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Btech cse (DS)

Ques - 6 :- distributed deadlock simulation :-

Given :- (wait for graph segments)

→ S1 :  $P_1 \rightarrow P_2$ ,  $P_3 \rightarrow P_4$

→ S2 :  $P_2 \rightarrow P_5$ ,  $P_5 \rightarrow P_6$

→ S3 :  $P_6 \rightarrow P_1$

(a) Global wait for graph.

Combine all edges :-

$P_1 \rightarrow P_2$

$P_3 \rightarrow P_4$

$P_2 \rightarrow P_5$

$P_5 \rightarrow P_6$

$P_6 \rightarrow P_1$

(b) Deadlock Detection :-

Look for cycle

$P1 \rightarrow P2 \rightarrow P5 \rightarrow P6 \rightarrow P1$

Date

$\rightarrow$  This forms a cycle  $\rightarrow$  DEADLOCK exists

Deadlock processes  $\rightarrow \{P1, P2, P5, P6\}$

(c) Chandy-Misra, Mass Algorithm  $\rightarrow$  (Suitable distributed Deadlock Algo) for detection

Uses probe messages to detect cycles in distributed wait for graphs without a central coordinator.

Ques 7  $\rightarrow$  given :-

$\rightarrow$  Local access = 5 ms

$\rightarrow$  Remote access = 25 ms

$\rightarrow$  probability (remote) = 0.3

$\rightarrow$  probability (local) = 0.7

(a) Expected file access time

$$E = (0.7 \times 5) + (0.3 \times 25)$$

$$E = 3.5 + 7.5$$

$$E = 11 \text{ ms.}$$

Uses "client side cache with LRU Eviction".

Reason :

- Frequently accessed remote files are cached locally  $\rightarrow$  reduces remote access.
- LRU ensures only least used files are removed.
- Minimises average access time and reduces network load.

Ques 8  $\rightarrow$  given :-

- $\rightarrow$  full checkpoint = 200 ms
- $\rightarrow$  incremental checkpoint = 50 ms.
- $\rightarrow$  must recover within 1 sec RPO.
- $\rightarrow$  Total period  $\rightarrow$  10 seconds

(a) optimal min : Use 1 full checkpoint for every 10 seconds.

- $\rightarrow$  incremental checkpoints every 1 second.
- So, over 10 seconds  $\rightarrow 10 \times 50 = 500$  ms.

(b) RPD = 1 second  $\rightarrow$  must not lose more than 1 second of progress.

- $\rightarrow$  incremental checkpoints every second satisfy RPD.

- full checkpoint every 10 sec prevents log chains and speed recovery.
- combination gives low overhead + fast no.

Ques 9 → (a) challenges :-

- 1) highly variable workload (flash sales, sudden spikes)
- 2) Ensuring requests are routed to the least loaded region.
- 3) Network latency and geographic distance
- 4) preventing overload of a single data centre

\* Suitable Algorithm :-

- Dynamic Load Balancing using "Least loaded sensor + Global Distributed load balancer".
- Each data centre reports load to a global scheduler.
- Requests routed to the region with minimum active load.
- Adapts in real time during flash sale spikes.



Use active-active Multi-Region Deployment Components :-

1) geo replicated databases (eg:- multimaster replication)

→ keeps  $RPO \approx 0$  (almost zero data loss)

2) Global Failover using anycast DNS/Traffic manager.

→ Redirects traffic to healthy region within seconds → low RTO.

3) Distributed message queues + stateless microservices.

→ Allows seamless services restart in other regions.

4) Regular checkpoints and log replication.

→ Ensures latest state recovery.

\* Result :-

→ Even if one regional data centre fails, a service continues from another region with minimal downtime (Low RTO) and minimal data loss (Low RPO).