

# Linear Programming: Linear Programming Formulations

Daniel Kane

Department of Computer Science and Engineering  
University of California, San Diego

Advanced Algorithms and Complexity  
Data Structures and Algorithms

# Learning Objectives

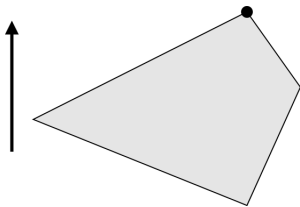
- Distinguish between the different types of linear programming problems.
- Use an algorithm that solves one formulation to solve another formulation.

# Formulations

Several different problem types that all go under the heading of “linear programming”.

# Full Optimization

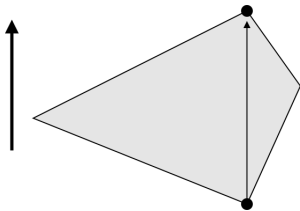
Minimize or maximize a linear function  
subject to a system of linear inequality  
constraints (or say that the constraints have  
no solution).



# Optimization from Starting Point

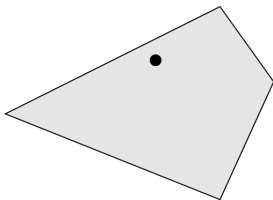
Given a system of linear inequalities and a **vertex of the polytope** they define, optimize a linear function with respect to these constraints.

Given a starting point



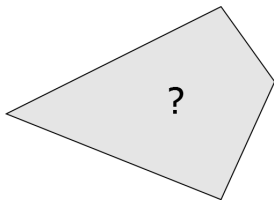
# Solution Finding

Given a system of linear inequalities, find  
some solution. NO objective given



# Satisfiability

Given a system of linear inequalities  
determine whether or not there is a solution.



# Equivalence

Actually, if you can solve any of these problems, you can solve any other!



# Full Optimization

Clearly capable of solving all the other versions.

- Start Opt: Ignore starting point.
- Solution Finding: Optimal is a solution.
- Satisfiability: See if finds a solution.

# Optimization from Starting Point

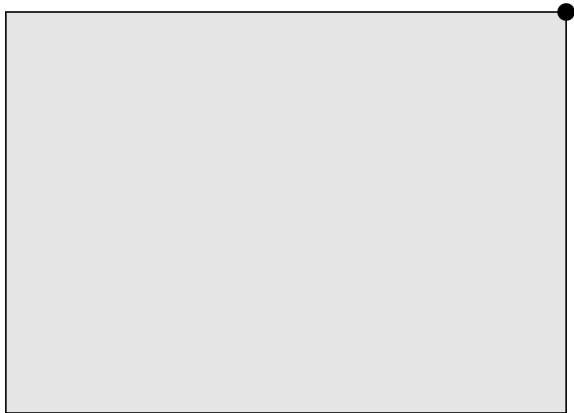
Going from other direction if you can only solve optimization from a starting point you can actually do the full optimization. Hence we have to first find starting point.

- How do you find starting point?

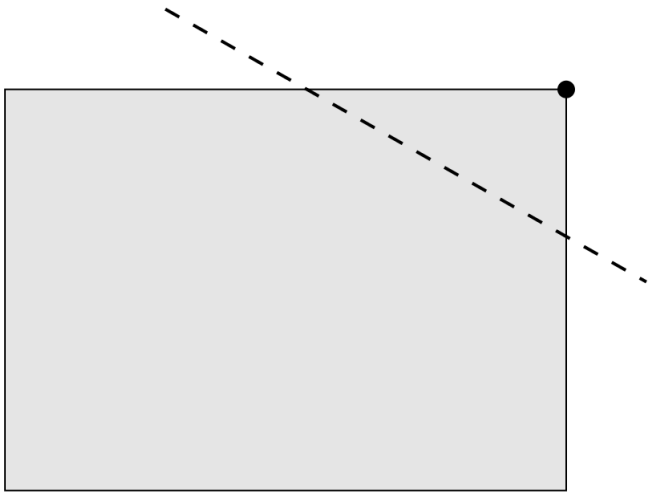
# Optimization from Starting Point

- How do you find starting point?
- Add equations one at a time.
- Optimize left hand side of next equation.

