Final Project

Isaac Schaal - CS 144

Question 1

Can you conclude that A = B if A, B, and C are sets such that:

- a. $A \cup C = B \cup C$?
- **b**. $A \cap C = B \cap C$?
- c. $A \cup C = B \cup C$ and $A \cap C = B \cap C$?

a)

Given $A \cup C = B \cup C$, we can't conclude that A = B. We will show this by providing a counterexample. Let $A = \{1, 2\}$, $B = \{1\}$, and $C = \{2, 3\}$. $A \cup C = \{1, 2, 3\} = B \cup C$. However, $A \neq B$.

b)

Given $A \cap C = B \cap C$, we can't conclude that A = B. We will show this by providing a counterexample. Let $A = \{1, 2, 3\}$, $B = \{3, 4, 5\}$, and $C = \{3, 6, 7\}$. $A \cap C = \{3\} = B \cap C$. However, $A \neq B$.

c)

Given $A \cup C = B \cup C$ and $A \cap C = B \cap C$, we can conclude that A = B. In order to prove that A = B, we must prove that $A \subseteq B$ and $B \subseteq A$. We will first prove that $A \subseteq B$. For any element $x \in A$, there are two cases, either $x \in C$ or $x \in C$. If $x \in C$, then $x \in A \cap C = B \cap C$, so $x \in B$. Now consider $x \in C$. Since $x \in A$, we have $x \in A \cup C = B \cup C$. Since $x \in C$ and $x \in B \cup C$, we have $x \in B$. Thus, for any $x \in A$, we have $x \in B$, which means that $A \subseteq B$. The proof that $B \subseteq A$ procedes in exactly the same way.

Extension

This question is an example of adding to our toolbox of set theory. If we are trying to prove set equality, we previously had to show that $A \subseteq B$ and $B \subseteq A$. However, we have just proved that if $A \cup C = B \cup C$ and $A \cap C = B \cap C$, we can conclude that A = B. If we are given two sets, A and B, and we know that $A \cup C = B \cup C$, we only have to prove that $A \cap C = B \cap C$ in order to prove set equality (and vice versa). Other similar problems in set theory give us similar tools, where we can use "shortcuts" to solve further problems.

In order to make this question more difficult, we can pose the following generalization.

Can you conclude that $A_1 = A_2 = ... = A_n$ if $A_1, A_2, ... A_n$ and C are sets such that:

d.
$$A_1 \cup C = A_2 \cup C = ... = A_n \cup C$$
?

e.
$$A_1 \cap C = A_2 \cap C = ... = A_n \cap C$$
?

f.
$$A_1 \cup C = A_2 \cup C = ... = A_n \cup C$$
 and $A_1 \cap C = A_2 \cap C = ... = A_n \cap C$?

d)

We can not conclude that $A_1=A_2=\ldots=A_n$ given $A_1\cup C=A_2\cup C=\ldots=A_n\cup C$. Let C be the universe. No matter what $A_1,\ A_2,\ \ldots A_n$ are, $A_1\cup C=A_2\cup C=\ldots=A_n\cup C$ will be equal, as C is the universe and any $A_i\cup C=C$. Thus, we can learn nothing about the equality of $A_1,\ A_2,\ \ldots A_n$.

e)

We can not conclude that $A_1=A_2=\ldots=A_n$ given $A_1\cap C=A_2\cap C=\ldots=A_n\cap C$. Let C be the empty set. No matter what $A_1,\ A_2,\ \ldots A_n$ are, $A_1\cap C=A_2\cap C=\ldots=A_n\cap C$ will be equal, as C is the empty set and any $A_i\cap C=C$. Thus, we can learn nothing about the equality of $A_1,\ A_2,\ \ldots A_n$.

f)

By applying our proof from **Question 1c)** multiple times, we can prove that $A_1 = A_2 = ... = A_n$ given that $A_1 \cup C = A_2 \cup C = ... = A_n \cup C$ and $A_1 \cap C = A_2 \cap C = ... = A_n \cap C$. For all possible pairs of A_i and A_j , we know that $A_i \cup C = A_j \cup C$ and $A_i \cap C = A_j \cap C$, and thus that $A_i = A_j$. This leads us to conclude that $A_1 = A_2 = ... = A_n$.

