

Q.1] a) True.

As the computers store the value of decimal numbers in binary format, there are some numbers which cannot be represented precisely.

e.g:-

In python,  $0.1 + 0.2$  will not give  $0.3$  but a value like  $0.30004$

b) True.

Reason:  $\Theta(g(n))$  is a tighter bound than  $\Omega(g(n))$

~~Q If  $\Theta(g(n))$~~

:  $\Theta(g(n))$ :  $c_1 g(n) \leq f(n) \leq c_2 g(n)$  for  $n \geq n_0$

$\Rightarrow f(n) \leq c'' g(n)$

$\left[ \begin{array}{l} \Theta(f(n)) \\ \Omega(g(n)) \end{array} \right]$

c) ~~True~~. False

Reason: The difference in iterations is too less to

'guarantee' that the value will be closer to its actual value.

d) False

Expected final state: [None, None, None, None, None, 5, 6, 7, 8, 9]

Reason:

Step 1: Enqueue 0-9:

$$\therefore Q = [0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7 \ 8 \ 9]$$

Step 2: Dequeue:

$$Q = [\text{None } \text{None } \text{None } \text{None } \text{None } 5 \ 6 \ 7 \ 8 \ 9]$$

Assumption: Q is not a circular queue.

e) False.

Selection sort is not adaptive, i.e., it does not speed up the sorting process if the array is already sorted.

Bubble Sort can be made adaptive by few modifications in its algorithm and thus will not necessarily speed up the sorting process.

f) False.

Adjacency matrix uses a space of  $O(n^2)$ . Although it can retrieve an edge, as well as delete or insert an edge in constant time. We can instead use an adjacency list which also has constant time for inserting or removing an edge.

g) True.

Even after multiplying the weights by a constant, all the weights will be scaled by the same factor. So shortest path will remain via same route as previous one.

h) False.

In BFS we can use queue <sup>or recursion</sup> and not stack.

i) True.

Reason:- In dynamic programming, most of the sub problems are dependent. Hence, most of them will be calculated multiple times if recursion is used. This is highly inefficient. Instead, the values of the smaller sub problems should be stored in dynamic programming methods so those values are not calculated again.

j) True.

Reason: `--add--` is a special method in python ~~for~~ for <sup>addition or concatenation</sup> adding two variables but of the same datatype (depends on the datatype).

Syntax:

`x.--add--(y)`

It takes 'y' as a parameter, and performs the operation on 'x'.

If x and y are strings, the output will be: xy

Q.2) a) All options are correct because the value of the bit generated is random, ~~but between~~, i.e., either 0 or 1.

Since all ~~POSSIBLE~~ options have either 0 or 1, all of them are correct

b) ii)  $\sim N \lg(N)$  where  $\lg N = \log_2 N$

Reason:

while ( $i < N$ ):

{  
     $i^* = 2$                            $\leftarrow \log_2 1 \lg N$   
}

for  $j$  in range ( $i$ ):                   $\leftarrow N \times \lg N$                    $= N \lg N$   
    Sum +=  $i$

Option (iv)

- Q.2) c) i) 101101, 010010  
ii) 100101, 011010

Reason: Parents : 101010 → 101101  
010101                                    010010 } Children  
    ↗  
    Crossing over event

101010 → 100101  
010101                                    011010

Q.3] a)  $\wedge t [\wedge t]^* \$$

b)  $\wedge [AEIOU]^* [AEIOVaeiou] [\wedge aeiou]^* [aeiou] [\wedge aeiou]^* [aeiou]$   
 $[\wedge aeiou]^* \$$

c)  $\wedge [ATC]^* TG^* [ATC]^* \$$

d)  $\wedge ([A-Za-z]) ([a-z]) ([a-Z]) ([a-z]) [a-z] \backslash \backslash \backslash 3 \backslash 2 \backslash 1 \$$

e)  $\wedge ([ATGC]) [ATGC] [ATGC] \backslash 1^* \$$

f)  $\wedge [\wedge a]^* (a [\wedge a]^* a [\wedge a]^*)^* \$$

g)  $\wedge [aeiouAEIOU] [a-z]^* [aeiou] \$$

h)  $\wedge [A-Za-z] \{ 20, \}$

}

17

Q.4] (a)  $t \cdot * t \$$

steepest

strategist

throughput

transcendent

b)  $\backslash \langle C.C.* \backslash 2 \backslash 1 \$ \rangle >$

addend a

etiquette

c)  $a \cdot * e \cdot * t$

strategist

transcendent

d)  $( [d_i n] ) @ C \cdot * \{ 1 \} \{ 2 \}$

addend a

transcendent

e) \<([tea]) c.\*\1.\*\1\>

## Anacenda

## automat a

## eigenvalue

## etiquette

Q.17]

