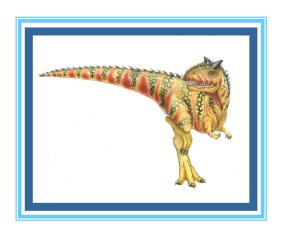
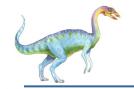
# Chapter 2a: Operating-System Services

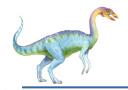




#### **Outline**

- Operating System Services
- User and Operating System-Interface
- System Calls





#### **Objectives**

- Identify services provided by an operating system
- Illustrate how system calls are used to provide operating system services



# Difference Between Kernel and Shell

Both of these are used in a computer system for establishing communication and operating the system. But there is a fundamental difference between shell and kernel. A shell is basically an interface present between the kernel and the user. It allows all of its users to establish communication with the kernel. A kernel is the very core of a typical OS. It functions to control all the tasks that come with a system. Let us know a bit more about each of these individually.





#### What is a Shell?

It refers to a special user program or an environment that provides a user with an interface for using the services of the operating system. A shell executes various programs on the basis of the input that a user provides.

#### What is a Kernel?

A kernel is basically the core and the heart of an OS (Operating system). It functions to manage the operations of the hardware and the computer. A kernel basically acts as a bridge between any user and the various resources offered by a system by accessing the various resources of a computer, such as the I/O devices, CPU, and various other resources.



# Difference Between Shell and Kernel

Parameters	Shell	Kernel	
Basics	A shell is basically an interface present between the kernel and the user.	A kernel is the very core of a typical OS.	
Meaning	A shell is a CLI (command-line interpreter).	A kernel is a type of low-level program that has its interfacing with the hardware on top of which all the applications run (disks, RAM, CPU, etc.).	
Uses and Purpose	A shell allows all of its users to establish communication with the kernel.	A kernel functions to control all the tasks that come with a system.	
Types	Korn Shell, C Shell, Bourne Shell, etc., are types of shells.	Hybrid kernel, Micro-kernel, Monolithic kernel, etc., are types of kernels.	
Functions	We can use shell commands such as mkdir, is, and many more for requesting the completion of the specific operation to the operating system (OS).	A kernel carries out the commands on a group of various files by specifying the pattern that can match.	
Management	A shell performs memory management.	A kernel performs process management.	
Layer of OS	The shell forms the outer layer of the operating system.	The kernel forms the inner layer of the operating system.	
Machine- Understandability	A shell interacts with all of its users and then interprets into a language that is understandable by the machine.	A kernel interacts with the hardware directly because it accepts the machine-understandable language from the available shell.	



#### **Operating System Services**

- Operating systems provide an environment for execution of programs and services to programs and users
- One set of operating-system services provides functions that are helpful to the user:
  - User interface Almost all operating systems have a user interface (UI).
    - Varies between Command-Line (CLI), Graphics User Interface (GUI), touch-screen
  - Program execution The system must be able to load a program into memory and to run that program, end execution, either normally or abnormally (indicating error)
  - I/O operations A running program may require I/O, which may involve a file or an I/O device
  - File-system manipulation The file system is of particular interest. Programs need to read and write files and directories, create and delete them, search them, list file Information, permission management.



#### **Operating System Services (Cont.)**

- One set of operating-system services provides functions that are helpful to the user (Cont.):
  - Communications Processes may exchange information, on the same computer or between computers over a network
    - Communications may be via shared memory or through message passing (packets moved by the OS)
  - Error detection OS needs to be constantly aware of possible errors
    - May occur in the CPU and memory hardware, in I/O devices, in user program
    - For each type of error, OS should take the appropriate action to ensure correct and consistent computing
    - Debugging facilities can greatly enhance the user's and programmer's abilities to efficiently use the system



