



Prevention of Street Harassment Through Constrained Shortest Path Algorithms



Isabel Mora
Report and
algorithms



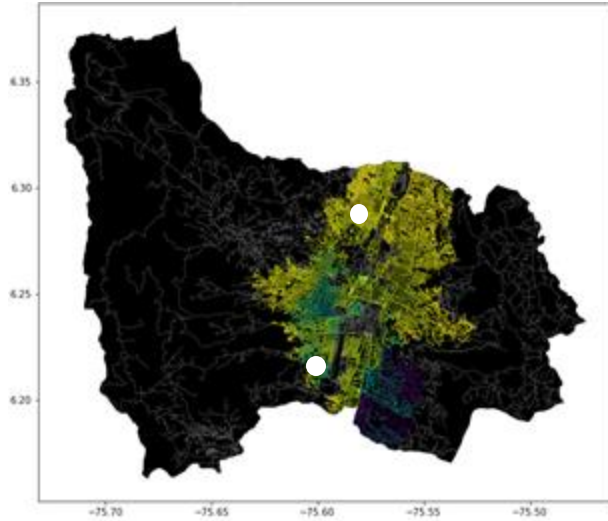
**Andrea
Serna**
Literature review



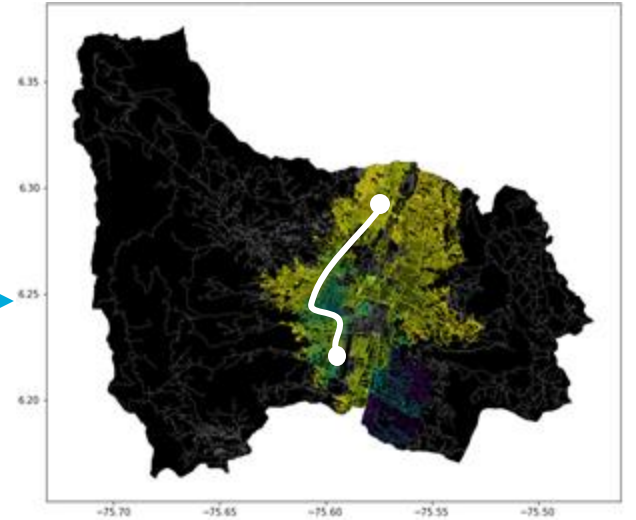
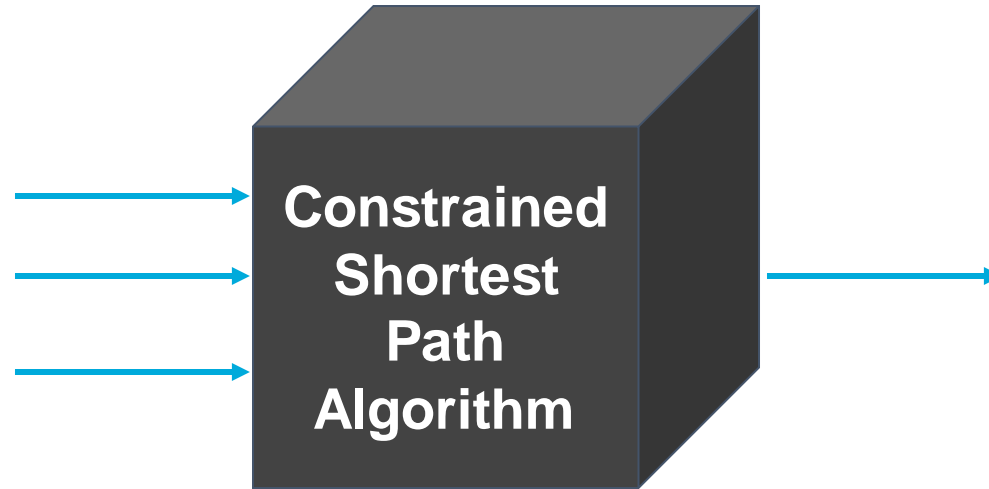
**Mauricio
Toro**
Data preparation