O	C	C	U	A	V	G	C	
O	0	2	4	6	8	10	12	14
C	2	0	2	4	6	8	10	12
U	4	20	1	2	4	6	8	10
A	6	4	3	2	2	4	6	8
V	8	6	5	3	3	2	4	6
G	10	8	7	5	4	4	2	4
G	12	10	9	7	6	5	4	3
C	14	12	10	9	8	7	6	(4)

Final Edit distance

Q.16]

A	B	C	D	E	F
<u>0</u>	∞	∞	∞	∞	∞

Start from A:

A	B	C	D	E	F
<u>0</u>	3	<u>10</u>	14	∞	∞

(Edge Relaxations = 3)

Now include C:

A	B	C	D	E	F
<u>0</u>	2	<u>1</u>	12	∞	3

(Edge relaxations = 3)

Now include B:

A	B	C	D	E	F
<u>0</u>	2	<u>1</u>	12	∞	3

(Edge relaxations = 0)

Now include F

A	B	C	D	E	F
<u>0</u>	2	<u>1</u>	12	15	3

(Edge relaxations = 1)

Now include D:

A	B	C	D	E	F
<u>0</u>	2	<u>1</u>	12	15	3

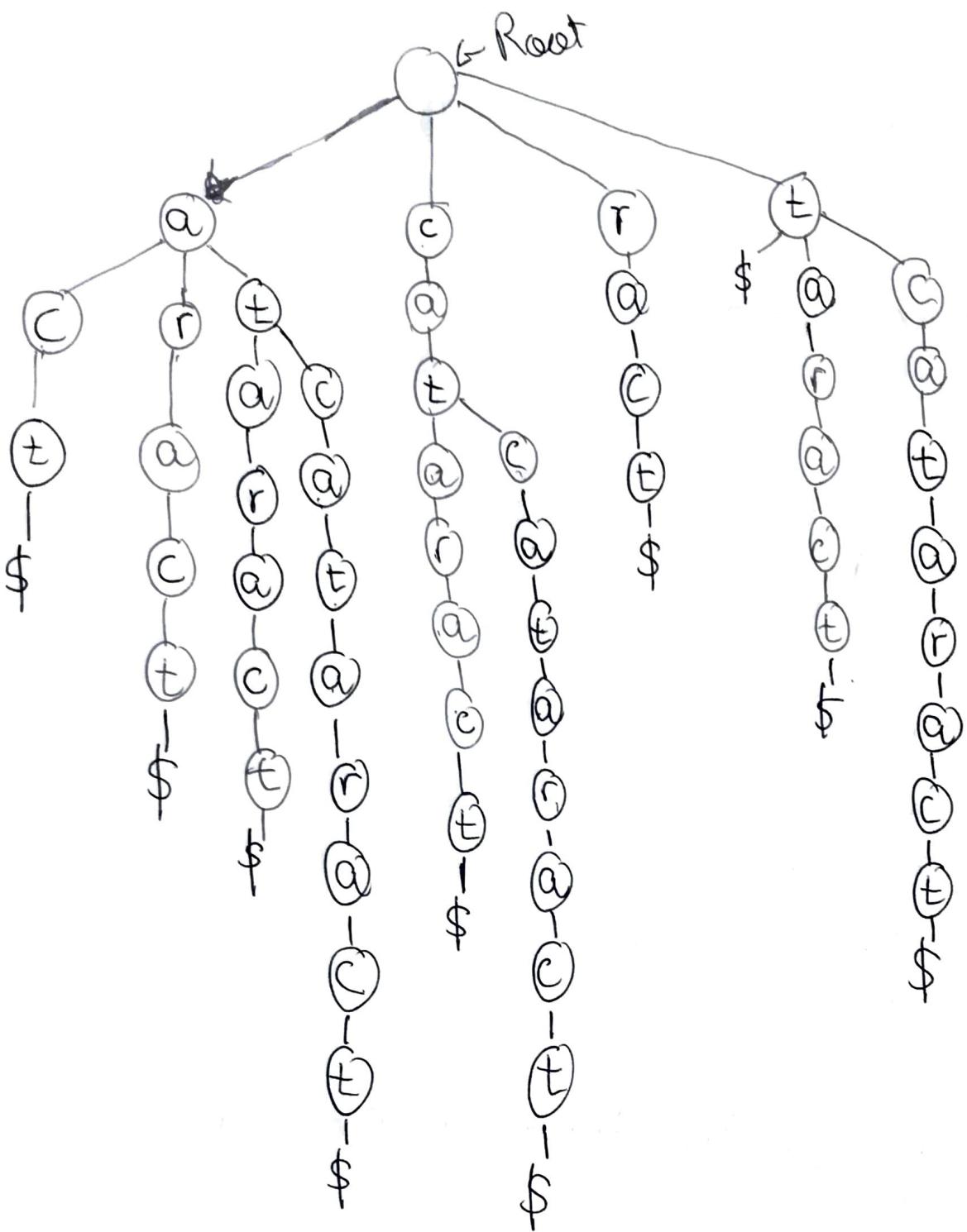
~Final distance values

Total edge relaxations = 37

Q.14) cataract

E  
ct  
@act  
ract  
Aract  
taract  
ataract  
cataract  
tcatara ct  
at catara ct  
cat catara ct

Suffix trie drawn on next page



Pattern : an

First, it will search amongst the ~~states~~ from root node to node 'a'. All nodes of node 'a'

node to node 'a'. Then it checks the child nodes of node 'a'.

Since no child node of 'a' has value 'n',  
the search will terminate and display that no  
pattern is found.

(Q.15) We can use a modified edit distance method to find shortest evolutionary path between a pair of sequences  $S_a$  and  $S_b$ .

