Issac Zheng	CS 492 Final	Bonelli	I pledge r	my honor that I have abided by the Stevens Honor System
<u>Tupabilitu</u>				<u>Typability Notes</u>
let f = proc (x: int) { proc (y: int) { x+y }} in (f 2)				-if 3 then 88 else 69 //not typable: 3 is not a bool
Typable. Take Γ to be $\{\}$ and t to be int -> int.				-let $x = 3$ in $(3 x)$ //not typable: 3 is not a function
{} - let f = proc (x:int) { proc (y: int) { x+y }} in (f 2) : int -> int				-proc(x) { (x 3) } //typable: x <u>has</u> to be a func ocam
T	Var TVar			-Γ is the typing environment
{x:int, y:int} - x: int {				$-\Gamma$ - e:t -> look into Γ and retrieve the type of e
{x:int, y:int}	TAdd			- Γ , id:t1 - e:t2 -> extend Γ with id of type t1
	TProc		TVar T	(extend_env id t1), then retrieve type of e
	{ x+y } : int -> int	{f:int->int->int} - f:int	->int->int {f:int->int->int} - 2: int	-{x: int} - x+2: int -> valid
				TApp -{x: bool} - x+2: int -> invalid
{} - proc (x:int) { proc (y: int) { x+y }} : int -> int -> int {f: int->int} - (f 2): int-> int				TLet -functions are type (example) t1->t2->t3, where t3 is
	$\{\}\ - \text{let f = proc (x:int) } \{ \text{ proc }$	(y: int) { x+y }} in (f 2) : in	-> int	the return type and t1 is the first param, t2 is
				second. t2 only exists if it's a nested function
proc (x: int) { pro	oc (y: bool) $\{if y then x$	else x-1} }		(whose purpose is to enable a second parameter).
• .	} and t to be int -> bool -> int.			
{} - proc (x: int) {proc	(y: bool) {if y then x else x-1}			-typing expr, env, type
		{x: int, y: bool}(x) = ir	nt /arTInt	
{x: int, u: bool}(u) = boo	ol {x: int, y: bool}(x) = int	·	t {x: int, y: bool} - 1: int	Dain lands are at at an
	TVarT			Pair Implementation
	ol {x: int, y: bool} - x: int		e e	Γ - e1: t1 Γ - e2: t2
	(v. int. u. book l- if	y then x else x-1: int	IIIE	TPair
		_	TProc	Γ - pair(e1, e2): t1*t2
$\{x: int\} \mid -proc(y: bool) \{if y then x else x-1\}: bool -> int$				Pair(e1,e2) ->
			TProc	chk expr e1 >>= fun t1 ->
$\{\}\ - proc(x: int) \{ proc(y: bool) \{ if y then x else x-1 \} \}: int -> bool -> int$				chk expre2 >>= funt2 ->
				return @@ PairType(t1,t2)
	{ proc (y: int) {x-y} } ir	n (f (3-6))		
• .	{} and f to be int->int->int			Γ - e1: t1*t2 Γ, id1:t1, id:t2 - e2:t
{} - let t=proc(x: int) { {x: int, y: bool}(x) = int	proc(y: int) {x-y} } in (f (3-6)):	int -> int		TUnpair
TVar				Γ - unpair(id1, id2) = e1 in e2: t
{x: int, y: int} - x: i				
	TSub		TintTin	t Unpair(id1,id2,e1,e2) ->
{x: int, y: int}		•	->i->i} - 3: i	chk_expr e1 >>= fun t ->
			TS	(material with
{x: int} - proc(y		i->i->i} - f: i->i->i		PairType(t1,t2) ->
		(f: i_si_si)		exteria_terivial (177)
{} - proc(x: int) { proc(y: int) {x-y} }: i->i->i			extend_tenv id2 t2 >>+	
	{} - let f=proc(x: int) { pro			CNK_expr e2
	() p. 55() (pro	-(3) (7. 9)) (1 (3	-//	> error "unpair: expected a pair")
Dahua				Debug Notes

<u>Debug</u>

Explicit-Refs

let a = 3

in let b = newref(if zero?(a) then 1 else 2)in debug(a)

Env := $[a \Rightarrow NumVal 3, b \Rightarrow RefVal 0]$ Store := [0 => NumVal 2]

let a = newref(2+1)in let b = newref(deref(a)+1)in debug(a)

 $Env := [a \Rightarrow RefVal 0, b \Rightarrow RefVal 1]$ Store:= $[0 \Rightarrow NumVal, 1 \Rightarrow NumVal 4]$ Implicit-Refs

let a = 2 in let b = 3

in begin

set a = b; debug(a)

end

Env := [a => RefVal 0, b => RefVal 1]Store := $[0 \rightarrow NumVal 3, 1 \rightarrow NumVal 3]$

let x=2 in let y=set x=3 in let z=proc(y) {proc(w) {set x=y+w}} in debug(3)

Env:= [x:=RefVal 0, y:=RefVal 1., z:=RefVal 2] Store:= [0:=NumVal 3, 1:=UnitVal $2:=ProcVal("y", proc(w) {set x=y+w}, [x:=RefVal]$ 0, y:=RefVal 1])]

Tupability Notes

- -if 3 then 88 else 69 //not typable: 3 is not a bool
- -let x = 3 in (3x) //not typable: 3 is not a function
- -proc(x) $\{(x 3)\}$ //typable: x has to be a func ocam
- $-\Gamma$ is the typing environment
 - $-\Gamma$ | e:t -> look into Γ and retrieve the type of e
 - - Γ , id:t1 |- e:t2 -> extend Γ with id of type t1 (extend_env id t1), then retrieve type of e
- $-\{x: int\} | -x+2: int -> valid$
- $-\{x: bool\} \mid -x+2: int -> invalid$
- -functions are type (example) t1->t2->t3, where t3 is the return type and t1 is the first param, t2 is second. t2 only exists if it's a nested function (whose purpose is to enable a second parameter).
- -typing expr, env, type

