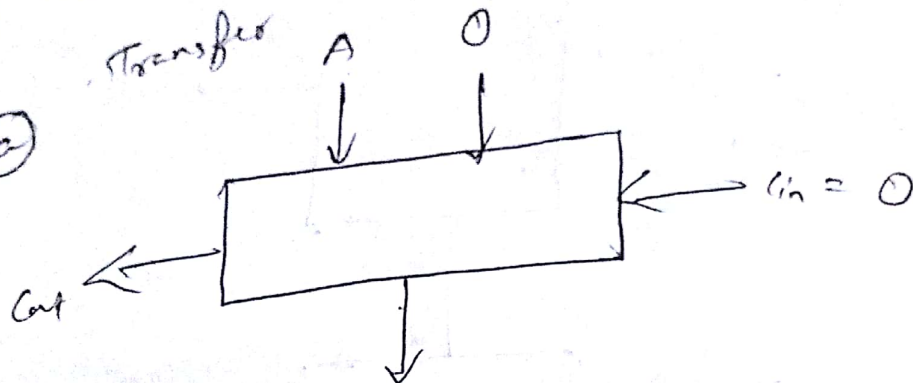
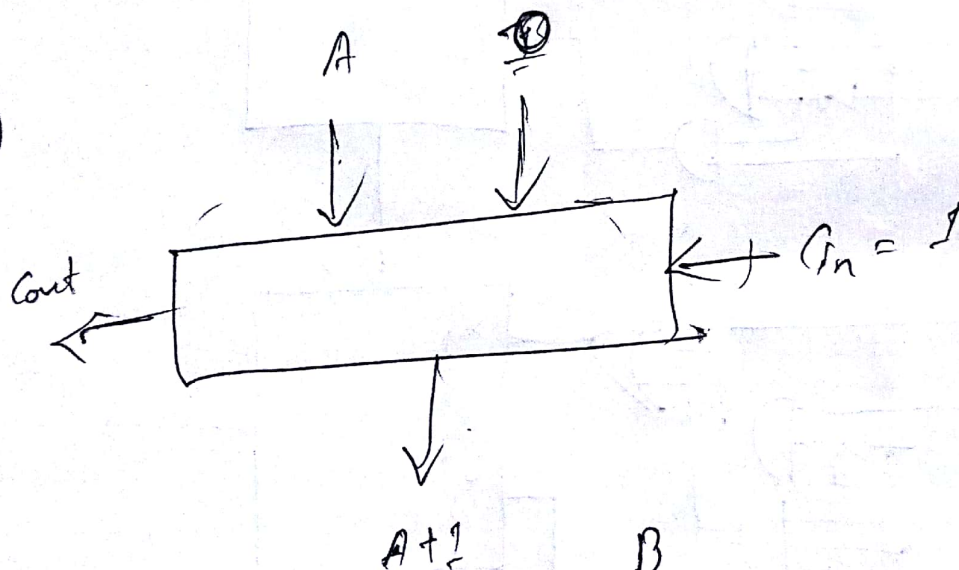


Arithmetic Part:

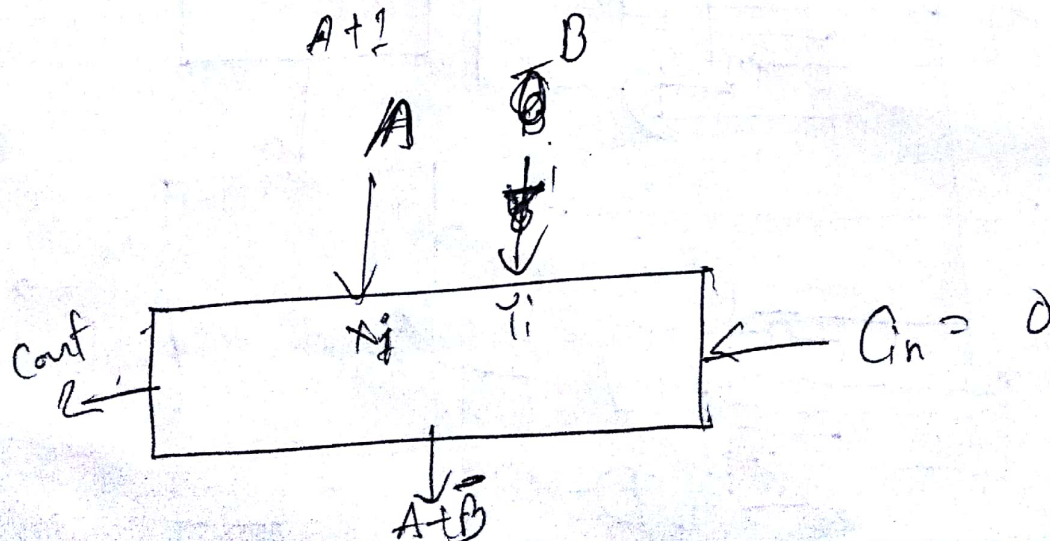
(a)



