# MPI (Message Passing Interface)

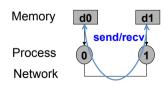
Part I

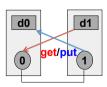
Jian-Jin LI

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# MPI Programming Model

- Communication model
  - Message passing (two side operation)
  - → Remote memory access (one side operation, MPI-2)

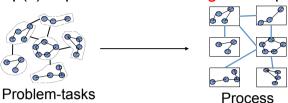




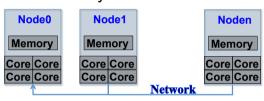
- Suitable architectures
  - ♦ Distributed memory MIMD
  - ♦ Cluster
  - Network of workstations

# MPI Programming Model

□ Group(s) of process to solve together a problem

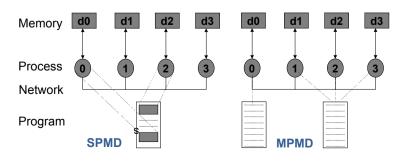


Distributed memory API



# MPI Programming Model

- □ SPMD / MPMD
  - ♦ Single / Multiple Program Multiple Data
  - ♦ Data distributed over process
  - Communication between process via message passing



### **MPI** Characteristics

- Message passing library, include
  - Environment management routines
  - ♦ Point to point communication
  - Datatypes management
  - Collective communications

to use with C, C++, Fortran

- Groups of process and communicators
- Process topologies
- ♦ Parallel I/O. RMA. dynamic process (MPI-2)
- Non-blocking collective communication, RMA improvement. parallel programming environment (MPI-3)
- Participants
  - ♦ Vendors: IBM, Intel, Meiko, Cray, ...
  - ♦ Libraries: PVM, Zipcode, Express, Linda, ...
  - Universities: San Francisco, Santa Barbara, ...

# MPI – Free Implementations

- MPICH
  - ♦ Leader: Argonne National Laboratory



- http://www.mpich.org
- Open MPI
  - ♦ û LAM (Ohio Supercomputer Center)
  - http://www.open-mpi.org



**Ohio Supercomputer Center** 

### References

- William Gropp, Ewing Lusk and Anthony Skiellum. Using MPI, 2 volume set, The MIT Press, 01/2000, ISBN: 0-262-57134-X.
- ♦ Peter Pachero, Parallel Programming with MPI. Morgan Kaufmann, 01/1997, ISBN: 1-55860-339-5
- http://www-unix.mcs.anl.gov/mpi
- http://www.mpi-forum.org/docs/docs.html
- http://www.idris.fr/data/cours/parallel/mpi/ choix doc.html

# **Environment Management Routines**

□ hello.c

```
#include "mpi.h" MPI's header file
#include <stdio.h>
int main(int argc, char **argv)
   int myrank, nbprocs;
   MPI Init( &argc, &argv ); Execution environment initialization
   MPI Comm rank ( MPI COMM WORLD, &myrank );
   MPI Comm size ( MPI COMM WORLD, &nbprocs);
   printf( " Hello from proc. %d of %d\n ",
            myrank, nbprocs);
   MPI Finalize(); End of MPI execution
   return 0:
```

# Compiling and Running MPI applications

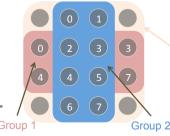
- Application implementation
  - → using C, C++ or Fortran and MPI library for process and communication management
- Compiling
  - ♦ mpicc hello.c -o hello
- Execution
  - ♦ 8 process on local node
    - mpiexec -n 8 ./hello
  - ♦ 16 process on nodes h1-4, 1 mpd run on each node mpiexec -hosts h1, h2, h3, h4 -n 16 ./hello

### Communicators

mpiexec -n 16 ./test

Communicators do not need to contain all processes in the system

Every process in a communicator has an ID called as "rank"



Can make copies of this

The same process might have different ranks in different communicators

program, there is one predefined communicator MPI COMM WORLD

When you start an MPI

communicator (same group of processes, but different "aliases")

### MPI's world

- □ Group of process and Communicator
  - ♦ MPI process are enrolled into groups
  - ♦ Group + context = Communicator
  - ♦ Default communicator: MPI COMM WORLD
  - ♦ Process identification (rank): 0, 1, ..., size-1
  - ♦ MPI Comm rank ( MPI COMM WORLD, &rank );
  - ♦ MPI Comm size( MPI COMM WORLD, &size );



MPI COMM WORLD

# Environment management routines

- MPI Abort
  - †int MPI Abort(MPI Comm comm, int errorcode);
  - terminates all process of communicator if exception: ex. malloc
- MPI Get processor name

  - return the processor name and its length
  - ♦ name buffer size: MPI MAX PROCESSOR NAME
- MPI Wtime / MPI Wtick
  - double MPI Wtime(void); double MPI Wtick();
  - return an ellapsed wall clock time in seconds / the number of seconds between successive clock ticks.