Pair Implementation

```
\Gamma |- e1: t1 \Gamma |- e2: t2
-----TPair
Γ |- pair(e1, e2): t1*t2
| Pair(e1,e2) ->
 chk_expr e1 >>= fun t1 ->
 chk_expr e2 >>= fun t2 ->
 return @@ PairType(t1,t2)
 Γ |- e1: t1*t2 Γ, id1:t1, id:t2 |- e2:t
-----TUnpair
 \Gamma |- unpair(id1, id2) = e1 in e2: t
| Unpair(id1,id2,e1,e2) ->
 chk_expr e1 >>= fun t ->
 (match t with
 | PairType(t1,t2) ->
   extend_tenv id1 t1 >>+
   extend tenv id2 t2 >>+
   chk_expr e2
 | _ -> error "unpair: expected a pair")
```

Debug Notes

- -newref evaluates the expression, then allocates it to the Store
 - -Store: 0 => ITE(IsZero?("a"), Int 1, Int 2) //wrong

Checked Implementation (* EXPLICIT-REFS *) | BeginEnd([]) -> return UnitType | BeginEnd(es) -> let rec check = function | [] -> return UnitType |[h] -> chk_expr h | h::t -> chk_expr h >>= fun _ -> check t in check es | NewRef(e) -> chk expr e >>= fun t -> return @@ RefTupe(t) | DeRef(e) -> chk expr e >>= fun t -> (match t with | RefType(t) -> return t -> error "deref: expected a reference type") | SetRef(e1,e2) -> chk expr e1 >>= fun t1 -> chk expr e2 >>= fun t2 -> (* check e1 is a ref, then check e2 = e1's type *) (match t1 with | RefType(t) -> if(t=t2)then return UnitType else error "setref: types of ref and expr don't match" |_ -> error "setref: expected a reference type") (* list *) | EmptyList(None) -> return @@ ListType(UnitType) | EmptyList(Some t) -> return @@ ListType(t) | Cons(h, t) -> chk expr h >>= fun t1 -> chk_expr t >>= fun t2 -> (match t2 with | ListType(t) \rightarrow if(t1=t) then return @@ ListType(t) else error "cons: type of head and tail don't match" -> error "cons: expected a list type") | IsEmptu(e) -> chk expr e >>= fun t -> (match t with | ListType(_) -> return BoolType | TreeType(_) -> return BoolType | _ -> error "isempty: expected a list type") | Hd(e) -> chk_expr e >>= fun te -> (match te with | ListType(t) -> return t -> error "hd: expected a list type" | Tl(e) -> chk expr e >>= fun te -> (match te with | ListType(t) -> return @@ ListType(t)

-> error "tl: expected a list type")

Checked Implementation

```
(* tree *)
| EmptyTree(None) ->
 return @@ TreeType(UnitType)
| EmptyTree(Some t) ->
 return @@ TreeType(t)
| Node(de, le, re) ->
 chk expr de >>= fun t1 ->
 chk_expr le >>= fun t2 ->
 chk expr re >>= fun t3 ->
 (* t2 and t3 must be trees of the same type as t1. *)
 if t2=TreeType(t1) && t3=TreeType(t1)
 then return @@ TreeType(t1)
 else error "node: types of don't match or subtrees aren't trees"
| CaseT(target,emptycase,id1,id2,id3,nodecase) ->
 chk expr target >>= fun tr ->
 chk_expr emptycase >>= fun tec ->
  (match tr with
  | TreeType(t) ->
     extend tenv id1 t >>+
     extend_tenv id2 (TreeType(t)) >>+
     extend tenv id3 (TreeType(t)) >>+
     chk_expr nodecase >>= fun tnc ->
     (* e2 and e3 must be of same type s *)
     if(tec=tnc)
     then return thc
     else error "caseT: types of emptycase and nodecase don't match"
 | _ -> error "caseT: expected a tree type"
```

the Store. *))

```
Record Example (Exam Implementation Guess)
 | Record(fs) ->
  sequence (List.map process_field fs) >>= fun evs ->
  return (RecordVal (addlds fs evs))
 | Proj(e,id) ->
  eval_expr e >>=
  fields_of_recordVal >>= fun f ->
  (match List.assoc opt id f with
  | None -> error "Field does not exist"
  | Some ( ,v) -> return v
    (*need ( ,v) bc now we have a true/false field for mutable*)
   (*List.assoc_opt still works bc f is a key(string)-value pair. Just that the value is now
     a tuple rather than an exp_val*)
 | SetField(e1,id,e2) ->
  eval_expr e1 >>=
  fields_of_recordVal >>= fun f ->
  eval_expr e2 >>= fun I ->
  (match List.assoc_opt id f with
  | None -> error "Field does not exist"
  | Some (b,v) ->
     match b with
     I false -> error "Error: Field not mutable"
     I true ->
      int_of_refVal v >>= fun t ->
      Store.set_ref g_store t | >>= fun _ -> return UnitVal
     (* mutable fields will be a RefVal with value of its value's address in Store. *)
     (* despite the pdf's example in 2.1, p.age will still return the RefVal value (the pointer value)
the actual NumVal. *)
     (* however, if you debug() after setting the ref, it'll show that the value has been corre
```

DO NOT PRINT THIS PAGE

Add:

- -type derivations
- -debug examples for explicit and implicit refs
- -Copy+Paste checked hw + guess final content and maybe do it? -> record, probably!!
- -figure out how to assign a bool to a var...
- -pair unpair code