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## 1. Monochromatic sets

## 1.1. Ramsey's theorem

**Notation 1.1** N denotes the set of positive integers,  $[n] = \{1, ..., n\}$ , and  $X^{(r)} = \{A \subseteq X : |A| = r\}$ . Elements of a set are written in ascending order, e.g.  $\{i, j\}$  means i < j. Write e.g. ijk to mean the set  $\{i, j, k\}$  with the ordering (unless otherwise stated) i < j < k.

**Definition 1.2** A k-colouring on  $A^{(r)}$  is a function  $c: A^{(r)} \to [k]$ .

## Example 1.3

- Colour  $\{i, j\} \in \mathbb{N}^{(2)}$  red if i + j is even and blue if i + j is odd. Then  $M = 2\mathbb{N}$  is a monochromatic subset.
- Colour  $\{i, j\} \in \mathbb{N}^{(2)}$  red if  $\max\{n \in \mathbb{N} : 2^n \mid (i+j)\}$  is even and blue otherwise.  $M = \{4^n : n \in \mathbb{N}\}$  is a monochromatic subset.
- Colour  $\{i, j\} \in \mathbb{N}^{(2)}$  red if i + j has an even number of distinct prime divisors and blue otherwise. No explicit monochromatic subset is known.

**Theorem 1.4** (Ramsey's Theorem for Pairs) Let  $\mathbb{N}^{(2)}$  are 2-coloured by  $c: \mathbb{N}^{(2)} \to \{1,2\}$ . Then there exists an infinite monochromatic subset M.

## Proof.

- Let  $a_1 \in A_0 := \mathbb{N}$ . There exists an infinite set  $A_1 \subseteq A_0$  such that  $c(a_1, i) = c_1$  for all  $i \in A_1$ .
- Let  $a_2 \in A_1$ . There exists infinite  $A_2 \subseteq A_1$  such that  $c(a_2,i) = c_2$  for all  $i \in A_2$ .
- Repeating this inductively gives a sequence  $a_1 < a_2 < \dots < a_k < \dots$  and  $A_1 \supseteq A_2 \supseteq \dots$  such that  $c(a_i,j) = c_i$  for all  $j \in A_i$ .

- One colour appears infinitely many times:  $c_{i_1} = c_{i_2} = \dots = c_{i_k} = \dots = c$ .
- $M = \{a_{i_1}, a_{i_2}, \ldots\}$  is a monochromatic set.

#### Remark 1.5

- The same proof works for any  $k \in \mathbb{N}$  colours.
- The proof is called a "2-pass proof".
- An alternative proof for k colours is split the k colours 1, ..., k into 2 colours: 1 and "2 or ... or k", and use induction.

Note 1.6 An infinite monochromatic set is very different from an arbitrarily large finite monochromatic set.

**Example 1.7** Let  $A_1 = \{1, 2\}$ ,  $A_2 = \{3, 4, 5\}$ , etc. Let  $\{i, j\}$  be red if  $i, j \in A_k$  for some k. There exist arbitrarily large monochromatic red sets but no infinite monochromatic red sets.

**Example 1.8** Colour  $\{i < j < k\}$  red iff  $i \mid (j + k)$ . A monochromatic subset  $M = \{2^n : n \in \mathbb{N}_0\}$  is a monochromatic set.

**Theorem 1.9** (Ramsey's Theorem for r-sets) Let  $\mathbb{N}^{(r)}$  be finitely coloured. Then there exists a monochromatic infinite set.

## Proof.

- r = 1: use pigeonhole principle.
- r = 2: Ramsey's theorem for pairs.
- For general r, use induction.
- Let  $c: \mathbb{N}^r \to [k]$  be a k-colouring. Let  $a_1 \in \mathbb{N}$ , and consider all r-1 sets of  $\mathbb{N} \setminus \{a_1\}$ , induce colouring  $c': (\mathbb{N} \setminus \{a_1\})^{(r-1)} \to [k]$  via  $c'(F) = c(F \cup \{a_1\})$ .
- By inductive hypothesis, there exists  $A_1 \subseteq \mathbb{N} \setminus \{a_1\}$  such that c' is constant on it (taking value  $c_1$ ).
- Now pick  $a_2 \in A_1$  and induce a colouring  $c': (A_1 \setminus \{a_2\})^{(r-1)} \to [k]$  such that  $c'(F) = c(F \cup \{a_2\})$ . By inductive hypothesis, there exists  $A_2 \subseteq A_1 \setminus \{a_2\}$  such that c' is constant on it (taking value  $c_2$ ).
- Repeating this gives  $a_1, a_2, \ldots$  and  $A_1, A_2, \ldots$  such that  $A_{i+1} \subseteq A_i \setminus \{a_{i+1}\}$  and  $c(F \cup \{a_i\}) = c_i$  for all  $F \subseteq A_{i+1}$ , for |F| = r 1.
- One colour must appear infinitely many times:  $c_{i_1} = c_{i_2} = \dots = c$ .
- $M = \{a_{i_1}, a_{i_2}, ...\}$  is a monochromatic set.

## 1.2. Applications of Ramsey's theorem

Example 1.10 In a totally ordered set, any sequence has monotonic subsequence.

Proof.

- Let  $(x_n)$  be a sequence, colour  $\{i,j\}$  red if  $x_i \leq x_j$  and blue otherwise.
- By Ramsey's theorem for pairs,  $M = \{i_1 < i_2 < \cdots \}$  is monochromatic. If M is red, then the subsequence  $x_{i_1}, x_{i_2}, \ldots$  is increasing, and is strictly decreasing otherwise.
- We can insist that  $(x_{i_j})$  is either concave or convex: 2-colour  $\mathbb{N}^{(3)}$  by colouring  $\{j < k < \ell\}$  red if  $(i, x_{i_j}), (j, x_{i_k}), (k, x_{i_\ell})$  form a convex triple, and blue if they form a concave triple. Then by Ramsey's theorem for r-sets, there is an infinite convex or concave subsequence.

**Theorem 1.11** (Finite Ramsey) Let  $r, m, k \in \mathbb{N}$ . There exists  $n \in \mathbb{N}$  such that whenever  $[n]^{(r)}$  is k-coloured, we can find a monochromatic set of size (at least) m.

## Proof.

- Assume not, i.e.  $\forall n \in \mathbb{N}$ , there exists colouring  $c_n : [n]^{(r)} \to [k]$  with no monochromatic m-sets.
- There are only finitely many (k) ways to k-colour  $[r]^{(r)}$ , so there are infinitely many of colourings  $c_r, c_{r+1}, \ldots$  that agree on  $[r]^{(r)}$ :  $c_i \mid_{[r]^{(r)}} = d_r$  for all i in some infinite set  $A_1$ , where  $d_r$  is a k-colouring of  $[r]^{(r)}$ .
- Similarly,  $[r+1]^{(r)}$  has only finitely many possible k-colourings. So there exists infinite  $A_2 \subseteq A_1$  such that for all  $i \in A_2$ ,  $c_i \mid_{[r+1]^{(r)}} = d_{r+1}$ , where  $d_{r+1}$  is a k-colouring of  $[r+1]^{(r)}$ .
- Continuing this process inductively, we obtain  $A_1 \supseteq A_2 \supseteq \cdots \supseteq A_n$ . There is no monochromatic m-set for any  $d_n : [n]^{(r)} \to [k]$  (because  $d_n = c_i|_{[n]^{(r)}}$  for some i).
- These  $d_n$ 's are nested:  $d_\ell|_{[n]^{(r)}} = d_n$  for  $\ell > n$ .

• Finally, we colour  $\mathbb{N}^{(r)}$  by the colouring  $c: \mathbb{N}^{(r)} \to [k], \ c(F) = d_n(F)$  where  $n = \max(F)$  (or in fact  $n \geq \max(F)$ , which is well-defined by above). So c has no monochromatic m-set (since M was a monochromatic m-set, then taking  $\ell = \max(M), \ d_\ell$  has a monochromatic m-set), which contradicts Ramsey's Theorem for r-sets.

#### Remark 1.12

- This proof gives no bound on n = n(k, m), there are other proofs that give a bound.
- It is a proof by compactness (essentially, we proved that  $\{0,1\}^{\mathbb{N}}$  with the product topology, i.e. the topology derived from the metric  $d(f,g) = \frac{1}{\min\{n \in \mathbb{N}: f(n) \neq g(n)\}}$ , is sequentially compact).

**Remark 1.13** Now consider a colouring  $c: \mathbb{N}^{(2)} \to X$  with X potentially infinite. This does not necessarily admit an infinite monochromatic set, as we could colour each edge a different colour. Such a colouring would be injective. We can't guarantee either the colouring being constant or injective though, as c(ij) = i satisfies neither.

**Theorem 1.14** (Canonical Ramsey) Let  $c: \mathbb{N}^{(2)} \to X$  be a colouring with X an arbitrary set. Then there exists an infinite set  $M \subseteq \mathbb{N}$  such that:

- 1. c is constant on  $M^{(2)}$ , or
- 2. c is injective on  $M^{(2)}$ , or
- 3. c(ij) = c(kl) iff i = k for all i < j and k < l,  $i, j, k, l \in M$ , or
- 4. c(ij) = c(kl) iff j = l for all i < j and  $k < l, i, j, k, l \in M$ .

