

# String Algorithms: Introduction

- This chapter focuses on efficient algorithms for string processing.
- Many string problems can be solved in  $O(n^2)$  time, but the challenge is to find algorithms that work in  $O(n)$  or  $O(n \log n)$  time.
- A fundamental string processing problem is the pattern matching problem: finding occurrences of a pattern in a string.
- We'll explore algorithms beyond the brute force approach, aiming for better time complexities.

# String Terminology

- **Zerobased indexing:** Strings are indexed starting from 0 (e.g., `s[0]` is the first character).
- **Alphabet:** The set of characters that can appear in a string (e.g., {A, B, ..., Z}).
- **Substring:** A sequence of consecutive characters within a string (e.g., 'AB' is a substring of 'ABC').
- **Subsequence:** A sequence of characters in a string, not necessarily consecutive, but maintaining their original order (e.g., 'AC' is a subsequence of 'ABC').
- **Prefix:** A substring starting at the beginning of a string (e.g., 'AB' is a prefix of 'ABC').
- **Suffix:** A substring ending at the end of a string (e.g., 'BC' is a suffix of 'ABC').

# Pattern Matching Problem

- **Goal:** Find all occurrences of a pattern (length  $m$ ) within a string (length  $n$ ).
- **Example:** Pattern 'ABC' occurs twice in the string 'ABABCBABC'.
- **Brute Force Approach:** Tests all possible positions for the pattern in the string, taking  $O(nm)$  time.
- **Challenge:** Find algorithms with better time complexities, ideally  $O(n + m)$ .

# Block Decomposition for Efficient Range Queries

- **Problem:** Given an array, efficiently calculate the sum of elements within a specified range.
- **Block Decomposition Approach:**
- Divide the array into blocks of size  $\sqrt{n}$ .
- Store the sum of elements in each block.
- To calculate the sum of a range, divide it into single elements and blocks, allowing for  $O(\sqrt{n})$  time complexity.

# Case Processing: Finding Closest Cells with Same Letter

- **Problem:** Given a grid of cells with letters, find two cells with the same letter that are closest to each other.
- **Approach:**
  - Consider each letter separately.
  - For each letter, find the minimum distance between two cells containing that letter.
- **Algorithms:**
  - **Algorithm 1 (Brute Force):** Check all pairs of cells with the same letter, taking  $O(k^2)$  time ( $k$  is the number of cells with that letter).
  - **Algorithm 2 (BreadthFirst Search):** Simultaneously start a breadthfirst search from each cell with the same letter, finding the minimum distance.