

Lecture 4.1

Network Layer

Address Mapping Protocols

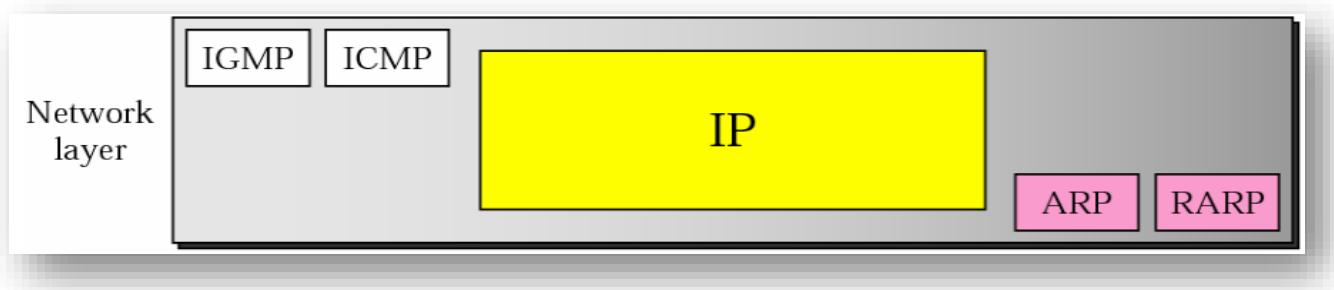
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INTRODUCTION

- IP (Internet Protocol) was designed as a **best-effort delivery protocol**, but it lacks **some features** such as **flow control** and **error control**.
- To make IP more **responsive** it takes help of other protocols.
- **Protocols** to create a **mapping a logical address to a physical address**: ARP (Address Resolution Protocol).
- Protocols to create a **reverse mapping** i.e. **mapping a physical address to a logical address**: RARP, BOOTP, and DHCP.
- Lack of **flow** and **error control** in the **Internet Protocol** has resulted in another protocol, **ICMP**.
- **ICMP reports congestion** and some types of **errors** in the network or destination host.

Address Mapping & Error Reporting Protocols



ADDRESS MAPPING

- An **internet** is made of a combination of **physical networks** connected by **internetworking devices** such as **routers**.
- A **packet** starting from a **source host** may pass through several different **physical networks** before finally reaching the **destination host**.
- The **hosts** and **routers** are **recognized** at the **network level** by their **logical (IP) addresses**, while at the **physical level**, they are **recognized** by their **physical (MAC) addresses**.
- Thus **delivery** of a **packet** to a **host** or a **router** requires **two levels of addressing**: **logical (IP)** and **physical (MAC)**.
- We need to be able to **map** a **logical address** to its **corresponding physical address** and **vice versa**.
- These can be done by using either **static** or **dynamic mapping**.

Static mapping

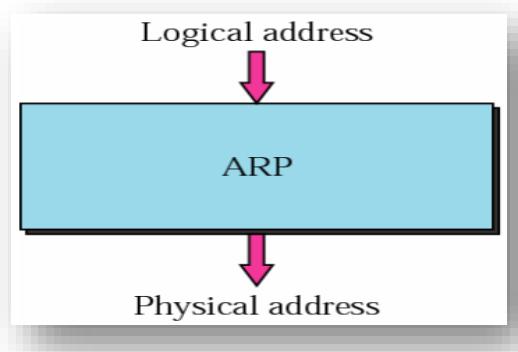
- Static mapping involves in the creation of a table that associates a logical(IP) address with a physical(MAC) address.
- This table is stored in each machine on the network.
- Static mapping has some limitations because physical addresses may change in the following ways:
 - A machine could change its NIC (Network Interface Card), resulting in a new physical address.
 - In some LANs, such as LocalTalk, the physical address changes every time the computer is turned on.

Dynamic mapping

In **Dynamic mapping** each time a machine knows one of the **two addresses** (**logical** or **physical**), it can use a **protocol** to find the other one.

Mapping Logical to Physical Address: ARP

- ARP stands for **Address Resolution Protocol** which is one of the most important protocols of the **Network layer** in the **TCP/IP model**.
- ARP finds the **physical address**, also known as **Media Access Control (MAC) address**, of a **host** from its **known IP address**.