### Point to Point Communication

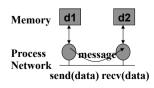
#### □ p2p.c

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### Point to Point Communication

- Communication between two process
  - Source and Destination
- ☐ Message = header + data
- Data conversion if necessary
- □ Transmission mechanism





#### Point to Point Communication

p2p.c (continue)

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#### Point to Point Communication

- Blocking communication
  - ♦ Blocking send and receive: MPI\_Send, MPI\_Recv
- □ Parameters of MPI Send and MPI Recv
  - ♦ Data address
  - Elements number of data
  - ♦ Type of data elements: MPI\_Datatype
  - ♦ Source or Destination of the message (MPI\_ANY\_SOURCE)
  - $\diamond$  Tag of message (MPI\_ANY\_TAG), may be used to indicate different type of message

  - Status: MPI\_Status (MPI\_SOURCE, MPI\_TAG,
     MPI\_ERROR)

# MPI Datatypes

□ Similar to C, examples:

MPI datatype	C datatype	
MPI_CHAR	signed char	
MPI_INT	signed int	
MPI_UNSIGNED	unsigned int	
MPI_FLOAT	float	
MPI_DOUBLE	double	
MPI_BYTE		

#### Complex datatypes

- ♦ MPI PACKED
- ♦ Derived types: structure, colonne of matrix ...

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# **Blocking Communication**

- Features
  - ♦ Completion of MPI\_Send means send variable can be reused
  - ♦ Completion of MPI\_Recv mean receive variable can be read
  - Cause synchronization -> Increase communication time
  - ♦ Affect the performance of parallel program
- Solution
  - ♦ Non-blocking communication

### Point to Point Communication

- Two side operation
- ⇒ a Send must be matched by a Recv
- □ Safe program in blocking communication

- 1

# Non-blocking Communication

Operation in 2 steps

- Completion: MPI\_WAIT(&request, &status);
- Test of completion:

```
MPI_TEST(&request, &flag, &status);
```

- Avoid dead lock (ex. T.16 with non-blocking send/recv)
- Allow communication / computation overlapping
- Persistent request can be used if many communication

# Non-blocking Communication

#### Example

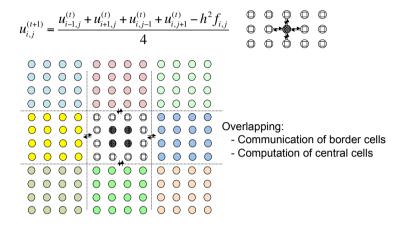
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# Message test routines

- MPI Probe/ MPI Iprobe
  - Availability test of message
  - Where is it from ?
  - What is its length ?

### Overlapping Communication/Computation

2D Poisson problem: Jacobi's algorithm



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# Design of Parallel Program

- □ Those who have to think
  - Launch of program with various number of process
  - Load balancing
  - Change of data over the time
  - Prefer local communication than distant one
  - ♦ Performance of parallel program, yet ?
- Another solutions
  - Collective communication
  - Decrease the communication number
    - Data grouping for communication

#### Collective Communication

- What is it?
  - ♦ Communication involving all processes of a group
- Objective
  - Increase the performance of parallel program
- □ How?
  - ♦ By reduce of idle processes ⇒ decrease the communication time
- Use cases
  - ♦ When I/O
  - → Parallel algorithms need collective communication

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### Collective Communication

□ Example - Broadcast of the dimension of a image

```
int myrank, size, dims[2];
int i, tag=30;

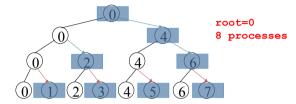
MPI_Comm_rank( MPI_COMM_WORLD, &myrank );
MPI_Comm_size( MPI_COMM_WORLD, &size );

if (myrank == 0) {
    /* Fill dims */
    for ( i=1; i<size; i++)
        MPI_Send( dims, 2, MPI_INT, i, tag, MPI_COMM_WORLD);
}
else
    MPI_Recv(dims, 2, MPI_INT, 0, tag, MPI_COMM_WORLD);</pre>
```

```
MPI_Bcast(dims, 2, MPI_INT, 0, MPI_COMM_WORLD);
Collective communication:
Steps - O(log,(size)) with binary tree
```

