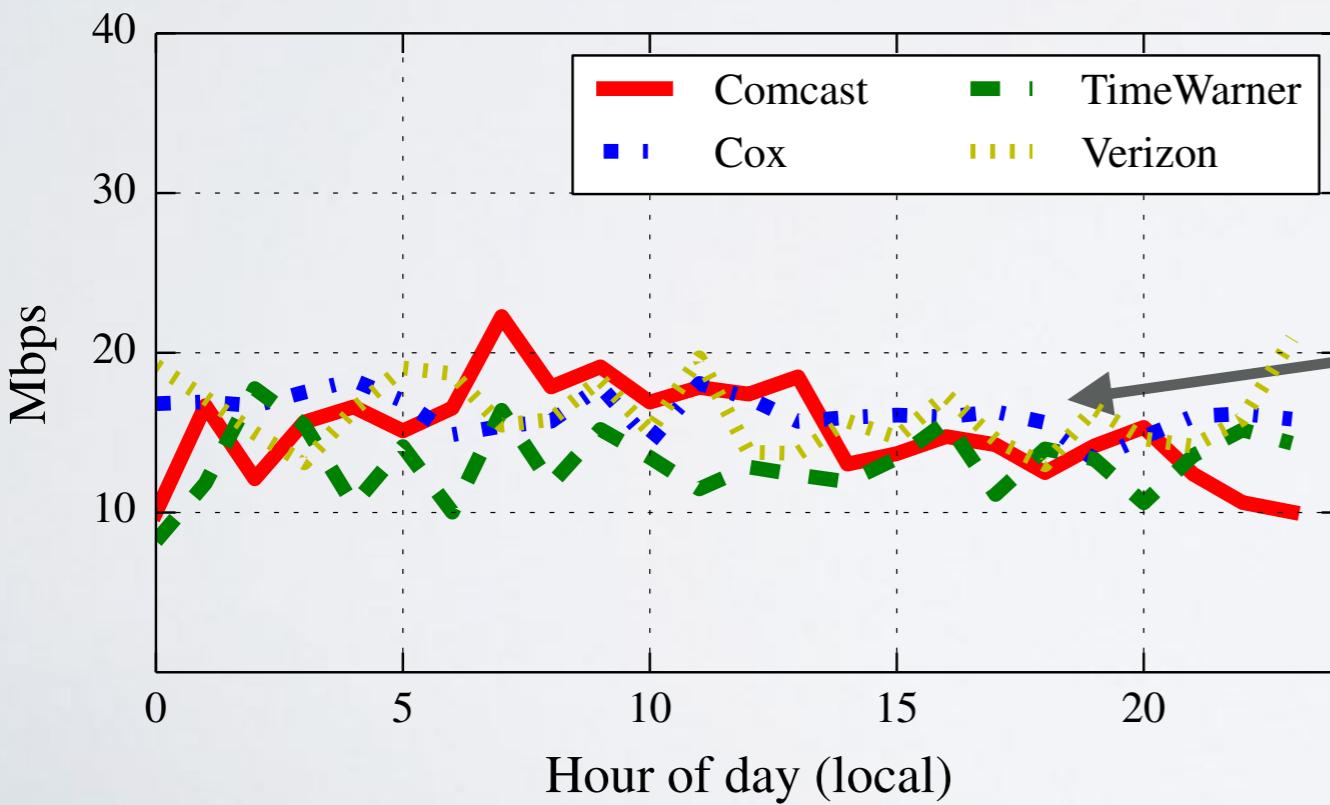


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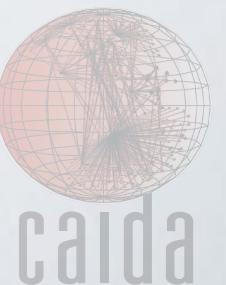


