

# Competitive Programming For Beginners

## Agenda

- Coding Problems
  - Approaches
  - Time Complexity
  - Space Complexity

## QUESTION

Find the element that is present once in an array where every other element is present twice?

Input:[1,3,5,6,6,3,1]

Output: 5

Input:[10,30,50,60,60,30,10]

Output: 50

## First Approach

**Using NESTED LOOPS: Time Complexity:  $O(n*n)$**

```
for(int i=0;i<arr.length;i++) {  
    for(int j=0;j<arr.length;j++) {  
        if(i != j && arr[i]==arr[j])  
            flag=1;  
    }  
    if(flag==0)  
    {  
        System.out.println(arr[i]);  
        break;  
    }  
    flag=0;  
}
```

## Second Approach

**Using SORTING: Time Complexity:  $O(n \cdot \log n)$**

```
Arrays.sort(arr);
for(int i=0;i<arr.length;i=i+2)
{
    if(arr[i]!=arr[i+1])
    {
        System.out.println(arr[i]);
        break;
    }
}
```

## Third Approach

**Using HASHMAP: Time Complexity:  $O(n)$  Space Complexity:  $O(n)$**

```
HashMap<Integer,Integer> map = new HashMap<Integer,Integer>();
    for(int i=0;i<arr.length;i++) {
        if(map.containsKey(arr[i]))
            map.put(arr[i],map.get(arr[i])+1);
        else
            map.put(arr[i],1);
    }
    System.out.println(map); //FOR DEBUGGING
    for(int x :map.keySet()) {
        if(map.get(x)==1){
            System.out.println(x);
            break;
        }
    }
```

## Fourth Approach

**Using for loop: Time Complexity:  $O(n)$  Space Complexity:  $O(1)$**

```
int res=0;
    for(int i=0;i<arr.length;i++) {
        {
            res= res^ arr[i];

        }
    }
    System.out.println(res);
```

## Bitwise XOR

INPUT X	INPUT Y	OUTPUT
0	0	0
0	1	1
1	0	1
1	1	0

So,  $8 \wedge 8 = 0$  And  $8 \wedge 0 = 8$

That is what we used:  $1 \wedge 1 \wedge 3 \wedge 3 \wedge 5 \wedge 6 \wedge 6 = 5 \wedge 0 = 5$



**What if the elements in array are sorted ?**

Input:[1,1,3,3,5,6,6]

Output: 5

**Using LOOPS: Time Complexity:  $O(n)$**

```
for(int i=0;i<arr.length;i=i+2)
{
    if(arr[i]!=arr[i+1])
    {
        System.out.println(arr[i]);
        break;
    }
}
```

## Second Approach

**Using Two Pointer: Time Complexity:  $O(n)$  Space Complexity:  $O(1)$**

```
int left_pointer = 0;
    int right_pointer = arr.length - 1;
    while (left_pointer < right_pointer) {
        if (left_pointer < right_pointer && arr[left_pointer] != arr[left_pointer + 1])
            return arr[left_pointer];
        if (left_pointer < right_pointer && arr[right_pointer] != arr[right_pointer - 1])
            return arr[right_pointer];

        left_pointer += 2;           //1,1,3,3,5,6,6
        right_pointer -= 2;
    }
    return arr[left_pointer];
```

## Third Approach

**Using BinarySearch: Time Complexity:  $O(\log n)$  Space Complexity:  $O(1)$**

Input array: 1,1,3,3,5,6,6]  
              0 1 2 3 4 5 6

First occurrence: Even indexes (0,2,4,.. )

Second occurrence: Odd indexes (1,3,5,..)

But that element (5)

First occurrence: Odd indexes (1,3,5,7,...)

Second occurrence: Even indexes (0,2,4,6...)

```
int low = 0;
int high = arr.length -1;
while(low < high){
    int mid = (low +high)/2;
    if(low==high)
    {
        System.out.println(arr[low]);
        break;
    }
    if( ( (mid & 1) == 0) && arr[mid] == arr[mid+1])        //1,1,2,2,3,3,5,6,6
        low = mid+2;
    else if( (mid & 1) == 0)
        high = mid;
    else if( ( (mid & 1) == 1) && arr[mid] == arr[mid-1])    //1,2,2,3,3,4,4,5,5
        low = mid+1;
    else if( (mid &1) == 1)
        high = mid-1;
}
```

## QUESTION

**Find the maximum sum subarray of size k?**

Input: [2, 1, 5, 1, 3, 2], k=3

Output: 9 //Subarray with maximum sum is [5, 1, 3].