Question 2

Data are transmitted over the Internet in datagrams, which are structured blocks of bits. Each datagram contains header information organized into a maximum of **14** different fields (specifying many things, including the source and destination addresses) and a data area that contains the actual data that are transmitted. One of the 14 header fields is the header length field (denoted by **HLEN**), which is specified by the protocol to be 4 bits long and that specifies the header length in terms of the 32 bit blocks of bits. For example, if HLEN = 0110, the header is made up of six 32-bit blocks. Another of the 14 header fields is the 16-bit long **total length field **(denoted by **TOTAL LENGTH**), which specifies the length in bits of the entire datagram, including both the header fields and the data area. The length of the data area is the total length of the datagram minus the length of the header.

a.** **The largest possible value of TOTAL LENGTH (which is 16 bits long) determines the maximum total length in octets (block of 8 bits) of an Internet datagram. What is this value?

b. ** **The largest possible value of HLEN (which is 4 bits long) determines the maximum total header length in 32-bit blocks. What is this value? What is the maximum total header length in octets?

c.** **The minimum (and most common) header length is 20 octets. What is the maximum total length in octets of the data area of an Internet datagram?

d.** **How many different strings of octets in the data area can be transmitted if the header length is 20 octets and the total length is as long as possible?

a)

TOTAL LENGTH is 16 bits long. The largest possible value would be given by the 16 bit string of all 1s, which has a value of $2^0 + 2^1 + 2^2 + ... + 2^{15} = 65535$. This means that the longest total length of an Internet datagram is 65535 octets.

b)

HLEN is 4 bits long. The largest possible value would be given by the 4 bit string of all 1s, which has a value of $2^0 + 2^1 + 2^2 + 2^3 = 15$. This means that the maximum total header length is 15 32-bit blocks. Accordingly, the maximum total header length in octets is $(15 \times 32)/8 = 60$.

c)

The maximum total length of the data area is equal to the largest possible value of TOTAL LENGTH - the minimum possible header length. Thus, the maximum total length is 65535 - 20 = 65515 octets.

d)

The maximum total length of the data area is 65515 octets. There are 2 choices for each of the 8 bits in the octet, so the total number of octets is $2^8 = 256$. Thus, for each of the 65515 octets, there are 256 options. Thus, there are $256^{65515} = 6.94 \times 10^{157775}$ possible strings of octets that can be transmitted in the data area.

Question 3

- a. Suppose that $\lim_{x\to c} f(x)$ exists. Prove that there exists a constant M and a $\delta>0$ such that |f(x)| < M for $0 < |x-c| < \delta$.
- b. Give an explicit example that illustrates the above statement. You should give an actual f, c, M, and δ . (Also include a well-labeled illustrative diagram.)

a)

We suppose that $\lim_{x\to c} f(x)$ exists. Assume that $\lim_{x\to c} f(x) = L$. This means that for every $\varepsilon > 0$, there exists a δ such that if $0 < |x-c| < \delta$, then $f(x) \in (L-\varepsilon, L+\varepsilon)$. Let $\varepsilon = 1$. Then, we know that there exists a δ such that if $0 < |x-c| < \delta$, we know that $f(x) \in (L-1, L+1)$. Let M = |L| + 1. Thus, $L+1 \le M$ and $L-1 \ge -M$. Thus, we know that $f(x) \in (-M, M)$ which means that |f(x)| < M.

b)

Let f(x) = 2x + 1 and c = 1. The $\lim_{x \to 1} f(x) = 3 = L$. We let M = |L| + 1 = 4. Let $\varepsilon = 1$ and $\delta = 1/2$. If 0 < |x - 1| < 1/2, then $f(x) \in (L - \varepsilon, L + \varepsilon) = (3 - 1, 3 + 1) = (2, 4)$. So, if 0 < |x - 1| < 1/2, then |f(x)| < 4 = M.