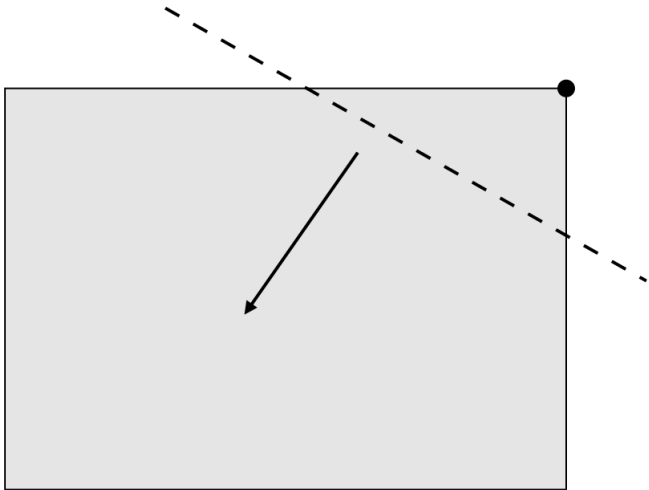
# Example



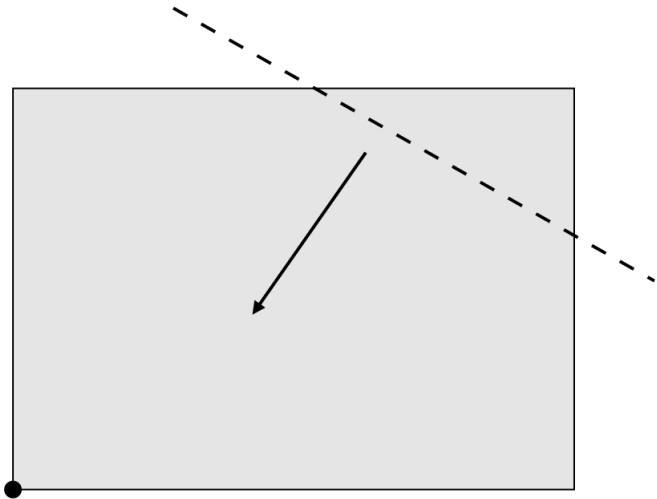
# Example



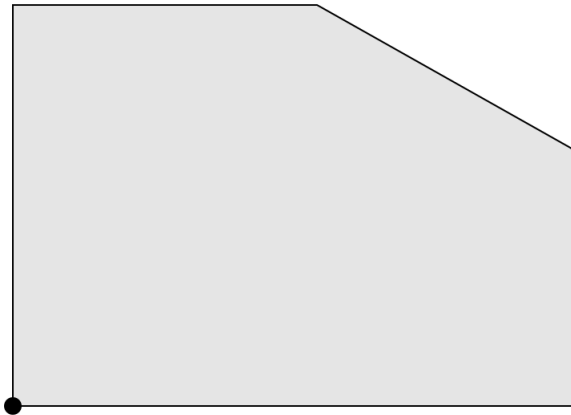
# Example



# Example

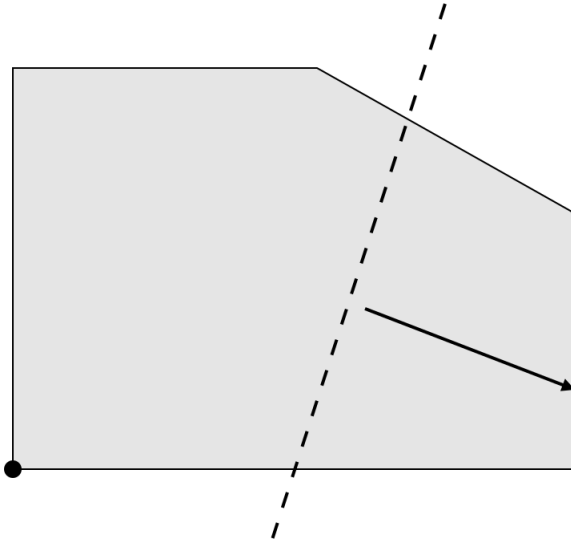


# Example

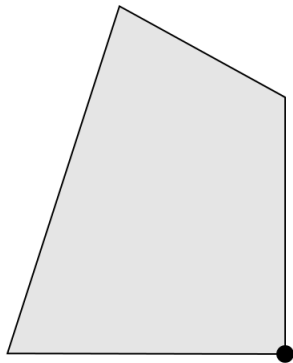




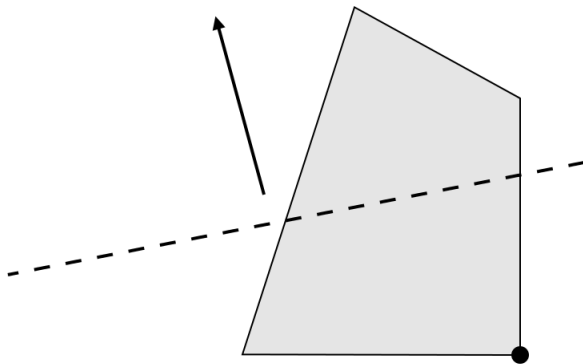
# Example



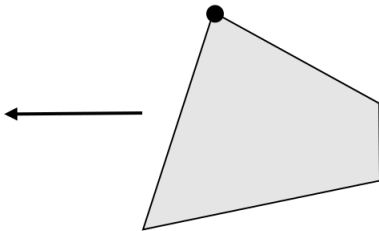
# Example



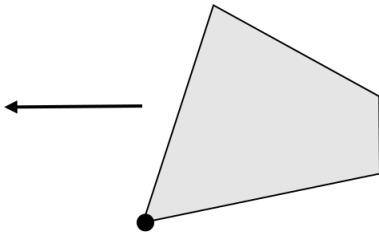
# Example



# Example



# Example



# Technical Point

Things are a bit messier if some intermediate systems don't have optima.

# Technical Point

Things are a bit messier if some intermediate systems don't have optima.

Fix: start with  $n$  constraints (gives a single vertex). Then while trying to add constraint  $v \cdot x \geq t$ , don't just maximize  $v \cdot x$ . Also add  $v \cdot x \leq t$  as a constraint (so that maximum will exist).

# Solution Finding

Q: How do we go from being able to find a solution to finding the best one?



