

## Bottom-up Chart Parsing: the CKY algorithm

Parsing  
ISCL-BA-06

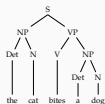
Çağrı Çöltekin  
ccolt@infs.uni-tuebingen.de

University of Tübingen  
Seminar für Sprachwissenschaft

Winter Semester 2020/21

Introduction CKY CKY

### Bottom-up parsing as search



S  $\rightarrow$  NP VP  
NP  $\rightarrow$  Det N  
VP  $\rightarrow$  V NP  
VP  $\rightarrow$  V  
Det  $\rightarrow$  a  
Det  $\rightarrow$  the  
N  $\rightarrow$  cat  
N  $\rightarrow$  dog  
V  $\rightarrow$  bites  
N  $\rightarrow$  bites

Ç. Çöltekin, INF | University of Tübingen

Winter Semester 2020/21 2 / 13

### Parsing so far

- Parsing is the task of automatic syntactic analysis
- For most practical purposes, context-free grammars are the most useful formalism for parsing
- We can formulate parsing as
  - Top-down: begin with the start symbol, try to produce the input string to be parsed
  - Bottom up: begin with the input, and try to reduce it to the start symbol
- Both strategies can be cast as search with backtracking
- Backtracking parsers are inefficient: they recompute sub-trees multiple times

Ç. Çöltekin, INF | University of Tübingen

Winter Semester 2020/21 3 / 13

Introduction CKY CKY

### Dealing with ambiguity

S  $\rightarrow$  NP VP  
NP  $\rightarrow$  Prn N  
NP  $\rightarrow$  Prn  
VP  $\rightarrow$  V NP  
VP  $\rightarrow$  V  
VP  $\rightarrow$  V S  
N  $\rightarrow$  duck  
V  $\rightarrow$  duck  
V  $\rightarrow$  saw  
Prn  $\rightarrow$  I  
Prn  $\rightarrow$  she  
Prn  $\rightarrow$  her

I saw her duck

1

Ç. Çöltekin, INF | University of Tübingen

Winter Semester 2020/21 4 / 13

Introduction CKY CKY

### Dealing with ambiguity

S  $\rightarrow$  NP VP  
NP  $\rightarrow$  Prn N  
NP  $\rightarrow$  Prn  
VP  $\rightarrow$  V NP  
VP  $\rightarrow$  V  
VP  $\rightarrow$  V S  
N  $\rightarrow$  duck  
V  $\rightarrow$  duck  
V  $\rightarrow$  saw  
Prn  $\rightarrow$  I  
Prn  $\rightarrow$  she  
Prn  $\rightarrow$  her

Prn

I saw her duck

2

Ç. Çöltekin, INF | University of Tübingen

Winter Semester 2020/21 5 / 13

Introduction CKY CKY

### Dealing with ambiguity

S  $\rightarrow$  NP VP  
NP  $\rightarrow$  Prn N  
NP  $\rightarrow$  Prn  
VP  $\rightarrow$  V NP  
VP  $\rightarrow$  V  
VP  $\rightarrow$  V S  
N  $\rightarrow$  duck  
V  $\rightarrow$  duck  
V  $\rightarrow$  saw  
Prn  $\rightarrow$  I  
Prn  $\rightarrow$  she  
Prn  $\rightarrow$  her

NP

Prn

I saw her duck

3

Ç. Çöltekin, INF | University of Tübingen

Winter Semester 2020/21 6 / 13

Introduction CKY CKY

### Dealing with ambiguity

S  $\rightarrow$  NP VP  
NP  $\rightarrow$  Prn N  
NP  $\rightarrow$  Prn  
VP  $\rightarrow$  V NP  
VP  $\rightarrow$  V  
VP  $\rightarrow$  V S  
N  $\rightarrow$  duck  
V  $\rightarrow$  duck  
V  $\rightarrow$  saw  
Prn  $\rightarrow$  I  
Prn  $\rightarrow$  she  
Prn  $\rightarrow$  her

NP

Prn

V

I saw her duck

4

Ç. Çöltekin, INF | University of Tübingen

Winter Semester 2020/21 7 / 13

Introduction CKY CKY

### Dealing with ambiguity

S  $\rightarrow$  NP VP  
NP  $\rightarrow$  Prn N  
NP  $\rightarrow$  Prn  
VP  $\rightarrow$  V NP  
VP  $\rightarrow$  V  
VP  $\rightarrow$  V S  
N  $\rightarrow$  duck  
V  $\rightarrow$  duck  
V  $\rightarrow$  saw  
Prn  $\rightarrow$  I  
Prn  $\rightarrow$  she  
Prn  $\rightarrow$  her

