

Implement a Heap Data Structure

You have been provided with a partial implementation of a heap data structure code. The code already has the following functions implemented:

- `initialize(int v[], int n)` – Initialize the heap to store vertices in array `v` with all their keys set to ∞ .
- `heapify(int i)` – This operation is required to re-store heap property for implementing certain operations, such as `removeMin` operation, and `updateKey` operation. This function is already implemented in given sample skeleton.
- `getKey(int data)` – This operation returns the `key` value of the given `data` in the heap.
- `empty()` – Returns `true` if the heap is empty, otherwise returns false.
- `printHeap()` - Prints the heap.

Your job is to complete the implementation that support the following heap operations:

- `buHeapify(int i)` – This operation is called when key value of an existing node is decreased or a new item is inserted at the end of the heap. For example, when the `key` value of a data item at node index `i` is decreased, the node may be required to move up the tree so that heap property (parent's `key` less than or equal to child's `key`) remains correct. This may be implemented as a recursive function or in iterative manner that starts swapping of nodes at node number `i`, then may call recursively at parent of `i`. In the worst case, the function will end at the root of the tree (all nodes from `i` up-to the root will be swapped).
- `void insertItem(int data, int key)` – This function inserts a new `(data, key)` pair in the heap. The new item is first inserted at the end of the heap which may violate heap property. So `buHeapify` operation is executed to ensure that heap property is restored. In the worst case, the inserted node may end up moving at the root node of the heap.
- `HeapItem removeMin()` – This operation will return the heap node that has the minimum key value. Must restore heap property by calling `heapify` after removal.
- `updateKey(int data, float key)` – This operation updates the `key` value of the given `data` to the given `key`. The function first searches for given `data` in the heap and then updates the `key` value. After update, it may happen that the new `key` violates min-heap property. Hence, after update, a call to `heapify` or `buheapify` is required.

Note the following-

-You must extend this class to implement your own Heap class. You cannot write your own class.

-Note that a special array *map* is kept in this *MinHeap* class to keep track of where each data value is currently stored inside the heap. This is required for searching in the *updateKey(data, key)* operation, because we need to know in which node the input data is stored.

-Without keeping this *map* array, you may find where *data* is stored by a linear search through all heap nodes which will require $O(n)$ time (n is the number of nodes in heap). In this case, *updateKey* will also require $O(n)$ time which is prohibited.

-However, keeping a *map* array does the data search in constant time, which gives the overall time of $O(\log n)$ for the *updateKey* operation. *This is very much important for certain algorithms such as Prim's MST algorithm or Dijkstra's shortest path algorithm to ensure their optimal running time.*

-The *map* array also helps us for *getKey* operation to return the *key* value of a vertex in constant time without keeping additional storage.