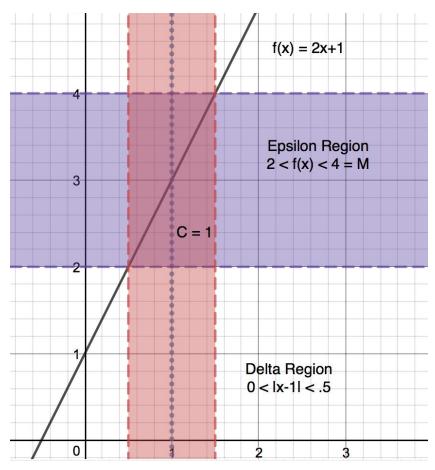


Figure 1: Epsilon and Delta regions

Question 4

(#algebra) Classify ALL groups of order 6 using the tools we have learned in this class. It is up to you to define what "classify" means, but it does NOT mean listing just the examples that we have seen in class. Hint: Break this problem up into cases. You must justify all your conclusions. You may use Sage if you think it is helpful. Note also that if you try to do background reading from other sources, you may encounter a lot of unfamiliar terminology! This is not necessary to do a thorough job on this task.

There are 2 groups of order 6. There can be multiple ways to represent them, but all representations can be classified into two cases. There is an abelian group of order 6, of which all representations are isomorphisms of the cyclic group Z_6 , and there is a non abelian group of order 6, of which all representations are isomorphisms of the symmetric group S_3 . In order to understand this, we can examine the cayley table for both groups. We can use the following code in Sage to create the below tables.

```
A = CyclicPermutationGroup(6)
print A.cayley_table()
B = SymmetricGroup(3)
print B.cayley_table()
```

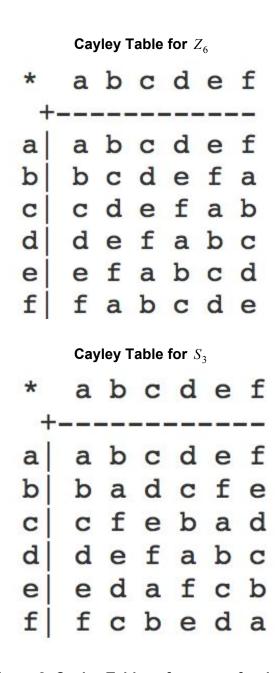


Figure 2: Cayley Tables of groups of order 6

We can first see that for any two elements α , β in the Cayley Table for Z_6 , $\alpha * \beta = \beta * \alpha$, and thus the group is abelian. For the Cayley Table for S_3 , on the other hand, we can see that $b*c=f\neq e=c*b$, and thus the group is non abelian.

As an example, we can choose a group of order 6 that is not represented by Z_6 or S_3 and show that it essentially the same as one of them. For two groups to be essentially the same, there is a bijection between the elements of the two groups such that both representations would have the same Cayley Table.

From **Corollary 6.11** in the Judson text, we know that the order of any element in a group of 6 elements must be 1,2,3 or 6, as they are the divisors of 6. If an element has order 6, it can be mapped to Z_6 , and if there are no elements of order 6 in the group, it can be mapped to S_3 .

We will choose the dihedral group D_3 . This group is the symmetries of a triangle, where the operators are a flip f and a rotation to the right r. There are 6 elements, the identity i, one flip f, one rotation r, a flip and a rotation f, two rotations r^2 , and a flip and two rotations, fr^2 . Note that three rotations or two flips is the same as the identity.

We can observe that this group is non abelian, as $f * r = fr \neq fr^2 = r * f$. This means that a rotation followed by a flip is not the same as a flip followed by a rotation.

By running the following Sage code,

```
C = DihedralGroup(3)
print C.cayley_table(["i", "f", "r", "fr", "r^2", "fr^2"])
```

We get the Cayley Table:

If we use the map $i \leftrightarrow a, f \leftrightarrow b, r \leftrightarrow c, fr \leftrightarrow d, r^2 \leftrightarrow e, fr^2 \leftrightarrow f$, we see that the Cayley Tables are the same.

We could do this same process for all other representations of groups of order 6, and could create a bijection between the representation and either Z_6 or S_3 .