NP

Prn

V

I saw her duck

5

Ç. Çöltekin, INF | University of Tübingen

Winter Semester 2020/21 8 / 13

Introduction CKY CKY

### Dealing with ambiguity

S  $\rightarrow$  NP VP  
NP  $\rightarrow$  Prn N  
NP  $\rightarrow$  Prn  
VP  $\rightarrow$  V NP  
VP  $\rightarrow$  V  
VP  $\rightarrow$  V S  
N  $\rightarrow$  duck  
V  $\rightarrow$  duck  
V  $\rightarrow$  saw  
Prn  $\rightarrow$  I  
Prn  $\rightarrow$  she  
Prn  $\rightarrow$  her

NP

Prn

V

I saw her duck

6

Ç. Çöltekin, INF | University of Tübingen

Winter Semester 2020/21 9 / 13

Introduction CKY CKY

### Dealing with ambiguity

S  $\rightarrow$  NP VP  
NP  $\rightarrow$  Prn N  
NP  $\rightarrow$  Prn  
VP  $\rightarrow$  V NP  
VP  $\rightarrow$  V  
VP  $\rightarrow$  V S  
N  $\rightarrow$  duck  
V  $\rightarrow$  duck  
V  $\rightarrow$  saw  
Prn  $\rightarrow$  I  
Prn  $\rightarrow$  she  
Prn  $\rightarrow$  her

NP

Prn

V

I saw her duck

7

Ç. Çöltekin, INF | University of Tübingen

Winter Semester 2020/21 10 / 13

Introduction CKY CKY

### Dealing with ambiguity

S  $\rightarrow$  NP VP  
NP  $\rightarrow$  Prn N  
NP  $\rightarrow$  Prn  
VP  $\rightarrow$  V NP  
VP  $\rightarrow$  V  
VP  $\rightarrow$  V S  
N  $\rightarrow$  duck  
V  $\rightarrow$  duck  
V  $\rightarrow$  saw  
Prn  $\rightarrow$  I  
Prn  $\rightarrow$  she  
Prn  $\rightarrow$  her

NP

Prn

V

I saw her duck

8

Ç. Çöltekin, INF | University of Tübingen

Winter Semester 2020/21 11 / 13

Introduction CKY CKY

### Dealing with ambiguity

S  $\rightarrow$  NP VP  
NP  $\rightarrow$  Prn N  
NP  $\rightarrow$  Prn  
VP  $\rightarrow$  V NP  
VP  $\rightarrow$  V  
VP  $\rightarrow$  V S  
N  $\rightarrow$  duck  
V  $\rightarrow$  duck  
V  $\rightarrow$  saw  
Prn  $\rightarrow$  I  
Prn  $\rightarrow$  she  
Prn  $\rightarrow$  her

NP

Prn

V

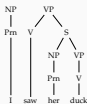
I saw her duck

9

Ç. Çöltekin, INF | University of Tübingen

Winter Semester 2020/21 12 / 13

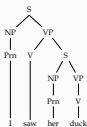
## Dealing with ambiguity



$S \rightarrow NP VP$   
 $NP \rightarrow Pm N$   
 $NP \rightarrow Pm$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $VP \rightarrow V S$   
 $N \rightarrow duck$   
 $V \rightarrow duck$   
 $V \rightarrow saw$   
 $Pm \rightarrow I$   
 $Pm \rightarrow she$   
 $Pm \rightarrow her$

10

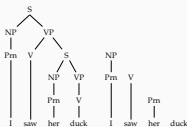
## Dealing with ambiguity



$S \rightarrow NP VP$   
 $NP \rightarrow Pm N$   
 $NP \rightarrow Pm$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $VP \rightarrow V S$   
 $N \rightarrow duck$   
 $V \rightarrow duck$   
 $V \rightarrow saw$   
 $Pm \rightarrow I$   
 $Pm \rightarrow she$   
 $Pm \rightarrow her$

11

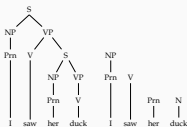
## Dealing with ambiguity



$S \rightarrow NP VP$   
 $NP \rightarrow Pm N$   
 $NP \rightarrow Pm$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $VP \rightarrow V S$   
 $N \rightarrow duck$   
 $V \rightarrow duck$   
 $V \rightarrow saw$   
 $Pm \rightarrow I$   
 $Pm \rightarrow she$   
 $Pm \rightarrow her$

12

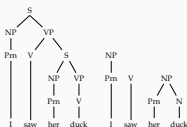
## Dealing with ambiguity



$S \rightarrow NP VP$   
 $NP \rightarrow Pm N$   
 $NP \rightarrow Pm$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $VP \rightarrow V S$   
 $N \rightarrow duck$   
 $V \rightarrow duck$   
 $V \rightarrow saw$   
 $Pm \rightarrow I$   
 $Pm \rightarrow she$   
 $Pm \rightarrow her$

