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C14-58.

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EXPERIMENT NO. 10.

Aim: write a program to eliminate left recursion from the given grammar.

Theory:

A production of grammar is said to have left recursion if the leftmost variable of its RHS is same as variable of its LHS.

A grammar G (V, T, P, S) is left recursive if it has a production in the form,

$$A \rightarrow A\alpha / \beta$$

the above grammar is left recursive because the left of production is only at first position on RHS.

we can replace left recursion by replacing a pair of production with

$$A \rightarrow \beta A'$$
$$A' \rightarrow \alpha A' / \epsilon$$

In left recursive grammar, expansion of A will generate $A\alpha$, $A\alpha\alpha$, $A\alpha\alpha\alpha$, on each side causing it to enter into an infinite loop.

Eg: $E \rightarrow E+T / T$

$T \rightarrow T * f / F$

$F \rightarrow (E) / id$

compare, $E \rightarrow E+T / T$

with $A \rightarrow A\alpha / \beta$

$\therefore A = E, \quad \alpha = +T, \quad \beta = T.$

$A \rightarrow A\alpha / \beta$ changes to $A' \rightarrow \alpha A' / \epsilon$

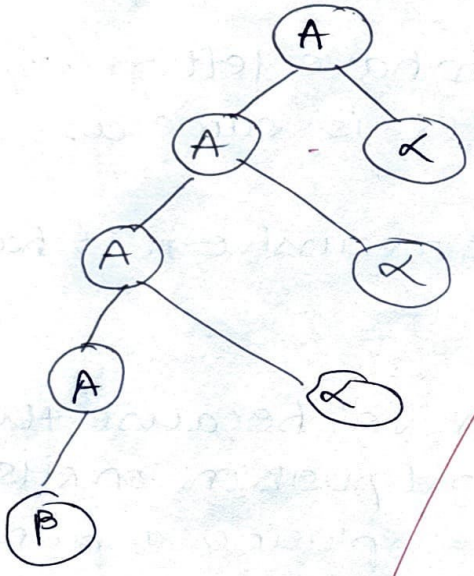
$A \rightarrow \beta A' \quad \therefore E = TE'$

$E \rightarrow E+T / T$

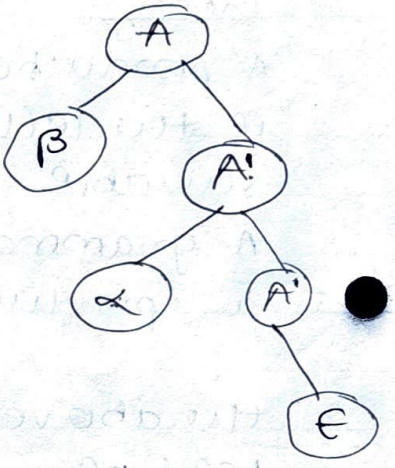
$E \rightarrow TE'$

$E' \rightarrow +TE' / \epsilon$

$A \rightarrow A\alpha \mid \beta$ $\xrightarrow{\text{removal of left recursion}}$ $A \rightarrow \beta A'$, $A' \rightarrow \alpha A' \mid \epsilon$



Removal of left recursion



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compare $T \rightarrow T*F \mid F$ with $A \rightarrow A\alpha \mid \beta$

$T \rightarrow FT'$

$T' \rightarrow *FT' \mid \epsilon$

∴ final set of production is.

$E \rightarrow TE'$

$E' \rightarrow +TE' \mid \epsilon$

$T \rightarrow FT'$

$T' \rightarrow +FT' \mid \epsilon$

$F \rightarrow (E) \mid id.$

18/3/24 (A-4)