

Section 1

Exercise 1.1.i

- (i) Let $f : x \rightarrow y$ be a morphism and $g, g' : y \rightarrow x$ be inverses. Then $fg = \text{id}_y$ and $gf = \text{id}_x$ (resp g'). Hence $gf = g'f$ so

$$\begin{aligned} gfg &= g'fg \\ g &= g' \end{aligned}$$

Therefore any morphism has at most one inverse. ///

- (ii) Suppose there are morphisms $g, h : y \rightrightarrows x$ such that $fg = \text{id}_y$ and $hf = \text{id}_x$. Then

$$\begin{aligned} hfg &= h \\ g &= h \end{aligned}$$

Hence $fg = \text{id}_y$ and $gf = \text{id}_x$ so f is an isomorphism. ///

Exercise 1.1.ii

Consider a category C and the subset of isomorphisms it contains. The subset contains the identity morphisms for each object because the identity is an isomorphism. The subset inherits associativity from C . Now pick two isomorphisms $f : x \rightarrow y$ and $g : y \rightarrow z$ and consider gf . We have $(f^{-1}g^{-1})gf = f^{-1}g^{-1}gf = f^{-1}f = \text{id}_x$ and $gf(f^{-1}g^{-1}) = \text{id}_y$ so gf is also an isomorphism. Hence C contains a maximal groupoid. ///

Exercise 1.1.iii

- (i) Observe that for any object $f : c \rightarrow x$ in c/C , the morphism $\text{id}_x : x \rightarrow x$ is the identity morphism for f . Now pick two morphisms in c/C , $h : x \rightarrow y$ and $k : y \rightarrow z$ and consider kh .

$$\begin{array}{ccccc} & & c & & \\ & f \swarrow & \downarrow g & \searrow b & \\ x & \xrightarrow{h} & y & \xrightarrow{k} & z \end{array}$$

Because $hf = g$ and $kg = b$ we have $khf = b$ so kh is a morphism in c/C which is enough to show it is a category. ///

- (ii) A similar argument as above shows that C/c is a category. ///

Section 2

Exercise 1.2.iii

Let $f : x \hookrightarrow y$ and $g : y \hookrightarrow z$ be monomorphisms and $h, k : w \rightarrow x$ be morphisms such that $gh = gfk$. Because g is monic, it follows that $fh = fk$. Similarly, because f is monic we have $h = k$. Hence gf is monic.

Now suppose f and g are simply morphisms and that gf is monic. BWO suppose that f is not monic. Then there exist morphisms h, k such that $fh = fk$ but $h \neq k$. Hence

we have $gfh = gfk$ this is a contradiction because gf being monic implies that $h = k$. Hence f is monic.

Now observe that the monomorphisms in a category form a subcategory. Every object has an identity morphism because the identity is monic. Associativity is inherited from the larger category and composition is preserved because the composition of two monomorphisms is also a monomorphism as shown above.

By duality the epimorphisms also form a subcategory. ///

Exercise 1.2.iv Recall that the only ideals in a field F are $\{0\}$ and F . Pick a morphism $f : X \rightarrow Y$ in **Field**. The kernel of f is an ideal so it is either $\{0\}$ or X . By the first isomorphism theorem, f is either 1-1 or the 0 map. Now pick g, h such that $fg = fh$. Then because f is 1-1, then it maps every element of its domain to a unique element in the codomain hence $h = g$. Therefore f is monic. Hence the monomorphisms in **Field** are the non-zero maps. ///

***Exercise 1.2.v** In **Ring**, consider the inclusion map $i : \mathbb{Z} \hookrightarrow \mathbb{Q}$. Now pick two morphisms, $f, g : R \rightrightarrows \mathbb{Z}$ such that $if = ig$ where R is some ring. Now BWOC suppose that $f \neq g$. Then there exists an $r \in R$ such that $f(r) \neq g(r)$. Hence $i(f(r)) \neq i(g(r))$ so $if \neq ig$ which is a contradiction. Hence the above implies that $f = g$ so i is monic.