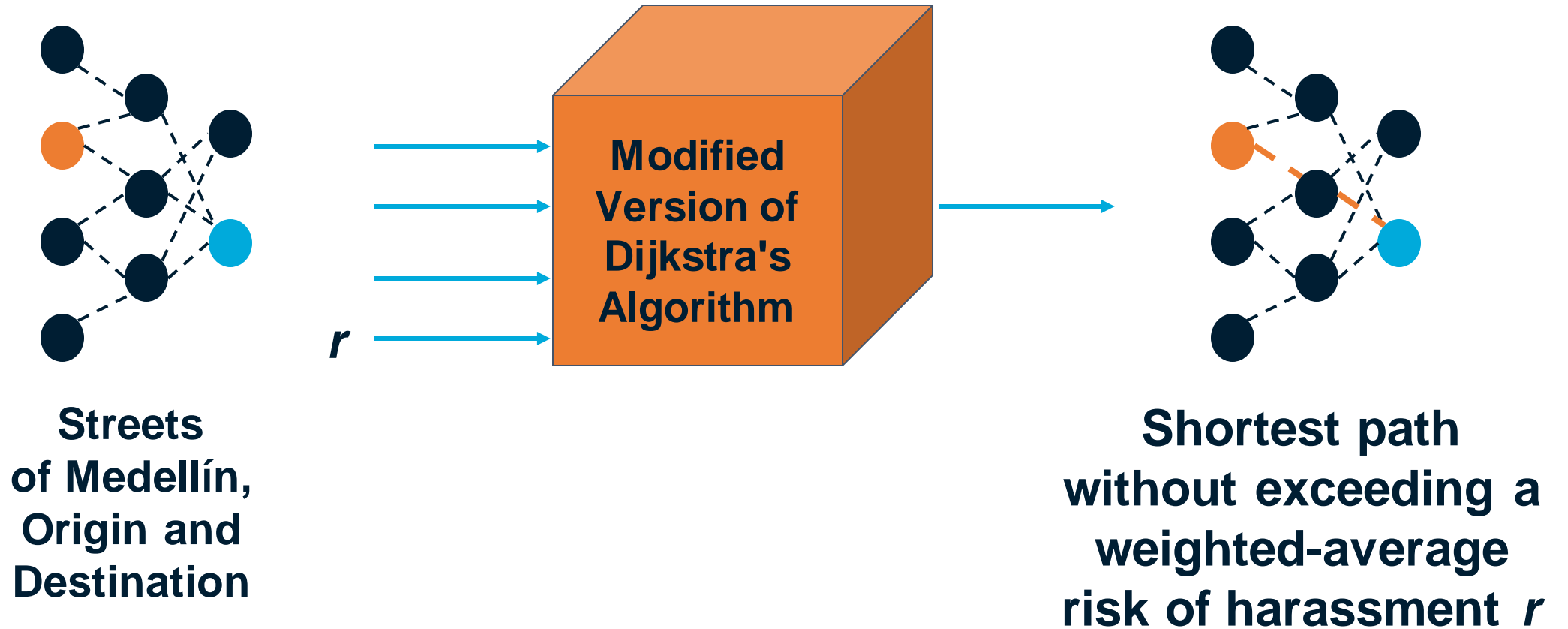
Problem Statement



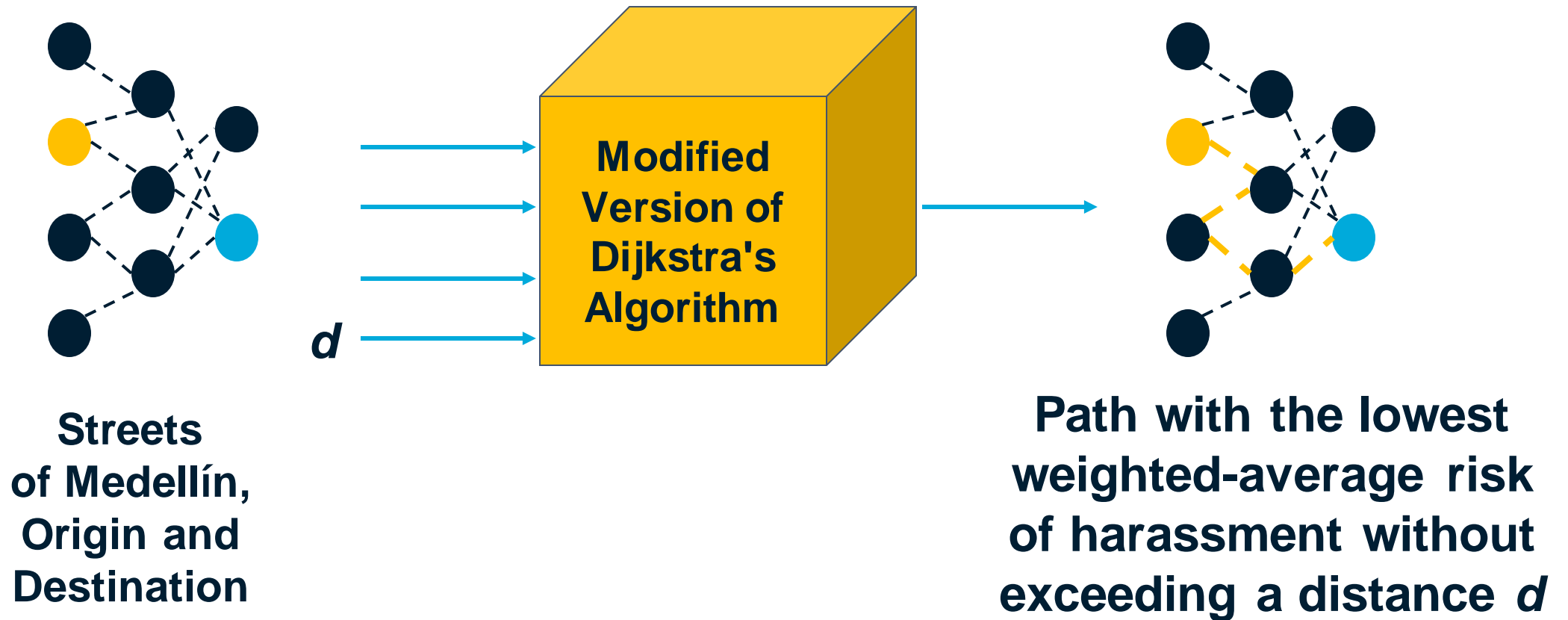
**Streets of Medellín:
Origin and
Destination**