#### Collective Communication

- Broadcast
- ♦ A process (root) has a message to send to others
- ♦ Possible implementation:

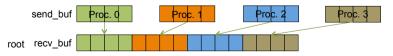


♦ MPI routine:

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#### Collective Communication

- Gather / Reduce
  - ♦ The processes of a group have data to merge together
  - ♦ MPI routines:



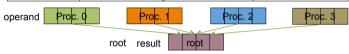
The length of recvbuf must be the one of sendbuf \*
number\_of\_processes, data in recvbuf are stored in the
order of processes

### Collective Communication

#### Reduce

int MPI\_Reduce( void \*operand, void \*result,
 int count, MPI\_Datatype datatype, MIP\_Op op,
 int root, MPI\_Comm comm );

Operation	Meaning	Operation	Meaning
MPI_MAX	Maximum	MPI_LOR	logical OR
MPI MIN	Minimum	MPI BOR	bitwise OR
MPI SUM	Sum	MPI LXOR	XOR
MPI PROD	Product	MPI BXOR	
MPI LAND	Logical AND	MPI MAXLOC	max or min +
MPI_BAND	bitwise AND	MPI_MINLOC	index



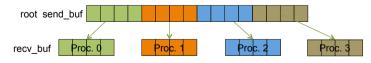
operand and result have the same length

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### **Data distribution**

#### MPI\_Scatter

♦ The root has data to distribute to the processes of the group



The length of sendbuf must be the one of recybuf \* number of processes.

#### Collective Communication

#### Example

```
int myrank, size;
double myresult, globalresult, *resultvec;

MPI_Comm_rank( MPI_COMM_WORLD, &myrank );
MPI_Comm_size( MPI_COMM_WORLD, &size );

/* computation of myresult */

resultvect = (double *)malloc(size*sizeof(double));
MPI_Gather( &myresult, 1, MPI_DOUBLE, resultvec, 1, MPI_DOUBLE, 0, MPI_COMM_WORLD );

MPI_Reduce( &myresult, &globalresult, 1, MPI_DOUBLE, MPI_SUM, 0, MPI_COMM_WORLD );
```

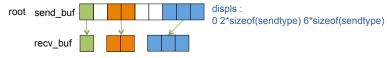
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#### Data distribution

#### MPI\_Gatherv / MPI\_Scatterv

More flexibility in message length of processes and message position in root memory

```
int MPI_Scatterv( void *sendbuf, int *sendcnts, int *displs, MPI_Datatype sendtype, void *recvbuf, int recvent, MPI_Datatype recvtype, int root, MPI_Comm comm );
```



### Collective Communication

- Barrier synchronization
  - Make a appointment for all processes

```
int MPI_Barrier( MIP_Comm comm );
```

- Use case: time measurement.
- Time measurement for each process

Execution time of program: the one of the slowest process

# Pack / Unpack

- Objective
  - Put the heterogeneous non-contiguous data together to be sent at one time
- Example

# Data grouping

- Need to send heterogeneous data
- Methods
  - ♦ MPI Pack/MPI Unpack
    - o Use of a buffer of bytes to send arbitrary data
  - Derived types
    - Corresponding MPI datatype for structure
- Choice of methods
  - ♦ MPI Pack/MPI Unpack
    - Heterogeneous data to send a few times
    - Communication of data of variable length
  - Derived types
    - Heterogeneous data to send repeatedly
    - $_{\circ}$  Non-contiguous data of the same type  $\Rightarrow$

```
MPI_Type_vector, MPI_Type_indexed
```

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# Pack / Unpack

■ Example (continue)

# Derived type

- Objective
  - ♦ Define heterogeneous data access
- General method
  - ♦ Build
  - Validate
  - Destroy

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# Derived types of matrix

Column of a matrix (in C): a derived type for data of a same type, evenly spaced

```
MPI_Type_vector( int blocnumber, int bloclength, int stride,
MPI_Datatype oldtype, MPI_Datatype *newtype );
```