Question 5

In this task, you demonstrate how to encode and decode your first and last name in the RSA cryptosystem by following the instructions in Judson, section 7.7. In particular, complete Judson 7.7: 1-4, except instead of sending the message "Math," send your first and last name. Be sure to explain how you deal with blocks of characters. For 4, you can use your name again. (Special note: Instead of using primes as specified in Judson, use two 200 digit primes that you find yourself. You must show how you found them, e.g., show any code you construct to find them.)

This problem is detailed in the following 10 pages of code, which was converted from a Jupyter Notebook running Sage.

```
## CONSTRUCTING A KEYPAIR FOR ALICE
# We use the random prime function to create a random prime
# with 200 digits.
# The random prime generates a prime between cutoff values
# which I chose specifically to generate a 200 digit number.
pa = random prime(10^200-1,False,10^199)
# We check that it is prime and that it has 200 digits
if is prime(p a) and len(str(p a)) == 200:
    print"p a is a prime number of 200 digits, it is: \n", p a
q a = random prime(10^200-1, False, 10^199)
# We check that it is prime, that it has 200 digits, and that
# it is not the same as p a
if is prime(q a) and len(str(q a)) == 200 and q a != p a:
    print"\np a is a prime number of 200 digits, it is: \n", q a
na = pa * qa
#We can't use m a = euler phi(n a) because our primes are to large
#and as such, more secure. Instead, we use the below formula.
m a = (p a - 1)*(q a - 1)
print "\nm a is : \n", m_a
p a is a prime number of 200 digits, it is:
38092642388946735939879406752157474086626160180141790937970175342361
85506156693422874984583229927501606375266480398744431409795798661950
5834237351512085135112198721704117498693487188021198066063134271
p a is a prime number of 200 digits, it is:
26028039689231637378592463899860479613091261654490757942831488745899
19041456093789454946895978556187009661033454140539282371780116127354
3655934894517867136372302064004948584811782695570894047608410061
```

m_a is :
99147680796721309777681971037944735318368903771662177699507406820940
70279994843940553979178470731516581819953827717173280262067158305893
46746773470316892643937864565058510336888942754644722081198566205315
03410497137129247169823364410081093292674709401569086093552418191221
47318497726163841642740553299589370483044524603461784989540193214680
51477705979970277071675375119005244022127550046174998756200

```
In [2]:
#It is not possible to factor m a, and thus
# we must find E a by randomally choosing numbers
# and checking if the gcd is = 1. Note that E is
# not neccesarily (and probably not) the lowest possible E.
E a = 0
while gcd(E a, m a) != 1:
    E_a = ZZ.random_element(m_a)
In [3]:
D a = inverse mod(E a, m a)
In [4]:
print "These are the values for Alice \n"
print " n = ", n_a
print " E = ", E_a
print " D = ", D a
These are the values for Alice
```

```
99147680796721309777681971037944735318368903771662177699507406
82094070279994843940553979178470731516581819953827717173280262067158
30589346746773470316892643937864565058510336888942754644722081198566
26943571618314974461094356888566205451065034858172656457166259961244
29576908597218959156989563588922161193000476498452974839942581019123
70958223723735932241761572461084185088749292011142138288670300531
      54593136901807186707304719835936886358120814272090599412690945
73613228262588510367817673653923723485888336577851064117890438254204
25758359638112179052373324792696051906510267417197543567963335991724
51376083831061261170007961509881545553468554061549974399893235423618
02984580962068659301667829563570775426982784340660326150244961011647
39467918555099177299995643014614405272004558543444795575101154007
      56629474870393970011029590160321712268036577276701298052678494
22335952886172051335172563089309909289435366678297091496040271524239
76645857184160763307641824901657630952744612082651764897919396206358
18635160939447379637646578374404796788523111305081304583344273670754
01498225817021500286083781751011442782743258588509736701255506828694
31799137502673404627286565839586179640805636796465069974406999143
```

In [5]:

```
print """We can assert that the encryption and decryption keys are multiplicativ
e inverses..."""
if inverse_mod(E_a, m_a) == D_a and inverse_mod(D_a, m_a) == E_a:
    print "They are multiplicative inverses"
```

We can assert that the encryption and decryption keys are multiplicative inverses...

