A Sound Type System for Physical Quantities, Units, and Measurements

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Abstract

We present a theory in Isabelle/HOL [7] that builds a formal model for both the *International System of Quantities* (ISQ) and the *International System of Units* (SI), which are both fundamental for physics and engineering [2]. Both the ISQ and the SI are deeply integrated into Isabelle's type system. Quantities are parameterised by *dimension types*, which correspond to base vectors, and thus only quantities of the same dimension can be equated. Since the underlying "algebra of quantities" from [2] induces congruences on quantity and SI types, specific tactic support is developed to capture these. Our construction is validated by a test-set of known equivalences between both quantities and SI units. Moreover, the presented theory can be used for type-safe conversions between the SI system and others, like the British Imperial System (BIS).

1 Introduction

Modern Physics is based on the concept of quantifiable properties of physical phenomena such as mass, length, time, current, etc. These phenomena, called *quantities*, are linked via an algebra of quantities to derived concepts such as speed, force, and energy. The latter allows for a dimensional analysis of physical equations, which had already been the backbone of Newtonian Physics. In parallel, physicians developed their own research field called "metrology" defined as a scientific study of the measurement of physical quantities.

The relevant international standard for quantities and measurements is distributed by the Bureau International des Poids et des Mesures (BIPM), which also provides the Vocabulaire International de Métrologie (VIM) [2]. The VIM actually defines two systems: the International System of Quantities (ISQ) and the International System of Units (SI, abbreviated from the French 'Système international d'unités'). The latter is also documented in the SI Brochure [3], a standard that is updated periodically, most recently in 2019. Finally, the VIM defines concrete reference measurement procedures as well as a terminology for measurement errors.

Conceived as a refinement of the ISQ, the SI comprises a coherent system of units of measurement built on seven base units, which are the metre, kilogram, second, ampere, kelvin, mole, candela, and a set of twenty prefixes to the unit names and unit symbols, such as milliand kilo-, that may be used when specifying multiples and fractions of the units. The system also specifies names for 22 derived units, such as lumen and watt, for other common physical

quantities. While there is still nowadays a wealth of different measuring systems such as the *British Imperial System* (BIS) and the *United States Customary System* (USC), the SI is more or less the de-facto reference behind all these systems. ¹

The present Isabelle theory builds a formal model for both the ISQ and the SI, together with a deep integration into Isabelle's order-sorted polymorphic type system [6]. Quantities and units are represented in a way that they have a quantity type as well as a unit type based on its base vectors and their magnitudes. Since the algebra of quantities induces congruences on quantity and SI types, specific tactic support has been developed to capture these. Our construction is validated by a test-set of known equivalences between both quantities and SI units. Moreover, the presented theory can be used for type-safe conversions between the SI system and others, like the British Imperial System (BIS).

As a result of our theory development², it is possible to express "4500.0 kilogram times metre per second squared" has the type $\mathbb{R}[kg \cdot m \cdot s^{-3}]$ This type means that the magnitude 4500.0 (which by lexical convention is considered as a real number) of the dimension $M \cdot L \cdot T^{-3}$ is a quantity intended to be measured in the SI-system, which means that it actually represents a force measured in Newtons. Via a type synonym, the above type expression gets the type \mathbb{R} newton.

In the example, the magnitude type part of this type is the real numbers \mathbb{R} . In general, however, magnitude types can be more general. If the term above is presented slightly differently as "4500 kilogram times metre per second squared", the inferred type will be ' $\alpha[kg \cdot m \cdot s^{-3}]$ where ' α is a magnitude belonging to the type-class numeral. This class comprises types like \mathbb{N} , \mathbb{Z} , 32 bit integers (32word), IEEE-754 floating-point numbers, as well as vectors belonging to the three-dimensional space \mathbb{R}^3 , etc. Thus, our type-system allows to capture both conceptual entities in physics as well as implementation issues in concrete physical calculations on a computer.