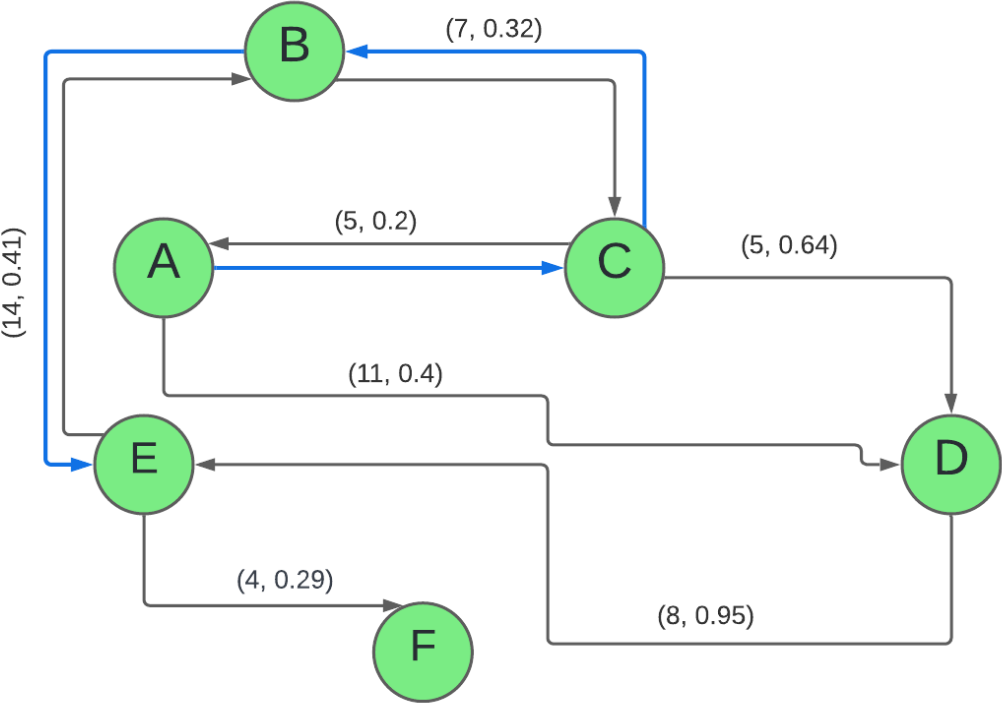
**Constrained
Shortest
Paths**



Second Algorithm



Algorithm Explanation



Example Execution:

