Tutorial-4

$$n \log_b^a = n \log_2^3$$

$$n\log_2^3 < n^2$$
 (case 3)

... according to master Theorem $T(n) = O(n^2)$.

Ans2

$$a = 4$$
, $b = 2$
 $n \log_{6}^{a} = n \log_{2}^{4} = n^{2} = f(n)$ (case2)

... according to master Theorem $T(n) = O(n^2 \log n)$

Ans3

$$T(n) = T(n/2) + 2^n$$

 $a = 1, b = 2$
 $n \log 2^1 = n^0 = 1$
 $1 < 2^n (case 3)$

... According to master Theorem T(n) = Q(2n)

Anster's Theorem is not applicable as

a is function of
$$n$$

Anss $T(n) = 16T(n/4) + n$
 $a = 16, b = 4$
 $n + \log b^a = n \log y^6 = n^2$
 $n^2 7 f(n)$
 $T(n) = 0 (n^2)$

Ans6 $T(n) = 2T(n/2) + n \log n$
 $a = 2, b = 2$
 $n \log b^a = n \log_2 2 = n$

Now $f(n) > n$

According to masters $T(n) = 0 (n \log n)$
 $a = 2, b = 2$
 $f(n) = n \log_2 2 = n$
 $f(n) = n \log_2 2 =$

Anso
$$T(n) = 2T(n/4) + no.51$$

$$a = 2, b = 4, f(n) = no.51$$

$$n \log b^{a} = n \log u^{2} = no.5$$

$$no.5 < T(n)$$

$$\therefore According to master's Theorem $T(n) = O(no.5)$$$

$$Anso of the end of$$

Ana12 T(n) = squt(n) + 1/2 + logn

Master's mot applicable as a is not constant.

And T(n) = 3T(n/2) + n $a = 3, b = 2 \qquad f(n) = n$ $n + \log_b a = n \log_2 a = n! \cdot 58$ $n! \cdot 58 + f(n)$

.: According to master's Theorem,

T(n) = 0(nlog23)

And $T(n) = 3T(n/3) + \sqrt{n}$ a = 3/b = 3 $n \log_b a = n \log_3^3 = n$

According to master Theorem, T(n) = Q(n)

And 15. T(n) = UT(n/3) + Cn $a = 4, b = 2 \qquad f(n) = C \times n$ $n \log_b a = n \log_b 4 = n^2$ $n^2 > C \times n$

According to masters Theorem, $T(n) = O(n^2)$

Anule

$$T(n) = 3T(n_{1}) + n \log n$$
 $a = 3$
 $b = 4$
 $f(n) = n \log n$
 $n \log b^{a} = n \log y^{3} = n^{0.79}$
 $n^{0.79} \ge n \log n$
 $n^{0.79} \ge n \log n$
 $f(n) = 0.79$
 $f(n) = 0 (n \log n)$

Anuly

 $T(n) = 3T(n_{3}) + n_{2}$
 $a = 3, b = 3$
 $f(n) = n_{2}$
 $n \log b^{a} = n \log 3^{3} = n$
 $g(n) = g(n_{2})$
 $f(n) = g(n_{2})$
 $f(n) = g(n_{3})$

Anuly

 $f(n) = g(n_{3})$
 $f(n) = g(n_{3})$

$$a=6$$

 $b=3$
 $f(n)=n^2\log n$ $n\log 6=\log 6=n^{1.63}$
 $n^{1.63} < n^2\log n$

: According to moster's Theorem T(n) = O(n2logn)

Anal9.

$$T(n) = 4T(n/2) + n \log n$$
 $a = 4, b = 2$
 $n \log 6^a = n^{\log 2^4} = n^2$
 $n^2 > n/\log n$

According to masters Theorem

 $T(n) = O(n^2)$

Ans20 T(n) = 64T(n/6)-n2 logn

Masters Theorem Is not applicable as f(n) not increasing function.

T(n) = 7T (n/3) +n2 $f(n) = n^2$ $n \log_b^a = n \log_{3}^7 = n^{1.7}$ $n^{1.7} < n^2$

According to masters, T(n) = O(n2) Ans22 T(n) = T(n/2) +n(2-cosn)

Master's theorem isn't applicable since regularity condition is isolated in ease3.