### Proof (Hints).

- First consider the 2-colouring  $c_1$  of  $\mathbb{N}^{(4)}$  where ijkl is coloured same if c(ij) = c(kl) and DIFF otherwise. Show that an infinite monochromatic set  $M_1 \subseteq \mathbb{N}$  (why does this exist?) coloured same leads to case 1.
- Assume  $M_1$  is coloured DIFF, consider the 2-colouring of  $M_1^{(4)}$ , which colours ijkl SAME if c(il) = c(jk) and DIFF otherwise. Show an infinite monochromatic  $M_2 \subseteq M_1$  (why does this exist?) must be coloured DIFF by contradiction.
- Consider the 2-colouring of  $M_2^{(4)}$  where ijkl is coloured SAME if c(ik) = c(jl) and DIFF otherwise. Show an infinite monochromatic set  $M_3 \subseteq M_2$  (why does this exist?) must be coloured DIFF by contradiction.
- 2-colour  $M_3^{(3)}$  by: ijk is coloured same if c(ij)=c(jk) and DIFF otherwise. Show an infinite monochromatic set  $M_4\subseteq M_3$  (why does this exist) must be coloured DIFF by contradiction.
- 2-colour  $M_4^{(3)}$  by the other two similar colourings to above, obtaining monochromatic  $M_6\subseteq M_5\subseteq M_4$ .
- Consider 4 combinations of these colourings on  $M_6$ , show 3 lead to one of the cases in the theorem, and the other leads to contradiction.

Proof.

- 2-colour  $\mathbb{N}^{(4)}$  by: ijkl is red if c(ij) = c(kl) and blue otherwise. By Ramsey's Theorem for 4-sets, there is an infinite monochromatic set  $M_1 \subseteq \mathbb{N}$  for this
- If  $M_1$  is red, then c is constant on  $M_1^{(2)}$ : for all pairs  $ij, i'j' \in M_1^{(2)}$ , pick m < nwith j, j' < m, then c(ij) = c(mn) = c(i'j').
- So assume  $M_1$  is blue.
- Colour  $M_1^{(4)}$  by giving ijkl colour green if c(il) = c(jk) and purple otherwise. By Ramsey's theorem for 4-sets, there exists an infinite monochromatic  $M_2 \subseteq M_1$  for this colouring.
- Assume  $M_2$  is coloured green: if  $i < j < k < l < m < n \in M_2$ , then c(jk) = c(in) = c(in)c(lm) (consider ijkn and ilmn): contradiction, since  $M_1$  is blue.
- Hence  $M_2$  is purple, i.e. for  $ijkl \in M_2^{(4)}$ ,  $c(il) \neq c(jk)$ .
- Colour  $M_2$  by: ijkl is orange if c(ik) = c(jl), and pink otherwise.
- By Ramsey's theorem for 4-sets, there exists infinite monochromatic  $M_3 \subseteq M_2$  for this colouring.
- Assume  $M_3$  is orange, then for  $i < j < k < l < m < n \in M_3$ , we have c(jm) =c(ln) (consider jlmn) and c(jm) = c(ik) (consider ijkm): contradiction, since  $M_3 \subseteq M_1$ .
- Hence  $M_3$  is pink, i.e. for ijkl,  $c(ik) \neq c(jl)$ .
- Colour  $M_3^{(3)}$  by: ijk is yellow if c(ij) = c(jk) and grey otherwise. By Ramsey's theorem for 3-sets, there exists infinite monochromatic  $M_4\subseteq M_3$  for this colouring.
- Assume  $M_4$  is yellow: then (considering  $ijkl \in M_4^{(4)}$ ) c(ij) = c(jk) = c(kl):
- contradiction, since  $M_4\subseteq M_1$ .

  So for any  $ijk\in M_4^{(3)},\ c(ij)\neq c(jk)$ .

  Finally, colour  $M_4^{(3)}$  by: ijk is gold if c(ij)=c(ik) and c(ik)=c(jk), silver if c(ij) = c(ik) and  $c(ik) \neq c(jk)$ , bronze if  $c(ij) \neq c(ik)$  and c(ik) = c(jk), and platinum if  $c(ij) \neq c(ik)$  and  $c(ik) \neq c(jk)$ .
- By Ramsey's theorem for 3-sets, there exists monochromatic  $M_5 \subseteq M_4$ .  $M_5$  cannot be gold, since then c(ij) = c(jk): contradiction, since  $M_5 \subseteq M_4$ . If silver, then we have case 3 in the theorem. If bronze, then we have case 4 in the theorem. If platinum, then we have case 2 in the theorem.

### Remark 1.15

- A more general result of the above theorem states: let  $\mathbb{N}^{(r)}$  be arbitrarily coloured. Then we can find an infinite M and  $I \subseteq [r]$  such that for all  $x_1...x_r \in M^{(r)}$  and  $y_1...y_r \in M^{(r)}, c(x_1...x_r) = c(y_1...y_r) \text{ iff } x_i = y_i \text{ for all } i \in I.$
- In canonical Ramsey,  $I = \emptyset$  is case 1,  $I = \{1, 2\}$  is case 2,  $I = \{1\}$  is case 3 and  $I = \{2\}$  is case 4.
- These  $2^r$  colourings are called the **canonical colourings** of  $\mathbb{N}^{(r)}$ .

## **Exercise 1.16** Prove the general statement.

## 1.3. Van der Waerden's theorem

Remark 1.17 We want to show that for any 2-colouring of  $\mathbb{N}$ , we can find a monochromatic arithmetic progression of length m for any  $m \in \mathbb{N}$ . By compactness, this is equivalent to showing that for all  $m \in \mathbb{N}$ , there exists  $n \in \mathbb{N}$  such that for any 2-colouring of [n], there exists a monochromatic arithmetic progression of length m. (If not, then for each  $n \in \mathbb{N}$ , there is a colouring  $c_n : [n] \to \{1,2\}$  with no monochromatic arithmetic progression of length m. Infinitely many of these colourings agree on [1], infinitely many of those agreeing in [1] agree on [2], and so on - we obtain a 2-colouring of  $\mathbb{N}$  with no monochromatic arithmetic progression of length m).

We will prove a slightly stronger result: whenever  $\mathbb{N}$  is k-coloured, there exists a length m monochromatic arithmetic progression, i.e. for any  $k, m \in \mathbb{N}$ , there exists  $n \in \mathbb{N}$  such that whenever [n] is k-coloured, we have a length m monochromatic progression.

**Definition 1.18** Let  $A_1, ..., A_k$  be length m arithmetic progressions:  $A_i = \{a_i, a_i + d_i, ..., a_i + (m-1)d_i\}$ .  $A_1, ..., A_k$  are **focussed** at f if  $a_i + md_i = f$  for all i.

**Example 1.19**  $\{4,8\}$  and  $\{6,9\}$  are focussed at 12.

**Definition 1.20** If length m arithmetic progressions  $A_1, ..., A_k$  are focused at f and are monochromatic, each with a different colour (for a given colouring), they are called **colour-focussed** at f.

**Remark 1.21** We use the idea that if  $A_1, ..., A_k$  are colour-focussed at f (for a k-colouring) and of length m-1, then some  $A_i \cup \{f\}$  is a length m monochromatic arithmetic progression.

**Theorem 1.22** Whenever  $\mathbb{N}$  is k-coloured, there exists a monochromatic arithmetic progression of length 3, i.e. for all  $k \in \mathbb{N}$ , there exists  $n \in \mathbb{N}$  such that any k-colouring of [n] admits a length 3 monochromatic progression.

Proof (Hints).

- Prove by induction the claim:  $\forall r \leq k, \exists n \in \mathbb{N}$  such that for any k-colouring of [n], there exists a monochromatic arithmetic progression of length 3, or r colour-focussed arithmetic progressions of length 2.
  - r = 1 case is straightforward.
  - Let claim be true for r-1 with witness n, let  $N=2n(k^{2n}+1)$ .
  - $\triangleright$  Partition N into blocks of equal size, show that two of these blocks must have the same colouring.
  - Using the inductive hypothesis, merge the r-1 colour-focussed arithmetic progressions from these two blocks into a new set of r-1 colour-focussed arithmetic progressions.
  - Find another length 2 monochromatic arithmetic progression, reason that this is of different colour.
- Reason that this claim implies the result.

Proof.