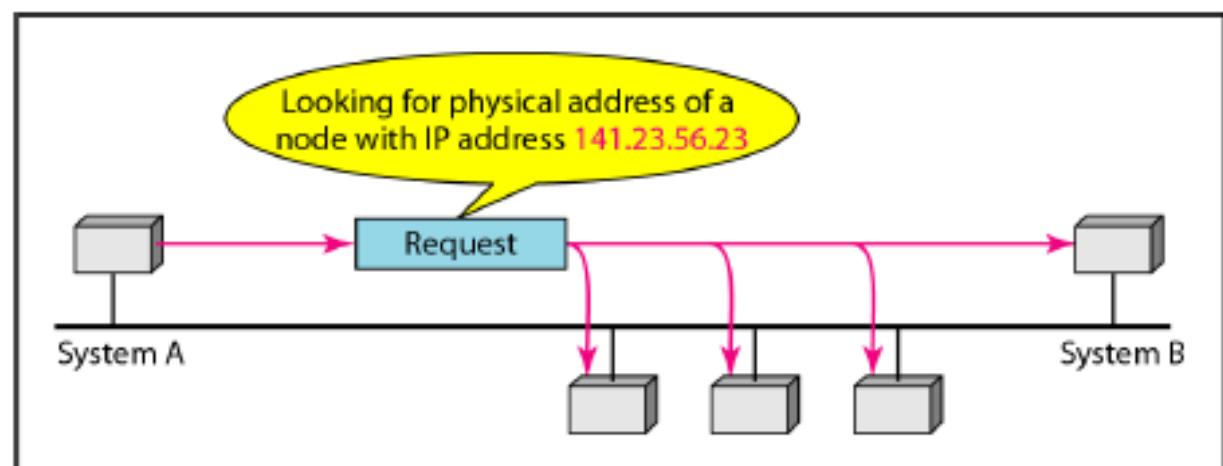


Mapping Logical to Physical Address: ARP

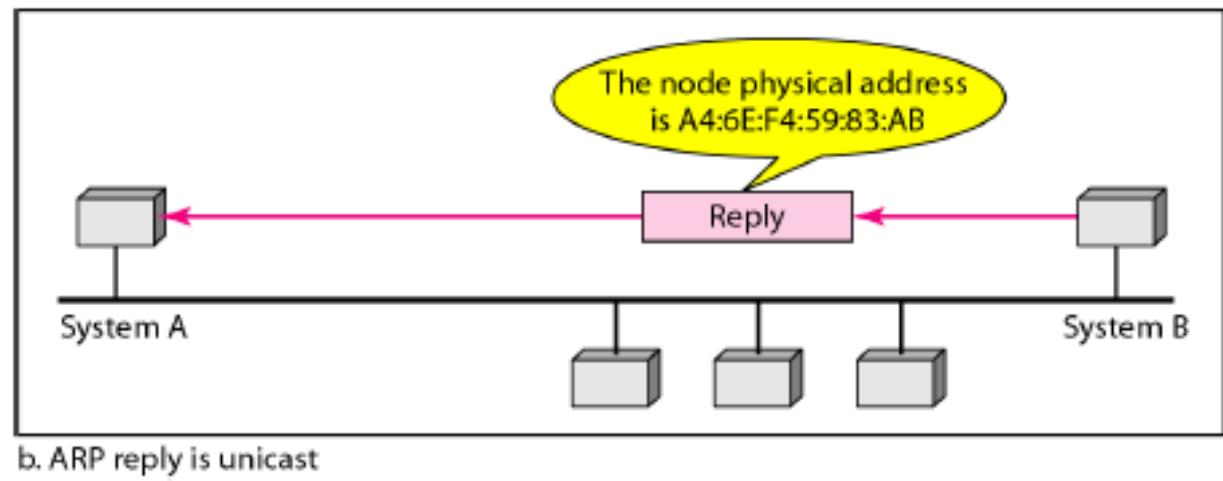
Following **steps** are involved in **logical to physical address mapping**:

- The **host** or the **router** sends an **ARP query packet**.
- The **ARP query packet** includes the **Physical** and **IP addresses** of the **sender** and the **IP address** of the **receiver**.
- As the **sender does not know** the **physical address** of the **receiver**, the **ARP query is broadcast** over the **network**.
- Every **host** or **router** on the **network** receives and processes the **ARP query packet**, but only the **intended recipient** recognizes its **IP address** and sends back an **ARP response packet**.
- The **ARP response packet** contains the **recipient's IP** and **physical addresses**.
- The **ARP response packet** is **unicast directly to the inquirer** (host/router) by using the **physical address** received in the **query packet**.

