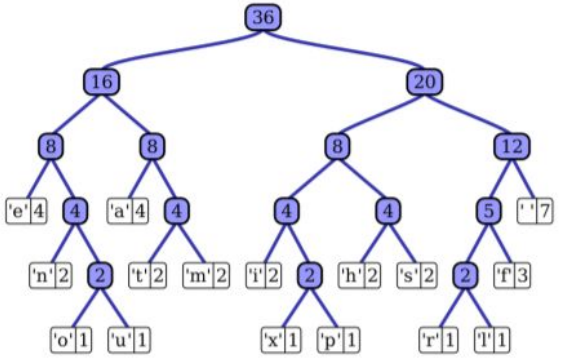
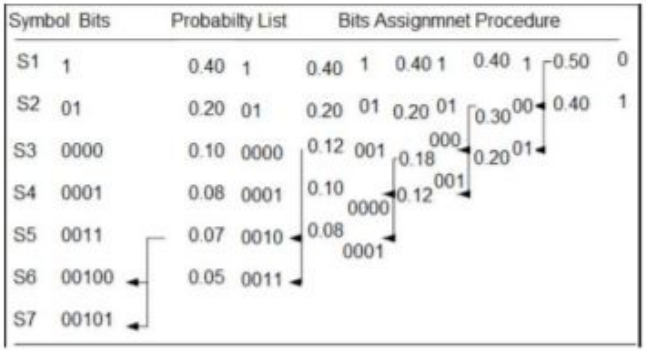


CO-203 Object Oriented Programming
Innovative Project Report

COMPRESSION USING HUFFMAN CODING

GUI built on Qt5 using C++



Ishan Khanna
2K19/EE/120

COMPRESSION USING HUFFMAN CODING

Huffman code is a technique that is commonly used for lossless data compression. The process of such compression proceeds through **Huffman coding.** This algorithm was developed by David A. Huffman who at the time was a Sc.D. student at MIT, USA. It was published in the 1952 paper which was titled "A Method for the Construction of Minimum-Redundancy Codes".

The output obtained from Huffman's algorithm can be seen as a variable-length code table for encoding a source symbol (such as a character in a file). The algorithm created this table from the frequency of occurrence (also known as *weight) for* each possible value of the source symbol. However, even though it is among the most optimal methods for encoding symbols separately, it is not always optimal among all compression methods.

Huffman Encoder utilises Huffman coding trees to reduce the count of bits that are needed to represent any symbol. It is a variable length encoding.

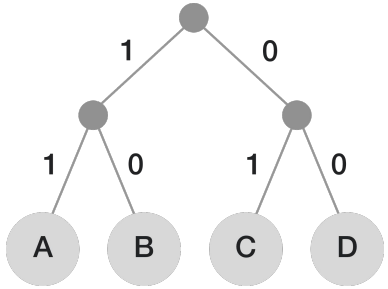
For Example :

- Message “CAB” is compressed to

011110.

- Otherwise the message would have been encoded using its ASCII representation as follows:

001000110010000100100010.



Algorithm

1) Find the frequency of each distinct symbol in the input.

2) Make a priority queue of single-node trees. Every node in the initial queue represents a distinct symbol from the set of possible symbols, and contains the frequency of that symbol in the message to be coded. Symbols with a frequency of zero are ignored.

3) Loop while there is more than 1 tree in the queue:

   3a. Remove the two trees from the queue that have the lowest frequency contained in their roots.

   3b. Create a new node that will be the root of a new tree. This new tree will have those two trees just removed in step 3a as left and right sub-trees. The frequency in the root of this new tree will be the sum of the frequency in the roots of its sub-trees.

   3d. Label the edge from this new root to its left sub-tree "0", and label the edge to its right sub-tree "1".

   3e. Insert this new tree in the queue, and go to 3.

4) Return the one tree in the queue as the Huffman code tree.