• We claim that for all  $r \leq k$ , there exists an  $n \in \mathbb{N}$  such that if [n] is k-coloured, then either:

- ► There exists a monochromatic arithmetic progression of length 3.
- ightharpoonup There exist r colour-focussed arithmetic progressions of length 2.
- This claim implies the result by the above remark.
- We prove the claim by induction on r:
  - r = 1: take n = k + 1, then by pigeonhole, some two elements of [n] have the same colour, so form a length two arithmetic progression.
  - Assume true for r-1 with witness n. We claim that  $N=2n(k^{2n}+1)$  works for r.
  - ▶ Let  $c : [2n(k^{2n} + 1)] \to [k]$  be a colouring. We partition [N] into  $k^{2n} + 1$  blocks of size 2n:  $B_i = \{2n(i-1) + 1, ..., 2ni\}$  for  $i = 1, ..., k^{2n} + 1$ .
  - Assume there is no length 3 monochromatic progression for c. By inductive hypothesis, each block  $B_i$  has r-1 colour-focussed arithmetic progressions of length 2.
  - Since  $|B_i| = 2n$ , each block also contains their focus. For a set M with |M| = 2n, there are  $k^{2n}$  ways to k-colour M. So by pigeonhole, there are blocks  $B_s$  and  $B_{s+t}$  that have the same colouring.
  - Let  $\{a_i, a_i + d_i\}$  be the r-1 arithmetic progressions in  $B_s$  colour-focussed at f, then  $\{a_i + 2nt, a_i + d_i + 2nt\}$  is the corresponding set of arithmetic progressions in  $B_{s+t}$ , each colour-focussed at f + 2nt.
  - Now  $\{a_i, a_i + d_i + 2nt\}$ ,  $i \in [r-1]$ , are r-1 arithmetic progressions colour-focused at f + 4nt. Also,  $\{f, f + 2nt\}$  is monochromatic of a different colour to the r-1 colours used (since there is no length 3 monochromatic progression for c). Hence, there are r arithmetic progressions of length 2 colour-focussed at f + 4nt.

Remark 1.23 The idea of looking at all possible colourings of a set is called a product argument.

**Definition 1.24** The **Van der Waerden** number W(k, m) is the smallest  $n \in \mathbb{N}$  such that for any k-colouring of [n], there exists a monochromatic arithmetic progression in [n] of length m.

**Remark 1.25** The above theorem gives a **tower-type** upper bound  $W(k,3) \le k^{k^{(\cdot)}k^{4k}}$ .

**Theorem 1.26** (Van der Waerden's Theorem) For all  $k, m \in \mathbb{N}$ , there exists  $n \in \mathbb{N}$  such that for any k-colouring of [n], there is a length m monochromatic arithmetic progression.

Proof (Hints).

• Use induction on m.

- Given induction hypothesis on m-1, prove the claim: for all  $r \leq k$ , there exists  $n \in \mathbb{N}$  such that for any k-colouring of [n], we have either a monochromatic length m arithmetic progression, or r colour-focussed arithmetic progressions of length m-1. Reason that this claim implies the result.
- Use induction on r. Give an explicit n for r = 1.
- Let n be the witness for r-1, let  $N=W(k^{2n},m-1)\cdot 2n$ . Assume a k-colouring of [N],  $c:[N]\to [k]$ , has no arithmetic progressions of length m.
- Partition [N] into the obvious choice of  $W(k^{2n}, m-1)$  blocks  $B_i$ , each of length 2n.
- Colour the indices  $1 \le i \le W(k^{2n}, m-1)$  of the blocks by

$$c'(i) = (c(2n(i-1)+1), c(2n(i-1)+2)...., c(2ni)) \\$$

- Reason that we can find monochromatic arithmetic progression s, s + t, ..., s + (m-2)t of length m-1 (w.r.t c'), and that this corresponds to sequence of blocks  $B_s, B_{s+t}, ..., B_{s+(m-2)t}$ , each identically coloured.
- Reason that  $B_s$  contains r-1 colour-focussed length m-1 arithmetic progressions  $A_i$  together with their focus f.
- Let  $A'_i$  be the same arithmetic progression but with common difference 2nt larger than that of  $A_i$ . Show the  $A'_i$  are colour-focussed at some focus in terms of f.
- Find another length m-1 arithmetic progression, show this must be monochromatic and of different colour to all  $A'_i$ . Show it also has same focus as all  $A'_i$ .

Proof.

• By induction on m. m = 1 is trivial, m = 2 is by pigeonhole principle. m = 3 is the statement of the previous theorem.

- Assume true for m-1 and all  $k \in \mathbb{N}$ .
- For fixed k, we prove the claim: for all  $r \leq k$ , there exists  $n \in \mathbb{N}$  such that for any k-colouring of [n], either:
  - $\rightarrow$  There is a monochromatic arithmetic progression of length m, or
  - ▶ There are r colour-focussed arithmetic progressions of length m-1.
- We will then be done (by considering the focus).
- To prove the claim, we use induction on r.
- r=1 is the claim of the first inductive hypothesis: take n=W(k,m-1).
- Assume the claim holds for r-1 with witness n, and assume there is no monochromatic arithmetic progression of length m. We will show that  $N = W(k^{2n}, m-1)2n$  is sufficient for r.
- Partition [N] into  $W(k^{2n},m-1)$  blocks of length 2n:  $B_i=\{2n(i-1)+1,...,2ni\}$  for  $i=1,...,W(k^{2n},m-1)$ .
- Each block has  $k^{2n}$  possible colourings. Colour the blocks as

$$c'(i) = (c(2n(i-1)+1), c(2n(i-1)+2)...., c(2ni)) \\$$

By definition of W, there exists a monochromatic arithmetic progression of length m-1 (w.r.t. to c'):  $\{\alpha, \alpha+t, ..., \alpha+(m-2)t\}$ . The repsective blocks  $B_{\alpha}, ..., B_{\alpha+(m-2)t}$  are identically coloured.

- $B_{\alpha}$  has length 2n, so by induction  $B_{\alpha}$  contains r-1 colour-focussed arithmetic progressions of length m-1, together with their focus (as length of block is 2n).
- Let  $A_1,...,A_{r-1},$   $A_i=\{a_i,a_i+d_i,...,a_i+(m-2)d_i\},$  be colour-focussed at f.
- Let  $A_i' = \{a_i, a_i + (d_i + 2nt), ..., a_i + (m-2)(d_i + 2nt)\}$  for i = 1, ..., r-1. The  $A_i'$  are monochromatic as the blocks are identically coloured and the  $A_i$  are monochromatic. Also,  $A_i$  and  $A_i'$  have the same colouring, and the  $A_i$  are colour-focussed, hence the  $A_i'$  have pairwise distinct colours.
- The  $A_i$  are focussed at f and the colour of f of different than the colour of all  $A_i$ .  $f = a_i + (m-1)d_i$  for all i.
- Now  $\{f, f+2nt, f+4nt, ..., f+2n(m-2)t\}$  is an arithmetic progression of length m-1, is monochromatic and of a different colour to all the  $A'_i$ .
- It is enough to show that  $a_i + (m-1)(d_i + 2nt) = f + 2n(m-1)t$  for all i, but this is equivalent to  $a_i + (m-1)d_i = f$ , which is true as all  $A_i$  were focussed at f.

Corollary 1.27 For any k-colouring of  $\mathbb{N}$ , there exists a colour class containing arbitrarily long arithmetic progressions.

Remark 1.28 We can't guarantee infinitely long arithmetic progressions, e.g.

- 2-colour  $\mathbb{N}$  by 1 red, 2, 3 blue, 4, 5, 6 red, etc.
- The set of infinite arithmetic progressions in  $\mathbb N$  is countable (since described by two integers: the start term and step). Enumerate them by  $(A_k)_{k \in \mathbb N}$ . Pick  $x_1 < y_1 \in A_1$ , colour  $x_1$  red and  $y_1$  blue. Then pick  $x_2, y_2 \in A_2$  with  $y_1 < x_2 < y_2$ , colour  $x_2$  red,  $y_2$  blue. Continue inductively.

**Theorem 1.29** (Strengthened Van der Waerden) Let  $m, k \in \mathbb{N}$ . There exists  $n \in \mathbb{N}$  such that for any k-colouring of [n], there exists a monochromatic length m arithmetic progression whose common difference is the same colour (i.e. there exists a, a + d, ..., a + (m - 1), d all of the same colour).

 $Proof\ (Hints).$ 

- Use induction on k.
- If n is the witness for k-1 colours, show that N=W(k,n(m-1)+1) is a witness for k colours, by considering n different multiples of the step of a suitable arithmetic progression.

Proof.

- Fix  $m \in \mathbb{N}$ . We use induction on k. k = 1 case is trivial.
- Let n be witness for k-1 colours.
- We will show that N = W(k, n(m-1) + 1) is suitable for k colours.
- If [N] is k-coloured, there exists a monochromatic (say red) arithemtic progression of length n(m-1) + 1: a, a + d, ..., a + n(m-1)d.

- If rd is red for any  $1 \le r \le n$ , then we are done (consider a, a + rd, ..., a + (m 1)rd).
- If not, then  $\{d, 2d, ..., nd\}$  is k-1-coloured, which induces a k-1 colouring on [n]. Therefore, there exists a monochromatic arithmetic progression b, b+s, ..., b+(m-1)s (with s the same colour) by induction, which translates to db, db+ds, ..., db+d(m-1)s and ds being monochromatic.

**Remark 1.30** The case m=2 of strengthened Van der Waerden is **Schur's theorem**: for any k-colouring of  $\mathbb{N}$ , there are monochromatic x, y, z such that x+y=z. This can be proved directly from Ramsey's theorem for pairs: let  $c: \mathbb{N} \to [k]$  be a k-colouring, then induce  $c': \mathbb{N}^{(2)} \to [k]$  by c'(ij) = c(j-i). By Ramsey, there exist i < j < k such that c'(ij) = c'(ik) = c'(jk), i.e. c(j-i) = c(k-i) = c(k-j). So take x = j-i, z = k-i, y = k-j.