□ Example Number of lines Number of columns

```
MPI_Type_vector( 4, 1, 6, MPI_DOUBLE, &type_column );
MPI_Type_Commit( &type_column ) ;
```



```
MPI_Type_free( &type_column );
```

# Derived type

□ General method (continue)

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# Distribution of square matrix

#### ■ MPI Scatter – 1 column per process

♦ Condition: processes number = order of matrix

# Distribution of square matrix

□ MPI\_Scattery – 1 column per process

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# Distribution of square matrix

□ MPI Scatter – several columns per process

```
/* Definition du type colonne (matrice globale) */
MPI Type vector(order mat, 1, order mat, MPI DOUBLE
MPI Type commit(&col);
MPI Type create resized(col, (MPI Aint)0,
  (MPI Aint) (1*sizeof(double)), &type col);
MPI Type commit(&type col);
/* Definition du type colonnes des processus */
Nb cols = order mat / size;
MPI Type vector (order mat, 1, nb cols, MPI DOUBLE, &colProc);
MPI Type commit(&colProc);
MPI Type create resized(colProc, (MPI Aint)0,
   (MPI Aint) (1*sizeof (double)), &type colProc);
MPI Type commit(&type colProc);
lA = (double *) malloc( nb cols * order mat * sizeof(double) );
MPI Scatter ( A, nb cols, type col, lA, nb cols,
             type colProc, 0, MPI COMM WORLD );
```

# Distribution of square matrix

- □ MPI Scatter 1 bloc of columns of matrix
  - ♦ Example: 2 columns per process

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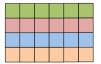
# Derived types of matrix

□ Line of a matrix (in C): a derived type for contiguous data of a same type

Example

```
Number of columns
```

```
MPI_Type_contiguous( 6, MPI_DOUBLE, &type_line );
MPI_Type_Commit( &type_line ) ;
```



```
MPI_Type_free( &type_line );
```

# MPI\_Type\_indexed

Derived type for data of a same type, not evenly spaced

Example

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### **Intra-Communicators**

- □ A intra-communicator is composed of
  - ♦ A group of p processes: identified by a unique rank (0, ..., p-1)
  - ♦ Predefined groups: MPI GROUP EMPTY
  - ♦ A context: a system-defined object, each context is exclusive
  - Attributes: topology
  - ♦ A minimal intra-communicator = a group + a context
- Group / Communicator
  - ♦ Groups and communicators are associated.
  - ♦ Groups/Communicators are dynamics.
  - A process may be in several groups/communicators.
     It has a unique rank within each group/communicator.
  - Creation from existing groups/ communicators



### **MPI Communicators**

- What is it?
  - Provide a separate communication space to subset of processes
  - System-defined object
  - Provide safe communications
  - ♦ Examples: MPI COMM WORLD, MPI COMM SELF
- Types of communicators
  - Intra-communicators
    - o A group of processes, each can send message to all other
    - Ease organization of task groups
    - Allow collective communication in a subset of processes
  - ♦ Inter-communicator
  - For sending message between processes belong to disjoint intracommunicators

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#### Intra-Communicators

Some defined functions

 MPI\_Comm\_group : get the processes group of a given communicator

```
int MPI_Comm_group( MPI_Comm comm, MPI_Group *group);
```

MPI\_Group\_incl : create un new group from a subset processes of a existing group, processes are reordered

MPI\_Comm\_create : create un new comm. from a group of processes, collective operation

MPI Comm free(&comm); MPI Group free(&group);

### Group / Intra-Communicator - Example

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# **MPI** Topologies

#### Characteristics

- → Topologies define different addressing scheme of processes
- Fit the communication pattern of parallel application to processes connection
- Topology is virtual in MPI (machine independent)
- Can match the network for performance (in theory)

#### □ Types of MPI topologies

- Cartesian topology (grid/torus)
- For Cartesian communication pattern
- Graph: general purpose case defined by
  - List of nodes-processes
  - Neighbours number
  - List of edges-connection between processes



### Intra-Communicators

#### Other constructors

- MPI\_Comm\_dup : create un new group from a subset processes of a existing group
- MPI\_Comm\_split: partition of a given communicator into disjoint sub-communicators. A sub-communicator is composed by processes with the same color.