## Sliding Window

**2** 1 5 1 3 2

**2** **1** 5 1 3 2

**2** 1 **5** 1 3 2

2 **1** **5** **1** 3 2

2 1 **5** **1** **3** 2

2 1 5 **1** **3** **2**

## Sliding Window

**Time Complexity:  $O(n)$  Space Complexity:  $O(1)$**

```
public static int function(int k, int[] arr) {  
    int sum = 0, maxSum = 0;  
    int windowStart = 0;  
    for (int windowEnd = 0; windowEnd < arr.length; windowEnd++) {  
        sum += arr[windowEnd];  
        if (windowEnd >= k - 1) {  
            maxSum = Math.max(maxSum, sum);  
            sum -= arr[windowStart];  
            windowStart++;  
        }  
    }  
    return maxSum;  
}
```

## QUESTION

**Find the maximum sum subarray ?**

Input: [-2, 1, 5, 1, 3,-2]

Output: 10    //Subarray with maximum sum is [1, 5, 1, 3].



## Kadane's Algorithm

-2 1 5 1 -3 4

-2 1 5 1 -3 4

-2 1 5 1 -3 4

-2 1 5 1 -3 4

-2 1 5 1 -3 4

-2 1 5 1 -3 4

Local

-2

-2 > global

0

1

1 > global

6

6 > 1

7 > 6

7

4 > 7

8 > 7

8

Global

Int.MinVal

-2

1

6

7

8

## Kadane's Algorithm

**Time Complexity:  $O(n)$  Space Complexity:  $O(1)$**

```
static int function(int arr[])
{
    int global = Integer.MIN_VALUE, local = 0;
    for (int i = 0; i < arr.length; i++)
    {
        local = local + arr[i];
        if (global < local)
            global = local;
        if (local < 0)
            local = 0;
    }
    return global;
}
```

## **Fibonacci Series**

Series: 0 1 1 2 3 5 8....

**What is its recurrence relation?**

## Fibonacci Series

```
static int fibS(int num)
{
    if (num <= 1)
        return num;

    return fibS(num-1) + fibS(num-2);
}
```

## **No of binary string without consecutive 1's**

Input:  $N = 2$

Output: 3 //3 strings are 00, 01, 10

Input:  $N = 3$

Output: 5 //5 strings are 000, 001, 010, 100, 101

## No of binary string without consecutive 1's

<b>0</b>	<b>1</b>	<b>1</b>	<b>2</b>	<b>3</b>	<b>5</b>	<b>8</b>
			n=1	n=2	n=3	n=4

```
static int function(int n)
{
    if(n==0||n==1)
        return n;
    return function(n-1)+function(n-2);
}
public static void main (String args[])
{
    int n=3;
    System.out.println(function(n+2));
}
```

## No of binary string without consecutive 1's

```
static int function(int n)
{
    if(n==0)
        return 0;
    if(n==1)
        return 2;
    if(n==2)
        return 3;
    return function(n-1)+function(n-2);
}
public static void main (String args[])
{
    int n=4;
    System.out.println(function(n));
}
```

**Count no of ways to reach the top. You can take either 1 or 2 step at a time**

**Input:**  $n = 1$

**Output:** 1 //only 1 stair to climb

**Input:**  $n = 2$

**Output:** 2 // (1, 1) and (2)

**Input:**  $n = 4$

**Output:** 5 //(1, 1, 1, 1), (1, 1, 2), (2, 1, 1), (1, 2, 1), (2, 2)



## **Count no of ways to reach the top. You can take either 1 or 2 step at a time**

```
static int fib(int n){
    if(n==0||n==1)
        return n;
    return fib(n-1)+fib(n-2);
}
public static void main (String args[]){
    int n=4;
    System.out.println(fib(n+1));
}
```

# Catalan Series

Series: 1 1 2 5 14 42 132.....

$$1. c(0) = c(1) = 1$$

$$2. c(2) = \begin{array}{ccc} 1 & & 1 \\ & \diagdown & \diagup \\ & 1 & 1 \end{array} = 1+1 = 2$$

$$3. c(3) = \begin{array}{ccc} 1 & & 1 \\ & \diagdown & \diagup \\ 1 & & 1 \\ & \diagdown & \diagup \\ 2 & & 2 \end{array} = 2+1+2 = 5$$

What is its recurrence relation?

