

# Chapter 8

# Stack

Stack is a one dimensional data structure.

Items are added and removed in last-in first-out manner.

The most recent item is called top of the stack (TOS).

# Declaring stack segment

`.STACK 100H`

When the program is assembled and loaded in memory,

- SS register contains the segment number of the stack segment.
- SP register is initialized to 100h for the preceding stack declaration. This indicates that the stack is empty. When the stack is not empty, SP contains the offset address of the TOS.

# PUSH and PUSHF

PUSH source (e.g. PUSH AX)

source = 16 bit register or memory word

Execution of PUSH causes the following to happen:

- SP is decreased by 2.
- A copy of the source content is pushed to the address specified by SS:SP. The source is unchanged

## PUSHF

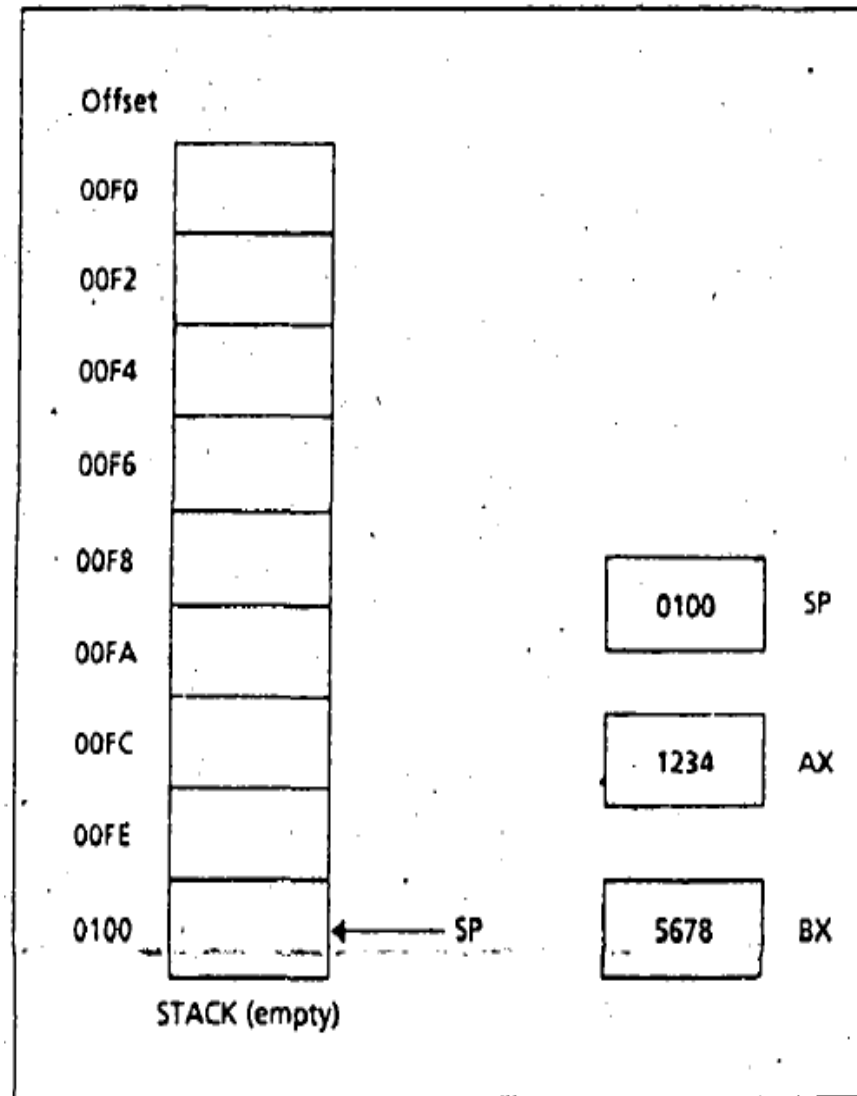
This instruction pushes the contents of the FLAGS register onto the stack.

There is no effect of PUSH/PUSHF on the flags.