Now pick morphisms $h, k : \mathbb{Q} \rightrightarrows R$ such that $hi = ki$. This means that for all integers z , $h(z) = k(z)$. Each rational q may be written canonically as nd^{-1} for $n, d \in \mathbb{Z}$. Hence $h(q) = h(nd^{-1}) = h(n)h(d^{-1})$. Further, because every element in \mathbb{Q} is a unit, $h(n)h(d^{-1}) = h(n)h(d)^{-1}$. Because h and k agree on the integers, we have $h(n)h(d)^{-1} = k(n)k(d)^{-1}$. Hence $h = k$ and i is an epimorphism.

\mathbb{Q} and \mathbb{Z} are not isomorphic in **Ring** because \mathbb{Q} is a field while \mathbb{Z} is not. Hence in **Ring**, monic and epic does not imply isomorphism. ///

Exercise 1.2.vii For a given subset of objects $A \subseteq P$, consider $M_A = \{x \in \text{obj}P \mid \forall a \in A, \exists f \in P(a, x)\}$, i.e. the set of all elements in P targeted by a morphism from every element in A . If it exists, the supremum of A is the least element in M_A w.r.t. \leq .

We define the infimum by duality as the greatest element in $m_A = \{x \in \text{obj}P \mid \forall a \in A, \exists f \in P(x, a)\}$.

Now suppose we have p, q both suprema of A . Then by definition $p \leq q$ and $q \leq p$ because each is a member of A . Hence by the antireflexivity of \leq , we have $p = q$. The dual proof shows uniqueness for the infimum. ///

Section 3

Exercise 1.3.i Given two groups G and H , a functor $BG \rightarrow BH$ is a group homomorphism. Given such a functor F , and elements $x, y \in G$, that is morphisms in BG , we have $F(x)F(y) = F(xy)$ which is exactly a group homomorphism. ///

Exercise 1.3.iii Pick a category C where $\text{obj}C = \{a, b, b', c\}$ and $|\text{obj}C| = 4$ such that the only morphisms included beyond the identities are $f : a \rightarrow b$ and $g : b' \rightarrow c$.

Now define a functor $F : C \rightarrow D$ for some category D where $\text{obj}D = \{Fa, Fb = Fb', Fc\}$ and $|\text{obj}D| = 3$. Let F preserve identities and $Ff : Fa \rightarrow Fb$ and $Fg : Fb \rightarrow Fc$. Observe that this is in fact a functor because it preserves identities and there are no nontrivial compositions in C .

Because $Fg \circ Ff$ is a morphism in D , but is not in the image of F , the image of F is not a sub category. ///

Exercise 1.3.iv First observe that $C(c, -)$ and $C(-, c)$ both preserve identities. Consider the action of $C(c, -)$ on $\text{id}_x : x \rightarrow x$. id_x maps to post composition by the identity on $C(c, x)$ which is the identity morphism for $C(c, x)$. Similar reason applies for $C(-, c)$.

Now pick two morphisms $f : x \rightarrow y$ and $g : y \rightarrow z$ in C . Then $C(c, -)g \cdot C(c, -)f = g_*f_* = (gf)_* = C(c, -)(gf)$. Hence $C(c, -)$ is in fact a functor. The same holds for $C(-, c)$ by duality. ///

Exercise 1.3.viii Consider the trivial Functor $\mathbf{0} : \text{Set} \rightarrow \text{Set}$ that maps every set to \emptyset and every morphism to id_\emptyset . Observe that this is indeed a functor because $\mathbf{0}(\text{id}_x) = \text{id}_{\mathbf{0}x} = \text{id}_\emptyset$ and $\mathbf{0}(gf) = \mathbf{0}(g)\mathbf{0}(f) = \text{id}_\emptyset$. Now pick any morphism in Set that is not an isomorphism. It is mapped to id_\emptyset which is an isomorphism. Hence functors do not reflect isomorphisms. ///