## Catalan Series

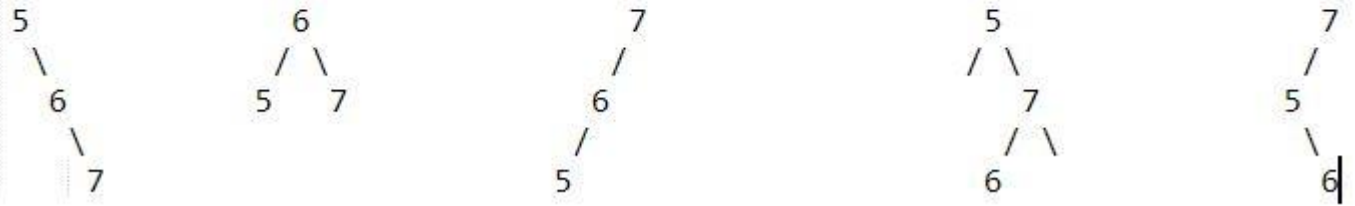
```
static int catS(int num) {  
    if (num <= 1)  
        return 1;  
    int res = 0;  
    for (int i=0; i<n; i++)  
        res += catS(i)*catS(n-i-1);  
    return res;  
}  
  
public static void main(String[] args) {  
    System.out.print(catS(5));  
}
```

# Number of BST's possible from n nodes ?

2 Nodes (5,7)



3 Nodes (5,6,7)



## QUESTION

**Find the Next Largest element / Next Greater element to right**

Input: [2, 1, 5, 1, 3, 2]

Output: [ 5,5,-1,3,-1,-1]

## First Approach

**Time Complexity:  $O(n*n)$  Space Complexity:  $O(1)$**

```
public static int function(int[] arr) {  
    int out=0;  
    for (int itr = 0; itr < arr.length; itr ++){  
        out = -1;  
        for( int jtr=itr+1; jtr<arr.length; jtr++) {  
            if( arr[itr]<arr[jtr]  
                {  
                    out = arr[jtr]; break;  
                }  
        }  
        System.out.println(out);  
    }  
}
```

## LOGIC

**2 1 5 1 3 2**



**STACK**

**5  
1  
3  
2**

Stack is not empty : `s.top()` [Greater the array element]

Stack is empty: -1

`s.pop()` [Less than the element] && stack is not empty

2 1 5 1 3 2



Stack

output

-1

2

2 1 5 1 3 2



pop()

-1

3

2 1 5 1 3 2



1

3

3



## Second Approach

**Time Complexity:  $O(n)$  Space Complexity:  $O(n)$**

```
public static int function(int[] arr) {  
    Stack<Integer> s = new Stack<>(); int new_arr[] = new int[arr.length], i=0;  
    for (int itr = arr.length-1; itr>=0; itr --) {  
        if(s.size()==0)    new_arr[i++]=-1;  
  
        else if( s.size()>0 && s.peek() >arr[itr])  
            new_arr[i++] = s.peek();  
  
        else if(s.size()>0 && s.peek() <= arr[itr]){  
            while(s.size()>0 && s.peek()<=arr[itr] ) s.pop();  
            if(s.size()==0)    new_arr[i++] = -1;  
            else                new_arr[i++] = s.peek();  
        }  
        s.push(arr[itr]);  
        // return the reversed new_arr  
    }  
}
```

## QUESTION

**Find the Next Greater element to left**

Input: [2, 1, 5, 1, 3, 2]

Output: [-1, 2, -1, 5, 5, 3]

## First Approach

**Time Complexity:  $O(n*n)$  Space Complexity:  $O(1)$**

```
for (int itr = 0; itr < arr.length; itr++) {  
    for( int jtr=itr-1; jtr>=0; jtr--) {  
    }  
}
```

## Second Approach

**Time Complexity:  $O(n)$  Space Complexity:  $O(n)$**

```
public static int function(int[] arr) {  
    Stack<Integer> s = new Stack<>(); int new_arr[] = new int[arr.length], i=0;  
    for (int itr = 0; itr<arr.length; itr++) {  
        if(s.size()==0)    new_arr[i++]=-1;  
  
        else if( s.size()>0 && s.peek() >arr[itr])  
            new_arr[i++] = s.peek();  
  
        else if(s.size()>0 && s.peek() <= arr[itr]){  
            while(s.size()>0 && s.peek()<=arr[itr] ) s.pop();  
            if(s.size()==0)    new_arr[i++] = -1;  
            else  
                new_arr[i++]=s.peek();  
        }  
        s.push(arr[itr]);  
    }  
}
```