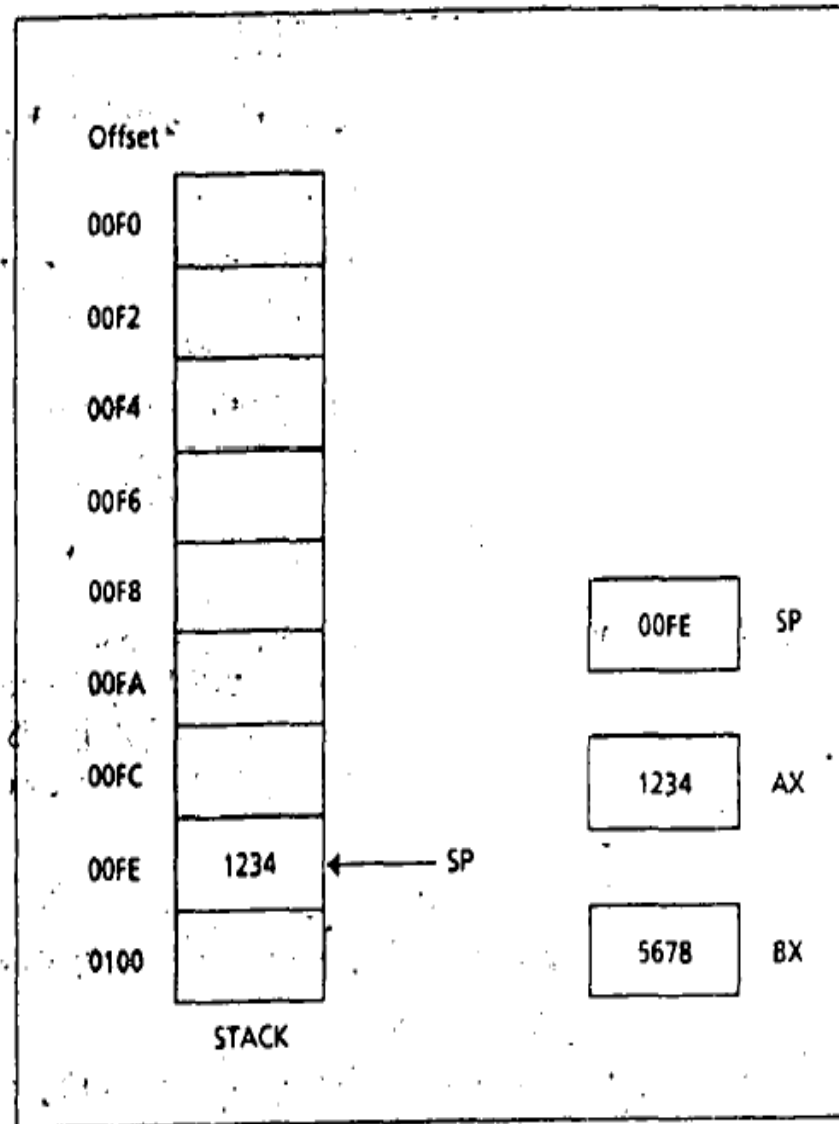
Illegal: PUSH DL (Byte instruction is illegal)

PUSH 2 (Push of immediate data is illegal)