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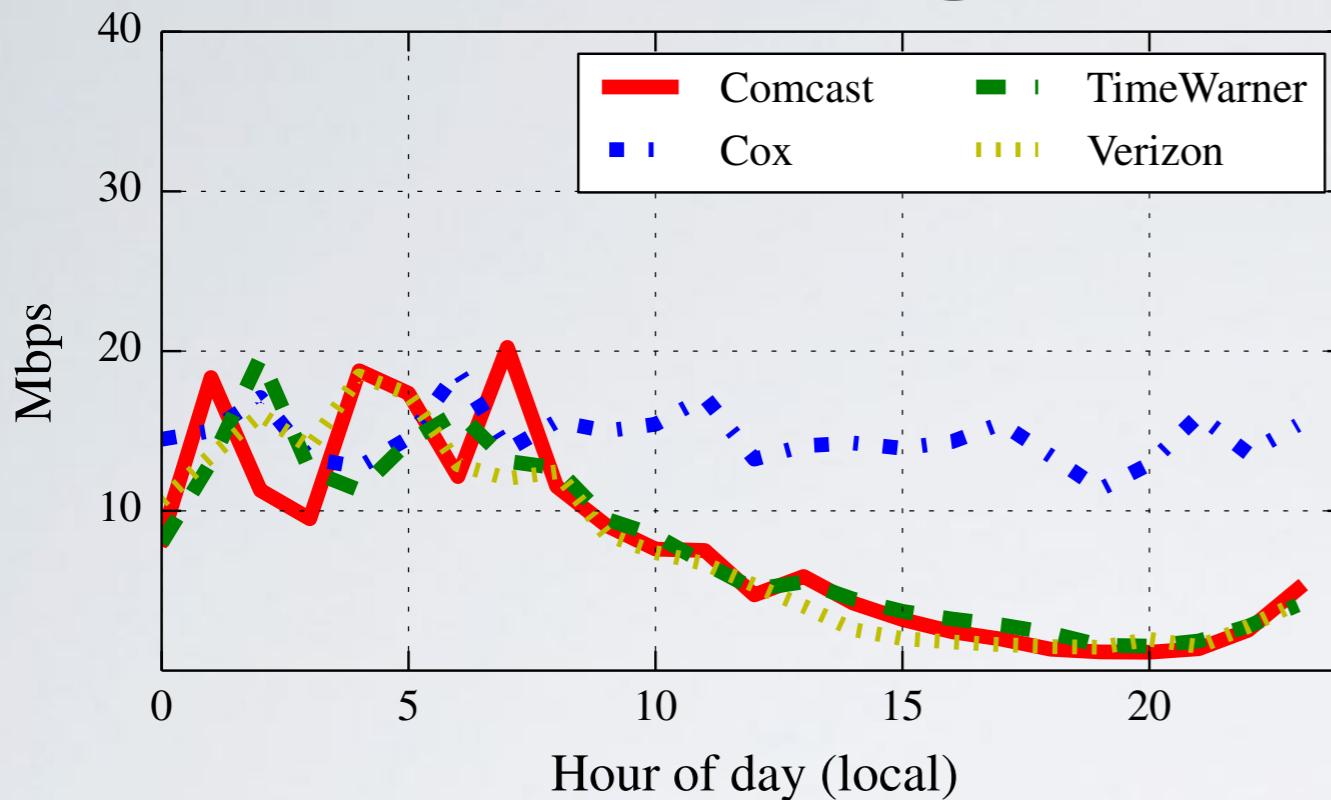
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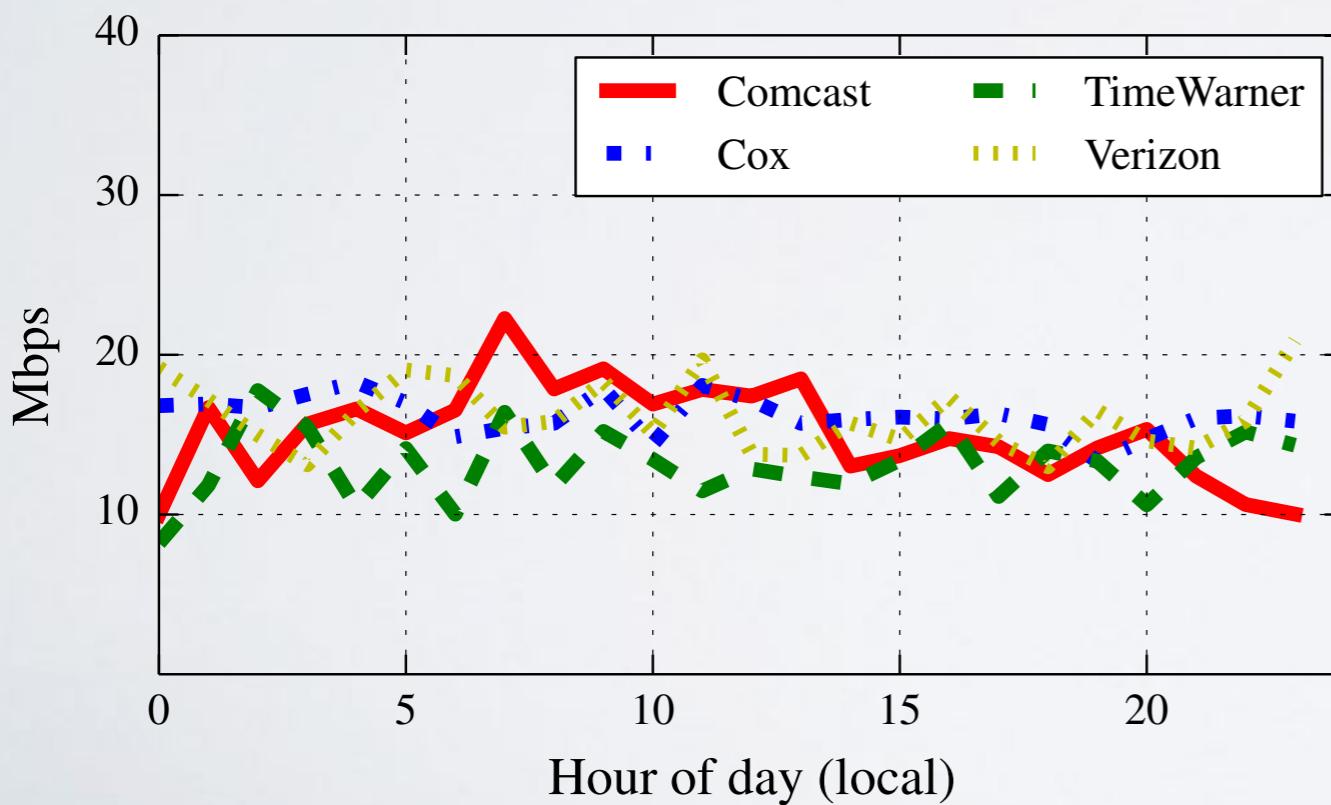
Interconnection dispute resolved; no diurnal effect