## QUESTION

**Find the Next Smallest element to left**

Input: [2, 1, 5, 1, 3, 2]

Output: [-1, -1, 1, -1, 1, 1]

## First Approach

**Time Complexity:  $O(n*n)$  Space Complexity:  $O(1)$**

```
for (int itr = 0; itr < arr.length; itr ++) {  
    for( int jtr=itr-1; jtr>0; jtr--) {  
  
    }  
}
```

## Second Approach

**Time Complexity:  $O(n)$  Space Complexity:  $O(n)$**

```
public static int function(int[] arr) {  
    Stack<Integer> s = new Stack<>(); int new_arr[] = new int[arr.length], i=0;  
    for (int itr = 0; itr<arr.length; itr++) {  
        if(s.size()==0)    new_arr[i++]=-1;  
  
        else if( s.size()>0 && s.peek() < arr[itr])  
            new_arr[i++] = s.peek();  
  
        else if(s.size()>0 && s.peek() >= arr[itr]){  
            while(s.size()>0 && s.peek()>=arr[itr] ) s.pop();  
            if(s.size()==0)    new_arr[i++] = -1;  
            else  
                new_arr[i++]=s.peek();  
        }  
        s.push(arr[itr]);  
    }  
}
```

## QUESTION

**Find the Next Smallest element to right**

Input: [2, 1, 5, 1, 3, 2]

Output: [ 1, -1 ,1, -1, 2, -1]



## First Approach

**Time Complexity:  $O(n*n)$  Space Complexity:  $O(1)$**

```
for (int itr = 0; itr < arr.length; itr++) {  
    for( int jtr=itr+1; jtr<arr.length; jtr++) {  
  
    }  
}
```

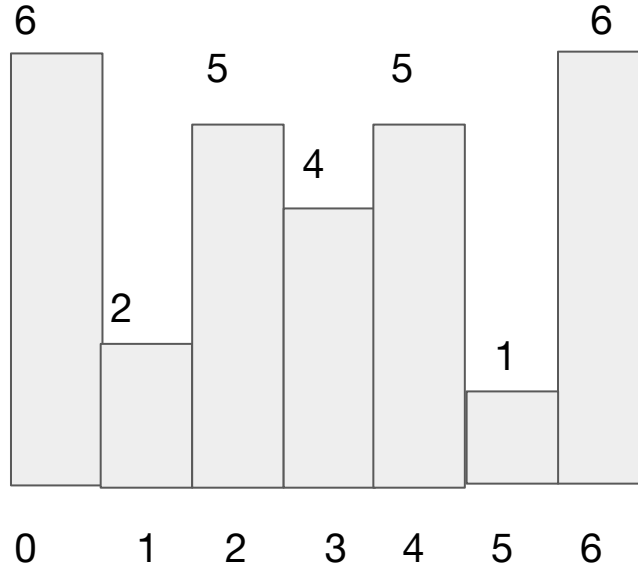
## Second Approach

**Time Complexity:  $O(n)$  Space Complexity:  $O(n)$**

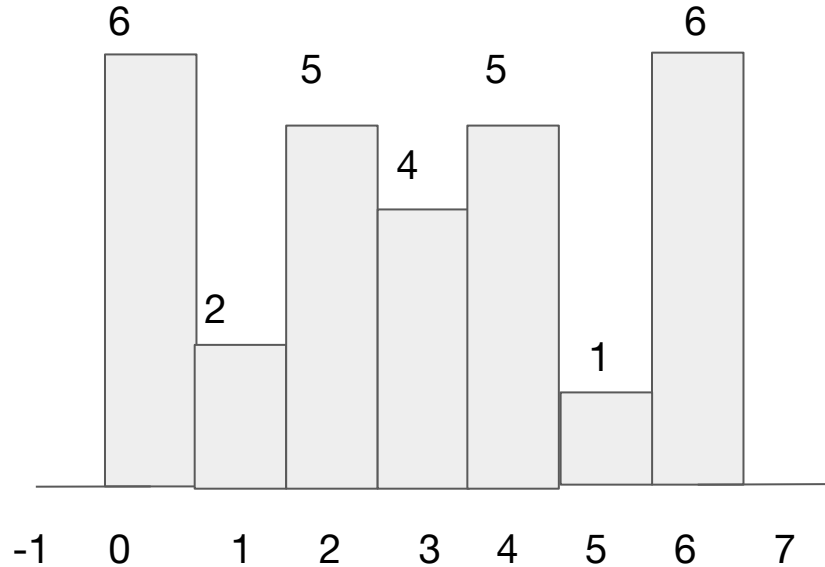
```
public static int function(int[] arr) {  
    Stack<Integer> s = new Stack<>(); int new_arr[] = new int[arr.length], i=0;  
    for (int itr =arr.length-1;itr>=0; itr --) {  
        if(s.size()==0)    new_arr[i++]=-1;  
  
        else if( s.size()>0 && s.peek() < arr[itr])  
            new_arr[i++] = s.peek();  
  
        else if(s.size()>0 && s.peek() >= arr[itr]){  
            while(s.size()>0 && s.peek()>=arr[itr] ) s.pop();  
            if(s.size()==0)    new_arr[i++] = -1;  
            else  
                new_arr[i++] =s.peek();  
        }  
        s.push(arr[itr]);  
        // return the reversed new_arr  
    }  
}
```