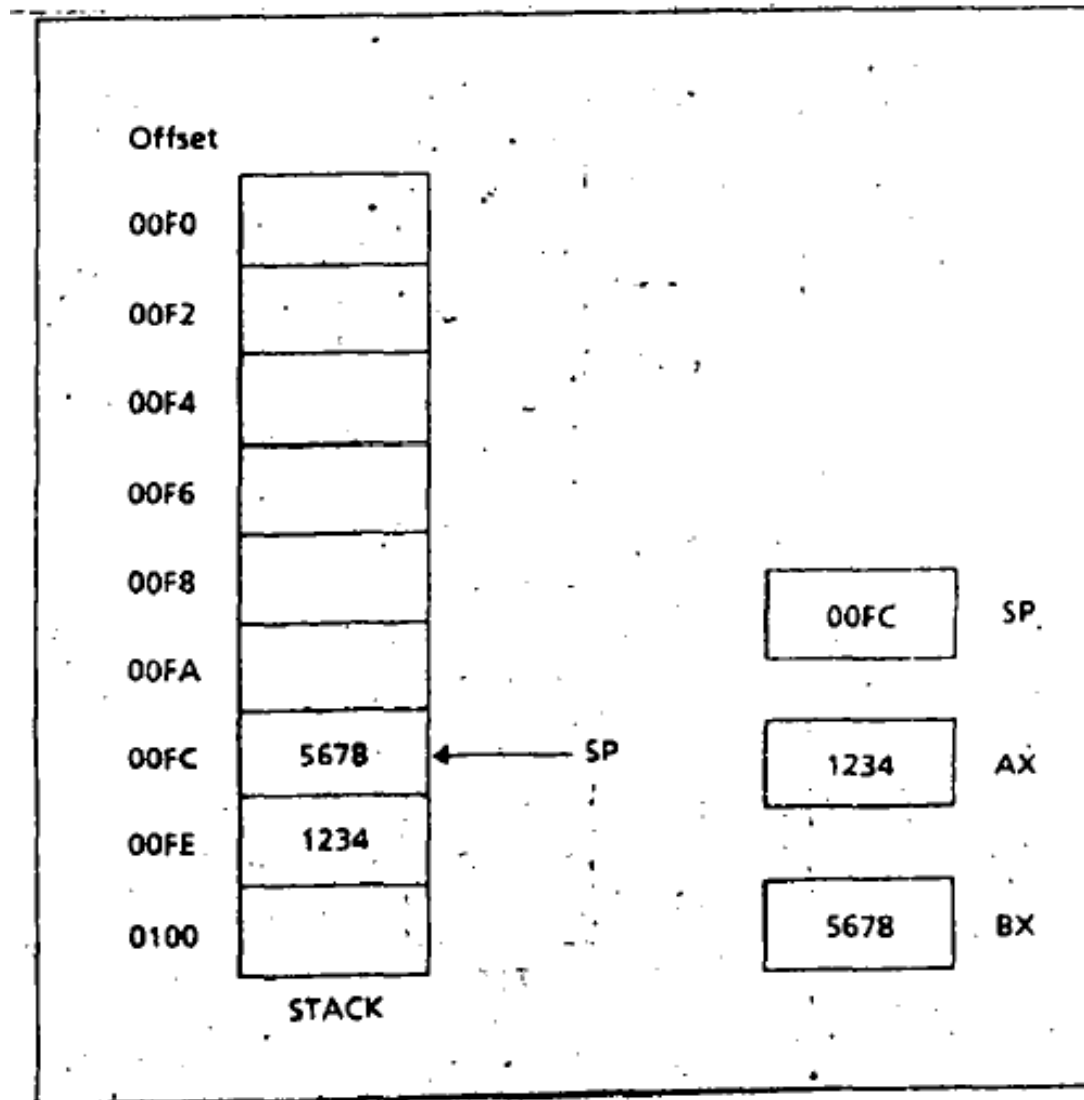
# Empty Stack



# After PUSH AX



# After PUSH BX



# POP and POPF

POP destination

destination = 16 bit register (except IP) or memory word

Execution of POP:

- The content of SS:SP (TOS) is moved to the destination.
- SP is increased by 2.

POPF

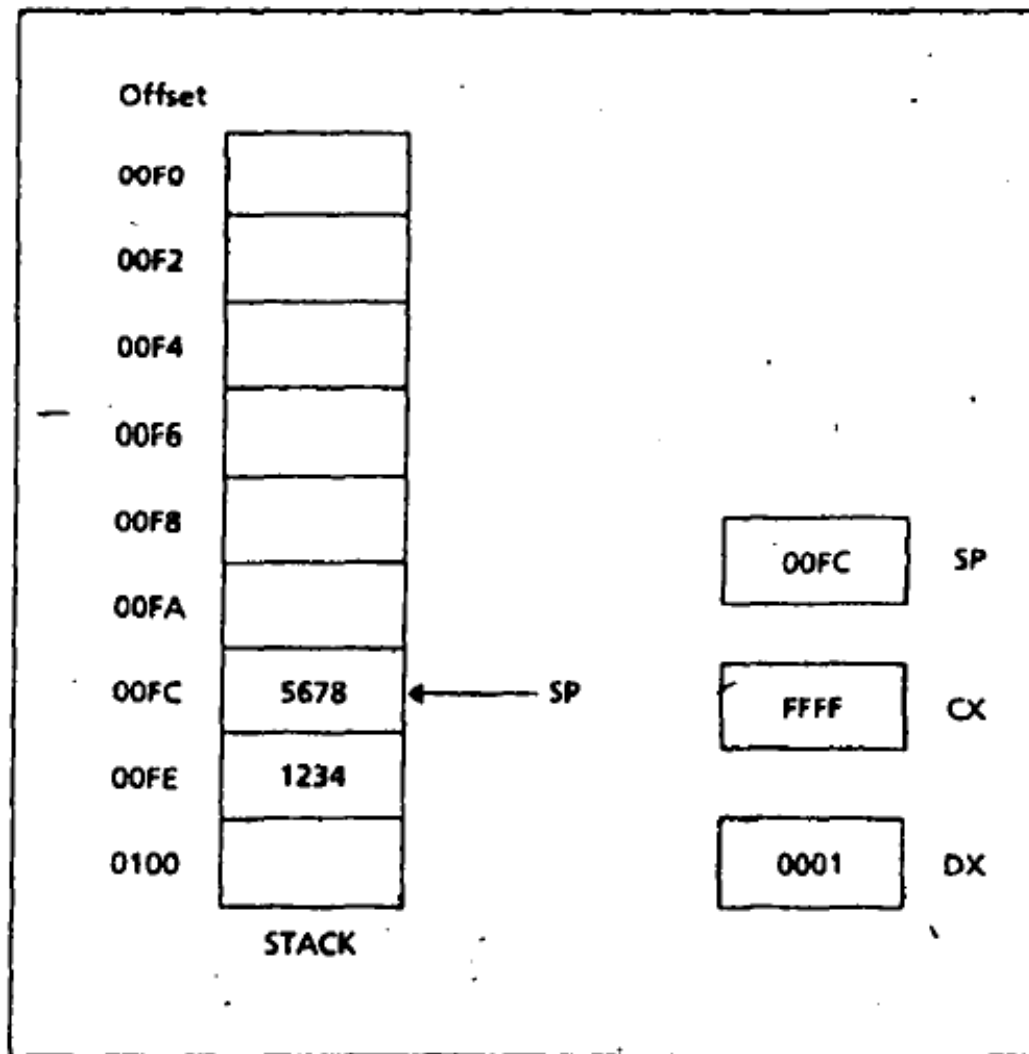
This instruction pops the TOS into the FLAGS register.