Source = A
Destination = E
Max Risk = 0.5

Shortest path from A to E
without exceeding av. risk of 0.5
A -> C -> B -> E

Vertex	Shortest Distance from A	Prev. Vertex	Average Weighted Risk
A	0	-1	0
B	∞ , 12	C	0.27
C	∞ , 5	A	0.2
D	∞ , 11, 10	A, C	0.4, 0.42
E	∞ , 26	B	0.34
F	∞		

Queue: {A, C, D, B, E}



Modified Dijkstra's Algorithm for the
Constrained Shortest Path Problem



<https://github.com/isabelmorar/ST0245-001>

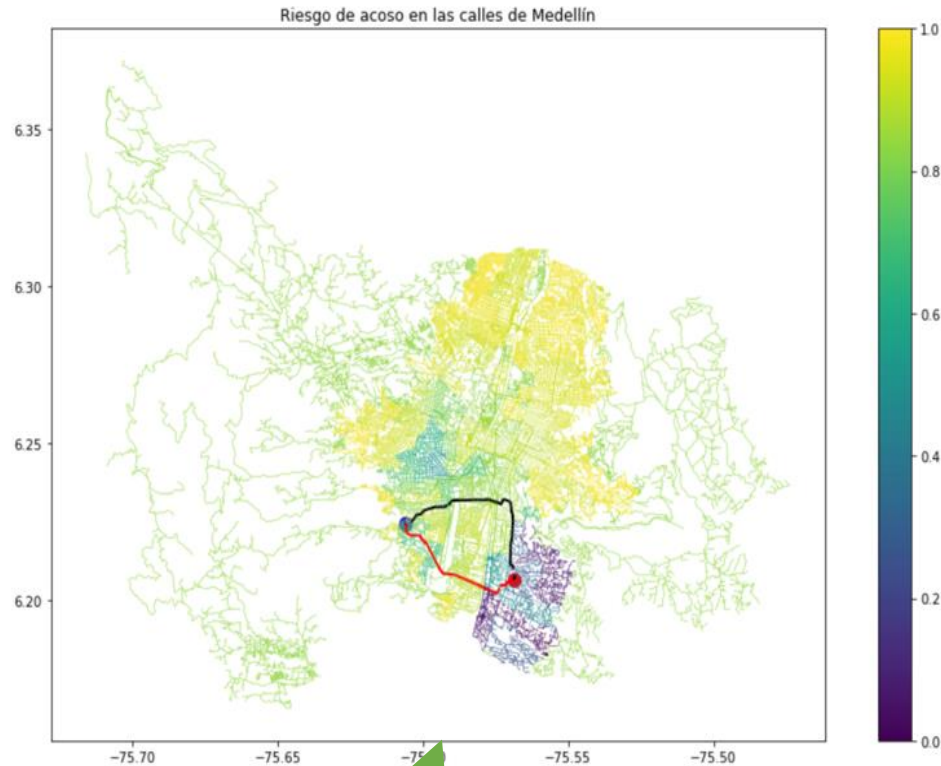


Example Output



Shortest path (red) –
Total Distance: 5903.91 meters
Average Risk: 0.688

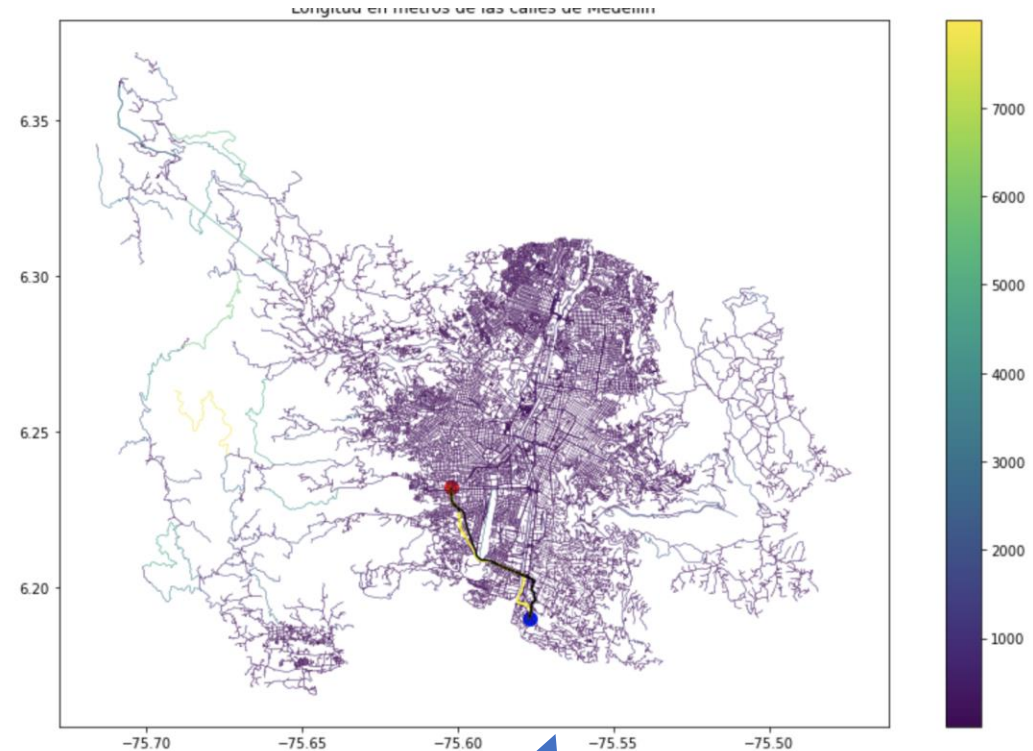
Shortest path without exceeding risk of 0.65 (black) –
Total Distance: 7644.081 meters
Average Risk: 0.647