## 1.4. The Hales-Jewett theorem

**Definition 1.31** Let X be finite set. We say  $X^n$  consists of words of length n on alphabet X.

**Definition 1.32** Let X be finite. A (combinatorial) line in  $X^n$  is a set  $L \subseteq X^n$  of the form

$$L = \left\{ (x_1,...,x_n) \in X^n : \forall i \not\in I, x_i = a_i \text{ and } \forall i,j \in I, x_i = x_j \right\}$$

for some non-empty set  $I \subseteq [n]$  and  $a_i \in X$  (for each  $i \notin I$ ). I is the set of **active** coordinates for L.

Note that a combinatorial line is invariant under permutations of X.

**Example 1.33** Let X = [3]. Some lines in  $X^2$  are:

- $I = \{1\}$ :  $\{(1,1),(2,1),(3,1)\}$  (with  $a_2 = 1$ ),  $\{(1,2),(2,2),(3,2)\}$  (with  $a_2 = 2$ ),  $\{(1,3),(2,3),(3,3)\}$  (with  $a_2 = 3$ ).
- $I=\{2\}$ :  $\{(1,1),(1,2),(1,3)\}$  (with  $a_1=1$ ),  $\{(2,1),(2,2),(2,3)\}$  (with  $a_1=2$ ),  $\{(3,1),(3,2),(3,3)\}$  (with  $a_1=3$ ).
- $I = \{1, 2\}: \{(1, 1), (2, 2), (3, 3)\}.$

Note that  $\{(1,3),(2,2),(3,1)\}$  is **not** a combinatorial line.

**Example 1.34** Some sets of lines in  $[3]^3$  are:

- $I = \{1\}$ :  $\{(1,2,3), (2,2,3), (3,2,3)\}$  (with  $a_2 = 2, a_3 = 3$ ).
- $I = \{1,3\}$ :  $\{(1,3,1), (2,3,2), (3,3,3)\}$  (with  $a_2 = 3$ ).

**Definition 1.35** In a line L, write  $L^-$  and  $L^+$  for the smallest and largest points in L (with respect to the ordering on  $[m]^n$  where  $x \leq y$  if  $x_i \leq y_i$  for all i).

**Definition 1.36** Lines  $L_1, ..., L_k$  are **focussed** at f if  $L_i^+ = f$  for all  $i \in [k]$ . They are **colour-focussed** if they are focussed and  $L_i \setminus \{L_i^+\}$  is monochromatic for all  $i \in [k]$ , with each  $L_i \setminus \{L_i^+\}$  a different colour.

**Theorem 1.37** (Hales-Jewett) Let  $m, k \in \mathbb{N}$  (we use alphabet X = [m]), then there exists  $n \in \mathbb{N}$  such that for any k-colouring of  $[m]^n$ , there exists a monochromatic combinatorial line.

**Notation 1.38** Denote the smallest such n by HJ(m, k).

 $Proof\ (Hints).$ 

- Induction on m. Prove by induction the claim that for all  $1 \le r \le k$ , there exists  $n \in \mathbb{N}$  such that for any k-colouring of  $[m]^n$ , we have either a monochromatic line, or r colour-focussed lines (reason that this claim implies the result).
- State why claim holds for r = 1.
- Let n be witness for r-1,  $n'=\mathrm{HJ}(m-1,k^{m^n})$ . Want to show that n+n' is witness for r.
- Write  $[m]^{n+n'} = [m]^n \times [m]^{n'}$ .
- For a colouring  $c:[m]^{n+n'} \to [k]$ , induce a suitable colouring  $c':[m]^{n'} \to [k]^{m^n}$  and consider what the definition of n' implies. Use this to induce a colouring  $c'':[m]^n \to [k]$ .
- Using the inductive hypothesis and the previous point, construct r-1 lines in  $[m]^{n+n'}$  which are colour-focussed. Find another line in  $[m]^{n+n'}$  (which should have first n coordinates constant) of different colour which has the same focus point.

*Proof.* By induction on m. The case m=1 is trivial as  $|[m]^n|=1$ . Assume that  $\mathrm{HJ}(m-1,k')$  exists for all  $k'\in\mathbb{N}$ . We claim that for all  $1\leq r\leq k$ , there exists  $n\in\mathbb{N}$  such that for any k-colouring of  $[m]^n$ , we have either:

- a monochromatic line, or
- r colour-focussed lines.

We can then take r = k and consider the focus.

We prove the claim by induction on r. For  $r=1,\ n=\mathrm{HJ}(m-1,k)$  suffices. Let n be a witness for r-1. Let  $n'=\mathrm{HJ}(m-1,k^{m^n})$ . We will show N=n+n' is a witness for r. Let  $c:[m]^N\to [k]$  be a k-colouring with no monochromatic lines. Writing  $[m]^N=[m]^n\times [m]^{n'}$ , colour  $[m]^{n'}$  by  $c':[m]^{n'}\to [k]^{m^n}$ ,  $c'(b)=(c(a_1,b),...,c(a_{m^n},b))$  (where  $[m]^n=\{a_1,...,a_{m^n}\}$ ). By the inductive hypothesis, there exists a line L in  $[m]^{n'}$  with active coordinates I such that

$$\forall a \in [m]^n, \forall b, b' \in L \setminus \{L^+\}, \quad c(a, b) = c(a, b').$$

But now this induces a (well-defined) colouring  $c'': [m]^n \to [k], c''(a) = c(a, b)$  for any  $b \in L \setminus \{L^+\}$ . By definition of n, there exist r-1 lines  $L_1, ..., L_{r-1}$  colour-focussed (w.r.t c'') at f, with active coordinates  $I_1, ..., I_{r-1}$ .

Finally, consider the r-1 lines  $L_i'$ ,  $1 \le i \le r-1$  in  $[m]^N$  that start at  $(L_i^-, L^-)$  with active coordinates  $I_i \cup I$ , and the line L' in  $[m]^N$  that starts at  $(f, L^-)$  with active coordinates I. By the construction of c'', the colour of each point in  $L_i'$  is determined by the first n coordinates which form a point lying in  $L_i$ . Hence, since the  $L_i$  are

colour-focussed, the  $L'_i$  are colour-focussed. As for L', the first n coordinates are constant (always equal to f), and so again by the construction of c'', the colour of each point in L' is equal to c''(f), which is a different colour to each colour of the  $L'_i$ . Hence all  $L'_1, ..., L'_{r-1}, L'$  colour-focussed at  $(f, L^+)$ , so we are done.

Corollary 1.39 Hales-Jewett implies Van der Waerden's theorem.

*Proof (Hints)*. For a colouring  $c: \mathbb{N} \to [k]$ , consider the induced colouring  $c'(x_1, ..., x_n) = c(x_1 + \cdots + x_n)$  of  $[m]^n$ .

*Proof.* Let c be a k-colouring of  $\mathbb{N}$ . For sufficiently large n (i.e.  $n \geq \mathrm{HJ}(m,k)$ ), induce a k-colouring c' of  $[m]^n$  by  $c'(x_1,...,x_n) = c(x_1+\cdots+x_n)$ . By Hales-Jewett, a monochromatic (with respect to c') combinatorial line L exists. This gives a monochromatic (with respect to c) length m arithmetic progression in  $\mathbb{N}$ . The step is equal to the number of active coordinates. The first term in the arithmetic progression corresponds to the point in L with all active coordinates equal to 1, the last term corresponds to the point in L with all active coordinates equal to m.

**Exercise 1.40** Show that the m-in-a-row noughts and crosses game cannot be a draw in sufficiently high dimensions, and that the first player can always win.

**Definition 1.41** A *d*-dimensional subspace (or *d*-point parameter set)  $S \subseteq X^n$  is a set such that there exist pairwise disjoint  $I_1, ..., I_d \subseteq [n]$  and  $a_i \in X$  for all  $i \in [n] - (I_1 \cup \cdots \cup I_d)$ , such that

$$S = \big\{x \in X^n : x_i = a_i \quad \forall i \in [n] - (I_1 \cup \dots \cup I_d),$$
 and  $x_i = x_j \quad \forall i, j \in I_k \text{ for some } k \in [d]\big\}.$ 

**Example 1.42** Two 2-dimensional subspaces in  $X^3$  are  $\{(x,y,2): x,y \in X\}$   $(I_1=\{1\},I_2=\{2\})$  and  $\{(x,x,y): x,y \in X\}$   $(I_1=\{1,2\},I_2=\{3\}).$ 

**Theorem 1.43** (Extended Hales-Jewett) For all  $m, k, d \in \mathbb{N}$ , there exists  $n \in \mathbb{N}$  such that for any colouring of  $[m]^n$ , there exists a monochromatic d-dimensional subspace.

*Proof (Hints)*. Use Hales-Jewett on  $m^d$  and k.

*Proof.* We can view  $X^{dn'}$  as  $\left(X^d\right)^{n'}$ . A line in  $\left(X^d\right)^{n'}$  (on alphabet  $Y=X^d$ ) corresponds to a d-dimensional subspace in  $X^{dn'}$  (on alphabet X). (Each inactive coordinate in the line corresponds to d adjacent inactive coordinates in the subspace, and each active coordinate in the line corresponds to d adjacent active coordinates in the subspace). Hence, we can take  $n=d\cdot \mathrm{HJ}(m^d,k)$ .