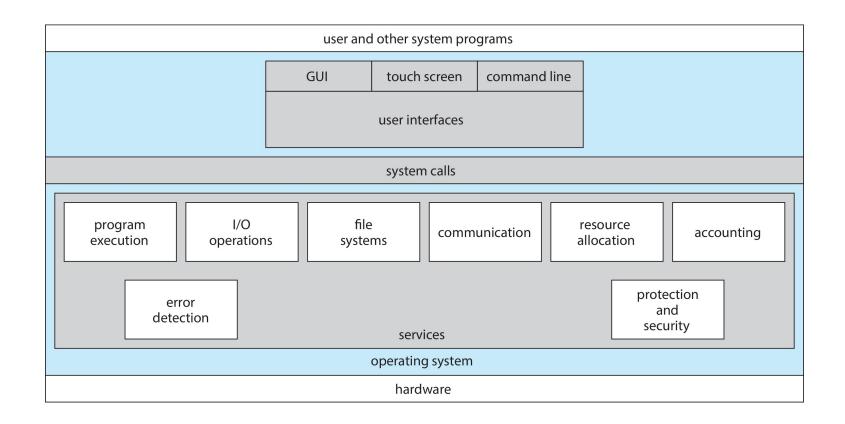


### **Operating System Services (Cont.)**

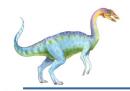
- Another set of OS functions exists for ensuring the efficient operation of the system itself via resource sharing
  - Resource allocation When multiple users or multiple jobs running concurrently, resources must be allocated to each of them
    - Many types of resources CPU cycles, main memory, file storage, I/O devices.
  - Logging To keep track of which users use how much and what kinds of computer resources
  - Protection and security The owners of information stored in a multiuser or networked computer system may want to control use of that information, concurrent processes should not interfere with each other
    - Protection involves ensuring that all access to system resources is controlled
    - Security of the system from outsiders requires user authentication, extends to defending external I/O devices from invalid access attempts



## **A View of Operating System Services**







#### **Command Line interpreter**

- CLI allows direct command entry
- Sometimes implemented in kernel, sometimes by systems program
- Sometimes multiple flavors implemented shells
- Primarily fetches a command from user and executes it
- Sometimes commands built-in, sometimes just names of programs
  - If the latter, adding new features doesn't require shell modification





#### **Bourne Shell Command Interpreter**

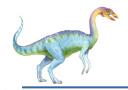
```
1. root@r6181-d5-us01:~ (ssh)
× root@r6181-d5-u... ● 第1 ×
                              ssh
                                      #2 × root@r6181-d5-us01... #3
Last login: Thu Jul 14 08:47:01 on ttys002
iMacPro:~ pbg$ ssh root@r6181-d5-us01
root@r6181-d5-us01's password:
Last login: Thu Jul 14 06:01:11 2016 from 172.16.16.162
[root@r6181-d5-us01 ~]# uptime
06:57:48 up 16 days, 10:52, 3 users, load average: 129.52, 80.33, 56.55
Froot@r6181-d5-us01 ~ 7# df -kh
Filesystem
                   Size Used Avail Use% Mounted on
/dev/mapper/vg_ks-lv_root
                        19G
                               28G 41% /
                    50G
tmpfs
                   127G 520K 127G
                                    1% /dev/shm
/dev/sda1
                   477M 71M
                             381M 16% /boot
                   1.0T 480G 545G 47% /dssd xfs
/dev/dssd0000
tcp://192.168.150.1:3334/orangefs
                    12T 5.7T 6.4T 47% /mnt/orangefs
/dev/apfs-test
                    23T 1.1T 22T
                                    5% /mnt/apfs
[root@r6181-d5-us01 ~]#
[root@r6181-d5-us01 ~]# ps aux | sort -nrk 3,3 | head -n 5
        97653 11.2 6.6 42665344 17520636 ? S<Ll Jul13 166:23 /usr/lpp/mmfs/bin/mmfsd
root
root
                                0?
                                               Jul12 181:54 [vpthread-1-1]
        69849 6.6 0.0
        69850 6.4 0.0 0 0? S Jul12 177:42 [vpthread-1-2]
root
         3829 3.0 0.0
                       0 0 ? S Jun27 730:04 [rp_thread 7:0]
root
                                0 ?
         3826 3.0 0.0
                           0
                                               Jun27 728:08 [rp_thread 6:0]
root
[root@r6181-d5-us01 ~]# ls -l /usr/lpp/mmfs/bin/mmfsd
-r-x---- 1 root root 20667161 Jun 3 2015 /usr/lpp/mmfs/bin/mmfsd
[root@r6181-d5-us01 ~]#
```