# Solution Finding

Q: How do we go from being able to find a solution to finding the best one?

A: Duality.

# Solution Finding

Q: How do we go from being able to find a solution to finding the best one?

A: Duality. Find a solution and a matching dual solution.

# Setup

Want to minimize  $x \cdot v$  subject to  $Ax \geq b$ .

# Setup

Want to minimize  $x \cdot v$  subject to  $Ax \geq b$ .

Instead find solution to:

$$Ax \geq b$$

$$y \geq 0$$

$$y^T A = v$$

$$x \cdot v = y \cdot b.$$

Will give optimal solution to original problem.

# Satisfiability

How does just **knowing** when you have a solution help?

# Satisfiability

How does just **knowing** when you have a solution help?

Can always find solution at a vertex. Means  $n$  equations are tight.

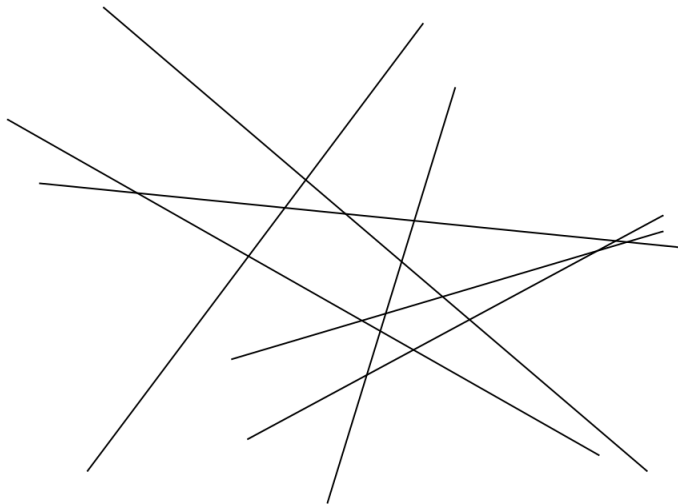
# Satisfiability

How does just **knowing** when you have a solution help?