There is no effect of POP/POPF on the flags.

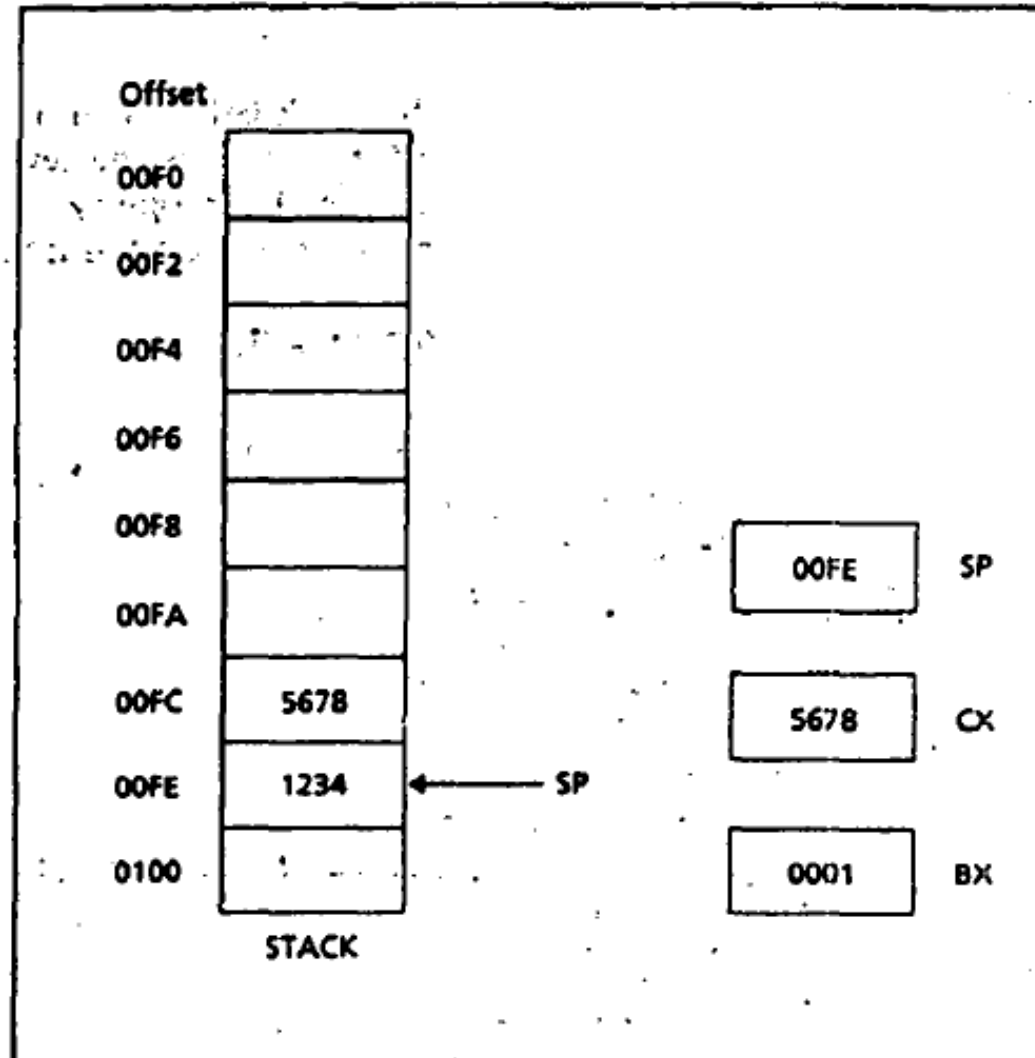
Illegal: POP DL (Byte instruction is illegal)