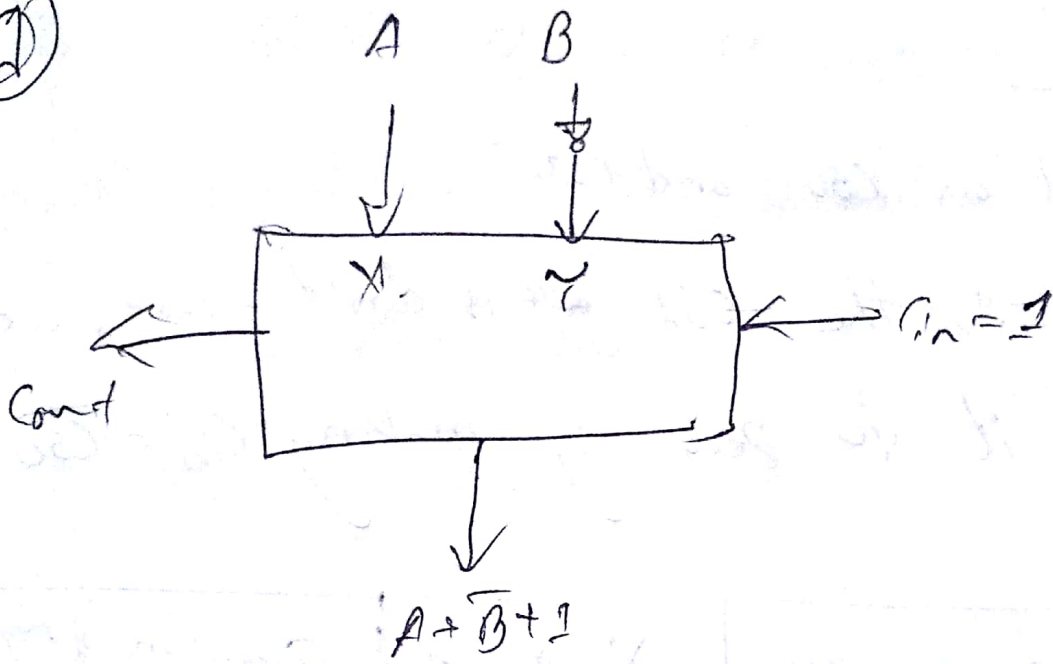
(b)



(c)



②



CS2	CS1	Y
0	0	0
0	1	0
1	0	\overline{B}
1	1	\overline{B}

$$Y_i = CS2 \overline{B} + CS2 CS1 \overline{B}$$

$$= CS2 \overline{B}_i$$

$$X_i = A_i$$

Logical Post:

Dependent on CS0 and CS1
 When CS0 = 1, the CS1 ~~and~~ is don't-care. We can force it to zero by making $C_{in} = \overline{CS0} CS1$.

Now,

CS2	CS1	CS0	X_i	Y_i	C_i	Operation	Required operation
0	0	X	A_i	0	0	$F_i = A_i$ (Transfer)	OR
1	X	1	A_i	B_i	0	$F_i = A_i \oplus B_i$ (Equivalence)	AND NOT

So, when $CS2 CS0 = 01$, we can OR B_i with A_i and the result will be $A + B$.

$$\therefore X_i = A_i + CS2 CS0 B_i$$

Again, when $CS_2 CS_0 = 11$, here, we have to get AND operation where the output is equivalence.

$$F_i = A_i \oplus B_i'$$

$$= A_i B_i + A_i B_i'$$

Let us investigate the possibility of ORing each input A_i with some boolean function K_i when $CS_2 CS_1 = 11$.

$$F_i = X_i \oplus Y_i$$

$$= (A_i + K_i) \oplus B_i'$$

$$= A_i B_i + K_i B_i + A_i' K_i B_i'$$

Taking $K_i = B_i'$, we get, $F_i = A_i B_i + B_i' B_i + A_i B_i B_i'$

$$= A_i B_i$$

$$\therefore X_i = A_i + CS_0 CS_2' B_i + CS_0 CS_2 B_i'$$

$$X_i = A_i + CS_0 \overline{CS_2} B_i + CS_0 CS_2 \overline{B_i}$$

$$Y_i = \overline{CS_2} B_i$$

$$C_{in} = \overline{CS_0} CS_1$$