**Definition 1.44** Let  $S \subseteq \mathbb{N}^d$  be finite. A **homothetic copy** of S is a set of the form  $a + \lambda S$  where  $a \in \mathbb{N}^d$  and  $\lambda \in \mathbb{N}$   $(l \neq 0)$ .

**Theorem 1.45** (Gallai) Let  $S \subseteq \mathbb{N}^d$  be finite. For every k-colouring of  $\mathbb{N}^d$ , there exists a monochromatic homothetic copy of S.

Proof (Hints). Let  $S = \{S_1, ..., S_m\}$ , consider colouring  $c' : [m]^n \to [k]$  (for suitable n) given by  $c'(x_1, ..., x_n) = c(S_{x_1}, ..., S_{x_m})$ .

Proof. Let  $S = \{S_1, ..., S_m\}$ . Let  $c: \mathbb{N}^d \to [k]$  be a k-colouring. For n large enough (i.e.  $n \geq \mathrm{HJ}(m,k)$ ), colour  $[m]^n$  by  $c'(x_1, ..., x_n) = c \left(S_{x_1} + \cdots + S_{x_m}\right)$ . By Hales-Jewett, there exists a monochromatic line (with respect to c') in  $[m]^n$  with active coordinates I. So  $c \left(\sum_{i \notin I} S_i + |I|S_j\right)$  is the same colour for all  $j \in [m]$ . So we are done, as  $\sum_{i \notin I} S_i + |I|S$  is a homothetic copy of S.

#### Remark 1.46

- Gallai's theorem can also be proven with a focusing + product colouring argument.
- For  $S = \{(x, y) \in \mathbb{N}^2 : x, y \in \{1, 2\}\}$ , Gallai's theorem proves the existence of a monochromatic square whereas extended Hales-Jewett only guarantees a monochromatic rectangle.

## 2. Partition regular systems

## 2.1. Rado's theorem

Strengthened Van der Waerden says that the system  $x_1 + x_2 = y_1, x_1 + 2x_2 = y_2, ..., x_1 + mx_2 = y_m$  has a monochromatic solution in  $x_1, x_2, y_1, ..., y_m$ . We want to find when a general system of equations is partition regular.

**Definition 2.1** Let  $A \in \mathbb{Q}^{m \times n}$  be a  $m \times n$  matrix. A is **partition regular (PR)** if for any finite colouring of  $\mathbb{N}$ , there exists a monochromatic  $x \in \mathbb{N}^n$  such that Ax = 0.

## Example 2.2

- Schur's theorem says that x + y = z has a monochromatic solution for any finite colouring of  $\mathbb{N}$ , and so that (1, 1, -1) is PR.
- Strengthened Van der Waerden states that

$$\begin{bmatrix} 1 & 1 & -1 & 0 & \dots & 0 \\ 1 & 2 & 0 & -1 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & m & 0 & 0 & \dots & -1 \end{bmatrix}$$

is PR.

- (a, b, -(a + b)) is PR for any a, b (a monochromatic solution is x = y = z).
- (2,-1) is not PR: colour N by n is red if  $\max\{m \in \mathbb{N} : 2^m \mid n\}$  is even, and blue otherwise. Then if 2x = y, x and y must have different colours.

**Definition 2.3** A rational matrix A with columns  $c_1, ..., c_n \in \mathbb{Q}^m$  has the **column property** (CP) if there exists a partition  $B_1 \sqcup \cdots \sqcup B_r$  of [n] such that:

- 1.  $\sum_{i \in B_1} c_i = 0$ .
- 2. For all  $s \in \{2, ..., r\}$ ,  $\sum_{i \in B_s} c_i \in \text{span}\{c_j : j \in B_1 \sqcup \cdots \sqcup B_{s-1}\}$  (note we can take the linear span over  $\mathbb{R}$  or over  $\mathbb{Q}$  here, as if a rational vector is a real linear combination of rational vectors, then it is also a rational linear combination of them).

#### Example 2.4

- (1,1,-1) has CP, with  $B_1=\{1,3\},\,B_2=\{2\}.$
- The matrix

$$\begin{bmatrix} 1 & 1 & -1 & 0 & \dots & 0 \\ 1 & 2 & 0 & -1 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & m & 0 & 0 & \dots & -1 \end{bmatrix}$$

from Strengthened Van der Waerden has CP, with  $B_1 = \{1, 3, ..., n\}$  and  $B_2 = \{2\}$ .

- (3, 4, -7) has CP with  $B_1 = \{1, 2, 3\}$ .
- $(\lambda, -1)$  has CP iff  $\lambda = 1$ .
- (3,4,-6) doesn't have CP.

### Example 2.5

$$\begin{bmatrix} 1 & -1 & 3 \\ 2 & -2 & a \\ 4 & -4 & b \end{bmatrix}$$

has CP iff (a, b) = (6, 12).

Remark 2.6  $x = (a_1, ..., a_n)$  is PR iff  $\lambda x$  is PR (for any  $\lambda \in \mathbb{Q}^{\times}$ ), so we can assume that each  $a_i \in \mathbb{Z}$ . Also, x has CP iff there exists  $\emptyset \neq I \subseteq [n]$  such that  $\sum_{i \in I} a_i = 0$ . We may also assume WLOG each  $a_i \neq 0$ . We will first show that if x is PR, then it has CP. Even in the  $1 \times n$  matrix case of Rado's theorem, neither direction is easy.

**Notation 2.7** For p prime and  $x = (a_k...a_0)_p \in \mathbb{N}$ , write e(x) for the rightmost non-zero digit in the base-p expansion of x, i.e.  $e(x) = a_{t(x)}$ , where  $t(x) = \min\{i : a_i \neq 0\}$ .

**Proposition 2.8** Let  $a_1, ..., a_n \in \mathbb{Q}^*$ . If  $(a_1, ..., a_n)$  is PR, then it has CP.

*Proof* (*Hints*). For p large enough (determine later a bound for p), colour  $\mathbb{N}$  by giving x colour e(x), and consider  $\min\{t(x_1),...,t(x_n)\}$ .

Proof. Let p be a large prime  $(p > \sum_{i=1}^n |a_i|)$ . Define a (p-1)-colouring of  $\mathbb N$  giving x colour e(x). By assumption, there are  $x_1,...,x_n$  of the same colour d such that  $\sum_{i=1}^n a_i x_i = 0$ . Let  $t = \min\{t(x_1),...,t(x_n)\}$ , and let  $I = \{i \in [n]: t(x_i) = t\}$  (note I is non-empty). So when summing  $\sum_{i=1}^n a_i x_i = 0$  and considering the last digit in the base p expansion, we have  $\sum_{i=1}^n a_i x_i = 0 \mod p^{t+1}$  and so obtain  $\sum_{i \in I} a_i d = 0 \mod p$ , so  $\sum_{i \in I} a_i = 0$  (since p is prime and was chosen large enough).

Remark 2.9 There is no other known proof of this proposition.

**Lemma 2.10** Let  $\lambda \in \mathbb{Q}$ . Then  $(1, \lambda, -1)$  is partition regular, i.e. for any finite colouring of  $\mathbb{N}$ , there exists monochromatic  $(x, y, z) \in \mathbb{N}^3$  such that  $x + \lambda y = z$ .

Proof (Hints).

- Reason that we can assume  $\lambda > 0$ . Write  $\lambda = r/s, r, s \in \mathbb{N}$ .
- Use induction on number of colours k: given n such that any (k-1)-colouring of [n] admits monochromatic solution, show that N = W(k, nr + 1)ns works for k colours, by considering the definition of W and isd for each  $i \in [n]$ .

*Proof.* The case  $\lambda = 0$  is trivial, and if  $\lambda < 0$ , we may rewrite the equation as  $z - \lambda y = x$ , so we may assume that  $\lambda > 0$ , so let  $\lambda = \frac{r}{s}$  for  $r, s \in \mathbb{N}$ . In fact, we show that for any k-colouring of [n] (for some n depending on k), there is a monochromatic solution.

We seek a monochromatic solution to  $x + \frac{r}{s}y = z$  for some finite colouring  $c : \mathbb{N} \to [k]$ . We use induction on the number of colours k. For k = 1,  $n = \max\{s, r+1\}$  is sufficient, with monochromatic solution (1, s, r+1). Assume n is a witness for k-1 colours. We will show N = nsW(k, nr+1) is suitable for k colours. By definition of W, given a k-colouring of [N], there is a monochromatic AP inside  $[W(k, nr+1)] \subseteq [N]$  of length nr+1: a, a+d, ..., a+nrd, coloured red.

Consider isd for each  $i \in [n]$ . Note that  $isd \leq nsW(k, nr + 1)$  so each isd does indeed have a colour. If some isd is also red, then (a, isd, a + ird) is a monochromatic solution. If no isd is red, then  $\{sd, ..., nsd\}$  is (k-1)-coloured, so by the inductive hypothesis, there exists  $i, j, k \in [n]$  such that  $\{isd, jsd, ksd\}$  is monochromatic and  $isd + \lambda jsd = ksd$ , so (isd, jsd, ksd) is a monochromatic solution.

#### Remark 2.11

- Note the similarity to the proof of Strengthened Van der Waerden.
- The case  $\lambda = 1$  is Schur's theorem, which can be proven directly by Ramsey's theorem; however, there is no known proof using Ramsey's theorem for general  $\lambda \in \mathbb{Q}$ .