Carrera 39 to Carrera 83

Path with lowest risk (yellow) –
Total Distance: 4302.909 meters
Average Risk: 0.727

Path with lowest risk without exceeding distance of 4100 meters (black) –
Total Distance: 4056.151 meters
Average Risk: 0.823



Calle 30A to Calle 16A

	Time Complexity	Memory Complexity
Modified Dijkstra's with Adjacency List	$O((V + E) \log V)$	$O(V + E)$

Time and memory complexity of the modified version of Dijkstra's algorithm, where V is the number of nodes and E is the number of edges in the graph. Specifically, V represents the intersections and E represents the streets in Medellin's map.



Shortest Path Results



Origin	Destination	Path	Maximum weighted-average risk of harassment
Universidad EAFIT	Universidad de Medellín	$d = 6142.57 \text{ m}$ $r = 0.758$	0.84
Universidad de Antioquia	Universidad Nacional	$d = 860.19 \text{ m}$ $r = 0.845$	0.85
Universidad Nacional	Universidad Luis Amigó	$d = 1910.13 \text{ m}$ $r = 0.842$	0.845

Path with shortest distance without exceeding a weighted average risk of harassment r .

Lowest Risk Results



Origin	Destination	Path	Maximum distance (meters)
Universidad EAFIT	Universidad de Medellín	$r = 0.719$ $d = 6183.71 \text{ m}$	7000
Universidad de Antioquia	Universidad Nacional	$r = 0.865$ $d = 815.44 \text{ m}$	820
Universidad Nacional	Universidad Luis Amigó	$r = 0.849$ $d = 1472.52$	1500

Path with lowest weighted-average risk of harassment without exceeding a distance d .

Algorithm Execution Times



Origin



Destination



Execution Time

**UNIVERSIDAD
EAFIT®**



11.787 seconds



**UNIVERSIDAD
DE ANTIOQUIA**



**UNIVERSIDAD
NACIONAL
DE COLOMBIA**



8.095 seconds



**UNIVERSIDAD
NACIONAL
DE COLOMBIA**



8.292 seconds

Web Development

- • • • •
- Graphical display of calculated paths

- • • • •
- Interactive user interface

Statistics

- • • • •
- Improve the numerical representation of sexual harassment risk

- • • • •
- Consider user demographics

Optimization

- • • • •
- Bi-objective optimization

- • • • •
- Calculate optimal path based on risk and distance simultaneously.

S & M 4

- • • • •
- Traffic Estimation

- • • • •
- Predict ideal paths through simulations of different scenarios



Mora, I., Serna, A., & Toro, M.
(2022, May 18). Prevention of
Street Harassment Through
Constrained Shortest Path
Algorithms. Universidad EAFIT.
<https://doi.org/10.31219/osf.io/9fr32>





THANK YOU!

This work was possible thanks to the help of the teacher assistants of Data Structures and Algorithms I. The first author is grateful to her parents for supporting her education at Universidad Eafit. All the authors would like to thank the "Vicerrectoría de Descubrimiento y Creación", of Universidad EAFIT, for their support on this research.