Mapping Logical to Physical Address: ARP



a. ARP request is broadcast

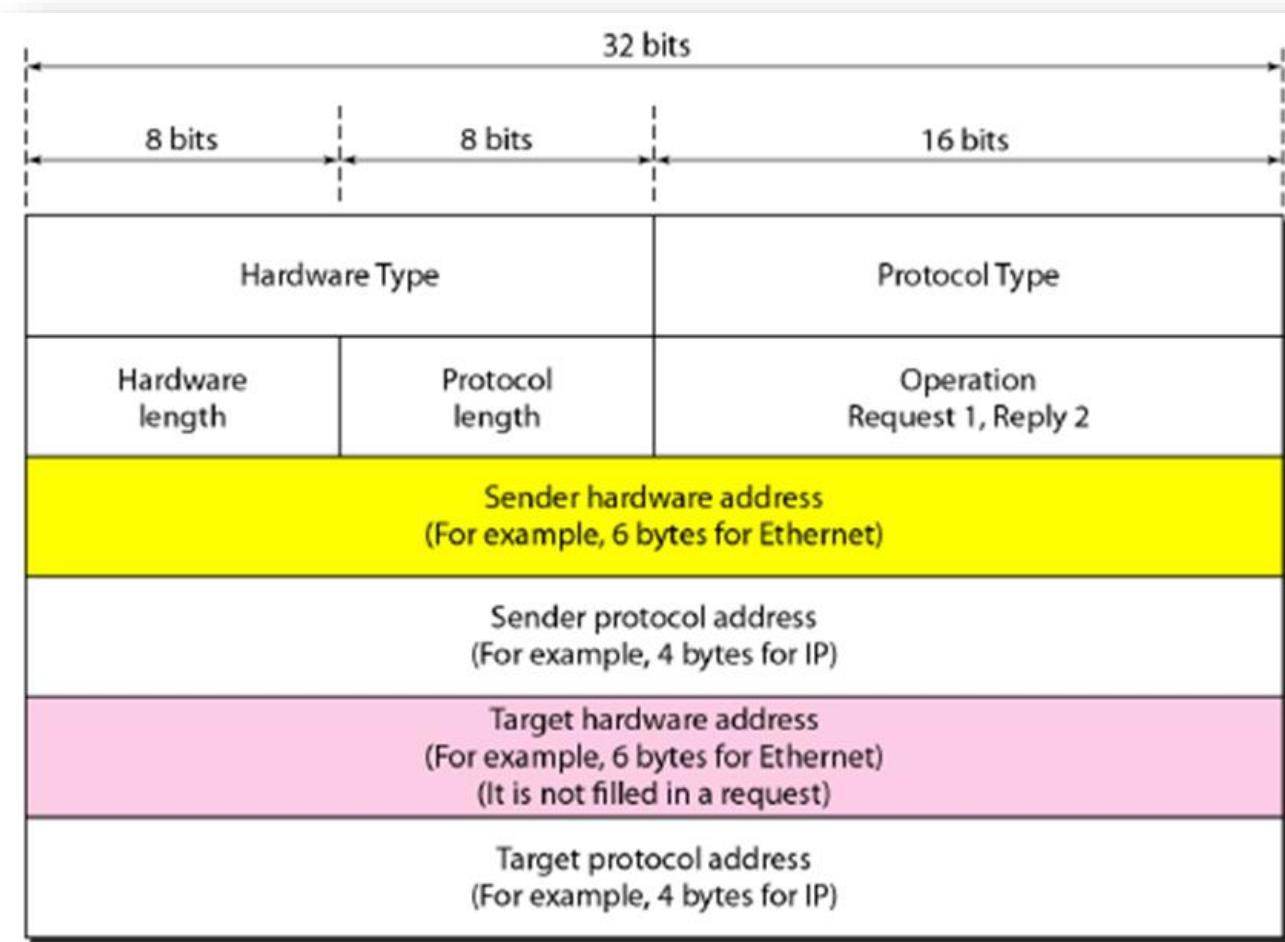


b. ARP reply is unicast

ARP Cache

- Using **ARP** is **inefficient** if **system A** needs to **broadcast** an **ARP request** for each IP **packet** it needs to send to **system B**.
- **ARP** can be useful if the **ARP reply is cached** (kept in **cache memory** for a while) because a system normally sends **several packets** to the **same destination**.
- A **system** that receives an **ARP reply** stores the **mapping** in the **cache memory** and keeps it for **20 to 30 minutes** unless the space in the **cache** is **exhausted**.
- Before sending an **ARP request**, the system first **checks its cache** to see if it can find the **mapping**.

ARP Packet Format



ARP Packet Format

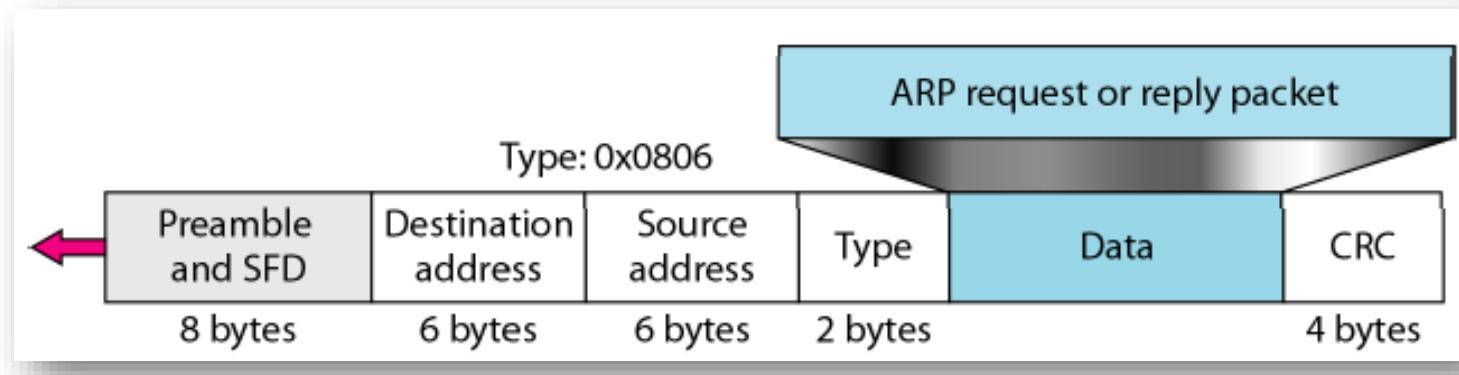
- a. **Hardware type.** This is a **16-bit** field **defining the type of the network** on which **ARP** is running. Each **LAN** has been assigned an **integer** based on its **type**. For example, **Ethernet is given type 1**.
- b. **Protocol type.** This is a **16-bit** field defining the **protocol**. For **example**, the value of this field for the **IPv4** protocol is **080016**.
- c. **Hardware length.** This is an **8-bit** field defining the **length** of the **physical address** in **bytes**. For **example**, for **Ethernet** the value is **6**.
- d. **Protocol length.** This is an **8-bit** field defining the **length** of the **logical address** in **bytes**. For **example**, for the **IPv4 protocol** the value is **4**.
- e. **Operation.** This is a **16-bit** field defining the **type of packet**. Two packet types are defined: **ARP request (1)** and **ARP reply (2)**.

ARP Packet Format

- f. **Sender hardware address.** This is a **variable-length field** defining the **physical address** of the **sender**. For **example**, for **Ethernet** this field is **6 bytes** long.
- g. **Sender protocol address.** This is a **variable-length field** defining the **logical (for example, IP) address** of the **sender**. For the **IPv4 protocol**, this field is **4 bytes** long.
- h. **Target hardware address.** This is a **variable-length field** defining the **physical address** of the **target**. For **example**, for **Ethernet** this field is **6 bytes** long. For an **ARP request** message, this field is **all 0s** because the **sender does not know** the **physical address** of the **target**.
- i. **Target protocol address.** This is a **variable-length field** defining the **logical (for example, IP) address** of the **target**. For the **IPv4 protocol**, this field is **4 bytes** long.

Encapsulation

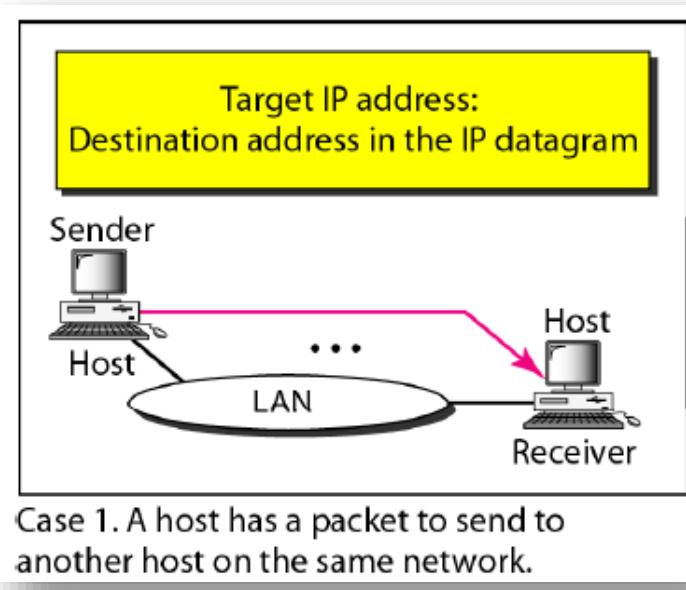
- An **ARP packet** is **encapsulated directly** into a **data link frame**.
- For **example**, in Figure an **ARP packet** is **encapsulated** in an **Ethernet frame**.
- Note that the **type field** indicates that the **data** carried by the **frame** are an **ARP packet**.