Let  $S_a$  be the source ~~string~~ RNA sequence

Let  $S_b$  be the target RNA sequence.

$$E(i,j) = \min \begin{cases} E(i-1, j) + 2 \\ E(i, j-1) + 2 \\ E(i-1, j-1) + 1 \end{cases}$$

But now we will apply an extra condition before calculating the distance.

The condition is:

$$E(i-1, j) \in S$$

$$E(i, j-1) \in S$$

$$E(i-1, j-1) \in S$$

Q.13]

~~def~~ import math

@

Q.13]

import math

from numpy import random

def func(x):

return math.sin(x)\*math.cos(x)

def area(func, a, b):

area = 0

points = []

for i in range(10000):

x = random.uniform(0, math.pi/2) (a, b)

y = random.uniform(0, 0.5) (0, 0.5)

points.append((x, y))

interior = 0

for i in range(len(points)):

if points[i][1] - func(points[i][0]) <= 0:

interior = interior + 1

b - a

area = (interior / 10000) \* (math.pi/2) \* 0.5

return area

# Actual area of curve through Riemann integration = 0.5

# To estimate error:

~~area =~~ or  
area - curve = area(func, 0, math.pi/2)

error = 0.5 - area - curve

print(error)

print(area - curve)

$$Q.2) \varphi = 1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \dots}}}$$

The above equation can be written as:

$$\varphi = 1 + \frac{1}{\varphi}$$

Python program:

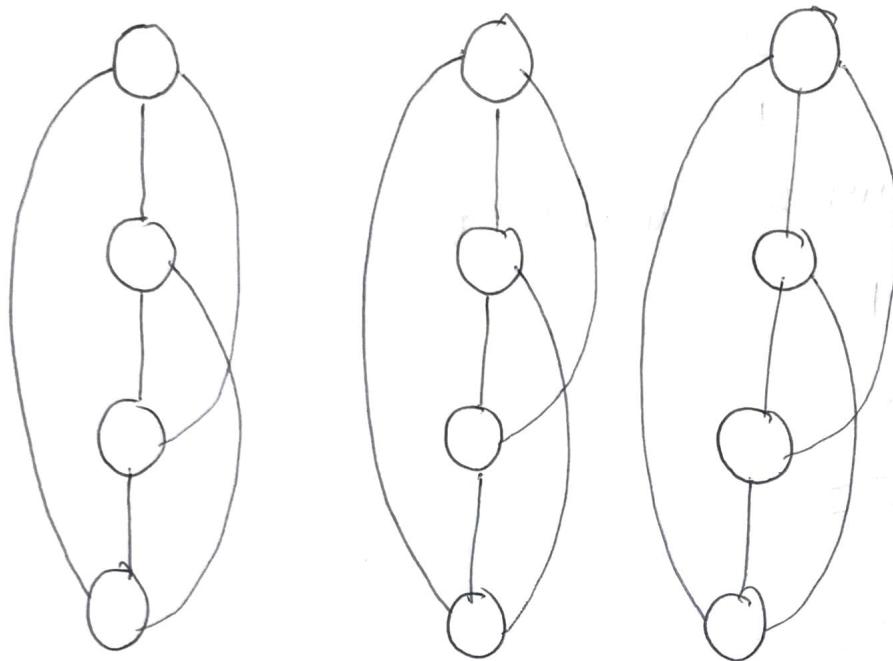
$$g = 1$$

for i in range(100000):

$$g = 1 + \frac{1}{g}$$

print(g)