Using the M-lab Data



Peak hour tests in
Jan/Feb 2014 are likely
“externally” congested



Off-peak tests in
Mar/Apr 2014 are likely
“self” congested



Noisy Data

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- Difficult to infer interdomain congestion using throughput*

*Sundaresan et al. "Challenges in Inferring Interdomain Congestion using Throughput Measurements", IMC 2017

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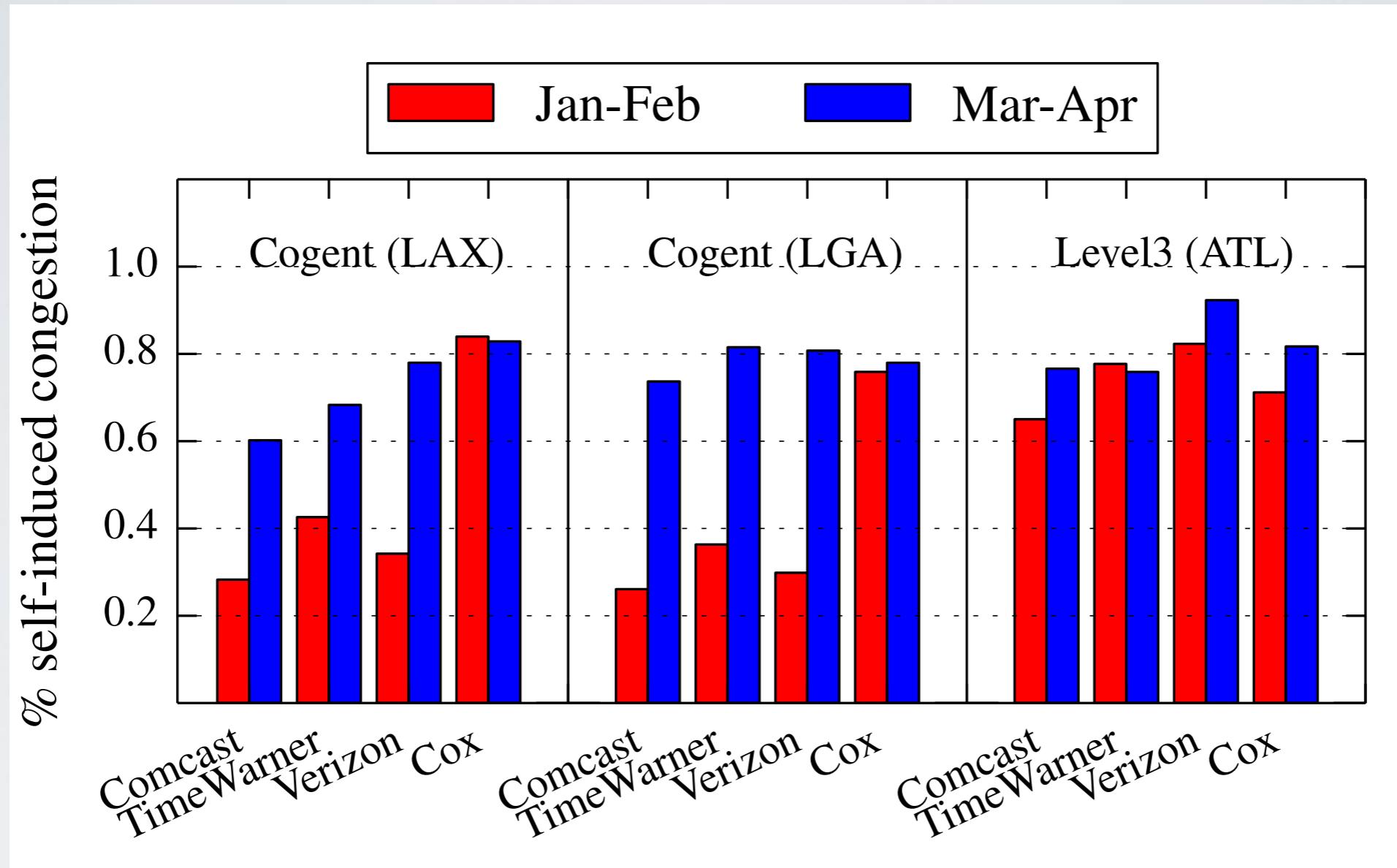
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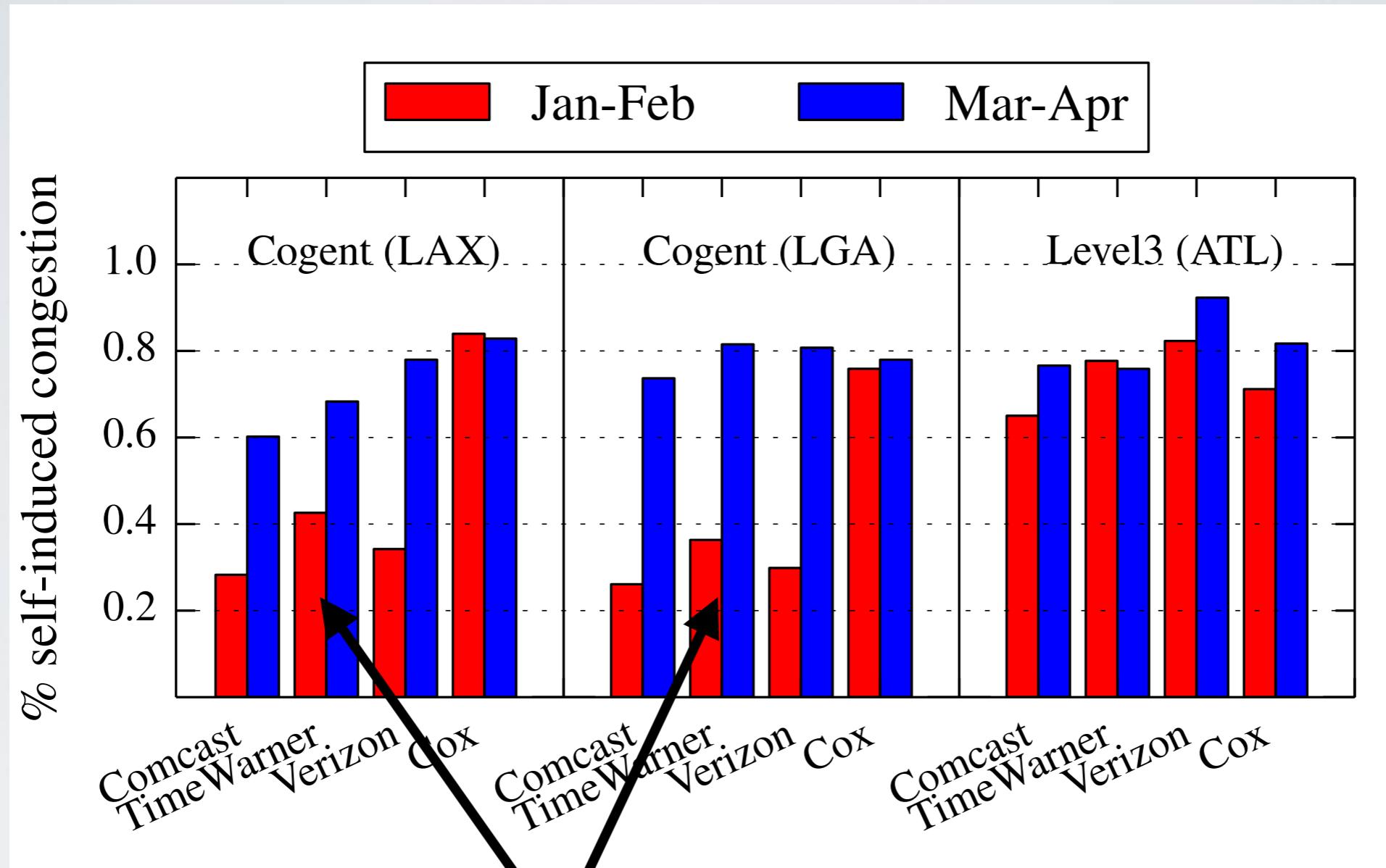
- Difficult to infer interdomain congestion using throughput*
- All tests labeled “external” may not have traversed congested interconnects
- We do not expect to identify all peak hour tests as externally congested, and vice versa
 - Looking for qualitative differences