Time and space complexity analysis for Huffman Coding Algorithms

Time Complexity:

Since efficient priority queue data structures require O(log *n*) time complexity, and a tree that has *n* leaves must have 2*n*−1 nodes. So, this algorithm operates in O(*n* log *n*) time, where *n* is the number of symbols.

Space Complexity:

If n is the count of distinct symbols in the input file, then we require O(n) space for building the tree. Hence, space complexity of Huffman coding is O(n).

Design and Implementation

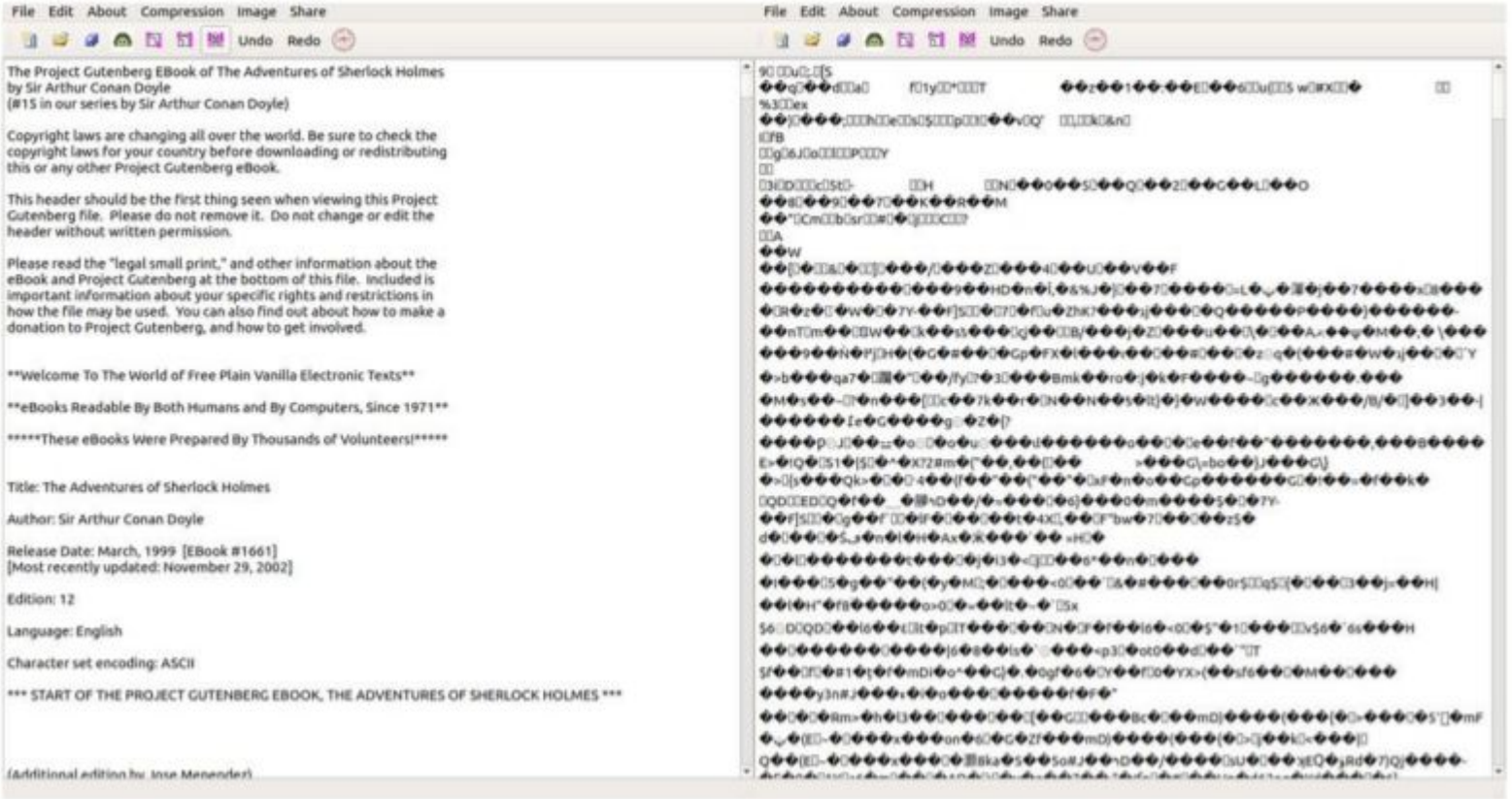
- I have built a GUI using Qt5. It includes a text view and an image view.
- The text view provides basic text editing features along with the options to compress and decompress the currently opened file.
- It also has provisions for sharing by e-mail option printing.
- The Image View includes an in-built image-display area and options to compress the file to JPEG.
- It also provides an option to print the image and to open an image for viewing.