**Theorem 2.12** (Rado's Theorem for Single Equations) Let  $a_1, ..., a_n \in \mathbb{Q} \setminus \{0\}$ .  $(a_1, ..., a_n)$  is PR iff it has CP.

*Proof (Hints)*. For  $\Leftarrow$ : for the obvious choice of  $I \subseteq [n]$ , fix  $i_0 \in I$ , and define  $x \in \mathbb{N}^n$  componentwise:

$$x_i = \begin{cases} x & \text{if } i = i_0 \\ y & \text{if } i \notin I \\ z & \text{if } i \in I \setminus \{i_0\} \end{cases}.$$

Show that x is a solution to  $\sum_{i=1}^{n} a_i x_i = 0$ .

*Proof.*  $\Longrightarrow$  is by Proposition 2.8. For  $\Longleftrightarrow$ : we have that  $\sum_{i \in I} a_i = 0$  for some  $\emptyset \neq I \subseteq [n]$ . Given a colouring  $c : \mathbb{N} \to [k]$ , we need to show that there are monochromatic  $x_1, ..., x_n$  such that  $\sum_{i=1}^n a_i x_i = 0$ .

Fix  $i_0 \in I$ . We construct the following vector  $x \in \mathbb{N}^n$  by defining its components:

$$x_i = \begin{cases} x & \text{if } i = i_0 \\ y & \text{if } i \notin I \\ z & \text{if } i \in I \setminus \{i_0\} \end{cases}$$

for some fixed suitable x, y, z. We need x, y, z to be monochromatic and

$$\begin{split} a_{i_0}x + \sum_{i \notin I} a_i y + \sum_{i \in I \backslash \{i_0\}} a_i z &= 0 \\ \iff a_{i_0}x - z a_{i_0} + \sum_{i \notin I} a_i y &= 0 \\ \iff x + \frac{\sum_{i \notin I} a_i}{a_{i_0}} y - z &= 0 \end{split}$$

and this holds, since x, y, z exist by the above lemma.

**Conjecture 2.13** (Rado's Boundedness Conjecture) Let A be an  $m \times n$  matrix that is not PR (so there exists a "bad" colouring, i.e. a k-colouring with no monochromatic solution to Ax = 0 for some  $k \in \mathbb{N}$ ). Is k bounded (for given m, n)?

This is known for  $1 \times 3$  matrices: 24 colours suffice.

**Proposition 2.14** Let  $A \in \mathbb{Q}^{m \times n}$ . If A is PR, then it has CP.

Proof (Hints).

- Let  $x \in \mathbb{N}^n$  be the monochromatic solution to Ax = 0. For fixed prime p, partition [n] into  $B_1, ..., B_r$  by grouping  $i, j \in [n]$  by  $t(x_i), t(x_j)$  (and preserving the ordering).
- Reason that the same partition exists for infinitely many p.
- Considering  $\sum_{i=1}^n x_i c_i = 0 \mod p$  for infinitely many p, show that  $\sum_{i \in B_1} c_i = 0$ , and

$$p^t \sum_{i \in B_k} \boldsymbol{c}_i + \sum_{i \in B_1, \dots, B_{k-1}} x_i d^{-1} \boldsymbol{c}_i \equiv \boldsymbol{0} \operatorname{mod} p^{t+1}.$$

• By taking the dot product with  $u \in \mathbb{N}^m$  for appropriate u, show by contradiction that  $\sum_{i \in B_k} c_i \in \text{span}\{c_i : i \in B_1, ..., B_{k-1}\}.$ 

*Proof.* Let  $c_1, ..., c_n \in \mathbb{Q}^m$  be the columns of A. For fixed prime p, colour  $\mathbb{N}$  as before by c(x) = e(x). By assumption, there exists a monochromatic  $x \in \mathbb{N}^n$  such that  $\sum_{i=1}^n x_i c_i = \mathbf{0}$ . We partition the columns (by partitioning  $[n] = B_1 \sqcup \cdots \sqcup B_r$ ) as follows:

- $i, j \in B_k$  iff  $t(x_i) = t(x_j)$ .
- $i \in B_k$ ,  $j \in B_\ell$  for  $k < \ell$  iff  $t(x_i) < t(x_j)$ .

We do this for infinitely many primes p. Since there are finitely many partitions of [n], for infinitely many p, we will have the same blocks  $B_1, ..., B_r$ .

Consider  $\sum_{i=1}^n x_i c_i = \mathbf{0}$  performed in base p. Each  $i \in [n]$  has the same colour  $d = e(x_i) \in [1, p-1]$ . So  $\sum_{i \in B_1} dc_i = 0 \mod p$  (by collecting the rightmost terms in base p), hence  $\sum_{i \in B_1} c_i = 0 \mod p$ . But this holds for infinitely many p, hence

$$\sum_{i \in B_1} c_i = 0.$$

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Now 
$$\sum_{i \in B_k} p^t d\mathbf{c}_i + \sum_{i \in B_1, \dots, B_{k-1}} x_i \mathbf{c}_i = \mathbf{0} \mod p^{t+1}$$
 for some  $t$ . So 
$$p^t \sum_{i \in B_k} \mathbf{c}_i + \sum_{i \in B_1, \dots, B_{k-1}} x_i d^{-1} \mathbf{c}_i \equiv \mathbf{0} \mod p^{t+1}.$$

We claim that  $\sum_{i \in B_k} c_i \in \operatorname{span}\{c_i : i \in B_1, ..., B_{k-1}\}$ . Suppose not, then there exists  $u \in \mathbb{N}^m$  such that  $u.c_i = 0$  for all  $i \in B_1, ..., B_{k-1}$ , but  $u.\left(\sum_{i \in B_k} c_i\right) \neq 0$ . Then dotting with u, we obtain  $p^t u.\left(\sum_{i \in B_k} c_i\right) \equiv 0 \operatorname{mod} p^{t+1}$ , so  $u.\sum_{i \in B_k} c_i \equiv 0 \operatorname{mod} p$ . But this holds for infinitely many p, so  $u.\sum_{i \in B_k} c_i = 0$ : contradiction.

**Definition 2.15** For  $m, p, c \in \mathbb{N}$ , an (m, p, c)-set  $S \subseteq \mathbb{N}$  with generators  $x_1, ..., x_m \in \mathbb{N}$  is of the form

$$S = \left\{ \sum_{i=1}^m \lambda_i x_i : \exists j \in [m] : \lambda_j = c, \lambda_i = 0 \ \forall i < j, \text{and} \ \lambda_k \in [-p,p] \ \forall k > j \right\}$$

where  $[-p, p] = \{-p, -(p-1), ..., p\}$ . So S consists of

$$\begin{split} cx_1 + \lambda_2 x_2 + \lambda_3 x_3 + \cdots + \lambda_m x_m, & \lambda_i \in [-p, p], \\ cx_2 + \lambda_3 x_3 + \cdots + \lambda_m x_m, & \lambda_i \in [-p, p], \\ & \vdots \\ cx_m. \end{split}$$

These are the **rows** of S. We can think of S as a "progression of progressions".

## Example 2.16

- A (2, p, 1)-set with generators  $x_1, x_2$  is of the form  $\{x_1 px_2, x_1 (p 1)x_2, ..., x_1 + px_2, x_2\}$ , so is an AP of length 2p + 1 together with its step.
- A (2, p, 3)-set with generators  $x_1, x_2$  is of the form  $\{3x_1 px_2, 3x_1 (p 1)x_2, ..., 3x_1, ..., 3x_1 + px_2, 3x_2\}$ , so is an AP of length 2p + 1, whose middle term is divisible by 3, together with three times its step.

**Theorem 2.17** Let  $m, p, c \in \mathbb{N}$ . For any finite colouring of  $\mathbb{N}$ , there exists a monochromatic (m, p, c)-set.

 $Proof\ (Hints).$ 

- Reason that an (m', p, c)-set contains an (m, p, c)-set for  $m' \ge m$ . With M = k(m-1) + 1, reason that if we can find an (M, p, c)-set with each row monochromatic, then we can find an monochromatic (m, p, c)-set.
- Let  $A_1 = \{c, 2c, ..., \lfloor n/c \rfloor c\}$ , reason that  $A_1$  contains a set of the form  $R_1 = \{cx_1 n_1d_1, cx_1 (n_1 1)d_1, ..., cx_1 + n_1d_1\}$  for some large  $n_1$ .
- Let  $B_1 = \left\{d_1, 2d_1, ..., \left\lfloor \frac{n_1}{(M-1)p} \right\rfloor d_1\right\}$ . We have  $cx_1 + \lambda_1 b_1 + \cdots + \lambda_{M-1} b_{M-1} \in R_1$ , explain why these are monochromatic.
- Inside  $B_1$ , define

$$A_2 = \bigg\{cd_1, 2cd_1, ..., \bigg\lfloor \frac{n_1}{(M-1)pc} \bigg\rfloor cd_1 \bigg\}.$$

and apply the argument as before, where the divisor in the  $\lfloor \cdot \rfloor$  expression in the new  $B_2$  is (M-2)p.

• Argue that after a certain number of steps, we have formed an (M, p, c)-set with each row monochromatic.

