


698. Partition to K Equal Sum Subsets

Solved 

Medium  Topics  Compares  Hint

Given an integer array `nums` and an integer `k`, return `true` if it is possible to divide this array into `k` non-empty subsets whose sums are all equal.

Example 1:

Input: `nums = [4,3,2,3,5,2,1]`, `k = 4`
Output: `true`
Explanation: It is possible to divide it into 4 subsets (5), (1, 4), (2,3), (2,3) with equal sums.

Example 2:

Input: `nums = [1,2,3,4]`, `k = 3`
Output: `false`

Constraints:

- $1 \leq k \leq \text{nums.length} \leq 16$
- $1 \leq \text{nums}[i] \leq 10^4$
- The frequency of each element is in the range $[1, 4]$.

Key Observations

We are given:

- `nums[]` with $n \leq 16$
- We must partition into `k` non-empty subsets
- Each subset must have the same sum

Step 1: Total sum check (mandatory)

Let:

`total = sum(nums)`

If:

- `total % k != 0`
- Impossible immediately \rightarrow return false

Let:

`target = total / k`

Every subset must sum to target.

We need to track:

- Which elements are already used
- How full the current subset is

Since $n \leq 16$, we can represent usage by a bitmask of length `n`:

- Bit `i` = 1 \rightarrow `nums[i]` already used
- Total states = $2^n \leq 65536$ \rightarrow totally manageable

So each DP state represents:

"Which elements have already been taken into some completed or current subset."

DP State

We use:

`dp[mask]` = current sum of the ongoing subset

Meaning:

- `mask` \rightarrow which elements are already used
- `dp[mask]` \rightarrow how much sum we have accumulated in the current group
- All previous groups are already completed with sum = target

Special meaning:

- If `dp[mask] == 0`, it means:
 - Either we are at the start
 - Or we just completed one subset and moved on to the next

Base case:

`dp[0] = 0`

No elements used, current subset sum = 0.

Final goal:

`mask == (1<<n)-1 AND dp[mask]==0`

All elements used and we just finished a group exactly.

Transitions

From a given `mask`, we try to add one unused element `i`:

Conditions:

- not used $\rightarrow (\text{mask} \& (1<<i)) == 0$
- Adding `nums[i]` does not overflow:

`dp[mask] + nums[i] <= target`

Then:

`newMask = mask | (1<<i)`
`newSum = (dp[mask] + nums[i]) % target`
`dp[newMask] = newSum`

Why % target?

If a subset reaches exactly target, we reset to 0 and start filling the next group.

Code

```
class Solution {
public:
    bool canPartitionKSubsets(vector<int>& nums, int k) {
        int n = nums.size();
        int total = 0;
        for (int x : nums) total += x;

        if (total % k != 0) return false;
        int target = total / k;

        // Optimization: sort descending (prunes faster)
        sort(nums.rbegin(), nums.rend());
        if (nums[0] > target) return false;

        int N = 1 << n;
        vector<int> dp(N, -1);
        dp[0] = 0;

        for (int mask = 0; mask < N; mask++) {
            if (dp[mask] == -1) continue;

            for (int i = 0; i < n; i++) {
                if (mask & (1 << i)) continue;

                if (dp[mask] + nums[i] <= target) {
                    int newMask = mask | (1 << i);
                    int newSum = (dp[mask] + nums[i]) % target;
                    dp[newMask] = newSum;
                }
            }
        }

        return dp[N - 1] == 0;
    }
};
```

At any moment:

- `mask` decides which numbers are already fixed permanently
- `dp[mask]` tracks how full the current bucket is
- When `dp == 0`, we just finished one bucket and move to the next
- We never exceed target
- Since we use exactly all numbers, we necessarily form `k` equal-sum subsets