Different Cases of ARP Operation

Case 1

- The **sender** is a **host** and wants to **send** a packet to **another host** on the **same network**.
- In this case, the **logical address** that must be **mapped** to a **physical address** is the **destination IP address** in the datagram header.

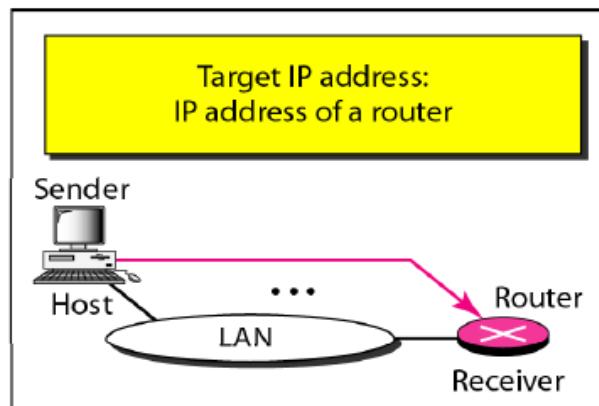


Different Cases of ARP Operation

Case 2

The **sender** is a **host** and wants to **send** a **packet** to **another host** on **another network**.

- In this case, the **host** looks at its **routing table** and finds the **IP address** of the **next hop (router)** for this **destination**.
- If it **does not** have a **routing table**, it looks for the **IP address** of the **default router**.
- The **IP address** of the **router** becomes the **logical address** that must be **mapped** to a **physical address**.



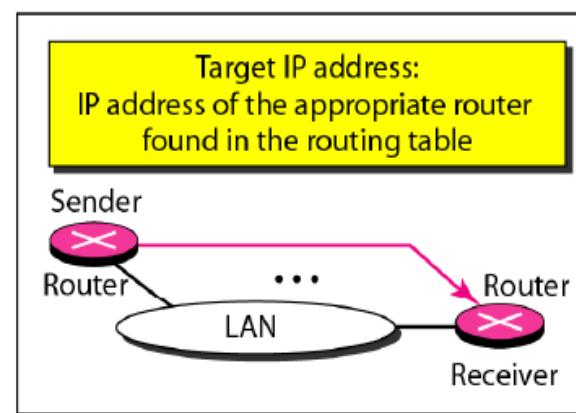
Case 2. A host wants to send a packet to another host on another network.
It must first be delivered to a router.

Different Cases of ARP Operation

Case 3

The **sender** is a **router** that has **received** a **datagram** destined for a **host** on **another network**.

- Router checks its **routing table** and finds the **IP address** of the **next router**.
- The **IP address** of the **next router** becomes the **logical address** that must be mapped to a **physical address**.



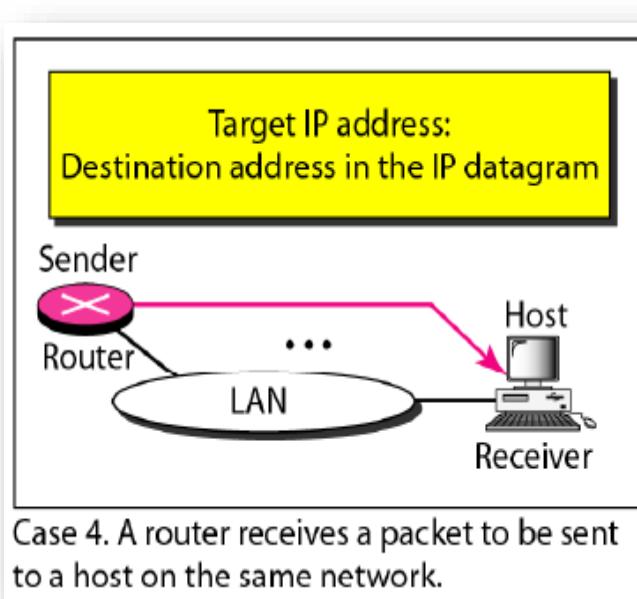
Case 3. A router receives a packet to be sent to a host on another network. It must first be delivered to the appropriate router.

Different Cases of ARP Operation

Case 4

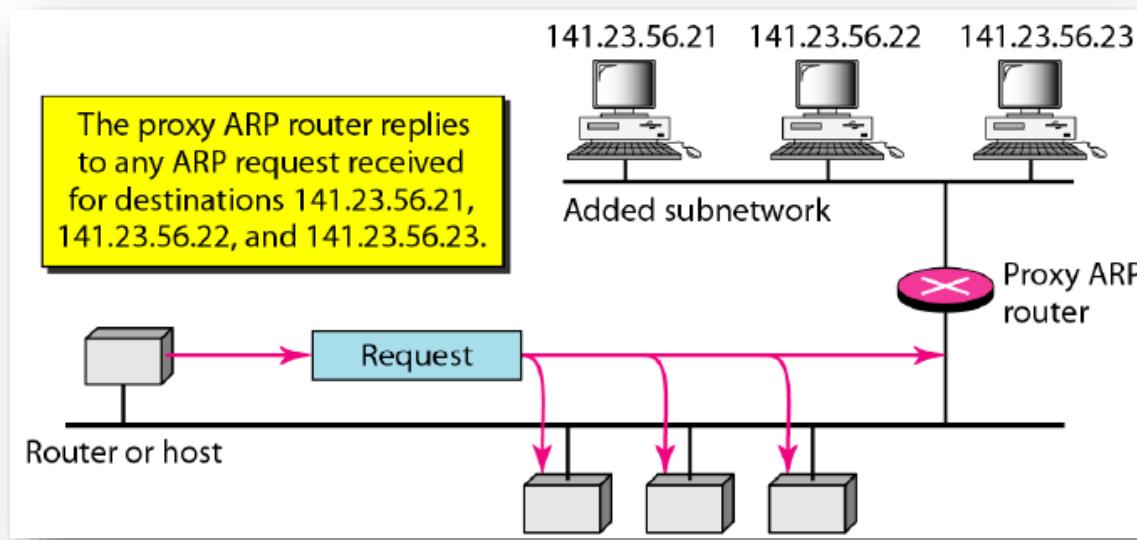
The **sender** is a **router** that has **received** a **datagram** destined for a **host** on the **same network**.