As mentioned before, it is a main objective of this work to support the quantity calculus of ISQ and the resulting equations on derived SI entities (cf. [3]), both from a type checking as well as a proof-checking perspective. Our design objectives are not easily reconciled, however, and so some substantial theory engineering is required. On the one hand, we want a deep integration of dimensions and units into the Isabelle type system. On the other, we need to do normal-form calculations on types, so that, for example, the units $\alpha[s \cdot m \cdot s^{-2}]$ and $\alpha[m \cdot s^{-1}]$ can be equated.

Isabelle's type system follows the Curry-style paradigm, which rules out the possibility of direct calculations on type-terms (in contrast to Coq-like systems). However, our semantic interpretation of ISQ and SI requires the foundation of the heterogeneous equivalence relation \cong_Q in semantic terms. This means that we can relate quantities with syntactically different dimension types, yet with same dimension semantics. This paves the way for derived rules that do computations of terms, which represent type computations indirectly. This principle is the basis for the tactic support, which allows for the dimensional type checking of key definitions of the SI system inside Isabelle/HOL, i. e. without making use of code-generated reflection. For example, the crucial definitions adapted from the SI Brochure that give the concrete definitions for the metre and the kilogram can be presented as follows:

 $^{^1{\}rm See}$ also https://en.wikipedia.org/wiki/International_System_of_Units.

²The sources can be found in the Isabelle Archive of Formal Proofs at https://www.isa-afp.org/entries/ Physical_Quantities.html

theorem metre-definition

- $\bullet \ \ (1{::}'a) *_Q metre \cong_Q \mathbf{c} \cdot ((299792458{::}'a) *_Q \mathbf{1})^{-\mathbf{1}} \cdot second$
- $(1::'a) *_{Q} metre \cong_{Q} (9192631770::'a) / (299792458::'a) *_{Q} \mathbf{c} \cdot ((9192631770::'a) *_{Q} second^{-1})^{-1}$

theorem kilogram-definition

 $\bullet \quad (1::'\alpha) *_Q kilogram \cong_Q \mathbf{h} \cdot ((662607015::'\alpha) \ / \ (10::'\alpha)^8 \cdot (1::'\alpha) \ / \ (10::'\alpha)^{34} *_Q \mathbf{1})^{-1} \cdot metre^{-\mathbf{2}} \cdot second$

These equations giving the concrete definitions for the metre and kilogram in terms of the physical constants \mathbf{c} (speed of light) and \mathbf{h} (Planck constant) can be proven directly using the tactic si-calc provided of our theory.

1.0.1 The Plan of the Theory Development

In the following we describe the overall theory architecture in more detail. Our ISQ model provides the following fundamental concepts:

- 1. dimensions represented by a type $(\mathbb{Z}, 'd)$ dimvec, i.e. a 'd-indexed vector space of integers representing the exponents of the dimension vector. 'd is constrained to be a dimension type later.
- 2. quantities represented by type (' α , 'd) Quantity, which are constructed as a vector space and a magnitude type ' α .
- 3. quantity calculus consisting of quantity equations allowing to infer that $L \cdot T^{-1} \cdot T^{-1} \cdot M$ is isomorphic to $M \cdot L \cdot T^{-2}$ is isomorphic to F (the left-hand-side equals mass times acceleration which is equal to force).
- 4. a kind of equivalence relation $=_Q$ on quantities, permitting to relate quantities of different dimension types.
- 5. base quantities for length, mass, time, electric current, temperature, amount of substance, and luminous intensity, serving as concrete instance of the vector instances, and for syntax a set of the constant symbols \mathbf{L} , \mathbf{M} , \mathbf{T} , \mathbf{I} , Θ , \mathbf{N} , \mathbf{J} corresponding to the above mentioned base vectors.
- 6. (Abstract) Measurement Systems represented by type (' α , 'd, 's) Measurement-System, which are a refinement of quantities. The refinement is modelled by a polymorphic record extensions; as a consequence, Abstract Measurement Systems inherit the algebraic properties of quantities.
- 7. derived dimensions types such as volume L^3 or energy $M \cdot L^2 \cdot T^{-2}$ corresponding to derived quantities.

