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Natural Language Processing: Assignment 1

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Question I

Question I, Part a)

So we want to show that:

$$abla^2 f(x) = egin{bmatrix}
abla (e_1
abla f(x)^T) \\ & \dots \\
abla (e_n
abla f(x)^T) \end{bmatrix}.$$

The LHS (Left Hand Side) can be written as:

$$\mathsf{LHS} = \nabla^2 f(\boldsymbol{x}) = \begin{bmatrix} f_{x_1 x_1} & f_{x_1 x_2} & \dots & f_{x_1 x_n} \\ f_{x_2} f_{x_1} & f_{x_2 x_2} & \dots & f_{x_2 x_n} \\ \dots & \dots & \dots & \dots \\ f_{x_n x_1} & f_{x_n x_2} & \dots & f_{x_n x_n} \end{bmatrix}.$$

I.e $[\nabla^2 f(x)]_{ij} = f_{x_i x_j}$ Now we consider the i^{th} row of the RHS (Right Hand Side):

$$\nabla (e_i \nabla f(x)^T)$$
.

Recall that $\nabla f(\mathbf{x}) = [f_{x_1}, f_{x_2}, ..., f_{x_n}] \implies e_i \nabla f(\mathbf{x})^T = f_{x_i}$

$$\implies i^{\text{ith}}$$
 row of RHS = ∇f_{x_i} .

So we can define $g(x) := f_{x_i}$ and it follows that:

$$\mathsf{RHS} = \nabla g(\boldsymbol{x}) = \begin{bmatrix} g_{\mathsf{x}_1}, & g_{\mathsf{x}_2}, & ..., & g_{\mathsf{x}_n} \end{bmatrix}.$$

$$\implies j^{\text{th}} \text{column of RHS} = g_{x_j} = \frac{\partial}{\partial x_j} (f_{x_i}) = f_{x_j x_i}.$$

So then if we assume continuity of the second order derivatives, we can apply *Schwarzes Theorem*, which gives symmetry of second order partial derivatives to see that

$$[\mathsf{LHS}]_{ij} = [\mathsf{LHS}]_{ji} = [\mathsf{RHS}]_{ij}.$$

QED.

It is obvious by the above mentioned symmetry of second order partial derivatives that the Hessian matrix is symmetric, thus it is clear that:

$$i^{\text{th}}$$
column of Hessian = $(\nabla(e_i \nabla f(x)^T))^T$.

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Question I, Part b)

From the identity above we can derive the following algorithm for computing the hessian:

- 1. **STEP I:** Use backpropagation on f to find it's **symbolic** first order derivatives.
- 2. **STEP 2:** For each f_{x_i} first order (symbolic) partial derivative, again apply the backpropagation algorithm (this time with forward propogation and numerical outputs) to get the i^{th} row of the hessian.

Since we have n first order partial derivatives, which we can evaluate in O(m) time for which we need to draw the corresponding computational graph and re-apply backpropogation (also in O(m) time), it follows that computing the Hessian with this method will have a time complexity of O(n*m). This is clearly more efficient than calculating the Hessian the "naive way" i.e entry by entry, which would clearly have a runtime complexity of $O(n^2*m)$.

Question I, Part c)

For the third order tensor, we can just apply the same idea, with

$$\nabla (\mathbf{e}_i \nabla (\mathbf{e}_i \nabla f(\mathbf{x})^T)^T).$$

If we visualize the 3rd order tensor as a 3D grid of values, then this will give us a vertical "column" of values, with

$$[\nabla (\mathbf{e}_i \nabla (\mathbf{e}_i \nabla f(\mathbf{x})^T)^T)]_k = f_{\mathsf{x}_i \mathsf{x}_i \mathsf{x}_k}.$$

Therefore per "vertical column" we will have a runtime of O(m) and there are n^2 vertical columns so therefore we will have a runtime complexity of $O(n^2 * m)$.

In general it is clear that if we apply this method to get the k^{th} order tensor then we will have a runtime of $O(n^{k-1}*m)$.

Question 2

https://colab.research.google.com/drive/100qHG-bIPm1uSX3v47s3x1rCIj46vc40?usp=sharing