- The **destination IP address** of the **datagram** becomes the **logical address** that must be **mapped** to a **physical address**.



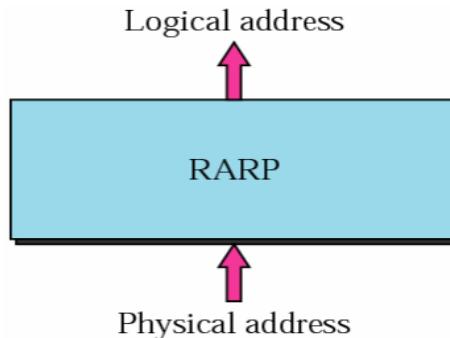
Proxy ARP

- A **proxy ARP** is an **ARP** that acts **on behalf** of a **set of hosts**.
- Whenever a **router** running a **proxy ARP** receives an **ARP request** looking for the **IP address** of one of these hosts, the **router sends** an **ARP reply** announcing **its own hardware (physical) address**.
- After the **router** receives the **actual IP packet**, it **sends** the **packet** to the **appropriate host or router**.



Mapping Physical Address to Logical Address

- There are **occasions** in which a **host** knows its **physical address**, but **needs** to know its **logical address**. This may happen in **two cases**:
- **Case 1:** A *diskless station* is just *booted*. The station can find its **physical address** by checking its **interface**, but it does **not know** its **IP address**.
- **Case 2:** An organization does **not have enough IP addresses** to assign to each **station**; it needs to **assign IP addresses on demand**. The station can **send** its **physical address** and ask for a **short time lease**.



Reverse Address Resolution Protocol (RARP)

- Reverse Address Resolution Protocol (RARP) finds the logical address for a machine that knows only its physical address.
- To create an IP datagram, a host or a router needs to know its own IP address .
- The IP address of a machine is usually read from its configuration file stored on a disk file.
- However, a diskless machine is usually booted from ROM, which has minimum booting information.
- The ROM is installed by the manufacturer.
- It cannot include the IP address because the IP addresses on a network are assigned by the network administrator.

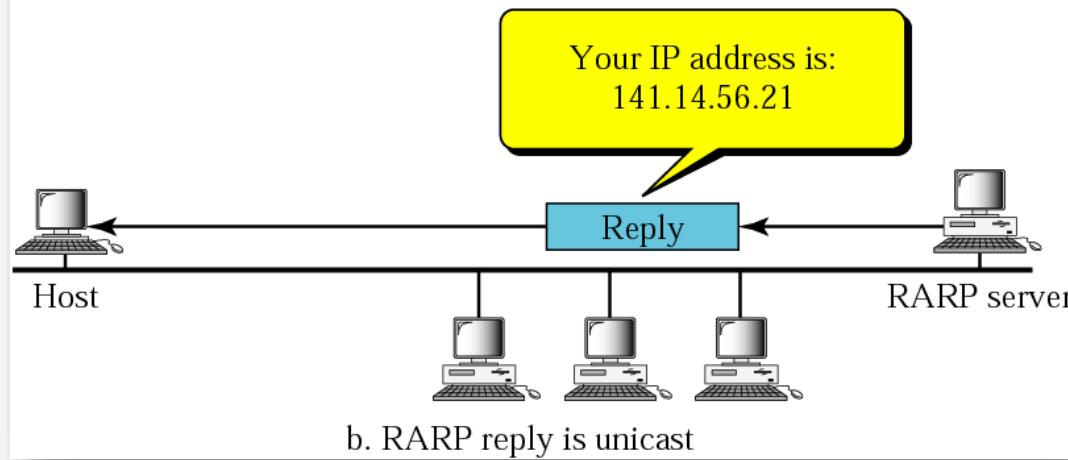
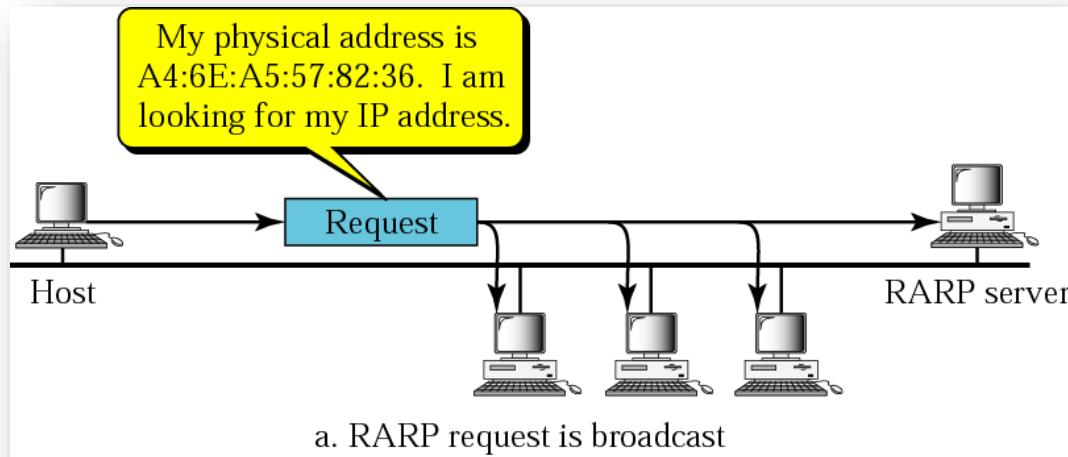
Reverse Address Resolution Protocol (RARP)

- The **machine** can get its **physical address** (by reading its **NIC**, for example), which is unique locally.
- It can then **use** the **physical address** to get the **logical address** by using the **RARP protocol**.