Then, through a fresh type-constructor SI, the abstract measurement systems are instantiated to the SI system — the *British Imperial System* (BIS) is constructed analogously. Technically, SI is a tag-type that represents the fact that the magnitude of a quantity is actually

a quantifiable entity in the sense of the SI system. In other words, this means that the magnitude 1::'a in quantity $(1::'a) *_Q metre$ actually refers to one metre intended to be measured according to the SI standard and has type $\mathbb{Z}[m]$. At this point, it becomes impossible, for example, to add one foot, in the sense of the BIS, to one metre in the SI without creating a type-inconsistency.

The theory of the SI is created by specialising the *Measurement-System*-type with the SI-tag-type and adding new infrastructure. The SI theory provides the following fundamental concepts:

- 1. measuring units and types corresponding to the ISQ base quantities such as *metre*, *kilogram*, *second*, *ampere*, *kelvin*, *mole* and *candela* (together with procedures how to measure a metre, for example, which are defined in accompanying standards);
- 2. a standardised set of symbols for units such as m, kg, s, A, K, mol, and cd;
- 3. a standardised set of symbols of SI prefixes for multiples of SI units, such as $giga \ (=10^9)$, $kilo \ (=10^3)$, $milli \ (=10^{-3})$, etc.; and a set of
- 4. unit equations and conversion equations such as $J = kg m^2 / s^2$ or 1 km/h = 1/3.6 m/s.

2 Background: Some Advanced Isabelle Constructs

This work uses a number of features of Isabelle/HOL and its meta-logic Isabelle/Pure, that are not necessarily available in another system of the LCF-Prover family and that needs therefore some explanation:

- Type-classes and order-sorted parametric polymorphism [5, 6]. Haskell-like type-classes allow for types depend on constants and represent therefore a restricted form of dependent types.
- The meta-logic Pure providing mechanisms to denote types inside the term-language: $'\alpha$ itself denotes an unspecified type and TYPE a constructor that injects the language of types into the language of terms.
- Code-generation: Reflection via eval
- The lifting package

3 Preliminary Algebraic Structures

At the core, the multiplicative operation on physical dimension results in additions of the exponents of base vectors:

$$\begin{array}{l} (M^{\alpha 1} \cdot L^{\alpha 2} \cdot T^{\alpha 3} \cdot I^{\alpha 4} \cdot \Theta^{\alpha 5} \cdot N^{\alpha 6} \cdot J^{\alpha 7}) * (M^{\beta 1} \cdot L^{\beta 2} \cdot T^{\beta 3} \cdot I^{\beta 4} \cdot \Theta^{\beta 5} \cdot N^{\beta 6} \cdot J^{\beta 7}) \\ = (M^{\alpha 1 + \beta 1} \cdot L^{\alpha 2 + \beta 2} \cdot T^{\alpha 3 + \beta 3} \cdot I^{\alpha 4 + \beta 4} \cdot \Theta^{\alpha 5 + \beta 5} \cdot N^{\alpha 6 + \beta 6} \cdot J^{\alpha 7 + \beta 7}) \end{array}$$

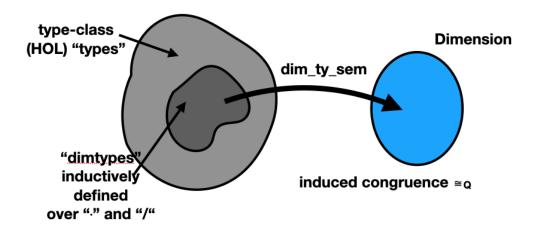


Figure 1: The "Inductive" Subset of dim-types interpreted in the Dimension-Type

This motivates type classes that represent this algebraic structure. We chose to represent it for the case of vectors of arbitrary length. We define the classes *group-mult* and the abelian multiplicative groups as follows:

```
notation times (infixl \cdot 70)

class group-mult = inverse + monoid-mult +
assumes left-inverse: inverse a \cdot a = 1
assumes multi-inverse-conv-div [simp]: a \cdot (inverse \ b) = a \ / \ b
...

class ab-group-mult = comm-monoid-mult + group-mult
...
abbreviation (input) npower :: '\alpha::{power,inverse} \Rightarrow nat \Rightarrow '\alpha ((---) [1000,999] 999)
where npower x \ n \equiv inverse \ (x \ \hat{} \ n)
```