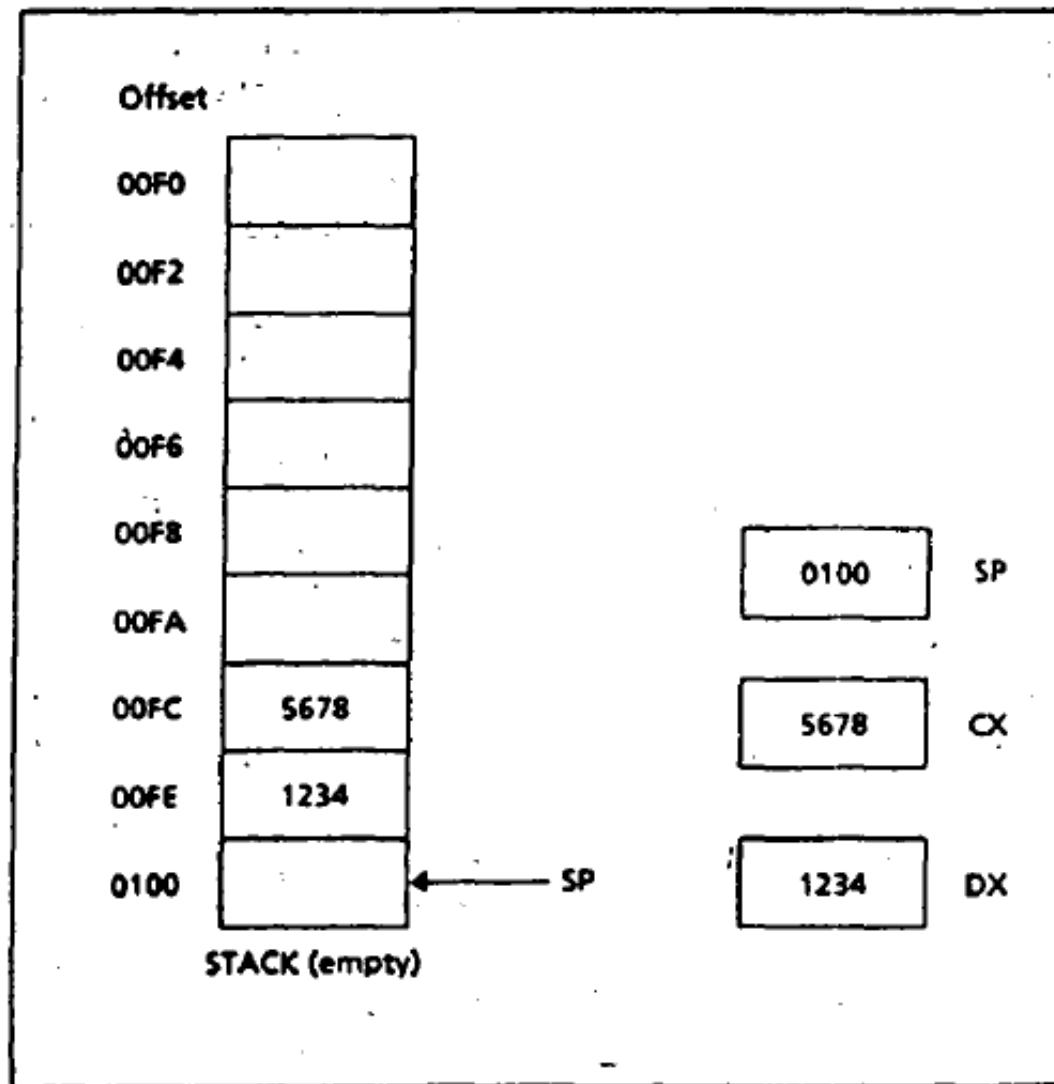
# Before POP



# After POP CX



# After POP DX



# Procedure declaration

```
name PROC type  
    ;body of the procedure
```

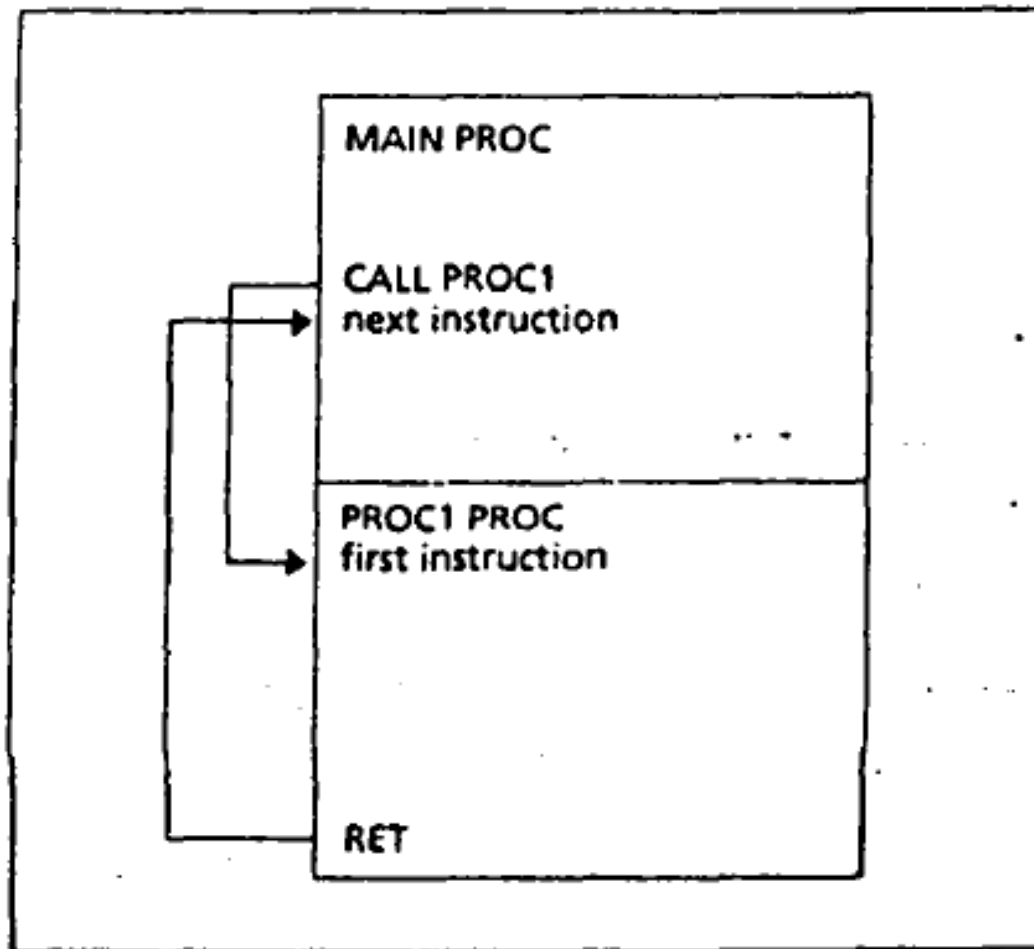
```
RET  
name ENDP
```

name = user defined name of the procedure  
type = (optional) NEAR or FAR

If type is omitted, NEAR is assumed.

NEAR means caller and called procedures are in the same segment. Far means caller and called procedures are in different segment.

