This different substring search algorithms testing shows dependencies of a string preprocessing and other methods of search optimizing

Processor: 3.4GHz Inter Core i5

Memory: 8 GB 1600 MHz DDR3

System: osx

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Brute Force algorithm associated simply comparing every substring char and pattern char until match

Brute Force is pretty stable alghorithm, works every time with $O(n^2)$ asymptotic

Best case:

0.01811002564288953 S

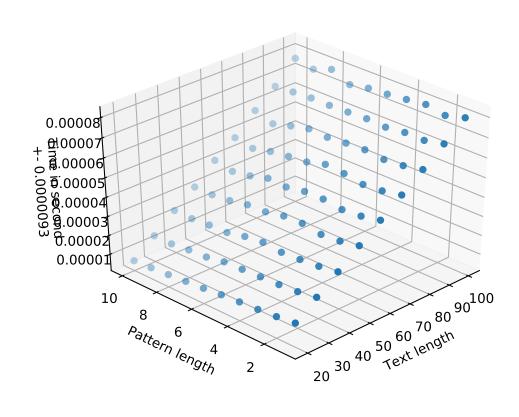
70.5625 B

pattern 'abc
Worst case:

0.021376972400094018 s

70.56640625 B

Brute Force



Hashing algorithm based on summarizing chars ordinals to simplify strings comparison

Hashes have a collision aspect that's why we can see an abrupt behavior change and time growthon a bad substring

pattern 'abc
Best case:

0.018540316711288245 s

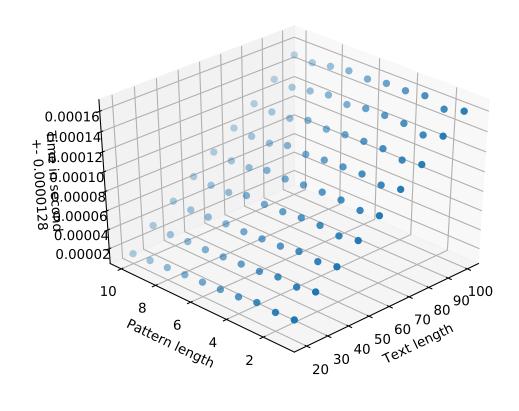
74.78125 B

pattern 'abc
Worst case:

0.03584653500014004 s

74.7265625 B

Hash linear



Hashing algorithm based on summarizing chars ordinals in square to simplify strings comparison

Hashes have a collision aspect that's why we can see an abrupt behavior change and time growthon a bad substring

pattern 'abc
Best case:

0.03533163470025708 s

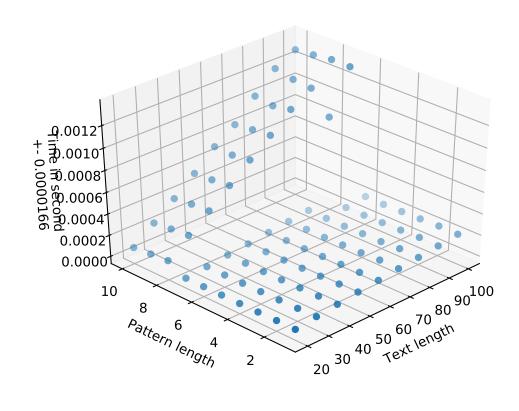
77.09765625 B

pattern 'abc
Worst case:

0.05473548321051088 s

77.09765625 B

Hash quad



Hashing algorithm based on summarizing chars ordinals with special Rabin Karph formula to simplify strings comparison

Hashes have a collision aspect that's why we can see an abrupt behavior change and time growthon a bad substring

pattern 'abc
Best case:

0.0285949156856077 s

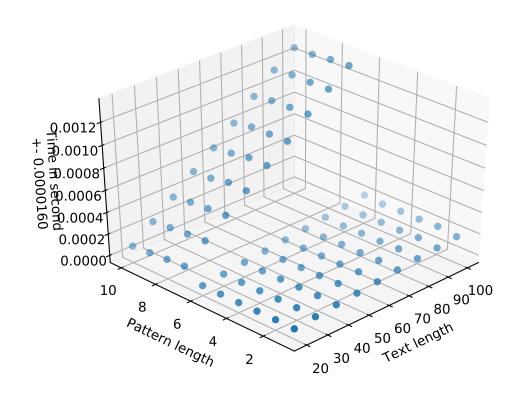
79.46875 B

pattern 'abc
Worst case:

0.03709414020710843 s

79.46875 B

Hash RK



Automate algorithm has a preprocessing component it builds a table of shifts out of pattern which is used during method execution

Automate has a preprocessing aspect that's why we can see time growth

Best case:

0.006662861530232183 S

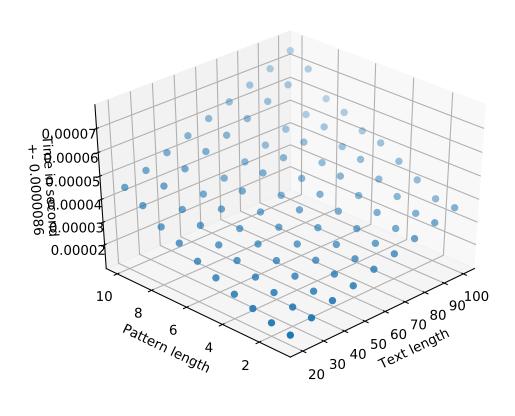
82.22265625 B

pattern 'abc
Worst case:

0.006932610507251608 S

82.22265625 B

Automate



Boyer Moore algorithm preprocess a pattern and builds shift tables for "bad char" anb "good suffix" heuristics

Boyer Moore algorithm works perfectly with prepared pattern on every text but as we can see there is recounting shift tables time delay

pattern 'ABCDABD

Best case:

0.003530097134952324 S

84.3984375 B

text 'ABCDAB ABCDABABCDAB ABCDABCDABABCDAB ABCDABCDABABCDAB ABCDABCDABABCABABC

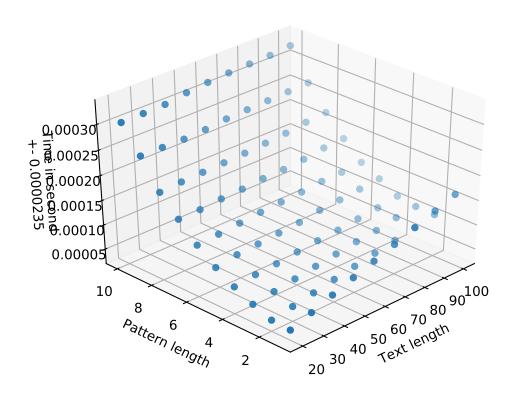
pattern 'ABCDABD

Worst case:

0.039851759153932034 s

84.3984375 B

Boyer Moore



KMP or Knuth-Morris-Pratt's algorithm builds shift table with least common subsequence method then uses it during search

KMP works fine on average string its pretty stable

pattern 'ABCDABD

Best case:

0.0025387278885903005 s

89.01953125 B

text 'ABC ABCDAB ABCDABCDABABC ABCDAB ABCDABCDABABC ABCDAB ABCDABCDABABC ABCDA

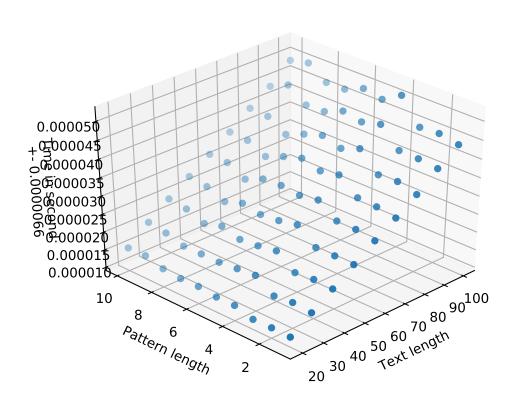
pattern 'ABCDABD

Worst case:

0.09061917733333757 s

89.01953125 B

KMP



Suffix Array algorithm preprocess a text string with making a suffix array then uses a binary search to find left and right answer borders Suffix Array has a text preprocessing that allows to work fast with O(n*log(n)) asymptotic on every pattern

Best case:

0.0060937853912933415 S

133.69921875 B

pattern 'b
Worst case:

0.0345027188695499 s

151.24609375 B

Suffix Array