Q. 8)



Q. 9) Genome Assembly Problem is a problem where we have short sequences (say  $n$  sequences), then we have to design an algorithm such that the ' $n$ ' sequences are aligned in such a way that it forms one single large <sup>genome</sup> sequence, in shortest time. In other words, we have to find a superstring or ' $S$ ' which contains all the <sup>short</sup> sequences ' $p_i$ ' as substrings. We can use hamiltonian method to solve this.

~~hamiltonian graph method~~  
We can use ~~hamiltonian graph method~~ to solve genome assembly problem. All the substrings represent vertices of the graph. We can create an edge between overlapping vertices and then see the path to find the path that covers all the vertices at least once. The order in which the vertices are passed can be noted and combine them to form the genome. 18

Q. 9] A drawback of hill climbing algorithms is it can get stuck at a local minima or a local maxima. At these points, the function cost function can move in any direction which is a problem. Method to overcome this: We can use simulated annealing.

We can use simulated annealing to overcome this issue. In this method, we don't reject moves that will increase the cost. We occasionally accept such moves with low probability. If this selected move then improves our cost ( $C_k$ ), then we will permanently accept this move.

Q.S] ~~k =~~ ~~int~~(  
    ~~input ("Enter value of k")~~)

seq = input ("Enter the sequence:")

∴ kmers = []

for i in range (0, len(seq) - ~~k + 1~~):  
    kmers.append (seq [i:i+k])

kmers\_sorted = sorted (kmers).

for i in kmers\_sorted:  
    print (i)

Q.10] Currently, all intervals lie between 0 and 1  
Let the value be  $s$ .

$$\therefore 0 < s < 1$$

$$\therefore \text{Here, } b = 12, a = 1$$

Multiplying by  $(b-a)$  on both sides

$$0 < s(11) < 11$$

Adding  $a$  on both sides

$$\therefore 1 < 11s + 1 < 12$$

$$\therefore 1 < \lceil 11s + 1 \rceil < 12 \quad \lceil \rceil \text{ :- a like math.ceil}$$

Now we use  $\lceil 11s + 1 \rceil$  for transformation  
where ' $s$ ' are the original probabilities

$$s = (0.8147, 0.9088, 0.1270, 0.9135, 0.6325, 0.0975)$$

$$0.2785, 0.5969, 0.9575, 0.9649)$$

∴ Applying the transformation  $\lceil 11s + 1 \rceil$ ,

New numbers: 10, 11, 3, 12, 8, 3, 5, 8, 12, 12

Q.12] Given: 1111000010 → let it be Chromosome A

$$\therefore x: 11110 \quad y: 00010$$

$$i.e. x=2$$

$$x=30$$

$$\therefore \text{Fitness function: } x^2 + y^3 - 5xy^2$$

$$\begin{aligned} \text{a) } & \text{Fitness of A: } 30^2 + 2^3 - 5 \times 30 \times 2^2 \\ & = \boxed{308} \end{aligned}$$

b) New chromosome: 0111000011

$$\therefore x = 01110$$

$$x=14$$

$$y = 00011$$

$$y=3$$

$$\begin{aligned} \text{Fitness of B} &= \textcircled{2} 14^2 + 3^3 - 5 \times 14 \times 3^2 \\ &= \boxed{-407} \end{aligned}$$

Q.11) First generate a random number within a sphere of radius, 1:

∴ The x, y, z coordinates can be calculated

$$x = r \cos \theta \sin \phi, y = r \sin \theta \cos \phi, z = r \cos \theta \sin \phi \\ \theta \in [0, 2\pi], \phi \in [0, \pi]$$

$$x = \text{sqrt}(\text{random}(0, 1)) * \cos(\text{random}(0, 2\pi)) \\ * \sin(\text{random}(0, \pi))$$

$$y = \text{sqrt}(\text{random}(0, 1)) * \sin(\text{random}(0, 2\pi)) \\ * \sin(\text{random}(0, \pi))$$

$$z = \text{sqrt}(\text{random}(0, 1)) * \cos(\text{random}(0, \pi))$$

Now, we can scale them to the ellipsoid:

$$\textcircled{Q} \quad x_{\text{ellipse}} = x * a$$

$$y_{\text{ellipse}} = y * b$$

$$z_{\text{ellipse}} = z * c$$