RARP Operation

- A **RARP request** is **created** and **broadcast** on the **local network**.
- Another **machine** on the **local network** that **knows all the IP addresses** will respond with a **RARP reply**.
- The **requesting machine** must be running a **RARP client** program; and the **responding machine** must be running a **RARP server** program.

RARP Operation



RARP Packet Format & Encapsulation

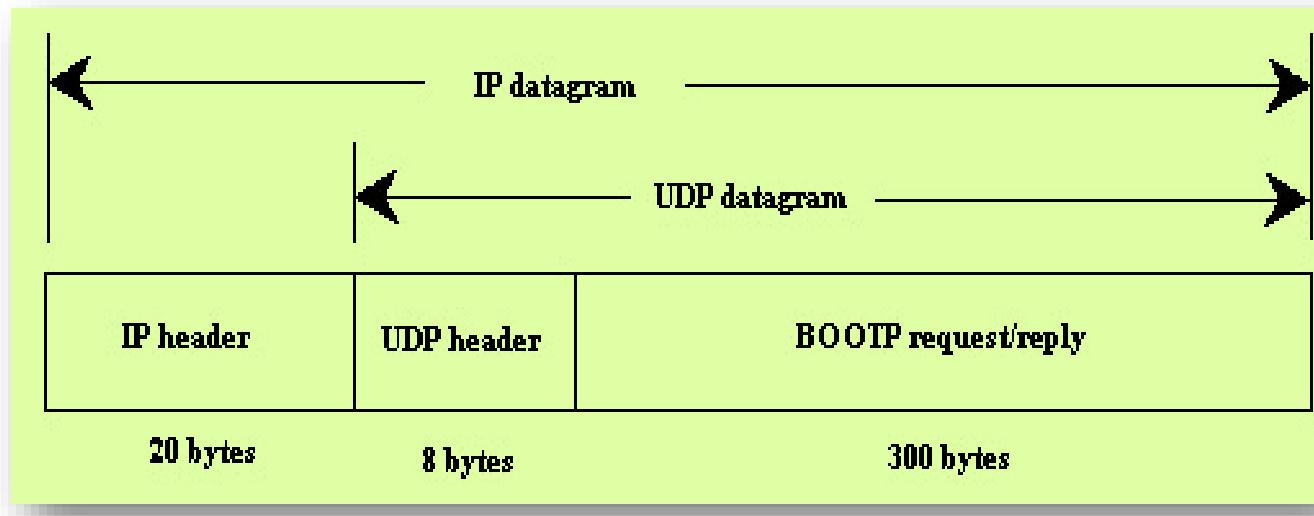
- The format of the **RARP packet** is the **same as the ARP packet** format, except that the **Operation field**.
- Its value is **3** for **RARP request** message and **4** for **RARP reply** message.
- An **RARP packet** is also **encapsulated** directly into a **data link frame** just like **ARP packet**.

Limitations of RARP

- As **broadcasting** is done at the **data link layer**.
- The **physical broadcast address**, all **1's** in the case of **Ethernet**, does not pass the boundaries of a network.
- This means that if an **administrator** has **several networks or several subnets**, it needs to **assign** a **RARP server** for **each network or subnet**.
- This is the **reason** that **RARP is almost obsolete**.
- Two protocols, **BOOTP** and **DHCP**, are replacing **RARP**.

Bootstrap Protocol (BOOTP)

- The **Bootstrap Protocol (BOOTP)** is a **Client/Server** based **protocol** at **Application layer**, designed to **provide physical address to logical address mapping**.
- The **administrator** may put the **client** and the **server** on the **same network** or on **different networks**.
- BOOTP** messages are **encapsulated** in a **UDP packet**, and the **UDP packet** itself is **encapsulated** in an **IP packet**.



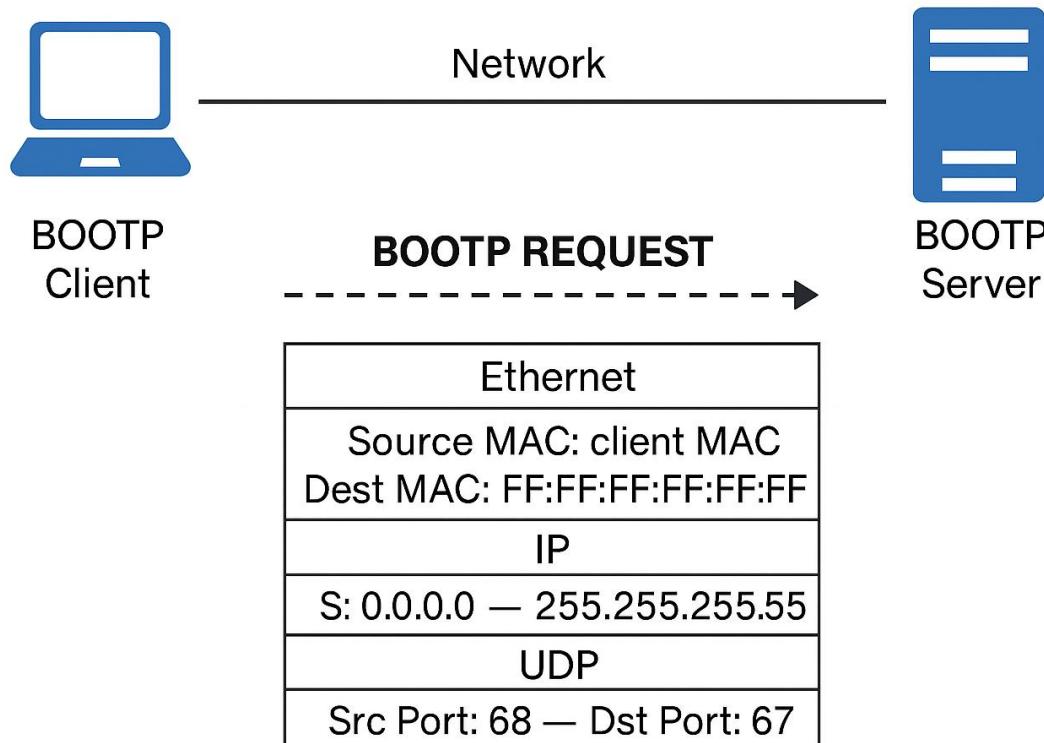
Bootstrap Protocol (BOOTP)

- *Case 1: Client and server on same network*
- When a **BOOTP client** is started, it has **no IP address**, so it **broadcasts** a **message** containing its **MAC address** onto the network.
 - A **BOOTP client** (e.g., diskless workstation or network device) **doesn't know** its **own IP address** when it starts.
 - It sets its **source IP address = 0.0.0.0** in the **request**, because it has no configured IP.
 - The **destination IP address = 255.255.255.255** (**broadcast**).
- This message is called a “**BOOTP request**”.