## **User Operating System Interface - GUI**

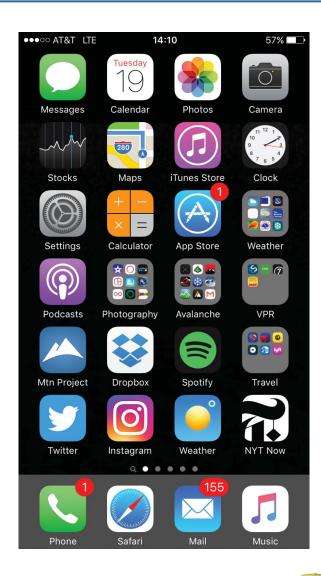
- User-friendly desktop metaphor interface
  - Usually mouse, keyboard, and monitor
  - Icons represent files, programs, actions, etc
  - Various mouse buttons over objects in the interface cause various actions (provide information, options, execute function, open directory (known as a folder)
- Many systems now include both CLI and GUI interfaces
  - Microsoft Windows is GUI with CLI "command" shell
  - Apple Mac OS X is "Aqua" GUI interface with UNIX kernel underneath and shells available
  - Unix and Linux have CLI with optional GUI interfaces (CDE, KDE, GNOME)





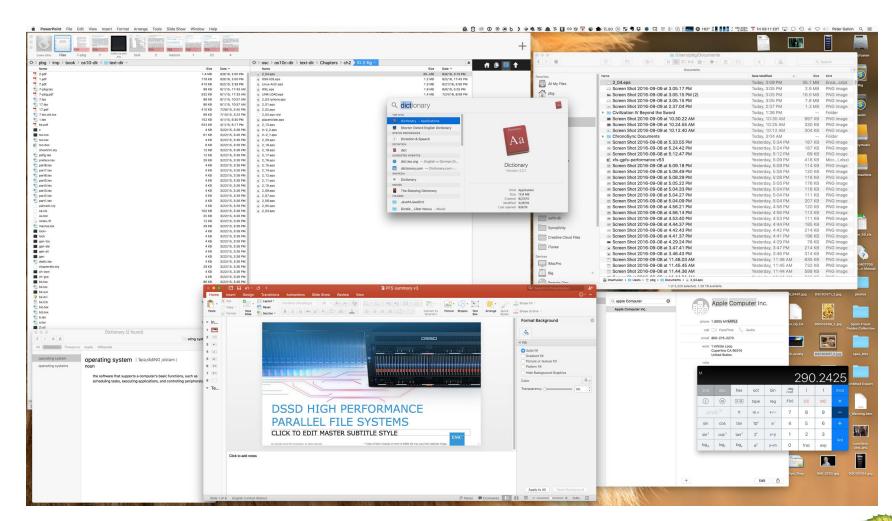
#### **Touchscreen Interfaces**

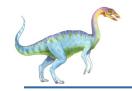
- Touchscreen devices require new interfaces
  - Mouse not possible or not desired
  - Actions and selection based on gestures
  - Virtual keyboard for text entry
- Voice commands





#### The Mac OS X GUI



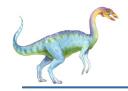


#### **System Calls**

- Programming interface to the services provided by the OS
- Typically written in a high-level language (C or C++)
- Mostly accessed by programs via a high-level Application Programming Interface (API) rather than direct system call use
- Three most common APIs are Win32 API for Windows, POSIX API for POSIX-based systems (including virtually all versions of UNIX, Linux, and Mac OS X), and Java API for the Java virtual machine (JVM)