# Procedure call and return



# CALL

## **Direct procedure call:**

CALL name

name = name of the procedure

## **Indirect procedure call:**

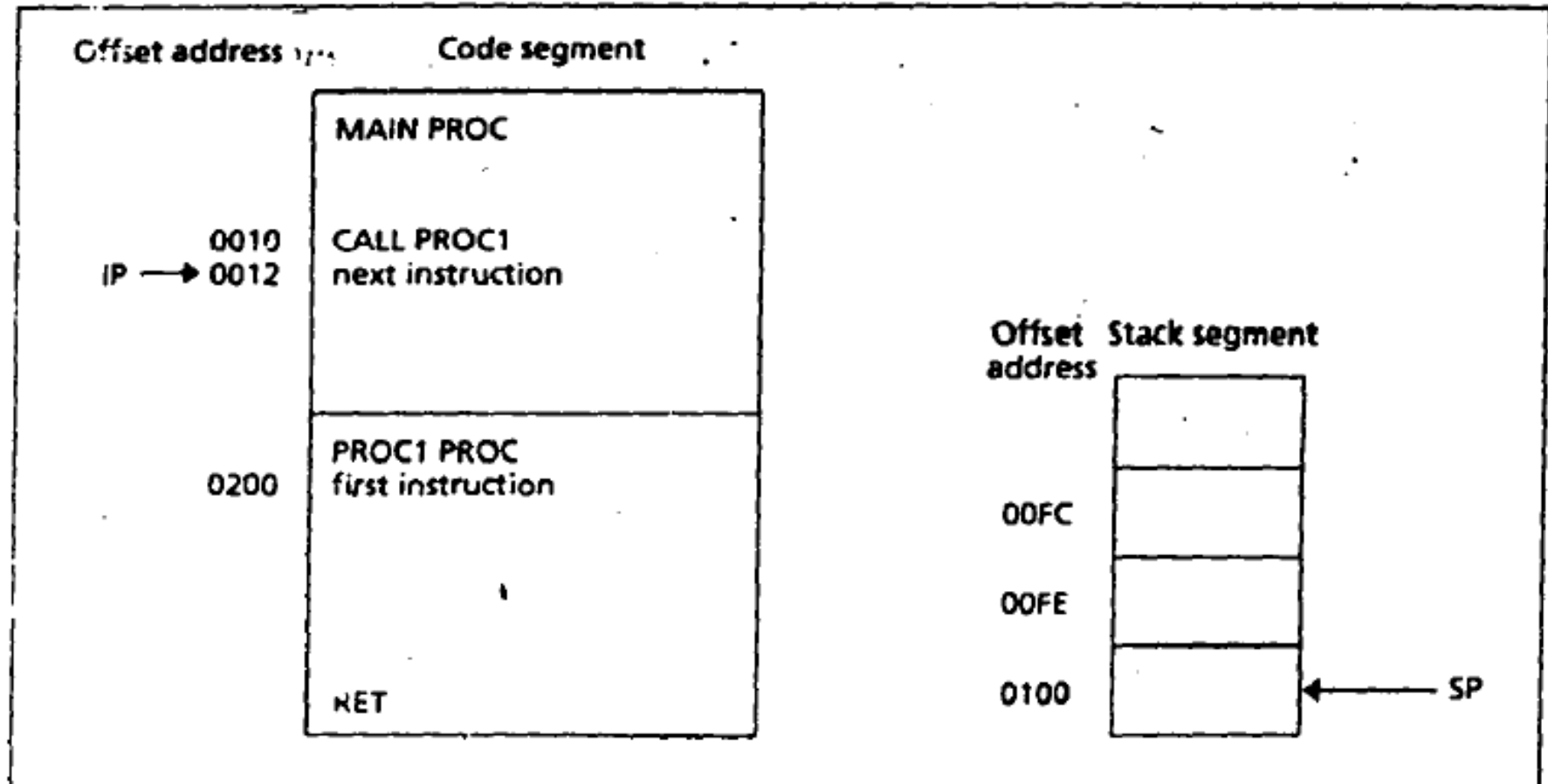
CALL address

address = a register or memory location containing the address of a procedure.

# Execution of CALL

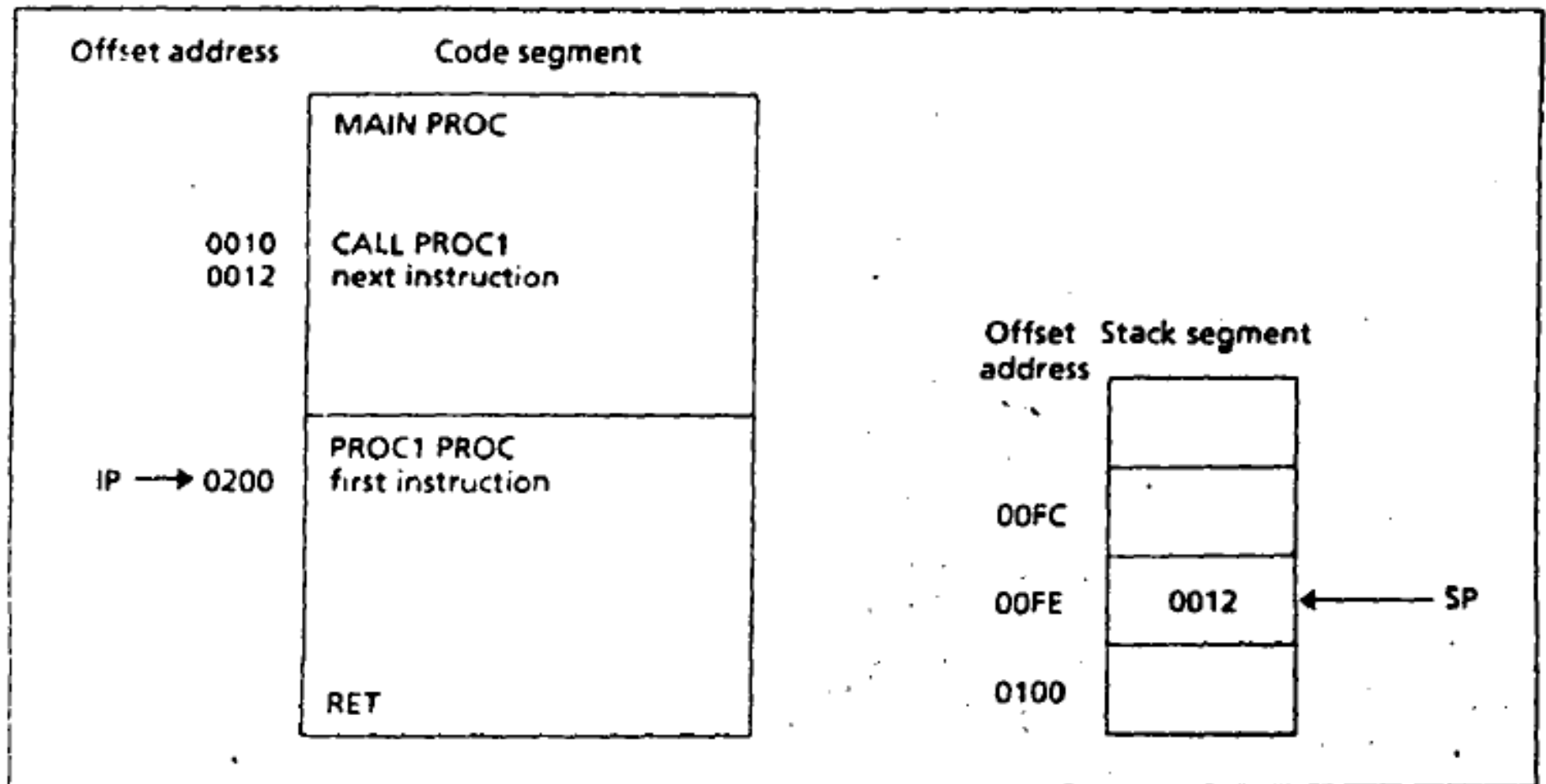
- Return address of the caller is saved on the stack.  
return address = offset of the next instruction after the CALL statement in the caller procedure.
- IP gets the offset address of the first instruction of called procedure.  
CS:IP = segment:offset of the first instruction of the called procedure.  
Control goes to the called procedure.