*Proof.* Let  $c: \mathbb{N} \to [k]$  be the colouring of  $\mathbb{N}$  with k colours. Note that an (m', p, c)-set with  $m' \geq m$  contains an (m, p, c)-set (by taking any m rows, and setting some suitable  $\lambda_i$  to 0). Let M = k(m-1) + 1. It is enough to find a (M, p, c)-set such that each row is monochromatic.

Let n be large (large enough to apply the argument that follows). Let  $A_1 = \{c, 2c, ..., \lfloor n/c \rfloor c\}$ . By Van der Waerden,  $A_1$  contains a monochromatic AP  $R_1$  of length  $2n_1 + 1$  where  $n_1$  is large enough:

$$R_1 = \{cx_1 - n_1d_1, cx_1 - (n_1 - 1)d_1, ..., cx_1 + n_1d_1\}.$$

has colour  $k_1$ . Now we restrict our attention to

$$B_{1} = \bigg\{d_{1}, 2d_{1}, ..., \bigg\lfloor \frac{n_{1}}{(M-1)p} \bigg\rfloor d_{1} \bigg\}.$$

Observe that

$$cx_1 + \lambda_1 b_1 + \dots + \lambda_{M-1} b_{M-1} \in R_1$$

for all  $\lambda_i \in [-p, p]$  and  $b_i \in B_1$ , so all these sums have colour  $k_1$ . Inside  $B_1$ , look at

$$A_2 = \bigg\{cd_1, 2cd_1, ..., \left\lfloor \frac{n_1}{(M-1)pc} \right\rfloor cd_1 \bigg\}.$$

By Van der Waerden,  $A_2$  contains a monochromatic AP  $R_2$  of length  $2n_2+1$  with colour  $k_2$ :

$$R_2=\{cx_2-n_2d_2,cx_2-(n_2-1)d_2,...,cx_2+n_2d_2\}.$$

Note that  $x_2 \subseteq B_1$ . Now we restrict our attention to

$$B_{2}=\left\{ d_{2},2d_{2},...,\left|\frac{n_{2}}{(M-2)p}\right|d_{2}\right\} .$$

Again, note that for all  $\lambda_i \in [-p, p]$  and  $b_i \in B_2$ , we have

$$cx_2+\lambda_1b_1+\cdots+\lambda_{M-2}b_{M-2}\in R_2$$

so has colour  $k_2$ .

We iterate this process M times, and obtain M generators  $x_1,...,x_M$  such that each row of the (M,p,c)-set generated by  $x_1,...,x_M$  is monochromatic. But now M=k(m-1)+1, so m of the rows have the same colour.

**Remark 2.18** Being extremely precise in this proofs (such as considering  $\lfloor \cdot \rfloor$ ) is much less important than the ideas in the proof. (Won't be penalised in the exam for small details like this).

Corollary 2.19 (Folkman's Theorem) Let  $m \in \mathbb{N}$  be fixed. For every finite colouring of  $\mathbb{N}$ , there exists  $x_1, ..., x_m \in \mathbb{N}$  such that

$$\mathrm{FS}(x_1,...,x_m) \coloneqq \left\{ \sum_{i \in I} x_i : \emptyset \neq I \subseteq [m] \right\}$$

is monochromatic.

*Proof (Hints)*. A specific case of Theorem 2.17.

*Proof.* By the (m, 1, 1) case of Theorem 2.17.

#### Remark 2.20

- The case n=2 of Folkman's theorem is Schur's theorem.
- For a colouring  $c: \mathbb{N} \to [k]$ , we induce a colouring  $c': \mathbb{N} \to [k]$  by  $c'(n) = c(2^n)$ . Then by Folkman's theorem for c', there exists  $x_1, ..., x_m$  such that

$$\mathrm{FP}(x_1,...,x_m) = \Bigg\{ \prod_{i \in I} x_i : \emptyset \neq I \subseteq [m] \Bigg\}.$$

• It is not known whether the same result holds for  $\mathrm{FS}(x_1,...,x_m) \cup \mathrm{FP}(x_1,...,x_m)$ . However, it does not hold for infinite sets  $\{x_n:n\in\mathbb{N}\}$ , and does hold for colourings of  $\mathbb{Q}$ .

**Proposition 2.21** Let A have CP. Then there exist  $m, p, c \in \mathbb{N}$  such that every (m, p, c)-set contains a solution y to Ay = 0, i.e. all  $y_i$  belong to the (m, p, c)-set.

*Proof.* Let  $c_1, ..., c_n$  be the columns of A. By assumption, there is a partition  $B_1 \sqcup ... \sqcup B_r$  of [n] such that  $\forall k \in [r]$ ,

$$\begin{split} &\sum_{i \in B_k} c_i \in \operatorname{span}\{c_i :\in B_1 \cup \dots \cup B_{k-1}\} \\ &\Longrightarrow \sum_{i \in B_k} c_i = \sum_{i \in B_1 \cup \dots \cup B_{k-1}} q_{ik} c_i \quad \text{for some } q_{ik} \in \mathbb{Q} \\ &\Longrightarrow \sum_{i=1}^n d_{ik} c_i = \mathbf{0} \end{split}$$

where

$$d_{ik} = \begin{cases} 0 & \text{if } i \notin B_1 \cup \dots \cup B_{k-1} \\ 1 & \text{if } i \in B_k \\ -q_{ik} & \text{if } i \in B_1 \cup \dots \cup B_{k-1} \end{cases}.$$

Take m=r. Let  $x_1,...,x_r\in\mathbb{N}$ , and let  $y_i=\sum_{k=1}^r d_{ik}x_k$  for each  $i\in[n]$ . Now  $\boldsymbol{y}=(y_1,...,y_n)$  is a solution to  $A\boldsymbol{y}=\boldsymbol{0}$ : we have

$$egin{aligned} \sum_{i=1}^n y_i c_i &= \sum_{i=1}^n \sum_{k=1}^r d_{ik} x_k c_i \ &= \sum_{k=1}^r x_k \sum_{i=1}^n d_{ik} c_i = \mathbf{0}. \end{aligned}$$

Let c be the LCD of all the  $q_{ik}$ . Now  $cy_i = \sum_{k=1}^n cd_{ik}x_k$  is an integral linear combination of the  $x_k$ , and cy is a solution since y is. Let p be c times maximum of the absolute values of the numberators of the  $q_{ik}$ . By definition of the  $d_{ik}$ , cy is in the (m, p, c)-set generated by  $x_1, ..., x_r$ .

**Theorem 2.22** (Rado)  $A \in \mathbb{Q}^{m \times n}$  is PR iff it has CP.

*Proof.*  $\Longrightarrow$  is by Proposition 2.14. For  $\longleftarrow$ , let  $c': \mathbb{N} \to [k]$  be a finite colouring of  $\mathbb{N}$ . Also, by the above proposition, since A has CP, there exists  $m, p, c \in \mathbb{N}$  such that  $Ax = \mathbf{0}$  has a solution x in any (m, p, c)-set by the above theorem. By Theorem 2.17, there is a monochromatic (m, p, c)-set with respect to c'. This gives a monochromatic solution x to  $Ax = \mathbf{0}$ .

**Remark 2.23** From the proof of Rado's Theorem, we obtain that if A is PR for the "mod p" colourings, then it is PR for any colouring. There is no proof of this fact that is more direct than using Rado's theorem.

**Theorem 2.24** (Consistency) Let A and B be rational PR matrices. Then the matrix

$$\begin{bmatrix} A & 0 \\ 0 & B \end{bmatrix}$$

is PR.

Proof (Hints). Rado's Theorem.

*Proof.* This is a trivial check of the CP given the CP of A and B, then we are done by Rado's Theorem.

Remark 2.25 The Consistency Theorem says that if we can find monochromatic solutions x and x' to Ax = 0 and By = 0, then we can find monochromatic solutions x' and y', of the same colour, to Ax' = 0 and By' = 0.

**Theorem 2.26** For any finite colouring of  $\mathbb{N}$ , some colour class contains solutions to all PR equations.

*Proof (Hints)*. Use the Consistency Theorem.

*Proof.* For a given k-colouring of  $\mathbb{N}$ , let  $\mathbb{N} = C_1 \sqcup \cdots \sqcup C_k$  be the colour classes. Assume the contrary, so for each  $1 \leq i \leq k$ , there exists a PR matrix  $A_i$  such that  $A_i \mathbf{x} = \mathbf{0}$  has no monochromatic solution of the same colour as  $C_i$ . But then by inductively applying the consistency theorem, the matrix

$$\begin{bmatrix} A_1 & 0 \\ & \ddots \\ 0 & A_k \end{bmatrix}$$

has a monochromatic solution of the same colour as some  $C_j$ . But then  $C_j x = 0$  has a solution x of the same colour as  $C_j$ : contradiction.

## 2.2. Ultrafilters

**Theorem 2.27** (Hindman) For any finite colouring of  $\mathbb{N}$ , there exists a sequence  $(x_n)_{n\in\mathbb{N}}$  such that

$$\mathrm{FS}(\{x_n:n\in\mathbb{N}\}) = \Bigg\{\sum_{i\in I}: I\subseteq\mathbb{N} \text{ finite}, I\neq\emptyset\Bigg\}.$$

**Definition 2.28** A filter on  $\mathbb{N}$  is a non-empty collection  $\mathcal{F}$  of subsets of  $\mathbb{N}$  such that:

- ∅ ∉ ℱ,
- If  $A \in \mathcal{F}$  and  $A \subseteq B$ , then  $B \in \mathcal{F}$ , i.e.  $\mathcal{F}$  is an **up-set**.
- If  $A, B \in \mathcal{F}$  then  $A \cap B \in \mathcal{F}$ , i.e.  $\mathcal{F}$  is closed under finite intersections.