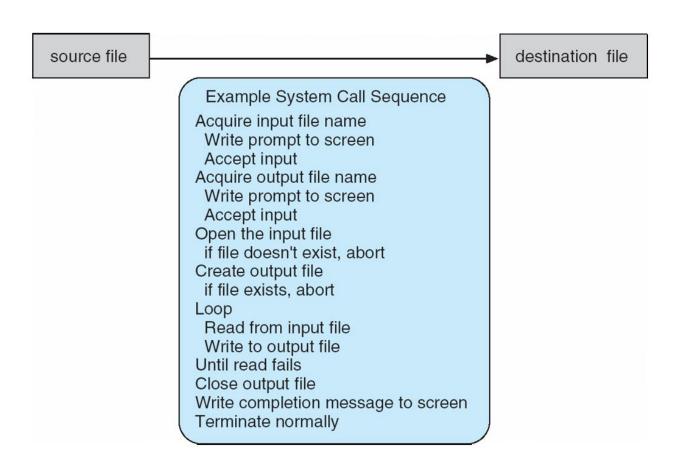
Note that the system-call names used throughout this text are generic



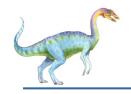


#### **Example of System Calls**

System call sequence to copy the contents of one file to another file







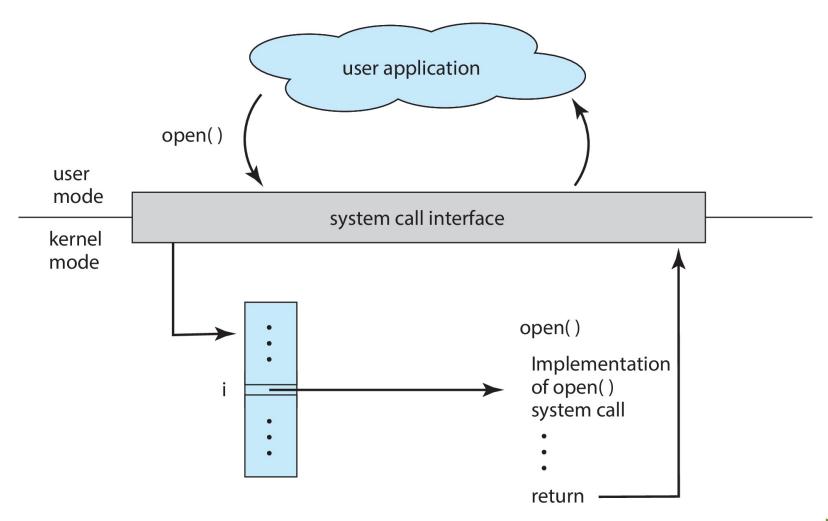
#### **System Call Implementation**

- Typically, a number is associated with each system call
  - System-call interface maintains a table indexed according to these numbers
- The system call interface invokes the intended system call in OS kernel and returns status of the system call and any return values
- The caller need know nothing about how the system call is implemented
  - Just needs to obey API and understand what OS will do as a result call
  - Most details of OS interface hidden from programmer by API
    - Managed by run-time support library (set of functions built into libraries included with compiler)



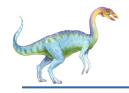


## API – System Call – OS Relationship



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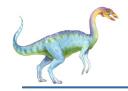