## Maximum Rectangular Area in A Histogram

Given : [ 6 2 5 4 5 1 6 ]



# Maximum Rectangular Area in A Histogram



## Approach

2 1 4 1 1 -1 -1      Next Smallest Right  
1 5 3 5 5 7 7      store indexes & reverse

-1 -1 2 2 4 -1 1      Next Smallest Left  
-1 -1 1 1 3 -1 5      store indexes

$\text{max} = \text{Math.max}(\text{max}, \text{arr}[i] * \{ r[i] - l[i] - 1 \})$

## Kth largest element

Input: [ 6 2 5 4 1 7]

K = 3

Output: 5

Input: [ 16 2 50 4 10 70]

K = 3

Output: 16

## First Approach

Sort the Input array: [ 6 2 5 4 1 7]

[ 1 2 4 5 6 7]

Then,  $\text{arr}[k-1]$  is the element

Output: 4

## First Approach



Sort the Input array: [ 6 2 5 4 1 7]

[ 1 2 4 5 6 7]

Then, `arr[arr.length - k]` is the element

Output: 5



## Second Approach



```
PriorityQueue<Integer> heap = new PriorityQueue<Integer>(k);
for(int i =0;i<arr.length;i++){
    heap.add(arr[i]);

    if(heap.size()>k){
        heap.poll();
    }
}
return heap.poll();
```

## Sort nearly sorted array

Input: [ 2 5 4 1 7]

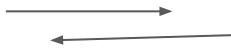
K = 3



Output: [ 1 2 4 5 7]

Input: [ 16 2 50 4 10 70]

K = 3



Output: [2 4 10 16 50 70]

## First Approach

Sort the Input array: [ 2 5 4 1 7]

Output: [ 1 2 4 5 7]

But 'k' was never used.

## Second Approach

```
PriorityQueue<Integer> heap = new PriorityQueue<>();
```

```
for(int i = 0; i < k + 1; i++)  
{  
    heap.add(arr[i]);  
}
```

```
int idx = 0;  
for(int i = k + 1; i < arr.length; i++)  
{  
    arr[idx++] = heap.poll();  
    heap.add(arr[i]);  
}
```

```
while (!heap.isEmpty()) {  
    arr[idx++] = heap.poll();  
}
```

## Minimize the cost of connecting ropes



Input: [ 2 5 4 1 7]

Output: 40

## Minimize the cost of connecting ropes



Input: [ 2 5 4 1 7 ]

Cost:  $3+7+ 12 + 19 = 41$

$$1+2 = 3$$

3 5 4 7

$$3+ 4 = 7$$

7 5 7

$$5+7 = 12$$

12 7

$$12+7$$

19

## Minimize the cost of connecting ropes



Input: [ 2 5 4 1 7]

If we do sorting:

1 2 4 5 7

$$1 + 2 = 3$$

3 4 5 7

$$3 + 4 = 7$$

5 7 7

$$5 + 7 = 12$$

7 12

$$7 + 12 = 19$$

cost : 41

## Minimize the cost of connecting ropes



```
PriorityQueue<Integer> heap = new PriorityQueue<>();
int min =0;
for(int i = 0; i<arr.length; i++)
{
    heap.add(arr[i]);
}

while (!heap.size()>1) {
    int temp = heap.remove()+ heap.remove();

    min += temp;
}

return min;
```



## Summary

- We learnt how to solve different coding problems with different approaches.
- How to reduce time complexity?
- How to reduce space complexity?

# Thank You