13

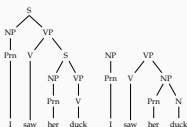
## Dealing with ambiguity



$S \rightarrow NP VP$   
 $NP \rightarrow Pm N$   
 $NP \rightarrow Pm$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $VP \rightarrow V S$   
 $N \rightarrow duck$   
 $V \rightarrow duck$   
 $V \rightarrow saw$   
 $Pm \rightarrow I$   
 $Pm \rightarrow she$   
 $Pm \rightarrow her$

14

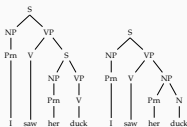
## Dealing with ambiguity



$S \rightarrow NP VP$   
 $NP \rightarrow Pm N$   
 $NP \rightarrow Pm$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $VP \rightarrow V S$   
 $N \rightarrow duck$   
 $V \rightarrow duck$   
 $V \rightarrow saw$   
 $Pm \rightarrow I$   
 $Pm \rightarrow she$   
 $Pm \rightarrow her$

15

## Dealing with ambiguity

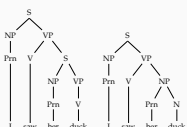


$S \rightarrow NP VP$   
 $NP \rightarrow Pm N$   
 $NP \rightarrow Pm$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $VP \rightarrow V S$   
 $N \rightarrow duck$   
 $V \rightarrow duck$   
 $V \rightarrow saw$   
 $Pm \rightarrow I$   
 $Pm \rightarrow she$   
 $Pm \rightarrow her$

16

## How to represent multiple parses

parse forest grammar



$S_{0,4} \rightarrow NP_{0,1} VP_{1,4}$   
 $NP_{0,1} \rightarrow Pm_{0,1}$   
 $NP_{0,1} \rightarrow I_{0,1}$   
 $VP_{1,4} \rightarrow V_{1,2} S_{2,4}$   
 $V_{1,2} \rightarrow saw_{1,2}$   
 $S_{2,4} \rightarrow Pm_{2,3} V_{3,4}$   
 $V_{3,4} \rightarrow duck_{3,4}$   
 $VP_{1,4} \rightarrow V_{1,2} NP_{2,4}$   
 $NP_{2,4} \rightarrow Pm_{2,3} N_{3,4}$

17

## CKY algorithm

- The CKY (Cocke-Kasami-Younger) parsing algorithm is a dynamic programming algorithm (Kasami 1965; Younger 1967; Cocke and Schwartz 1970)
- It processes the input *bottom up*, and saves the intermediate results on a *chart*
- Time complexity for *recognition* is  $O(n^3)$
- Space complexity is  $O(n^2)$
- It requires the CFG to be in *Chomsky normal form* (CNF) (can somewhat be relaxed, but not common)

## Chomsky normal form (CNF)

- A CFG is in CNF, if the rewrite rules are in one of the following forms
  - $A \rightarrow BC$
  - $A \rightarrow a$
 where  $A, B, C$  are non-terminals and  $a$  is a terminal
- Any CFG can be converted to CNF
- Resulting grammar is *isotypically equivalent* to the original grammar:
  - it generates/accepts the same language
  - but the derivations are different

## Converting to CNF: example

$S \rightarrow NP VP$   
 $S \rightarrow Aux NP VP$   
 $NP \rightarrow the N$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V$   
 $N \rightarrow cat$   
 $N \rightarrow dog$   
 $V \rightarrow bites$   
 $N \rightarrow bites$

$S \rightarrow Aux NP VP$   
 $S \rightarrow Aux NP VP \Rightarrow S \rightarrow Aux X$   
 $X \rightarrow NP VP$   
 $NP \rightarrow the N$   
 $NP \rightarrow the N \Rightarrow NP \rightarrow XN$   
 $X \rightarrow the$   
 $VP \rightarrow V$   
 $VP \rightarrow V \Rightarrow VP \rightarrow bites$

## Converting to CNF

- Eliminate the  $\epsilon$  rules; if  $A \rightarrow \epsilon$  is in the grammar
  - replace any rule  $B \rightarrow \alpha A \beta$  with two rules
    - $B \rightarrow \alpha \beta$
    - $B \rightarrow A' \beta$
  - add  $A' \rightarrow \alpha$  for all  $\alpha$  (except  $\epsilon$ ) whose LHS is  $A$
  - repeat the process for newly created  $\epsilon$  rules
  - remove the rules with  $\epsilon$  on the RHS (except  $S \rightarrow \epsilon$ )
- Eliminate unit rules; for a rule  $A \rightarrow B$ 
  - Replace the rule with  $A \rightarrow \alpha_1 | \dots | \alpha_n$ , where  $\alpha_1, \dots, \alpha_n$  are all RHS of rule  $B$
  - Remove the rule  $A \rightarrow B$
  - Repeat the process until no unit rules remain
- Binarize the non-binary rules with non-terminal on the RHS; for a rule  $A \rightarrow X_1 X_2 \dots X_n$ :
  - Replace the rule with  $A \rightarrow A_1 X_2 \dots X_n$ , and add  $A_1 \rightarrow X_1 X_2$
  - Repeat the process until all new rules are binary

