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Btech CSE (DS)

Ques :- 6 :- Distributed Deadlock simulation :-

Given :- (wait for graph segments)

$\rightarrow S_1 : P_1 \rightarrow P_2, P_3 \rightarrow P_4$

$\rightarrow S_2 : P_2 \rightarrow P_5, P_5 \rightarrow P_6$

$\rightarrow S_3 : P_6 \rightarrow P_1$

(a) Global wait for graph.

Combine all edges :-

$P_1 \rightarrow P_2$

$P_3 \rightarrow P_4$

$P_2 \rightarrow P_5$

$P_5 \rightarrow P_6$

$P_6 \rightarrow P_1$

(b) Deadlock detection :-

Look for cycle

$P_1 \rightarrow P_2 \rightarrow P_5 \rightarrow P_6 \rightarrow P_1$

→ This forms a cycle → DEADLOCK exists

Deadlock processes $\rightarrow \{P_1, P_2, P_5, P_6\}$

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(c) Chandy-Misra-Hoag Algorithm \rightarrow (Suitable distributed Deadlock Algo) for detection

Uses probe messages to detect cycles in distributed wait for graphs without a central coordinator.

Ques 1 → Given :-

→ Local access = 5 ms

→ Remote access = 25 ms

→ probability (remote) = 0.3

→ probability (local) = 0.7

(a) Expected file access time

$$E = (0.7 \times 5) + (0.3 \times 25)$$

$$E = 3.5 + 7.5$$

$$E \approx 11 \text{ ms.}$$

Uses "client side cache with LRU Eviction".

Reason :-

- Frequently accessed remote files are cached locally \rightarrow reduces remote access.
- LRU ensures only least used files are removed.
- Minimises average access time and reduces network load.

Ques 8 \rightarrow given :-

\rightarrow full checkpoint = 200 ms

\rightarrow incremental checkpoint = 50 ms.

\rightarrow must recover within 1 sec RPO.

\rightarrow Total period \rightarrow 10 seconds

(a) optimal min : Use 1 full checkpoint for every 10 seconds.

\rightarrow incremental checkpoints every 1 second.

So, over 10 seconds $\rightarrow 10 \times 50 = 500$ ms.

(b) RPD = 1 second \rightarrow must not lose more than 1 second of progress.

\rightarrow incremental checkpoints every second satisfy RPD.

- full checkpoint every 10 sec prevents log chains and speed recovery.
- combination gives low overhead + fast rec.

Ques 9 → (a) challenges :-

- 1) highly variable workload (flash sales, sudden spikes).
- 2) Ensuring requests are routed to the least loaded region.
- 3) Network latency and geographic distance.
- 4) preventing overload of a single data centre.

* Suitable Algorithm :-

- Dynamic Load Balancing using "Least loaded sensor + Global Distributed load balancer".
- Each data centre reports load to a global scheduler.
- Requests routed to the region with minimum active load.
- Adapts in real time during flash sale spikes.

Use active-active Multi-Region Deployment
Components :-

- 1) geo replicated databases (eg:- multimaster replication)
→ keeps $RPO \approx 0$ (almost zero data loss)
- 2) global failover using anycast DNS/Traffic manager.
→ Redirects traffic to healthy region within seconds → low RTO.
- 3) distributed message queues + stateless microservices.
→ Allows seamless services restart in other regions.
- 4) Regular checkpoints and log replication.
→ Ensures latest state recovery.

* Result :-

- even if one regional data centre fails, a service continues from another region with minimal downtime (low RTO) and minimal data loss (low RPO).