# Before CALL





# After CALL



# RET

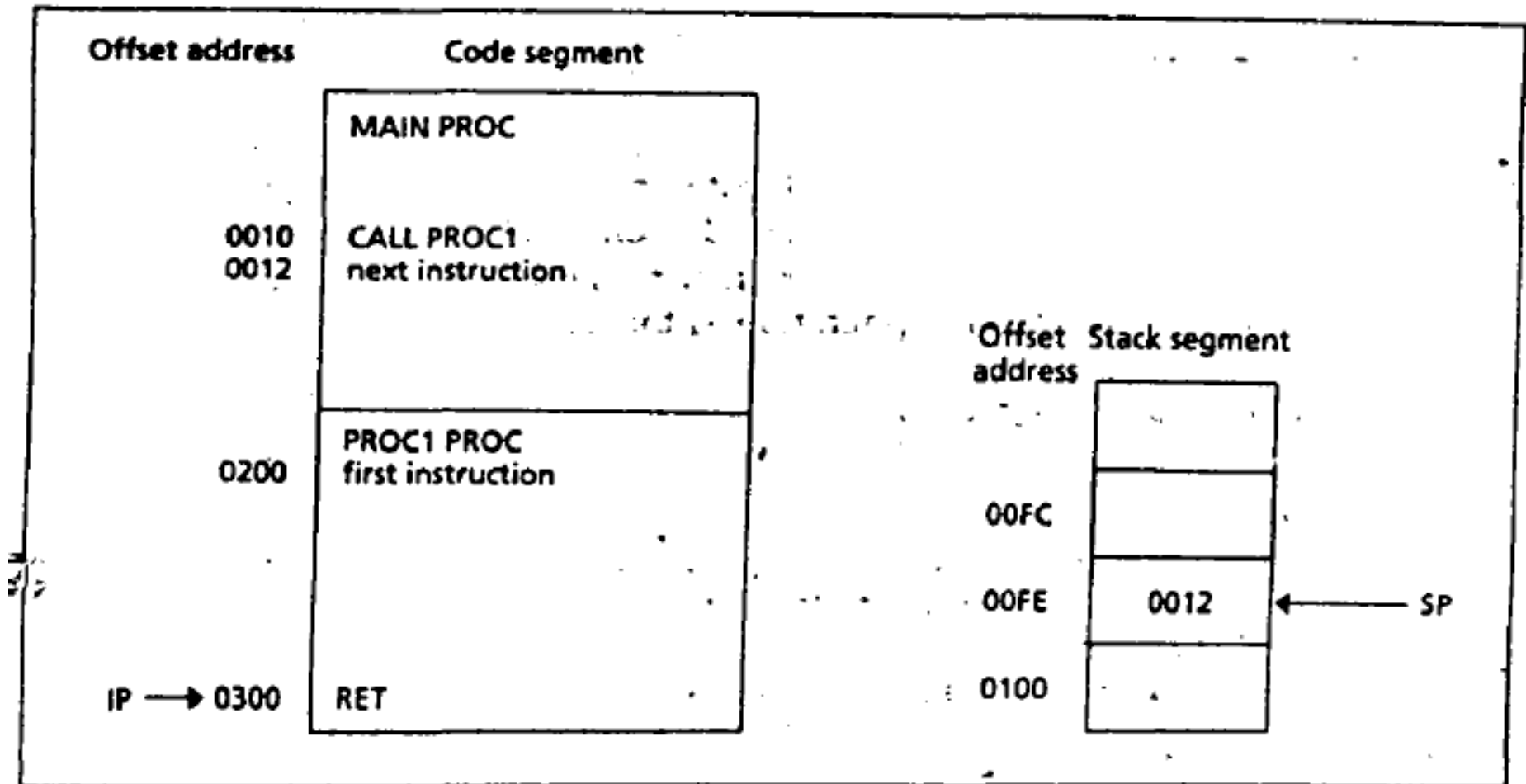
RET pop\_value

pop\_value = (optional)

Execution of RET:

- IP gets the value of TOS. So, CS:IP = segment:offset of the return address.  
Thus, control goes back to the caller program

# Before RET



# After RET

