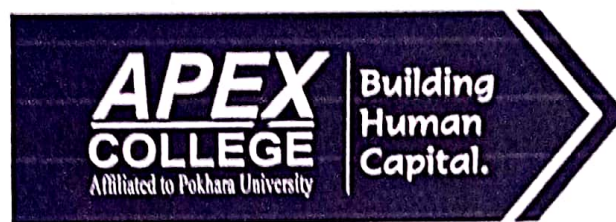


# Apex College

## BCIS Program

Affiliated to Pokhara University



### Data Structure & Algorithms Lab Report

*3. Implementation of Circular Queue*

Date: \_ \_ \_ \_ \_

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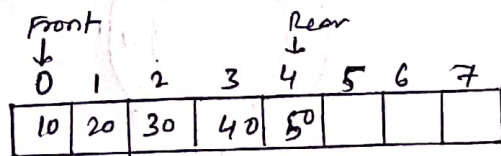
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## # Lab 3 Objectives

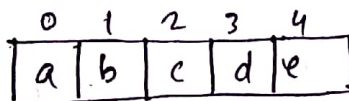
- To understand circular queue and implementation of its operations.
- To create and execute algorithms of enqueue and dequeue operations of circular queue.

## # Introduction : Circular Queue

Circular queue is a linear data structure in which the operations are performed based on FIFO sequence and last position is connected back to the first position of queue to make circle. It is also called 'Ring Buffer'.



In above normal queue, we can insert element until queue becomes full, but we cannot insert the next element even if there is a space in front of queue



$F = 0, 1, 2, 3, 4$

$R = -1, 0, 1, 2, 3$

Now,  $F > R$

then, if we try to insert, it shows full state

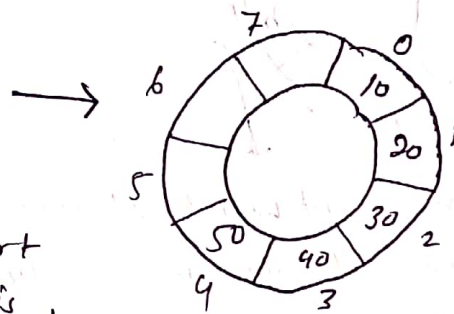
and

if we try to ~~insert~~ delete, it shows empty

that is called. Absurd Condition,

So, to solve this problem, we use Circular queue

where we have to sacrifice either one element or one space in queue



### Applications of Circular Queue

- CPU scheduling
- Memory Management
- Traffic Management

In Circular Queue, we use Modulo Operation to find the value of Rear and Front pointer,

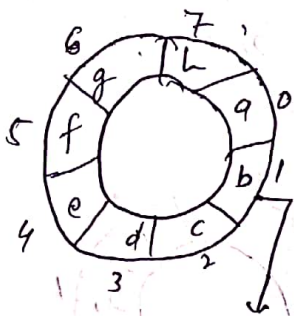
$$\text{Front} = (\text{Front} + 1) \% \text{MAX}$$

$$\text{Rear} = (\text{Rear} + 1) \% \text{MAX}$$

We can define Full and Empty condition after sacrifice of either a element or a space in a circular queue.

\* One Element Sacrifices Method

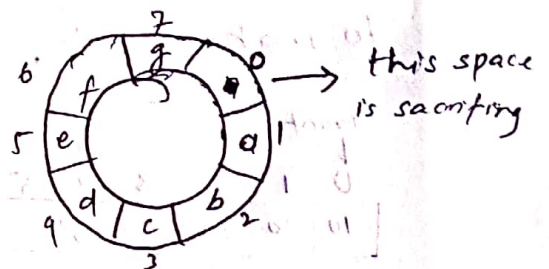
$$R = F$$



This element won't get service, because it shows 'queue is full' but 'queue is empty' at some time

\* One Space Sacrifices Method

$$R = 0; F = 1$$



Now, we can define

$(R + 1) \% \text{MAX} == F$  (i.e.  $R = 1, F = 1$ ) as a empty queue, and.

$(R + 2) \% \text{MAX} == F$  (i.e.  $R = 0, F = 1$ ) as a full queue.

$\therefore$  sacrificing one space is better than sacrificing a resource from queue, Hence, we use one space sacrifice method.