Bootstrap Protocol (BOOTP)

- *Case 1: Client and server on same network*
- Encapsulation of **BOOTP request** at data link layer
 - Since the **client** doesn't know the **server's MAC**, it uses **Ethernet broadcast**.
 - **Destination MAC** = **FF:FF:FF:FF:FF:FF**
 - **Source MAC** = client's **own MAC address** (unique and known).

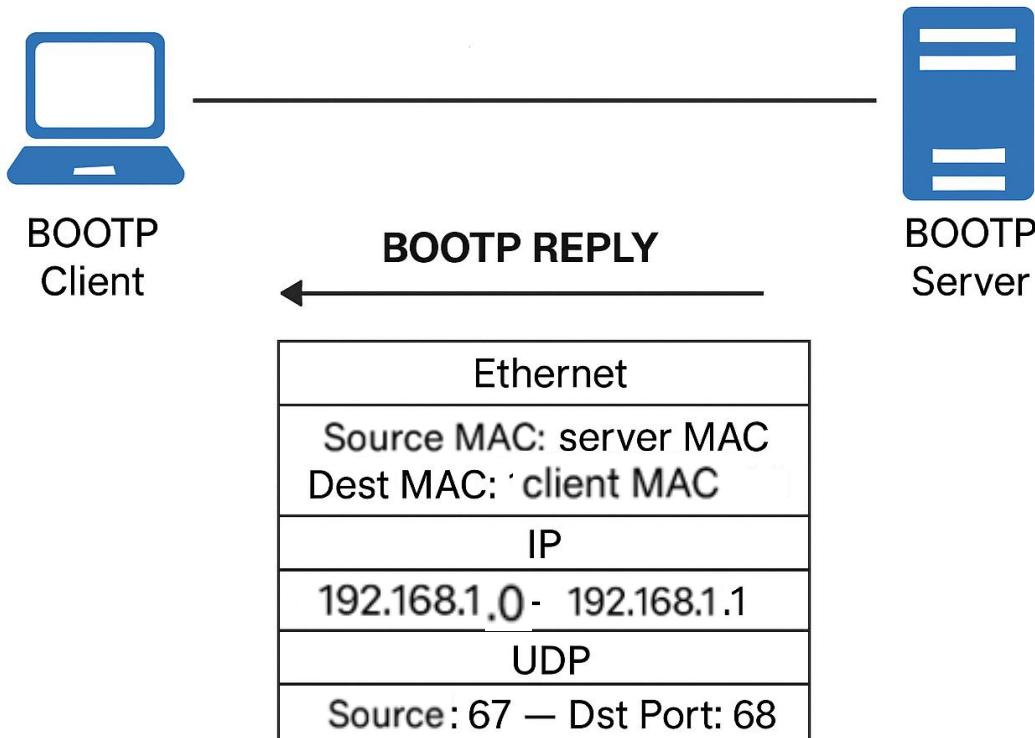
BOOTP client and server on the same network



Bootstrap Protocol (BOOTP)

- *Case 1: Client and server on same network*
- “**BOOTP request**,” is picked up by the **BOOTP server**, which **replies** to the client with the following information that the client needs:
 1. The **client’s IP address**, **subnet mask**, and **default gateway address**.
 2. The **IP address** and **host name** of the **BOOTP server**.
- When the **client receives** this information from the **BOOTP server**, it **configures** and **initializes** its **TCP/IP protocol stack**, and then **connects** to the **server**.

Bootstrap Protocol (BOOTP)