*Sundaresan et al. “Challenges in Inferring Interdomain Congestion using Throughput Measurements”, IMC 2017

Applying the Model to M-lab data

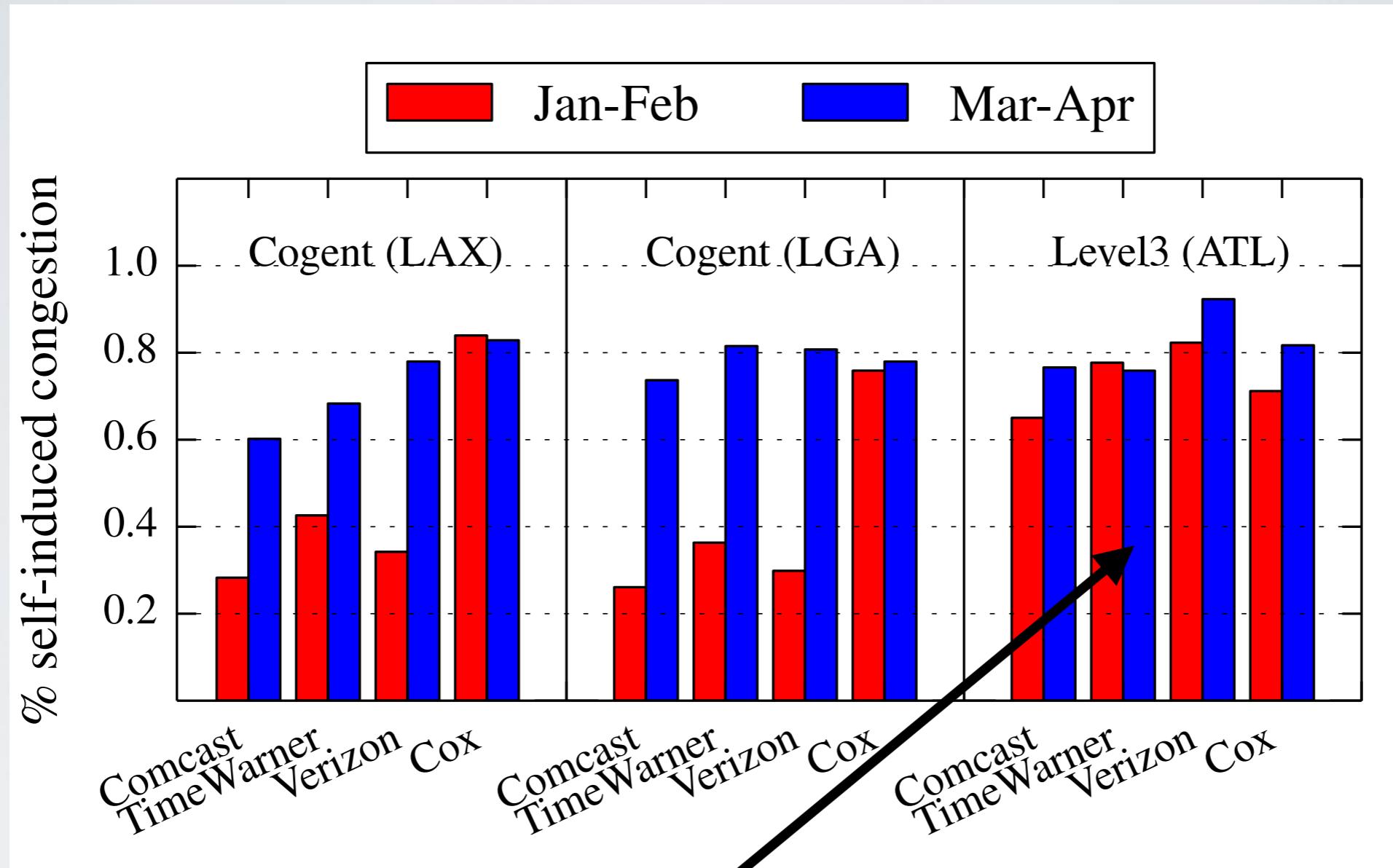


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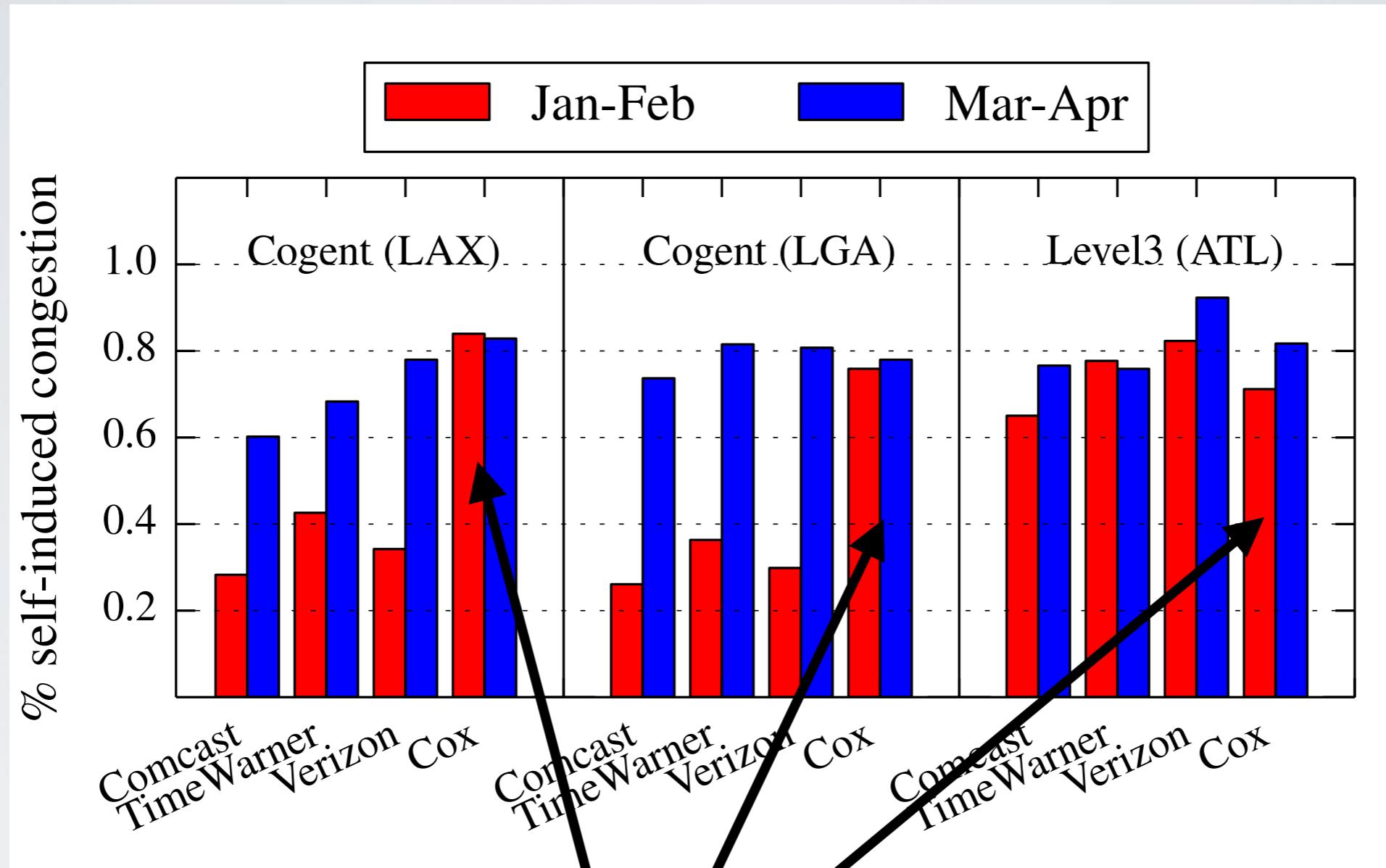
Much lower incidences of self-induced congestion for Cogent in Jan/Feb 2014 as compared to Mar/Apr