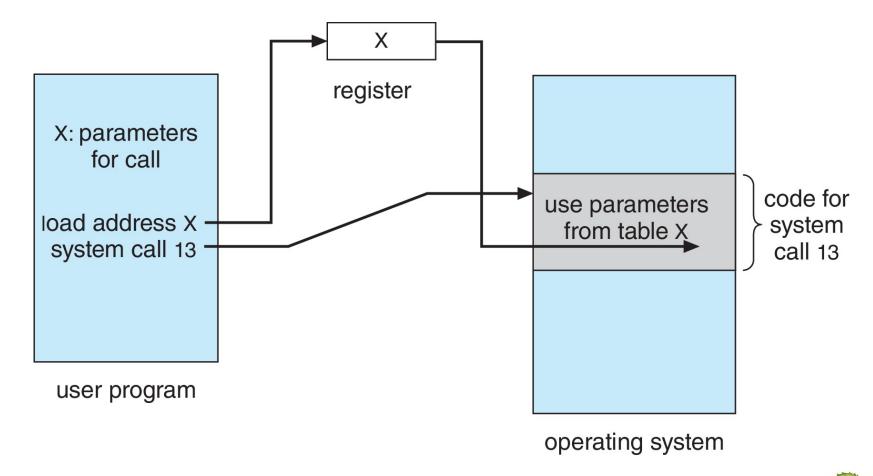
#### System Call Parameter Passing

- Often, more information is required than simply identity of desired system call
  - Exact type and amount of information vary according to OS and call
- Three general methods used to pass parameters to the OS
  - Simplest: pass the parameters in registers
    - In some cases, may be more parameters than registers
  - Parameters stored in a block, or table, in memory, and address of block passed as a parameter in a register
    - This approach taken by Linux and Solaris
  - Parameters placed, or pushed, onto the stack by the program and popped off the stack by the operating system
  - Block and stack methods do not limit the number or length of parameters being passed

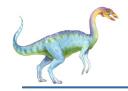




#### Parameter Passing via Table



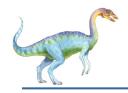




#### **Types of System Calls**

- Process control
  - create process, terminate process
  - end, abort
  - load, execute
  - get process attributes, set process attributes
  - wait for time
  - wait event, signal event
  - allocate and free memory
  - Dump memory if error
  - Debugger for determining bugs, single step execution
  - Locks for managing access to shared data between processes



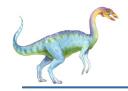


#### **Types of System Calls (Cont.)**

- File management
  - create file, delete file
  - open, close file
  - read, write, reposition
  - get and set file attributes
- Device management
  - request device, release device
  - read, write, reposition
  - get device attributes, set device attributes
  - logically attach or detach devices



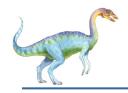
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#### **Types of System Calls (Cont.)**

- Information maintenance
  - get time or date, set time or date
  - get system data, set system data
  - get and set process, file, or device attributes
- Communications
  - create, delete communication connection
  - send, receive messages if message passing model to host name or process name
    - From client to server
  - Shared-memory model create and gain access to memory regions
  - transfer status information
  - attach and detach remote devices

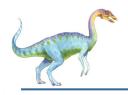




## **Types of System Calls (Cont.)**

- Protection
  - Control access to resources
  - Get and set permissions
  - Allow and deny user access





#### **Examples of Windows and Unix System Calls**

#### **EXAMPLES OF WINDOWS AND UNIX SYSTEM CALLS**

The following illustrates various equivalent system calls for Windows and UNIX operating systems.

	Windows	Unix
Process control	<pre>CreateProcess() ExitProcess() WaitForSingleObject()</pre>	<pre>fork() exit() wait()</pre>
File management	<pre>CreateFile() ReadFile() WriteFile() CloseHandle()</pre>	<pre>open() read() write() close()</pre>
Device management	<pre>SetConsoleMode() ReadConsole() WriteConsole()</pre>	<pre>ioctl() read() write()</pre>
Information maintenance	<pre>GetCurrentProcessID() SetTimer() Sleep()</pre>	<pre>getpid() alarm() sleep()</pre>
Communications	<pre>CreatePipe() CreateFileMapping() MapViewOfFile()</pre>	<pre>pipe() shm_open() mmap()</pre>
Protection	<pre>SetFileSecurity() InitlializeSecurityDescriptor() SetSecurityDescriptorGroup()</pre>	chmod() umask() chown()



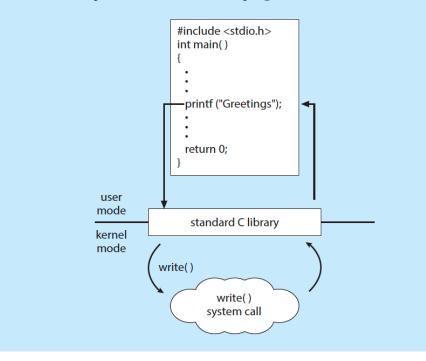


#### **Standard C Library Example**

C program invoking printf() library call, which calls write() system call

#### THE STANDARD C LIBRARY

The standard C library provides a portion of the system-call interface for many versions of UNIX and Linux. As an example, let's assume a C program invokes the printf() statement. The C library intercepts this call and invokes the necessary system call (or calls) in the operating system—in this instance, the write() system call. The C library takes the value returned by write() and passes it back to the user program:





# **End of Chapter 2a**