... and derive the respective properties:

On this basis we define *dimension vectors* of arbitrary size via a type definition as follows:

```
typedef ('\beta, '\nu) dimvec = UNIV :: ('\nu::enum \Rightarrow '\beta) set morphisms dim-nth dim-lambda ..
```

Here, the functions dim-nth and dim-lambda represent the usual function pair that establish the isomorphism between the defined type $('\beta, '\nu)$ dimvec and an implementing domain, in this case the universal set of type $('\nu \Rightarrow '\beta)$ set. Note that the index-type $'\nu$ is restricted to be enumerable by type class enum.

Via a number of intermediate lemmas over types, we can finally establish the desired result in Isabelle compactly as follows:

```
instance dimvec :: (ab-group-add, enum) ab-group-mult by (<proof ommitted>)
```

If ' β is an abelian additive group, and if the index type ' ν is enumerable, (' β , ' ν) dimvec is an abelian multiplicative group.

4 The Domain: ISQ Dimension Terms and Calculations

In the following, we will construct a concrete semantic domain as instance of (β, ν) dimvec. This is where the general model of the dimension vector space of section 3 becomes a specific instance of the current ISQ standard as defined [2]; should physicians discover one day a new physical quantity, this would just imply a change of the following enumeration. Moreover, we will define the ISQ standards dimensions as base vectors in this vector space; historically, there had been alternative proposals of a quantity system that boil down to the choice of another eigen-vector set in this vector space.

The definition of an enumeration and the proof that it can be accommodated to the required infrastructure of the *enum*-class is straight-forward, and the construction of our domain *Dimension* follows immediately:

```
\begin & \textbf{datatype} \ sdim = Length \mid Mass \mid Time \mid Current \mid Temperature \mid Amount \mid Intensity \\ \begin & \textbf{definition} \ enum-sdim = [Length, Mass, Time, Current, Temperature, Amount, Intensity] \\ \begin & \textbf{definition} \ enum-sdim = [Length, Mass, Time, Current, Temperature, Amount, Intensity] \\ \begin & \textbf{definition} \ enum-sdim \ P \ \longleftrightarrow P \ Length \land P \ Mass \land P \ Time \land \dots \\ \begin & \textbf{definition} \ enum-ex-sdim \ P \ \longleftrightarrow P \ Length \lor P \ Mass \lor P \ Time \lor \dots \\ \begin & \textbf{instance} < \textbf{proof} \ ommitted > \\ \begin & \textbf{end} \\ \begin & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sdim) \ dimvec \\ \end & \textbf{type-synonym} \ Dimension = (\mathbb{Z}, sd
```

Note that the *enum*-class stems from the Isabelle/HOL library and is intended to present sufficient infrastructure for the code-generator. Note, further, that [2] discusses also the possibility of rational exponents, but finally defines them as integer numbers \mathbb{Z} .

A base dimension is a dimension where precisely one component has power 1: it is the dimension of a base quantity. Here we define the seven base dimensions. For the concrete definition of the seven base vectors we define a constructor:

```
definition mk-BaseDim :: sdim \Rightarrow Dimension where mk-BaseDim d = dim-lambda (\lambda i. if (i = d) then 1 else 0)
```

which lets us achieve a first major milestone on our journey: a *term* representation of base vectors together with the capability to prove and to compute dimension-algebraic equivalences. We introduce the ISQ dimension symbols defined in [2]:

```
abbreviation LengthBD (L) where \mathbf{L} \equiv mk\text{-}BaseDim\ Length abbreviation MassBD (M) where \mathbf{M} \equiv mk\text{-}BaseDim\ Mass} ... abbreviation BaseDimensions \equiv \{\mathbf{L},\,\mathbf{M},\,\mathbf{T},\,\mathbf{I},\,\Theta,\,\mathbf{N},\,\mathbf{J}\} lemma BD\text{-}mk\text{-}dimvec\ [si\text{-}def]: \mathbf{L} = mk\text{-}dimvec\ [1,\ 0,\ 0,\ 0,\ 0,\ 0] \mathbf{M} = mk\text{-}dimvec\ [0,\ 1,\ 0,\ 0,\ 0,\ 0]
```