Can always find solution at a vertex. Means  $n$  equations are tight.

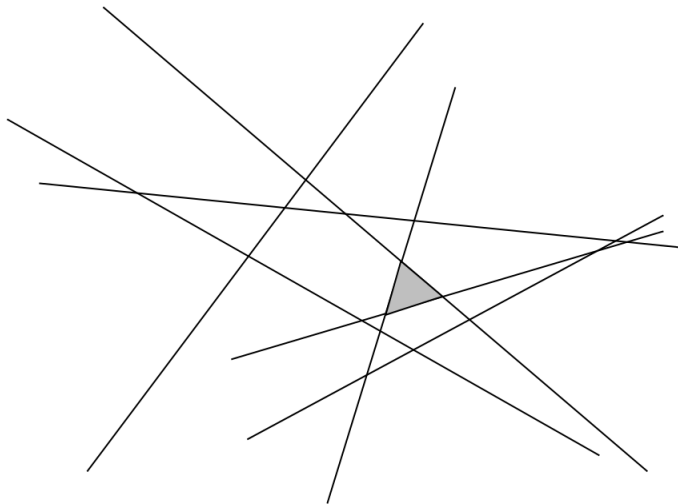
Figure out which equations to use.

# Example

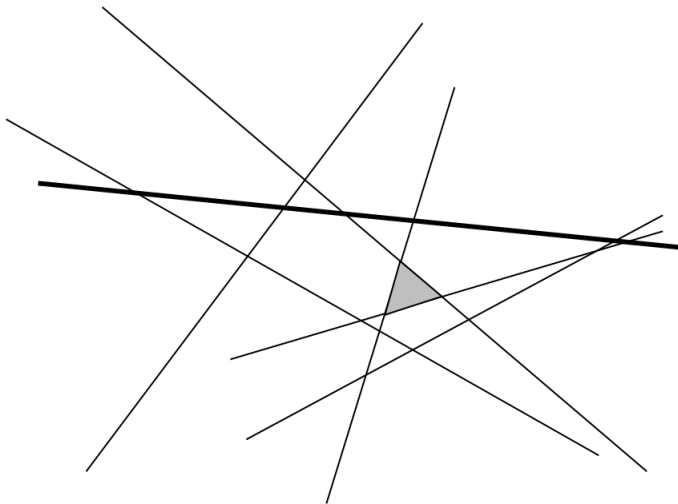




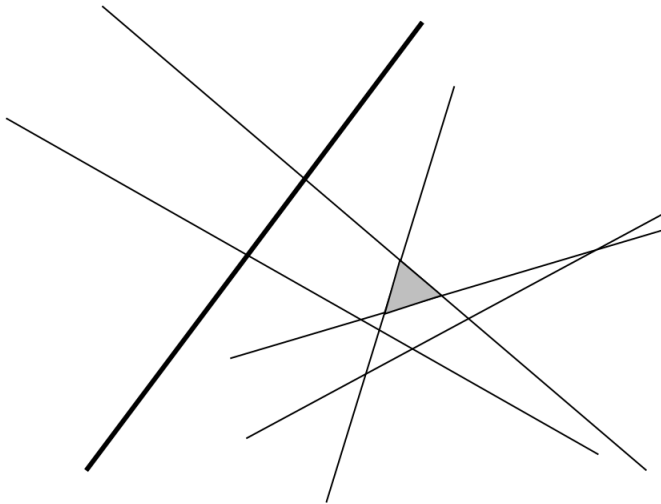
# Example



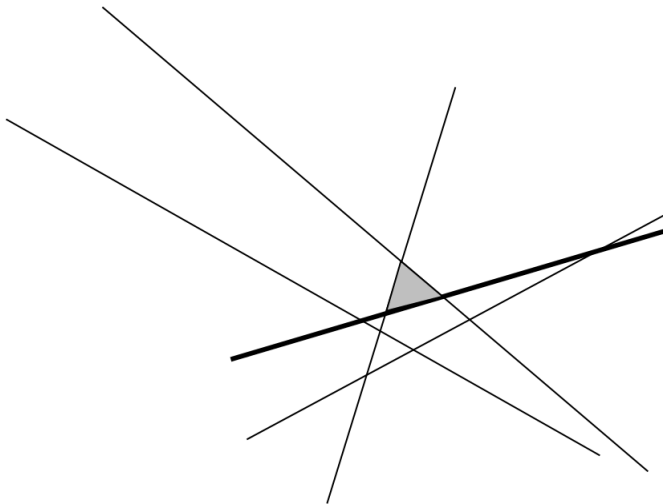
# Example



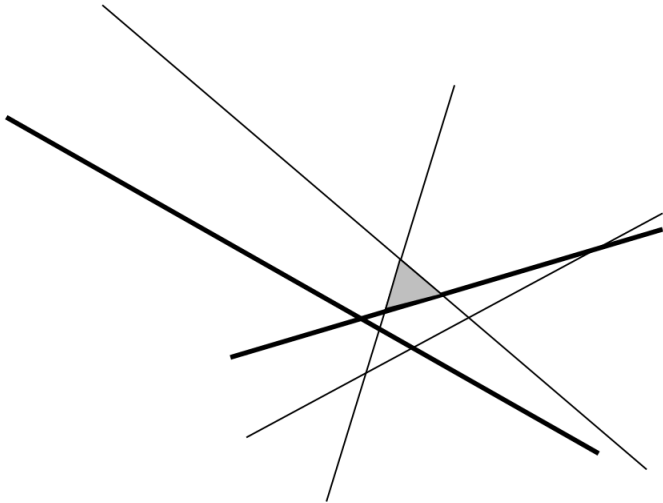
# Example



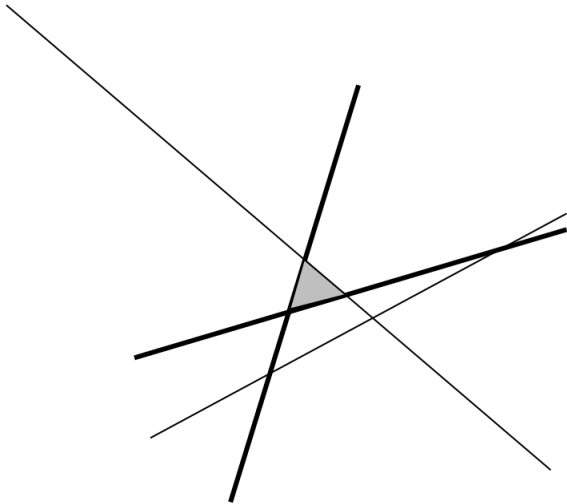
# Example



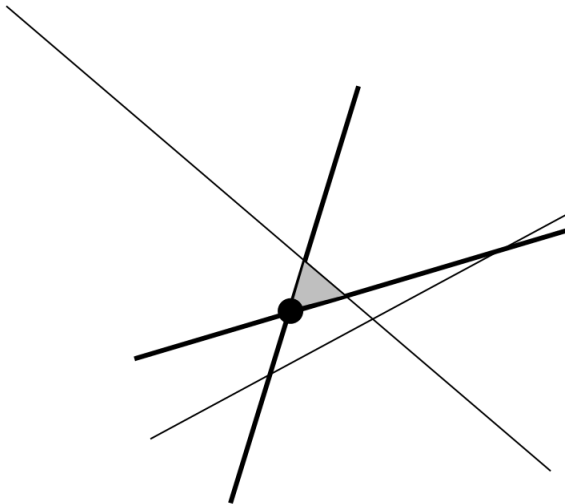
# Example



# Example



# Example



# Problem

In order to find a solution to a linear program with  $m$  equations in  $n$  variables, how many times would one have to call a satisfiability algorithm?



# Solution

In order to find a solution to a linear program with  $m$  equations in  $n$  variables, how many times would one have to call a satisfiability algorithm?

$m$  times. You need to test each equation once, keeping the ones that work.