They are multiplicative inverses

```
In [6]:
## NOW FOR BOB
p b = random prime(10^200-1, False, 10^199)
if is prime(p b) and len(str(p b)) == 200:
    print"p_b is a prime number of 200 digits, it is: \n", p_b
q b = random prime(10^200-1, False, 10^199)
if is prime(q b) and len(str(q b)) == 200 and q b != p b:
    print"\np b is a prime number of 200 digits, it is: \n", q b
n_b = p_b * q_b
m b = (p b -1)*(q b-1)
E b = 0
while gcd(E b, m b) != 1:
    E b = ZZ.random element(m b)
D b = inverse mod(E b, m b)
print "\n These are the values for Bob \n"
print " n = ", n_b
print " E = ", E_b
print " D = ", D_b
```

print """\nWe can assert that the encryption and decryption keys are multiplicat

if inverse mod(E b, m b) == D b and inverse mod(D b, m b) == E b:

print "They are multiplicative inverses"

ive inverses..."""

p_b is a prime number of 200 digits, it is:
48300024253958051679091274870598542327015093633381379473315222422414
33273921719077552238101930883346212194345031900067864404935969145963
1609607201616769388870124453108533815515886384216748693300734819

p_b is a prime number of 200 digits, it is:
20547157568939371780290655252793604200664164763754120123115895297750
13912308054424170157659135080672576605115372794626085257809807481324
7168573055631244405587374828387036171059603796088461190579548621

These are the values for Bob

- $\begin{array}{lll} n &=& 99242820892966941528409779863077841097029738263330163534249854\\ 19445555599260174946501533967046574331244042600490855371483778472305\\ 25077140900774008171932565432617437869594270385420196322801374437990\\ 97083423883092914922204156831378016887648304785656320517847804367203\\ 41105054355925064036133888383476532497172682772590791733407354876013\\ 63240379865911350812085344306297612778707593824269167412638134599 \end{array}$
- $\begin{array}{lll} E &=& 81281740855071145590232591513700767207565638549445012177148478\\ 69769130178607820751601255889017928164699193663951289143936383091250\\ 32019534700020742707159998325054674600517041273422798579760679322483\\ 43489908735128686090213165829522890175177985011391707774756870970347\\ 23576524342172399256484732430752187124870399594476829035325526140764\\ 88049036971217050108634401370104443606735515454920512321090854613 \end{array}$

We can assert that the encryption and decryption keys are multiplica tive inverses...

They are multiplicative inverses

```
In [7]:
### ENCODING MY FIRST AND LAST NAME
# I now encode "Isaac Schaal"...
# however, I know I need to send strings that
# are divisible by 4. I thus will encode and send
# the string "IsaacxSchaal", with the "x" representing
# a space. I will then break this string into 3 strings
# of length 4.
def split by len(text, chunksize):
    return [text[i*chunksize:(i*chunksize+chunksize)] for i in range(len(text)/c
hunksize)]
string = "IsaacxSchaal"
string = split by len(string, 4)
print string
['Isaa', 'cxSc', 'haal']
In [8]:
# We now must encrypt the message
digits list = []
messages = []
```

The message is: [205027785, 208993379, 228094184]

messages.append(ZZ(digits, 128))

print "The message is :", messages

digits_list.append([ord(letter) for letter in word])

for word in string:

for digits in digits list:

```
signed_list = []
for message in messages:
    signed_list.append(power_mod( message, D_a, n_a))
print "The signed message is : ", signed_list
```

The signed message is: [254164326502664313598874651600280660381579 61702043069714825, 7556741169340472409842847444709190329052688910565 0032224443, 45083820833899571128426451969546685057129008985711515893 20535283005725514625109353754520146150791554867604775995974698515941707]

In [10]:

In [9]:

```
encrypted_list = []
for signed in signed_list:
    encrypted_list.append(power_mod(signed, E_b, n_b))
print "The encrypted message is : ", encrypted_list
```