- Features:

1. Text editor – Cut, Copy, Paste, Undo, Redo
2. Text Compression
3. Text Decompression
4. Image Viewer
5. Image Compression
6. Sharing by Mail
7. Print

Results

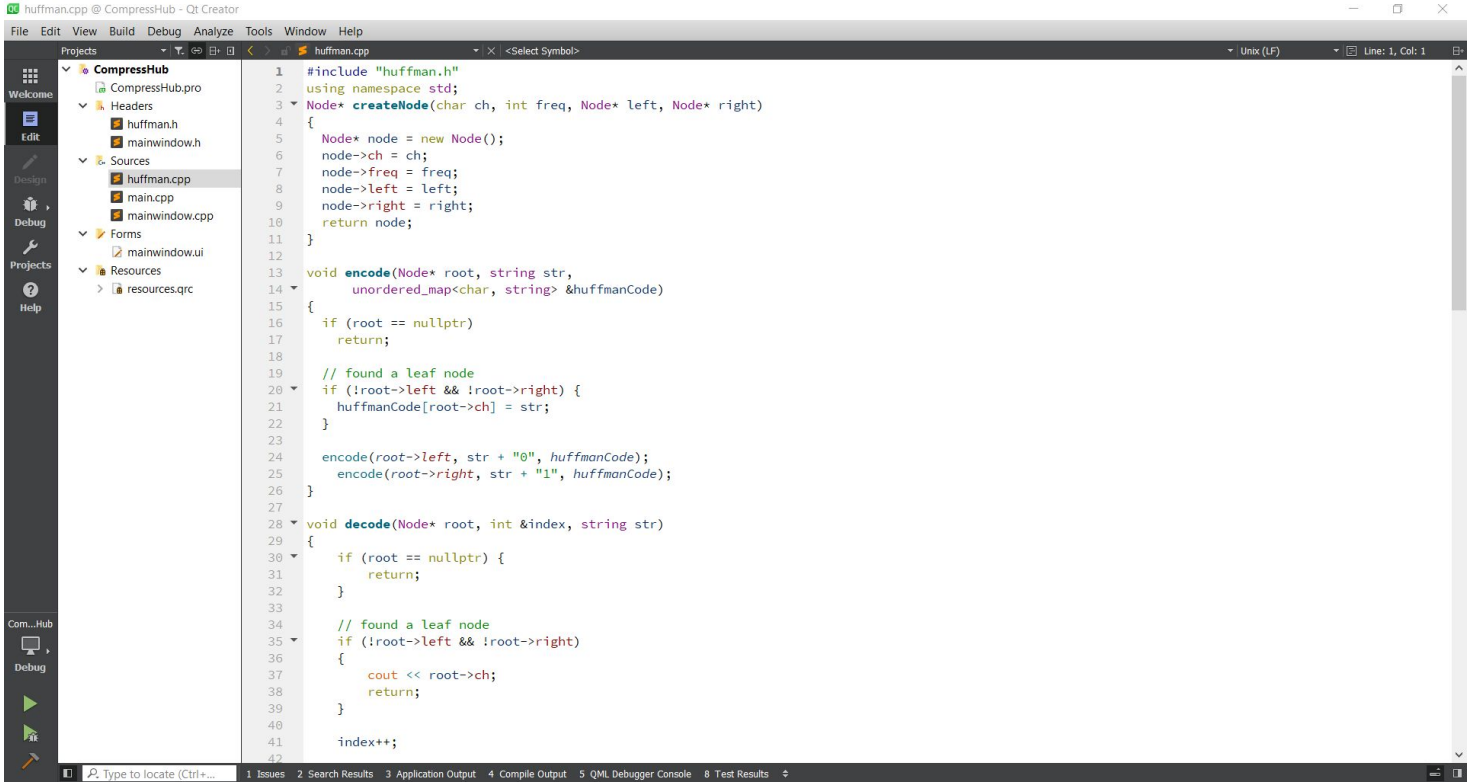
Text Compression

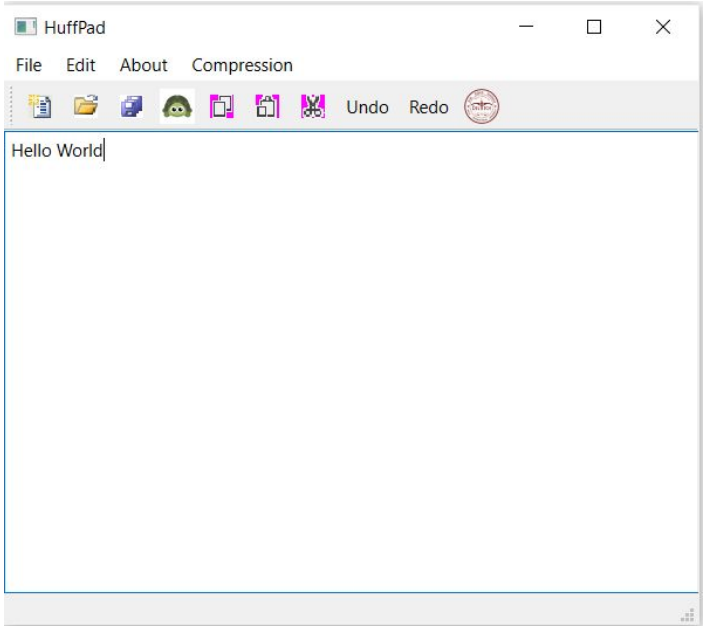


Results

Image Compression

Results


huffman.cpp @ CompressHub - Qt Creator
File Edit View Build Debug Analyze Tools Window Help
Projects
CompressHub
CompressHub.pro
Headers
huffman.h
mainwindow.h
Sources
huffman.cpp
main.cpp
mainwindow.cpp
Forms
mainwindow.ui
Resources
resources.qrc
Welcome
Edit
Design
Debug
Projects
Help
Com...Hub
Debug
#include "huffman.h"
using namespace std;
Node* createNode(char ch, int freq, Node* left, Node* right)
{
Node* node = new Node();
node->ch = ch;
node->freq = freq;
node->left = left;
node->right = right;
return node;
}
void encode(Node* root, string str,
unordered_map<char, string> &huffmanCode)
{
if (root == nullptr)
return;
// found a leaf node
if (!root->left && !root->right) {
huffmanCode[root->ch] = str;
}
encode(root->left, str + "0", huffmanCode);
encode(root->right, str + "1", huffmanCode);
}
void decode(Node* root, int &index, string str)
{
if (root == nullptr) {
return;
}
// found a leaf node
if (!root->left && !root->right)
{
cout << root->ch;
return;
}
index++;
Type to locate (Ctrl+...
1 Issues
2 Search Results
3 Application Output
4 Compile Output
5 QML Debugger Console
8 Test Results



Results

```
32      }
Application Output
CompressHub2
Hello World

Encoded string is :
00011111010110111000111001010011

Decoded string is:
Hello World02:51:35: C:\Users\Ishan's PC\Desktop\QT Projects\build-CompressHub-Desktop_Qt_6_1_2_MinGW_64_bit-Debug\debug\CompressHub.exe exited with code 0
1 Issues  2 Search Results  3 Application Output  4 Compile Output  5 QML Debugger Console  8 Test Results
```

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