+ LINEAR SEARCH \* \* searching: finding given value position in a list of values. \* Linear/sequential search : Basic & simple, compares the target value meth elements of list to identify \* linear clarch me one by one value of pe jeake chek kiya pata hai w.s.t underlying budition. Time Complexity \* Best case - No. of cheeks well the loop make is best case, the element found at othe luder evined so one comparison ie done that is best case. lek array me jitne kau check karne pade har ex position pe wo best case hai \*worst case - 100p will go through every element & they it can be say element not found. code = (oudition (y-dse) loop literation (\* practice questions from github) \* Algorithm = 1) Traversal of array. 2) Compare seil éléments one by one with targeted key. 2) It élément mateines key, écétur current index. 
 0
 1
 2
 3
 4
 5
 6

 10
 20
 40
 30
 15
 9
 35

-> BINARY SCARCH	
* Soutedarrays ascending forders  decendending	
* search sorted array which we us half.	peatedly divide the slatch interval
* idea of binary search to reduce time complexity.	
<b>y ·</b>	
" ALMORITHM "	start index + end index
	Can lange romas index - start + (end-star
(ii) Take see middle shewent	Farget element > middle - seavels
reptuepeating. until search is done array	farget element > middle -> search sight part target element < middle -> search left target = = middle -> oms wer
done	target = = middle -> omswes
I if start index ) end widers	
-> To cheek whether areas	, is sorted in ascending of
(descending.	a wall in and index wall
any oall start we	dus value to end judes value.
ORDER \	-) In out of ( )
BINARY SEARCH else	, decreasing/descending
pool for discussion tosted	allay
farget element	> middle -> search tift part middle -> search sight part
target element &	(middle ) search sight par
A Time Company.	( ) solvi somme ( ) sin angeogo)
Bust case: o(1)	
worst case: ollongn	1 - log Africe of
Total no.	of = log_ Offize of in worst care
	大学を1~7季×変性♥ N - 1777 1777

Pseudocode for Binary Search // whether array is sorted in ascending of descending boolean check Array Order = arr [start] < arr [end]; if this condition is true, so ascending and if false then descending >/function int binaryseasch (int are, int start, write (stard <=end) { calculate mid term intend, int key) { if (mid == ky) } int mid = (start tend)/2; // calculate return mid; mid = (start + lend-start))/2) if (check Array Order) } if (key < mid) { Of Economic spicosod programs & & end = mid -1' (ascend. of (mid == key) { (11 start = same) sectures mid; else } thend = Same; if (check Array Order) { if [mid > key) } ueturn binary search (arr, start, mid-1, key); y (key>mid)} end = mid -1; (Istart = saine) Justind else { setuen binary Search (arr, midt), Start = mid+1' elses Hend = same; 14 (key > mid) { return binary Search (arr, start, mid-1, key) return binarylearch (arr, midtl,

\* when to use sinary search

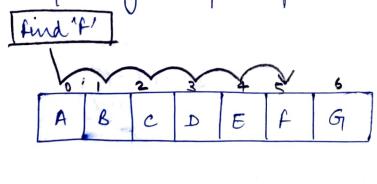
- sorted array

nehen tere problem stated is in sequential particular partiers/way so that if follows a source sequential pattern to get answer.

-> Difference 6/w linear & Binary Search

## LINEAR SCARCH

- \* Input data need not to be sorted.
- \* It does sequential access
- \* Tune complainty -> =0(n)
- \* equality no. of comparisons



## BINARY, SEARCH

- \* Tuput data needs to be Sorted.
- \* It access data randomly \* Time complexity -> O(logn)
- \* performs order of compail-