## CKY demonstration

an ambiguous example



### CKY demonstration

an ambiguous example



### CKY demonstration

an ambiguous example



### CKY demonstration

an ambiguous example



### CKY demonstration

an ambiguous example



### CKY demonstration

an ambiguous example

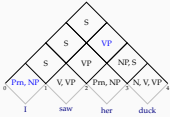


## CKY demonstration

an ambiguous example

 $S \rightarrow NP VP$ 

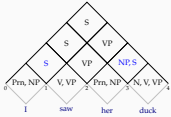
$S \rightarrow NP VP$   
 $NP \rightarrow Prn N$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V S$   
 $N \rightarrow duck$   
 $VP \rightarrow duck | saw$   
 $V \rightarrow duck | saw$   
 $Prn \rightarrow I | she | her$   
 $NP \rightarrow I | she | her$



## CKY demonstration

an ambiguous example

$S \rightarrow NP VP$   
 $NP \rightarrow Prn N$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V S$   
 $N \rightarrow duck$   
 $VP \rightarrow duck | saw$   
 $V \rightarrow duck | saw$   
 $Prn \rightarrow I | she | her$   
 $NP \rightarrow I | she | her$

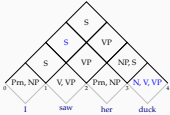


## CKY demonstration

an ambiguous example

Introduction CKY CKY

$S \rightarrow NP VP$   
 $NP \rightarrow Prn N$   
 $VP \rightarrow V NP$   
 $VP \rightarrow V S$   
 $N \rightarrow duck$   
 $VP \rightarrow duck | saw$   
 $V \rightarrow duck | saw$   
 $Prn \rightarrow I | she | her$   
 $NP \rightarrow I | she | her$



## CKY demonstration: the chart

our chart is a 2D array

NP, Prn	S	S	S					
	V, VP	VP	VP					
		Prn	NP, S					
			V, N, NP					
0	she	1	saw	2	her	3	duck	4

Space complexity is  $O(n^2)$ .

Space complexity is  $O(n^2)$ .

## CKY demonstration: the chart

our chart is a 2D array – this is more convenient for programming

S								
S	VP							
S	VP	NP, S						
NP, Prn	V, VP	Prn, NP	V, N, NP					
0	she	1	saw	2	her	3	duck	4

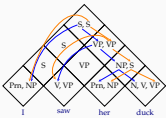
Space complexity is  $O(n^2)$ .

## Parsing vs. recognition

- We went through a recognition example
- Note that the algorithm is not directional: it takes the complete input
- Recognition accepts or rejects a sentence based on a grammar
- For parsing, we want to know the derivations that yielded a correct parse
- To recover parse trees, we
  - we follow the same procedure as recognition
  - add back links to keep track of the derivations

## Chart parsing example (CKY parsing)

Introduction CKY CKY

The chart stores a *parse forest* efficiently.

## Summary

- + CKY avoids re-computing the analyses by storing the earlier analyses (of sub-spans) in a table
  - It still computes lower level constituents that are not allowed by the grammar
  - CKY requires the grammar to be in CNF
  - CKY has  $O(n^3)$  recognition complexity
  - For parsing we need to keep track of backlinks
  - CKY can efficiently store all possible parses in a chart
  - Enumerating all possible parses have exponential complexity (worst case)
- Next:
- Top-down chart parsing: Earley algorithm
  - Suggested reading: Grune and Jacobs (2007, section 7.2)

## Acknowledgments, references, additional reading material

- Gries, John and J. T. Schwartz (1975). *Programming languages and their compiles: preliminary notes*. Tech. rep. Courant Institute of Mathematical Sciences, NYC.  
 Greif, David and Carol L. R. Jacobs (2007). *Parsing Techniques: A Practical Guide*. second. *Monographs in Computer Science*. The first edition is available at [http://linguistics.berkeley.edu/~greif/2007\\_monographs\\_in\\_computer\\_science.pdf](http://linguistics.berkeley.edu/~greif/2007_monographs_in_computer_science.pdf). Springer New York, isbn: 9780387085658.  
 Kuroda, Shizuo (1965). *An Efficient Recognition and Syntax Analysis Algorithm for Context-Free Languages*. Tech. rep. MIT, Document 2.  
 Younger, David H. (1967). "Recognition and parsing of context-free languages in time  $n^3$ ". In *Information and control* 10.2, pp. 105–132.