The encrypted message is: [401925809566064050975869074026045151585 71352418905441508646, 2279108591140747084730939355803877498618107901 7681121210277, 94711587906195612238276766839171629325828761250654544

```
In [11]:
```

```
### DECODING
# The message is sent to Bob. He decodes it.
decrypted list = []
for encrypted in encrypted list:
    decrypted list.append(power mod(encrypted, D b, n b))
print "The decrypted message is : ",decrypted list
received list = []
for decrypted in decrypted list:
    received list.append(power mod(decrypted, E_a, n_a))
print "The unsigned message (from Alice) is : ", received list
digits list = []
for received in received list:
    digits list.append(received.digits(base=128))
letters list = []
for digits in digits list:
    letters list.append([chr(ascii) for ascii in digits])
lis = []
for letter in letters list:
    lis.append("".join(letter))
answer = ''.join(lis)
print "The final answer is : ", answer
```

```
The decrypted message is: [254164326502664313598874651600280660381
57922194211070839961729895792864409699002015416580841519456565600599
79393612715104625515791470952269676047158579189228307913858599527924
77772925872727837100434345157878025258843808093072140162691277758533
44941365731521877791330952422685819779099683015129261070441220048572
56871119074984076957512963068361424731625698420264601552205926969517
74961702043069714825, 7556741169340472409842847444709190329052688910
56581713296269843936827734798687325927792555989207389803895891878350
28124778000746427622135420326600721917116903378692599261361215559480
13424478613058180577392185195301112651878800619638997840421506390574
59884006513105632024786321758375335174805280477178253919823940897454
92924480411417488593655385997615568956268639618572238526310652993622
7810032224443, 45083820833899571128426451969546685057129008985711515
89320535283005725514625109353754520146150791554867604775995974698515
94190319029670142152867757019540281450892262397271081680470925673916
83824372612459120960284314567986340235015063042876430332758499750545
16940870667628973354650391034574750615582778496225803332489198708168
76088619298725790695222596159490394562330116731666661051426925465166
2707071
The unsigned message (from Alice) is: [205027785, 208993379, 22809
4184]
The final answer is : IsaacxSchaal
```

In [12]:

```
### TAMPERED MESSAGE
# We first display the sent message again.
print "The encrypted message is : ", encrypted_list
```

The encrypted message is: [401925809566064050975869074026045151585 71352418905441508646, 2279108591140747084730939355803877498618107901 7681121210277, 94711587906195612238276766839171629325828761250654544

```
In [13]:
```

```
#We then add 1 to each segment of the encrypted message
# to emulate it being tampered with.
tampered_list = []
for i in range(len(encrypted_list)):
    tampered_list.append(encrypted_list[i]+1)
print "The tampered message is : ", tampered_list
```

The tampered message is : [4019258095660640509758690740260451515857 1352418905441508647, 22791085911407470847309393558038774986181079016 681121210278, 947115879061956122382767668391716293258287612506545443

```
In [14]:
```

```
# We now see what happens if Bob trys to decrypt the message
decrypted list = []
for tampered in tampered list:
    decrypted list.append(power mod(tampered, D b, n b))
received list = []
for decrypted in decrypted list:
    received_list.append(power_mod(decrypted, E_a, n_a))
digits list = []
for received in received list:
    digits list.append(received.digits(base=128))
letters list = []
for digits in digits list:
    letters list.append([chr(ascii) for ascii in digits])
lis = []
for letter in letters list:
    lis.append("".join(letter))
answer = ''.join(lis)
print "The final answer is : ", answer
The final answer is: h h"VuxeJ0{; | $[ngp<(m?7]]?P{JF{ bAq3?i0I<gZRm}
sOSEe\{\#B[>'-5]
{3Ruw?##2W!b`$)6kC^AF
IMnR9K<hwnQrD(v{B1%5*qIpKPk*Trp\tw.{gz"J[<iV$1}`MZw</pre>
uyPCPskYudO7C0jPe*N"xm1p@5B=Bf|,I/= w)4CTgD,r:
;BqQ -Lw-TSi ,b?xlGW0}4s,=&D\HyXV-,b
. HNwRT
i6TD<PPhxJV"bl#tpArjou2=::oUDymji lm(@
Q<i\%iw+j;v^WU<L-\sim Qt\%sAWC:iGlL = (Z
@wrfE ~u<,E(]2E48Y</pre>
In [1]:
#This is cleary gibberish, and shows us how sensitive
```

#RSA is to small amounts of tampering with the message