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# MPI Topologies

#### Information for grid/torus creation

- Number of dimensions
- Order of dimensions
- ♦ Wrap on (torus) or no (grid)

#### Predefined functions

→ Grid constructor

♦ Transformation rank <-> coordinates

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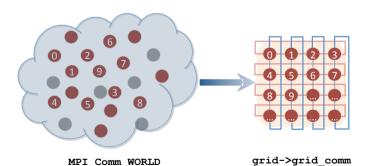
# Topology and intra-communicator

- □ Torus creation
  - Data structure

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# Topology and intra-communicator

□ Torus creation – torus 2D communicator



P.S.: grid->grid comm = MPI COMM NULL for process not in torus

# Topology and intra-communicator

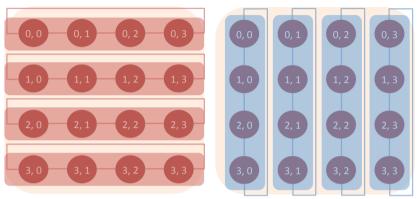
#### ■ Torus creation

# Topology and intra-communicator

#### Torus creation

# Topology and intra-communicator

□ Torus creation – row/column communicators



row communicators

column communicators

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# MxM – Algorithm of Fox

Parallel algorithm

/\* Process (i, j) computes  $C_{ij} = \sum_{k=0}^{q-1} A_{ik} B_{kj} */$  for ( k=0; k<q; k++) {

- 1. Select a block of A for each row of the grid
- 2. Broadcast of the chosen block in each line for each process
  - 3. Multiply the block of A with the block of B in each process

$$A_{i,u}$$
 avec  $u = (i+k) \mod q$  for the row  $i */$  /\* example: bleu, rouge, vert, orange if  $q=4 */$ 

$A_{00}$	$A_{01}$	$A_{02}$	$A_{03}$
$A_{10}$		$A_{12}$	
$A_{20}$		$A_{22}$	
$A_{30}$		$A_{32}$	$A_{33}$

$B_{00}$	$B_{01}$	$B_{02}$	$B_{03}$
		$B_{12}$	
		$B_{22}$	
${\rm B}_{30}$		$B_{32}$	

# MxM – Algorithm of Fox

#### Assumptions

- ♦ A, B: n-by-n matrices
- $\Rightarrow$  N: the number of processes and N=q<sup>2</sup>, n'=n/q
- ♦ Input data of Process (i, j):

$$A_{ij} = \begin{pmatrix} a_{i} * n', j * n' & \dots & a_{(i} * n', (j+1) * n'-1 \\ \vdots & \dots & \vdots \\ a_{(i+1)} * n'-1, j * n' & \dots & a_{(i+1)} * n'-1, (j+1) * n'-1 \end{pmatrix} \begin{pmatrix} A_{00} A_{01} | A_{02} A_{0} \\ A_{10} A_{11} | A_{12} A_{1} \\ A_{20} A_{21} | A_{22} A_{2} \\ A_{30} A_{31} | A_{32} A_{3} \end{pmatrix}$$

$$B_{ij} = \begin{pmatrix} b_{i} * n', j * n' & \dots & b_{i} * n', (j+1) * n'-1 \\ \vdots & \dots & \vdots \\ b_{(i+1)} * n'-1, j * n' & \dots & b_{(i+1)} * n'-1, (j+1) * n'-1 \end{pmatrix} \begin{pmatrix} B_{00} B_{01} | B_{02} B_{0} \\ B_{10} B_{11} | B_{12} B_{1} \\ B_{20} B_{21} B_{22} B_{2} B_{2} B_{2} B_{22} B_{22} B_{22} B_{23} B_{23} B_{24} B$$

**Principle:**  $A_{ij}$  remain in each P(i,j),  $B_{ij}$  move

Data distribution with q=4

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# MxM – Algorithm of Fox

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# MxM – Algorithm of Fox

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### Distribution of blocks of matrix

## Summary

#### Processes group and communicator

- ♦ A process is identified by a rank in a group
- A communicator is composed by a group of processes and a context.
   It may have some attributes (ex. topology).

#### Communication types

- ♦ One-to-one
- Collective communication
- → Blocking / non-blocking communication

#### Advanced data types

- ♦ Pack/Unpack
- Derived data type

#### Performance of parallel program

- Collective communication
- non-blocking communication -> overlapping communicationcomputation

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### Gather of blocks of matrix