A demonstration of a computation ³ and a proof is shown in the example below:

```
value \mathbf{L} \cdot \mathbf{M} \cdot \mathbf{T}^{-2}
lemma \mathbf{L} \cdot M \cdot \mathbf{T}^{-2} = mk-dimvec [1, 1, -2, 0, 0, 0, 0] by (simp\ add:\ si\text{-def})
```

Note that the multiplication operation (\cdot) is inherited from the fact that the *Dimension*-type is a proven instance of the *ab-group-mult*-class. So far, the language of dimensions is represented by a shallow embedding in the *Dimension* type.

5 Dimension Types and their Semantics in Terms of the *Dimension*-Type

The next section on our road is the construction of a sub-language of type-expressions. To this end, we define a *type class* by those type-terms for which we have an interpretation function dim-ty-sem into the values of the Dimension-type. For our construction it suffices that the type-symbols of this class have a unitary, i. e., one-elementary, carrier-set.

```
class dim\text{-}type = unitary +

fixes dim\text{-}ty\text{-}sem :: '\alpha \ itself \Rightarrow Dimension

class basedim\text{-}type = dim\text{-}type +

assumes is\text{-}BaseDim: is\text{-}BaseDim \ (dim\text{-}ty\text{-}sem \ (TYPE('\alpha)))
```

³The command value compiles the argument to SML code and executes it

Recall that the type constructor ' α itself from Isabelle/Pure denotes an unspecified type and TYPE a constructor that injects the language of types into the language of terms. We also introduce a sub-type-class basedim-type for base-dimensions.

The definition of the basic dimension type constructors is straightforward via a oneelementary set, *unit set*. The latter is adequate since we need just an abstract syntax for type expressions, so just one value for the dimension-type symbols. We define types for each of the seven base dimensions, and also for dimensionless quantities.

```
\begin{array}{lll} \textbf{type-def} & Length & = UNIV :: unit \ set \ .. \ \textbf{setup-lifting} \ \ type-definition-Length \\ \textbf{type-synonym} \ \ L & = Length \\ \textbf{type-def} & Mass & = UNIV :: unit \ set \ .. \ \textbf{setup-lifting} \ \ type-definition-Mass \\ \textbf{type-synonym} \ \ M & = Mass \\ \dots \end{array}
```

The following instantiation proof places the freshly constructed type symbol L in the class basedim-type by setting its semantic interpretation to the corresponding value in the Dimension-type.

```
\label{eq:continuity} \begin{array}{l} \textbf{instantiation} \ \ Length :: basedim-type \\ \textbf{begin} \\ \textbf{definition} \ [si\text{-}eq] : dim\text{-}ty\text{-}sem\text{-}Length \ (\alpha :: Length \ itself) = L \\ \textbf{instance} < \textbf{proof} \ ommitted > \\ \textbf{end} \end{array}
```

Note that Isabelle enforces a convention to name the definition of an operation assumed in the interface of the class to be the concatenation of the interface name (e.g. dim-ty-sem) and the name of the class intantiation (e.g. Length). For the other 6 base-types we proceed analogously.