# A program to demonstrate operations in Circular Queue.

```
#include <stdio.h>
```

```
#include <stdlib.h>
```

```
#define MAX 10 or 8
```

```
struct queue {
```

```
    int items[MAX];
```

```
    int rear, front;
```

```
};  
typedef struct queue queue;
```



```

void enqueue(queue *q, int item) {
    if((q->rear + 1) % MAX == q->front)
        printf("Circular queue is full. \n");
    else {
        q->rear = (q->rear + 1) % MAX;
        q->itemx[q->rear] = item;
    }
}

```

```

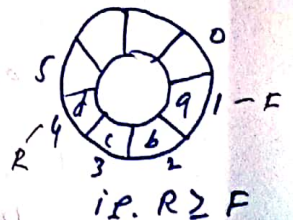
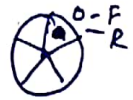
int dequeue(queue *q) {
    int item = -1;
    if((q->rear + 1) % MAX == q->front)
        printf("Circular queue is empty. \n");
    else {
        item = q->itemx[q->front];
        q->front = (q->front + 1) % MAX;
    }
    return item;
}

```

```

void display(queue *q) {
    int i;
    printf("Elements in Circular Queue: ");
    if((q->rear + 1) % MAX == q->front)
        printf(":- is empty. \n");
    else if(q->front <= q->rear) {
        for(i = q->front; i <= q->rear; i++)
            printf("%d \n", q->itemx[i]);
    }
    else {
        for(i = q->front; i < MAX; i++) // 1
            printf("%d \n", q->itemx[i]);
    }
}

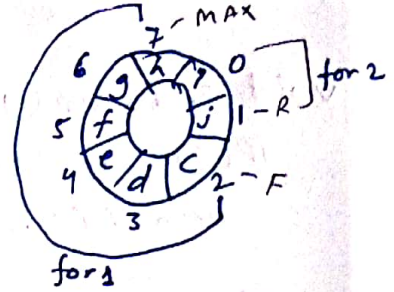
```



```

for (i = 0; i <= q->rear; i++) //2
    printf("%d\n", q->items[i]);
}

```



```

void main () {

```

```

    queue q;
    int ch, item;

```

```

    q.front = 0;

```

```

    q.rear = -1;

```

```

    while (1) {

```

```

        printf("Enter elements in Circular Queue:\n");

```

```

        printf("1. Enqueue. \t 2. Dequeue. \t 3. Display \t 4. Exit\n");

```

```

        scanf("%d", &ch);

```

```

        switch (ch) {

```

```

            case 1:

```

```

                printf("Enter data:");

```

```

                scanf("%d", &item);

```

```

                enqueue(&q, item);

```

```

                break;

```

```

            case 2:

```

```

                item = dequeue(&q);

```

```

                printf("%d is dequeued.\n", item);

```

```

                break;

```

```

            case 3:

```

```

                display(&q);

```

```

                break;

```

```

            case 4:

```

```

                exit(0);

```

```

            default:

```

```

                printf("Invalid choice.\n");

```

```

        }
    }
}

```



## # Activities

### ① Dequeue

- dequeue operation in circular queue is similar to normal queue's dequeue operation but it used modulo operations to check conditions

• Check whether the queue is empty or not

i.e. if  $(R+1) \% MAX == F$

printf("Queue is empty.") & exit

• Otherwise, queue is not ~~full~~ empty

- set, element =  $q \rightarrow \text{items}[q \rightarrow F]$ ;

- set, value of front

i.e.  $q \rightarrow \text{front} = (q \rightarrow \text{front} + 1) \% MAX$ ;

↳ return item or element,

### ② Enqueue

- similarly, enqueue also used modulo operations

• check whether queue is full or not

i.e. if  $((q \rightarrow \text{rear} + 1) \% MAX == F)$

printf("Queue is full.in") & exit

• otherwise, queue is not full

- set, value of rear

$\text{Rear} = (\text{rear} + 1) \% MAX$

- set value at rear pos

$\text{items}[\text{rear}] = \text{item}$

### ③ Display

- In display function, we have tried a condition which can traverse and ~~check~~ display value till MAX pos of rear but it generates error if rear pointer goes beyond MAX pos in circular queue.

// Function with error generating condition

:

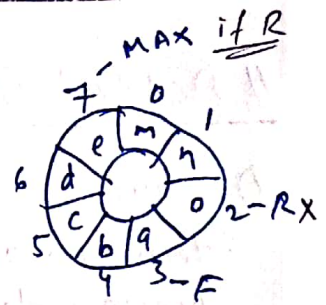
else {

for (i = q -> rear; i >= q -> front; i--)

printf("%d\n", q -> item[i]);

}

;



→ This condition can traverse only items ind between F and MAX if R = MAX.  
Now, R is less than MAX, so, that it returns error.

To overcome this issue, we have used different & multiple method, which is done in program section.

④ Exit

- exit(0);

## # Conclusion

- We have learned about circular queue, difference between normal & circular queue with their own advantages and disadvantages and their implementations.