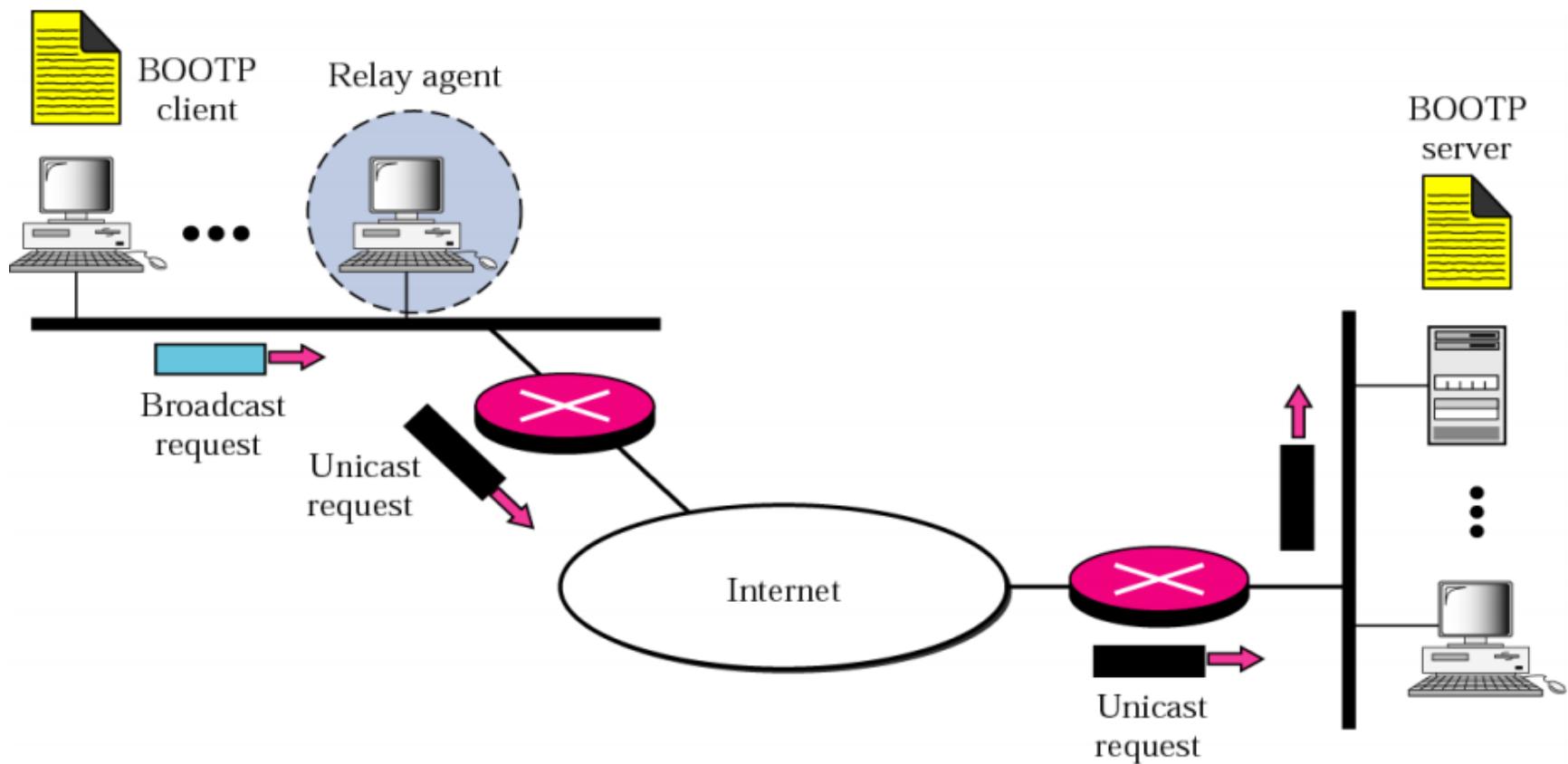
Bootstrap Protocol (BOOTP)

- *Case 2 : Client and server on different networks*
- The **BOOTP request** is **broadcasted** because the **client** does not know the **IP address** of the **server**.
- The client simply uses **all 0's** as the **source address** and **all 1's** as the **destination address**.
- But a **broadcast** IP datagram **cannot pass through any router**.
- To **solve the problem**, there is a need for an **intermediary**.
- **One of the hosts in local network** (or a **router** that can be configured to operate at the **application layer**) can be used as a **relay** .
- The **host** in this case is called a **Relay Agent**.

Bootstrap Protocol (BOOTP)

- *Case 2 : Client and server on different networks*
- The **relay agent** knows the **unicast address** of a **BOOTP server**.
- When it **receives** this type of packet, it **encapsulates** the message in a **unicast datagram** and **sends** the **request** to the **BOOTP server**.
- The packet, carrying a **unicast destination address**, is **routed** by any router and reaches the **BOOTP server**.
- The **BOOTP server** knows the message comes from a **relay agent** because one of the fields in the request message defines the **IP address** of the **relay agent**.
- **BOOTP server** sends a **BOOTP reply** message to the **relay agent**.
- The **Relay Agent**, after receiving the **reply**, **sends it** to the **BOOTP client**.

BOOTP client and server on different network



Limitations of BOOTP

- **BOOTP is not a Dynamic configuration protocol.**
- When a **client** requests its **IP address**, the **BOOTP server** consults a **table** that matches the **physical address** of the **client** with its **IP address**.
- This implies that the **binding** between the **physical address** and the **IP address** of the client **already exists**.
- The **binding is predetermined** i.e. **Static**.
- However, what if a **host moves** from **one physical network** to **another?**
- What if a **host wants** a **temporary IP address?**
- **BOOTP cannot handle these situations** because the **binding** between the **physical** and **IP addresses** is **static and fixed** in a **table** until changed by the **administrator**.
- Thus **BOOTP is a Static configuration protocol**.

Dynamic Host Configuration Protocol (DHCP)

- The **Dynamic Host Configuration Protocol (DHCP)** has been devised to provide **Static** and **Dynamic Address Allocation** that can be manual or automatic as required.

Static Address Allocation

- In this capacity **DHCP acts as BOOTP** does.
- It is **backward compatible** with **BOOTP**, which means a **host** running the **BOOTP client** can request a **static address** from a **DHCP server**.
- A **DHCP server** has a **static database** that **Statically binds Physical addresses to IP addresses**.

Dynamic Host Configuration Protocol (DHCP)

Dynamic Address Allocation

- DHCP has a **Second database** with a pool of **available IP addresses**.
- This **second database** makes **DHCP dynamic**.
- When a **DHCP client** requests a **temporary IP address**, the **DHCP server** goes to the **pool of available (unused) IP addresses** and **assigns an IP address** for a **negotiable period of time**.
- When a **DHCP client** sends a **DHCP request** to a **DHCP server**, the server **first checks its Static database**.
- If an **entry** with the requested **physical address** exists in the **Static database**, the permanent **IP address** of the **client** is assigned.

Dynamic Host Configuration Protocol (DHCP)

- On the other hand, if the entry **does not exist** in the **Static database**, the server selects an **IP address** from the **available pool**, **assigns** the **address** to the client, and **adds** the entry to the **Dynamic database**.
- The **Dynamic** aspect of **DHCP** is **needed** when a **host moves from network to network** or is connected and disconnected from a network.
- **DHCP** provides **temporary** IP addresses for a **limited time**.
- The **DHCP server** issues a **lease for a specific time**.
- When the **lease expires**, the **client** must either **stop** using the **IP address** or **renew** the **lease**.
- The **server** has the **option** to **agree** or **disagree** with the **renewal**.
- If the **server disagrees**, the **client stops** using the **address**.