Dimension type expressions can be constructed by multiplication and division of the base dimension types above. Consequently, we need to define multiplication and inverse operators at the type level as well. On the class of dimension types (in which we have already inserted the base dimension types), the definitions of the type constructors for inner product and inverse is straightforward.

```
\textbf{typedef}\ ('\alpha :: dim-type,\ '\beta :: dim-type)\ DimTimes\ (\textbf{infixl}\cdot 69) = UNIV :: unit\ set\ .. \textbf{setup-lifting}\ type-definition-DimTimes
```

The type $'\alpha \cdot '\beta$ is parameterised by two types, $'\alpha$ and $'\beta$ that must both be elements of the dim-type class. As with the base dimensions, it is a unitary type as its purpose is to represent dimension type expressions. We instantiate dim-type with this type, where the semantics of a product dimension expression is the product of the underlying dimensions. This means that multiplication of two dimension types yields a dimension type.

```
instantiation DimTimes :: (dim-type, dim-type) dim-type
```

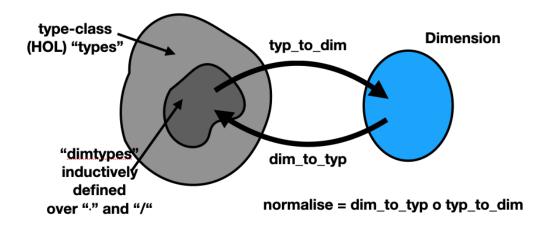


Figure 2: The "Inductive" subset of dim-types interpreted in SML Lists

```
\begin{array}{l} \textbf{begin} \\ \textbf{definition} \ dim\text{-}ty\text{-}sem\text{-}DimTimes :: ('\alpha \cdot '\beta) \ itself \Rightarrow Dimension \ \textbf{where} \\ [si\text{-}eq] : \ dim\text{-}ty\text{-}sem\text{-}DimTimes \ x = (dim\text{-}ty\text{-}sem \ TYPE('\alpha)) \cdot (dim\text{-}ty\text{-}sem \ TYPE('\beta)) \\ \textbf{instance by} \ (intro\text{-}classes, \ simp\text{-}all \ add: \ dim\text{-}ty\text{-}sem\text{-}DimTimes\text{-}def, \ (transfer, \ simp)+) \\ \textbf{end} \end{array}
```

Thus, the semantic interpretation of the product of two *dim-type*'s is a homomorphism over the product of two dimensions. Similarly, we define inversion of dimension types and prove that dimension types are closed under this.

```
typedef '\alpha DimInv ((--1) [999] 999) = UNIV :: unit set .. setup-lifting type-definition-DimInv instantiation DimInv :: (dim-type) dim-type begin definition dim-ty-sem-DimInv :: ('-1) itself \Rightarrow Dimension where [si-eq]: dim-ty-sem-DimInv x= inverse (dim-ty-sem TYPE('\alpha)) instance by (intro-classes, simp-all add: dim-ty-sem-DimInv-def, (transfer, simp)+) end
```

Finally, we introduce some syntactic sugar for such as ' α^4 for ' $\alpha \cdot '\alpha \cdot '\alpha \cdot '\alpha$ or ' α^{-4}

By the way, we also implemented two morphisms on the SML-level underlying Isabelle, which is straight-forward and ommitted here. These functions yield for:

```
\begin{array}{l} \mathbf{ML} (\\ \textit{Dimension-Type.typ-to-dim} \ @\{\textit{typ} \ L^{-2} \cdot M^{-1} \cdot T^4 \cdot I^2 \cdot M\}; \\ \textit{Dimension-Type.dim-to-typ} \ [\textit{1,2,0,0,0,3,0}]; \\ \textit{Dimension-Type.normalise} \ @\{\textit{typ} \ L^{-2} \cdot M^{-1} \cdot T^4 \cdot I^2 \cdot M\} \\ \end{pmatrix}
```

the system output:

$$\begin{array}{l} val~it=[^\sim 2,~0,~4,~2,~0,~0,~0]\colon int~list\\ val~it=L\cdot M^2\cdot N^3\colon \mathbf{typ}\\ val~it=L^{-2}\cdot T^4\cdot I^2\colon \mathbf{typ} \end{array}$$

6 ISQ Quantity and SI Types

7 Validation by the VIM and the 'Brochure'

8 Related Work and Conclusion

This work has drawn inspiration from some previous formalisations of the ISQ and SI, notably Hayes and Mahoney's formalisation in Z[4] and Aragon's algebraic structure for physical quantities[1]. To the best of our knowledge, our mechanisation represents the most comprehensive account of ISQ and SI in a theory prover.

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