Exercise 1.3.vi Consider the map $\text{dom} : F \downarrow G \rightarrow D$ that takes an object $(d, e, f) \mapsto d$ and a morphism $(h, k) \mapsto h$. Now pick an identity morphism $(\text{id}_d, \text{id}_e)$; clearly this functor preserves identity. Now pick a composable pair of morphisms (h, k) and (h', k') . We have $\text{dom}((h', k')(h, k)) = \text{dom}(h'h, k'k) = h'h = \text{dom}(h', k')\text{dom}(h, k)$. Hence dom is a functor. Define cod analogously. ///

Exercise 1.3.x Let $\text{cl}(g)$ be the conjugacy class of g and $\text{Cl}(G)$ be the set of all conjugacy classes of G (I don't know what the standard notation for this is). Consider the following mapping:

$$\begin{aligned} \text{Conj} : \text{Group} &\longrightarrow \text{Set} \\ G &\longmapsto \text{Cl}(G) \\ (f : G \rightarrow H) &\longmapsto (f' : \text{Cl}(G) \rightarrow \text{Cl}(H)) \end{aligned}$$

Where f' maps $\text{cl}(x)$ to $\text{cl}(f(x))$. Note that this is indeed a well defined map: pick an element b in G . Then $f'(b) = \text{cl}(f(b))$. For any other element $a \in \text{cl}(b)$, we have $f'(a) = \text{cl}(f(a)) = \text{cl}(f(g)f(b)f(g)^{-1}) = \text{cl}(f(b))$.

Observe that Conj preserves identities. Now pick $h : G \rightarrow H$ and $k : H \rightarrow I$. Then $\text{Conj}(kh) = f' : \text{Cl}(G) \rightarrow \text{Cl}(I) = \text{Conj}(k)\text{Conj}(h)$. Hence Conj is a functor.

Now pick a pair of groups such that the cardinalities of their sets of conjugacy classes differ. If there was an isomorphism between the two groups, say $k : G \rightarrow H$, then

$\text{Conj}k : \text{Cl}(G) \rightarrow \text{Cl}(H)$ would be an isomorphism as well. But this cannot be because there are no isomorphisms between sets with different cardinalities. ///

Section 1.4

Exercise 1.4.i Consider the inverse of a component of the natural isomorphism $\alpha : F \Rightarrow G$ such as $\alpha_c^{-1} : Gc \rightarrow Fc$. It remains to show that the following diagram commutes:

$$\begin{array}{ccc} Gc & \xrightarrow{\alpha_c^{-1}} & Fc \\ G\phi \downarrow & & \downarrow F\phi \\ Gc' & \xrightarrow{\alpha_{c'}^{-1}} & Fc' \end{array}$$

By the naturality of α we have that $G\phi \cdot \alpha_c = \alpha_{c'} \cdot F\phi$. Hence by composing each side with α_c^{-1} and $\alpha_{c'}^{-1}$ we get $\alpha_{c'}^{-1} \cdot G\phi = F\phi \cdot \alpha_c^{-1}$. Hence $\alpha^{-1} : G \Rightarrow F$ is a natural isomorphism. ///

Exercise 1.4.ii Pick two groups H, K and two functors $F, G : BH \Rightarrow BK$. By exercise 1.3.i, F and G are group homomorphisms. Now pick a natural transformation $\alpha : F \Rightarrow G$. α and an element $a \in H$. Examining the naturality square for α_a shows that $\alpha_a \circ Fa = Ga \circ \alpha_a$.

Exercise 1.4.iv Observe that f and g define the same natural transformation exactly when their components are equal. The naturality condition for f_* is that $h^*f_* = f_*h^*$ and for g_* it is $h^*g_* = g_*h^*$. Hence because f and g are distinct morphisms, we do not have $h^*f_* = h^*g_*$, for example if h^* is the identity. ///

Exercise 1.4.vi The components of a natural transformation are defined to be morphisms in the shared target category of two functors. Hence the definition of natural (or extranatural) transformation does not allow for different target categories because the same notion of morphism does not apply. ///

Questions: Why is A included in the definition extranatural transformation? Why not just B and B^{op} ?

Are there notions of morphism between objects in different categories? (Not functors, but a morphism from an object $x \in C$ to an object $y \in D$).