Applying the Model to M-lab data



Level3 does not show significant differences, was not affected by interconnection disputes

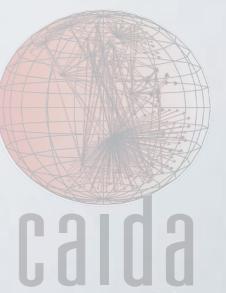
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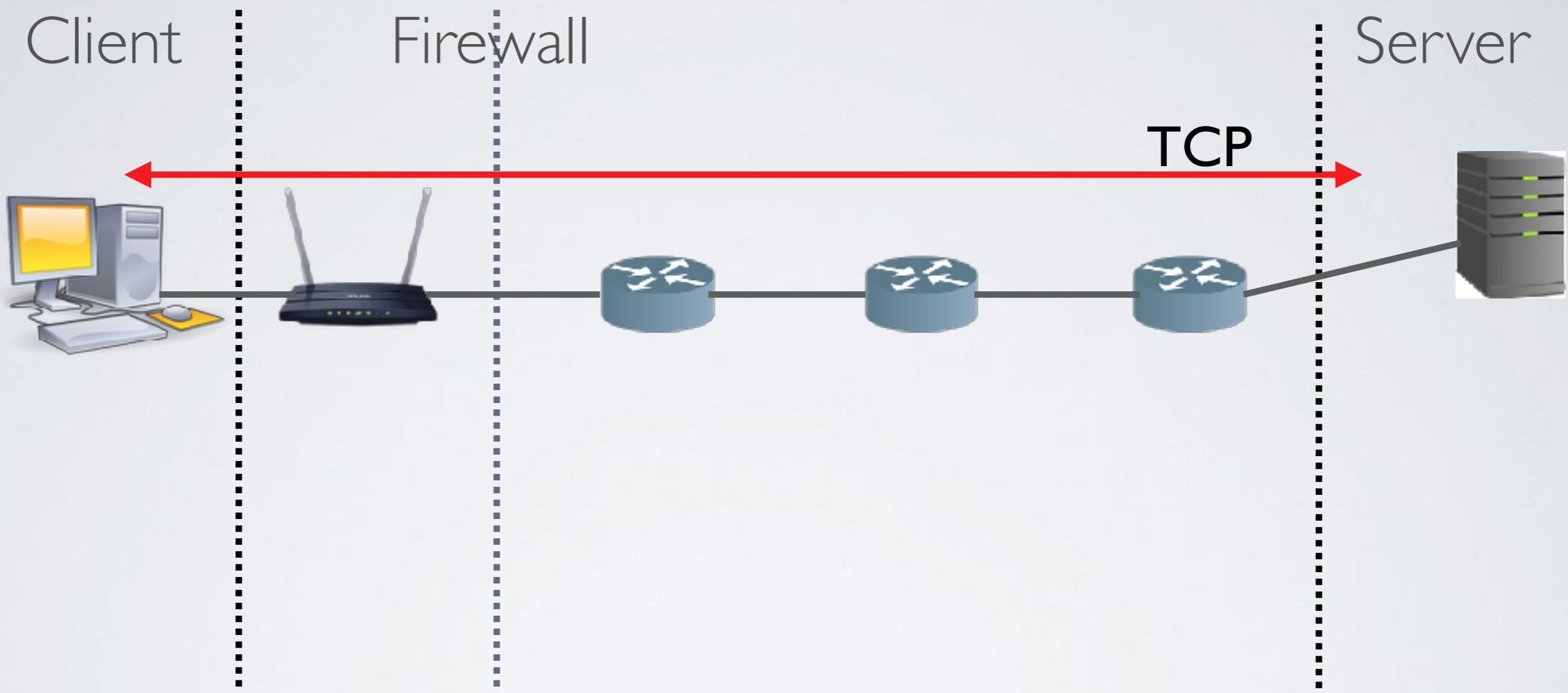
Cox does not show significant differences,
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In-band Measurements