A filter is a notion of "large" subsets of  $\mathbb{N}$ .

## Example 2.29

- $\mathcal{F}_1 = \{A \subseteq \mathbb{N} : 1 \in A\}$  is a filter.
- $\mathcal{F}_2 = \{A \subseteq \mathbb{N} : 1, 2 \in A\}$  is a filter.
- $\mathcal{F}_3 = \{A \subseteq \mathbb{N} : A^c \text{ finite}\}\$ is a filter, called the **cofinite filter**.
- $\mathcal{F}_4 = \{A \subseteq \mathbb{N} : A \text{ infinite}\}\$ is not a filter, since it contains  $2\mathbb{N}$  and  $2\mathbb{N}+1$  but not  $\emptyset = (2\mathbb{N}) \cap (2\mathbb{N}+1)$ .
- $\mathcal{F}_5 = \{A \subseteq \mathbb{N} : 2\mathbb{N} \setminus A \text{ finite}\}$  is a filter.

**Definition 2.30** An ultrafilter is a maximal filter.

**Definition 2.31** For  $x \in \mathbb{N}$ , the principal ultrafilter at x is

$$\mathcal{U}_x\coloneqq \{A\subseteq \mathbb{N}: x\in A\}.$$

**Proposition 2.32** The principal ultrafilter at x is indeed an ultrafilter.

*Proof (Hints)*. Straightforward.

*Proof.* If  $B \notin \mathcal{U}_x$ , then  $x \in B^c$  so  $B^c \in \mathcal{U}_x$ , but  $B^c \cap B = \emptyset$ , so  $\mathcal{U}_x \cup \{B\}$  is not a filter.

#### Example 2.33

- $\mathcal{F}_1 = \{A \subseteq \mathbb{N} : 1 \in A\}$  is an ultrafilter.
- $\mathcal{F}_2=\{A\subseteq\mathbb{N}:1,2\in A\}$  is not an ultrafilter as  $\mathcal{F}_1$  extends it.
- $\mathcal{F}_3 = \{A \subseteq \mathbb{N} : A^c \text{ finite}\}\$ is not an ultra filter, as  $\mathcal{F}_5$  extends it.
- $\mathcal{F}_5 = \{A \subseteq \mathbb{N} : 2\mathbb{N} \setminus A \text{ finite}\}\$  is not an ultrafilter, as  $\{A \subseteq \mathbb{N} : 4\mathbb{N} \setminus A \text{ finite}\}\$  extends it.

**Proposition 2.34** A filter  $\mathcal{F}$  is an ultrafilter iff for all  $A \subseteq \mathbb{N}$ , either  $A \in \mathcal{F}$  or  $A^c \in \mathcal{F}$  $\mathcal{F}$ . *Proof* (*Hints*).  $\Leftarrow$ : straightforward.  $\Rightarrow$ : show if  $A \notin \mathcal{F}$ , then  $\exists B \in \mathcal{F}$  such that  $A \cap$  $B = \emptyset$ . *Proof.*  $\Leftarrow$ : since  $A \cap A^c = \emptyset \notin \mathcal{F}$ .  $\implies$ : let  $\mathcal{F}$  is an ultrafilter. We cannot have  $A, A^c \in \mathcal{F}$  as  $A \cap A^c = \emptyset \notin \mathcal{F}$ . Suppose there is  $A \subseteq \mathbb{N}$  such that  $A, A^c \notin \mathcal{F}$ . By maximality of  $\mathcal{F}$ , since  $A \notin \mathcal{F}$ , then  $\exists B \in \mathcal{F}$ such that  $A \cap B = \emptyset$  (suppose not, then  $\mathcal{F}' = \{S \subseteq \mathbb{N} : S \supseteq A \cap B \text{ for some } B \in \mathcal{F}\}$ extends  $\mathcal{F}$ ). Similarly,  $\exists C \in \mathcal{F}$  such that  $A^c \cap C = \emptyset$ . So we have  $C \subseteq A$ , so  $B \cap C = \emptyset$  $\emptyset \notin \mathcal{F}$ : contradiction (or also  $C \subseteq A \Longrightarrow A \in \mathcal{F}$ : contradiction). Corollary 2.35 Let  $\mathcal{U}$  be an ultrafilter and  $A = B \cup C \in \mathcal{U}$ . Then  $B \in U$  or  $C \in \mathcal{U}$ . *Proof (Hints)*. Straightforward. *Proof.* If not, then  $B^c, C^c \in \mathcal{U}$  by Proposition 2.34, hence  $B^c \cap C^c = (B \cup C)^c =$  $A^c \in \mathcal{U}$ : contradiction. **Proposition 2.36** Every filter is contained in an ultrafilter. *Proof (Hints)*. Use Zorn's Lemma. *Proof.* By Zorn's Lemma, it is enough to show that every non-empty chain of filters has an upper bound. Let  $\{\mathcal{F}_i : i \in I\}$  be a chain of filters, and set  $\mathcal{F} = \bigcup_{i \in I} \mathcal{F}_i$ . •  $\emptyset \notin \mathcal{F}$  since  $\emptyset \notin \mathcal{F}_i$  for each  $i \in I$ . • If  $A \in \mathcal{F}$  and  $A \subseteq B$ , then  $A \in \mathcal{F}_i$  for some  $i \in I$ , so  $B \in \mathcal{F}_i$ , so  $B \in \mathcal{F}$ . • Let  $A,B\in\mathcal{F},$  so  $A\in\mathcal{F}_i$  and  $B\in\mathcal{F}_j$  for some i,j. WLOG,  $\mathcal{F}_i\subseteq\mathcal{F}_j,$  so  $A\cap B\in\mathcal{F}_i$  $\mathcal{F}_i$ , so  $A \cap B \in \mathcal{F}$ .  $\mathcal{F}$  is an upper bound for the chain, so we are done. **Proposition 2.37** Let  $\mathcal{U}$  be an ultrafilter. Then  $\mathcal{U}$  is non-principal iff  $\mathcal{U}$  extends the cofinite filter  $\mathcal{F}_C$ .  $Proof\ (Hints). \iff$  straightforward.  $\implies$  use Corollary 2.35. *Proof.*  $\Leftarrow$ : if  $\mathcal{U} = \mathcal{U}_x$  is principal, then we have  $\{x\} \in \mathcal{U}$  so  $\{x\}^c \notin \mathcal{U}$  by <u>Proposition 2.34</u>, but also  $\{x\}^c \in \mathcal{F}_C$ : contradiction.

**Notation 2.38** Let  $\beta \mathbb{N}$  denote the set of all ultrafilters on  $\mathbb{N}$ .

is in  $\mathcal{U}$ , so  $\mathcal{U}$  is principal: contradiction.

**Definition 2.39** Define a topology on  $\beta\mathbb{N}$  by its base (basis), which consists of

 $\implies$ : let  $A \in \mathcal{F}_C$ , so  $A^c = \{a_1, ..., a_k\}$  is finite. Assume  $A \notin \mathcal{U}$ , then  $A^c \in \mathcal{U}$ , so by Corollary 2.35, some  $a_i \in \mathcal{U}$ . But then by definition of a filter, each set containing  $a_i$ 

$$C_{\Lambda} := \{ \mathcal{U} \in \beta \mathbb{N} : A \in \mathcal{U} \}$$

for each  $A\subseteq\mathbb{N}$ . The sets above indeed form a base: we have  $\bigcup_{A\subseteq\mathbb{N}}C_A=\beta\mathbb{N}$ , and  $C_A\cap C_B=C_{A\cap B}$ , since  $A\cap B\in\mathcal{U}$  iff  $A,B\in\mathcal{U}$ . The open sets are of the form  $\bigcup_{i\in I}C_{A_i}$  and the closed sets are of the form  $\bigcap_{i\in I}C_{A_i}$ .

Remark 2.40  $\beta \mathbb{N} \setminus C_A = C_{A^c}$ , since  $A \notin \mathcal{U}$  iff  $A^c \in \mathcal{U}$ . We can view  $\mathbb{N}$  as being embedded in  $\beta \mathbb{N}$  by identifying  $n \in \mathbb{N}$  with  $\tilde{n} := \mathcal{U}_n$ , the principal ultrafilter at n. Each point in  $\mathbb{N}$  under this correspondence is isolated in  $\beta \mathbb{N}$ , since  $C_{\{n\}} = \{\tilde{n}\}$  is an open neighbourhood of  $\tilde{n}$ . Also,  $\mathbb{N}$  is dense in  $\beta \mathbb{N}$ , since for every  $n \in A$ ,  $\tilde{n} \in C_A$ , so every non-empty open set in  $\beta \mathbb{N}$  intersects  $\mathbb{N}$ .

**Theorem 2.41**  $\beta \mathbb{N}$  is a compact Hausdorff topological space.

# 3. Euclidean Ramsey theory