

<p>C. Classy Numbers</p> <p>time limit per test: 3 seconds memory limit per test: 256 megabytes</p> <p>Let's call some positive integer <i>classy</i> if its decimal representation contains no more than 3 non-zero digits. For example, numbers 4, 200000, 10203 are <i>classy</i> and numbers 4231, 102306, 7277420000 are not.</p> <p>You are given a segment $[L; R]$. Count the number of <i>classy</i> integers x such that $L \leq x \leq R$.</p> <p>Each testcase contains several segments, for each of them you are required to solve the problem separately.</p> <p>Input</p> <p>The first line contains a single integer T ($1 \leq T \leq 10^4$) — the number of segments in a testcase.</p> <p>Each of the next T lines contains two integers L_i and R_i ($1 \leq L_i \leq R_i \leq 10^{18}$).</p> <p>Output</p> <p>Print T lines — the i-th line should contain the number of <i>classy</i> integers on a segment $[L_i; R_i]$.</p>	<p>– Problem tags</p> <p>combinatorics dp *1900</p> <p>No tag edit access</p>
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Problem Breakdown

Lets understand the problem first

A number is *classy* if:

- It has at most 3 non-zero digits (0, 1, 2, or 3 non-zero digits allowed)
- Examples:
 - 00000 - 1 non-zero - *classy*
 - 102000 - 2 non-zero - *classy*
 - 10203 - 3 non-zero - *classy*
 - 4231 - 4 non-zero - NOT *classy*
 - 7277420000 - many non-zero - NOT *classy*

We need to answer up to $T = 10,000$ queries, each with:

L, R ($1 \leq L \leq R \leq 10^{18}$)

Goal:

Count *classy* numbers in $[L, R]$.

This means we need something fast — brute force is impossible.

Digit DP is perfect.

Digit DP State

We design:

```
dfs(pos, nonZeroCount, tight, leadingZero)
```

Where:

- pos : digit index (0..18)
- nonZeroCount : how many non-zero digits we have placed so far (must be ≤ 3)
- tight : whether we match the prefix of X
- leadingZero : whether number hasn't started yet (still 000.. prefix)

Why leadingZero?

Because leading zeros do not count as non-zero digits.

Base Case

```
If pos == len:
    return (nonZeroCount <= 3) ? 1 : 0
```

Transition

At pos, digit can be:

```
for d = 0 .. (tight ? digits[pos] : 9):
    nextTight = tight && (d == limit)
    nextLeadingZero = leadingZero && (d == 0)
    nextNonZero = nonZeroCount + (d != 0 ? 1 : 0)
    if nextNonZero <= 3:
        ans += dfs(...)
```

Code

```
#include <bits/stdc++.h>
using namespace std;
using ll = long long;

vector<int> digits;
ll dp[20][4][2]; // pos, nonZeroCount, tight, leadingZero
int len;

ll dfs(int pos, int nz, int tight, int leadingZero) {
    if (pos == len) {
        return (nz <= 3) ? 1 : 0;
    }

    ll &res = dp[pos][nz][tight][leadingZero];
    if (res != -1) return res;

    res = 0;
    int limit = tight ? digits[pos] : 9;

    for (int d = 0; d <= limit; d++) {
        int nextTight = tight && (d == limit);
        int nextLZ = leadingZero && (d == 0);
        int nextNZ = nz + ((d == 0 && nextLZ) ? 0 : (d != 0));
        // simpler: if (leadingZero && d==0) --> nextNZ = nz
        if (leadingZero && d == 0) nextNZ = nz;

        if (nextNZ <= 3) {
            res += dfs(pos + 1, nextNZ, nextTight, nextLZ);
        }
    }
    return res;
}

ll solveUpTo(ll x) {
    if (x < 0) return 0;
    digits.clear();

    if (x == 0) digits.push_back(0);
    while (x > 0) {
        digits.push_back(x % 10);
        x /= 10;
    }
    reverse(digits.begin(), digits.end());
    len = digits.size();
}

memset(dp, -1, sizeof(dp));
return dfs(0, 0, 1, 1);
}

int main() {
    ios::sync_with_stdio(false);
    cin.tie(nullptr);

    int T;
    cin >> T;

    while (T--) {
        ll L, R;
        cin >> L >> R;
        cout << solveUpTo(R) - solveUpTo(L - 1) << "\n";
    }
}
```