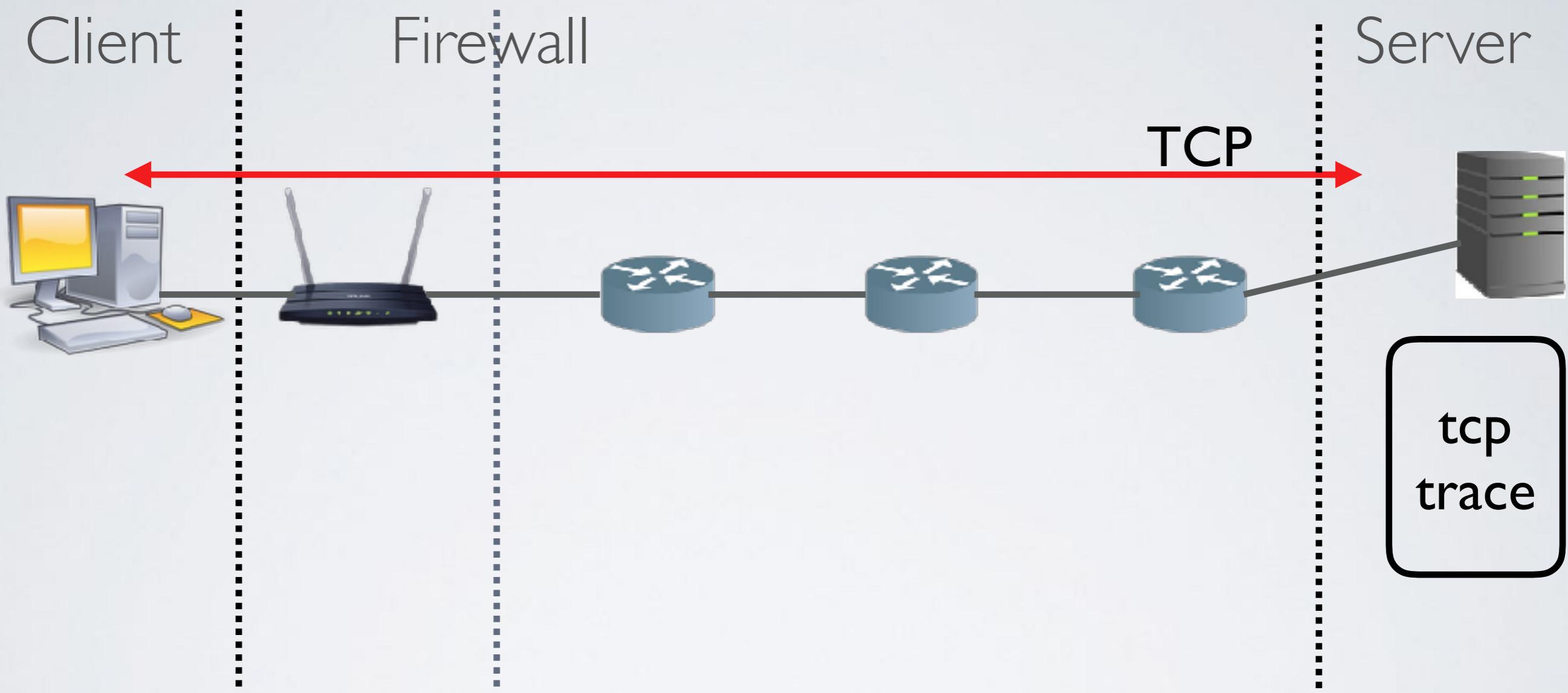
- Can we leverage an ongoing TCP connection for path measurements?
- e.g., where is the TCP flow bottlenecked? is the client's wireless access network the bottleneck? What is the capacity/available bandwidth of the path?
- Why in-band? No need to send external flows (which may be treated differently than the application)
- TCP flow has already punched a hole in the NAT at the client side



