Stacks & Queue

- godo any consiguous - west to each other stack - lastin lastoct

queue - lastin furtant

linked lists - only accessor a particular order

Stack--top

- capacity - size -head items

-make constructors à destructors -make, vouse, destroy : 4

(heisenby "
-deprods on time à concomstance

Java performance depends:
- memory dean up only happer when memory runs out

Opregre data typee - any elements in it

malloc sometimes faits - Check : F NULL

realloc, reallocates aremony -con fail he sure to cheeks, prob naturalse

delete Queve free (gn ->Q); #free gute free (g); It recycle in deut

- new head is successor to old head

Queve is full it successor of head is tast

Priority Queve -every element has a priority assigned - has a specialize order -element of highest priority degreed frot -ext heap -manutained ut log steps -big C

Discrete Event Sim here are good for that

Dynamic c numery alloc
- malloc gives you down beyond data section of unso
- no order no stack discipline
- not freezy depletes system resources (womany lenk)

The Herp

- aranguard menory

- needs poolities

- pile of date @ end of program

calloc better than malloc sme it puts O:

Realloc can change your location. Your pointers may no longer work