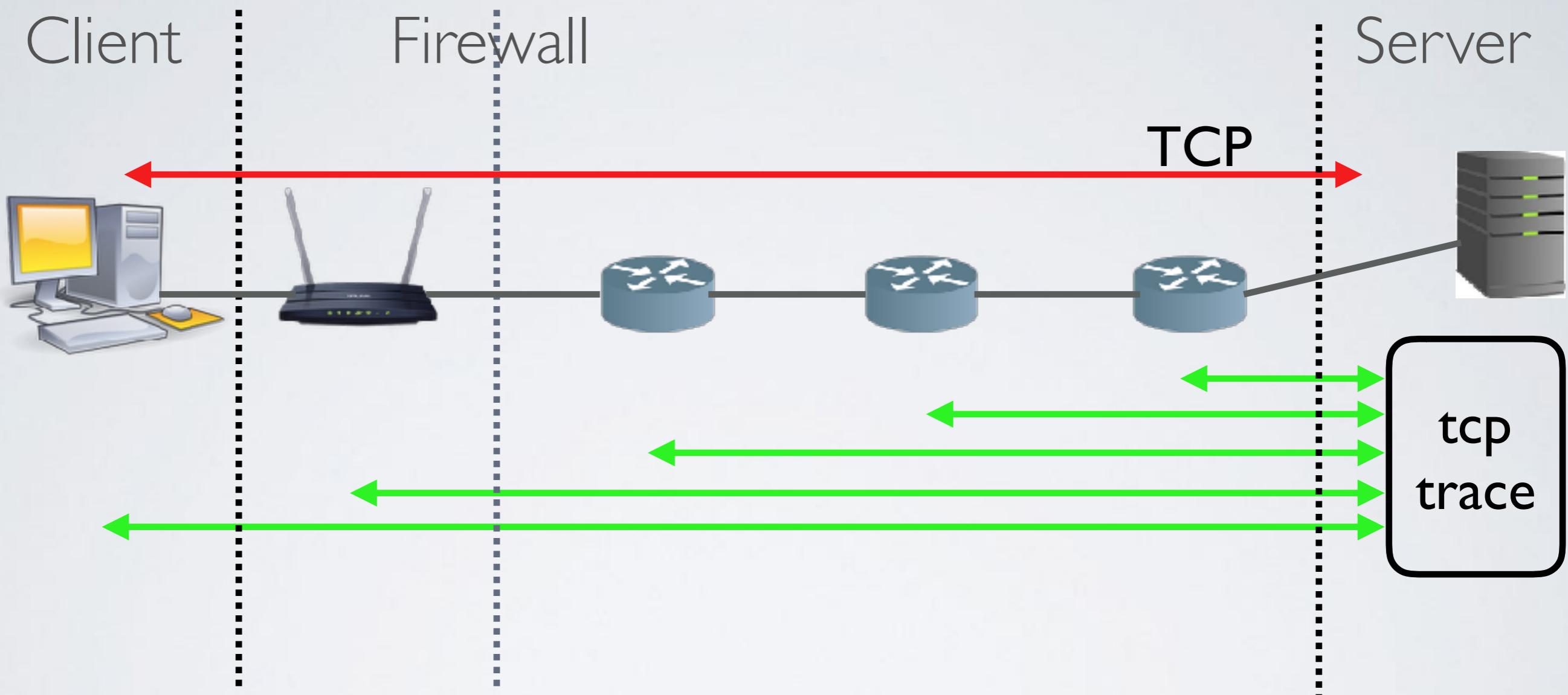
tcptrace



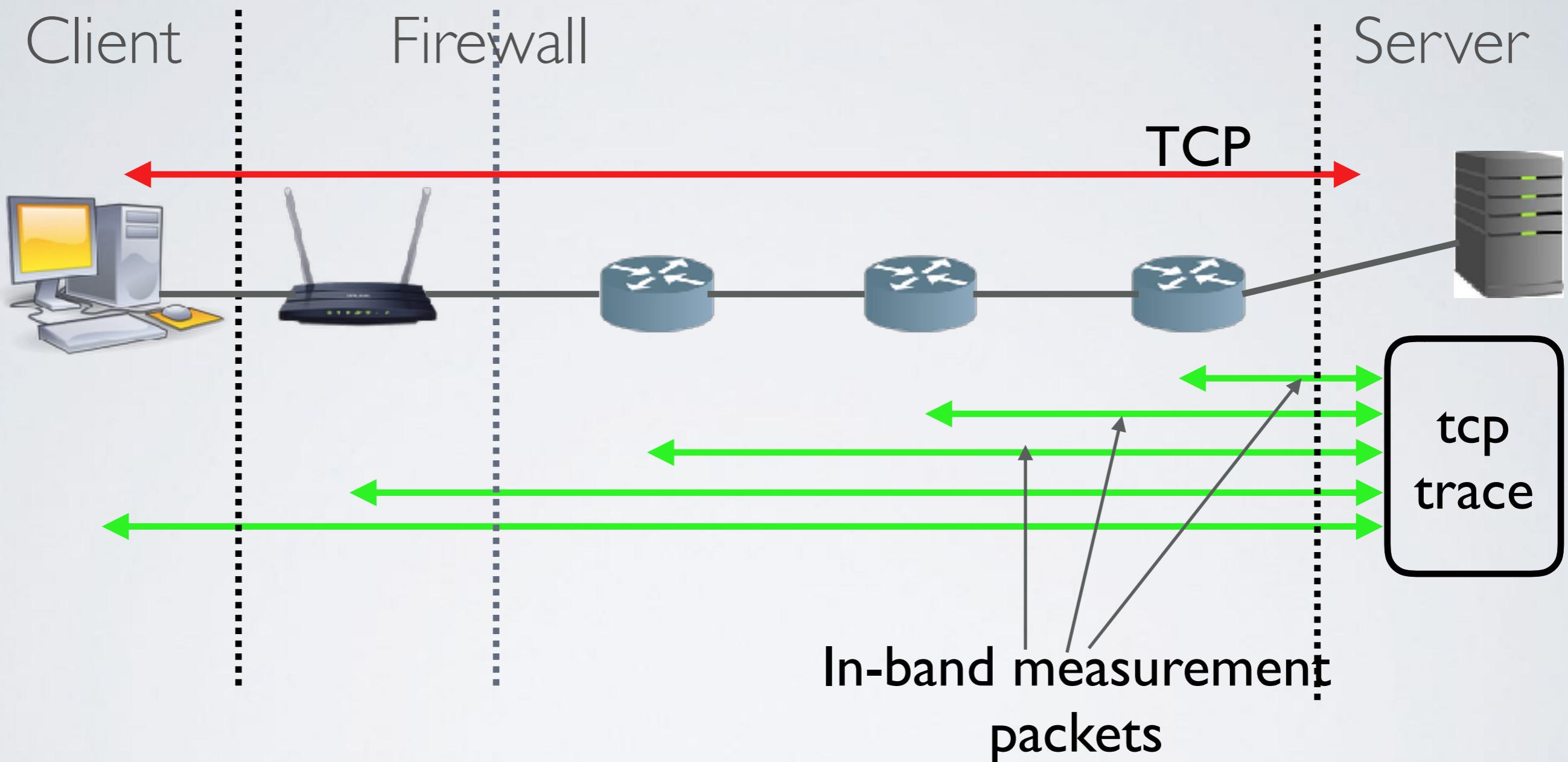
tcptrace



tcptrace



tcptrace



Tracetcp

- **In-band high-frequency traceroute:** injecting empty TTL-limited packets in the ongoing TCP flow
- Ability to observe the buffer building up at bottleneck
- Can measure to the client past the NAT, and observe wireless delays
- prototype at: <https://github.com/ssundaresan/tracetcp>
(ask me for access)



Tracetcp

(more measurements in the works)

- Packet-pair and packet-train techniques to measure per-hop capacity and available bandwidth (in-band pathneck)
- Per-hop loss rate
- Main challenge: how to utilize packets from the TCP stream, and smartly insert measurement packets without affecting the ongoing flow



Thanks!
Questions?
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