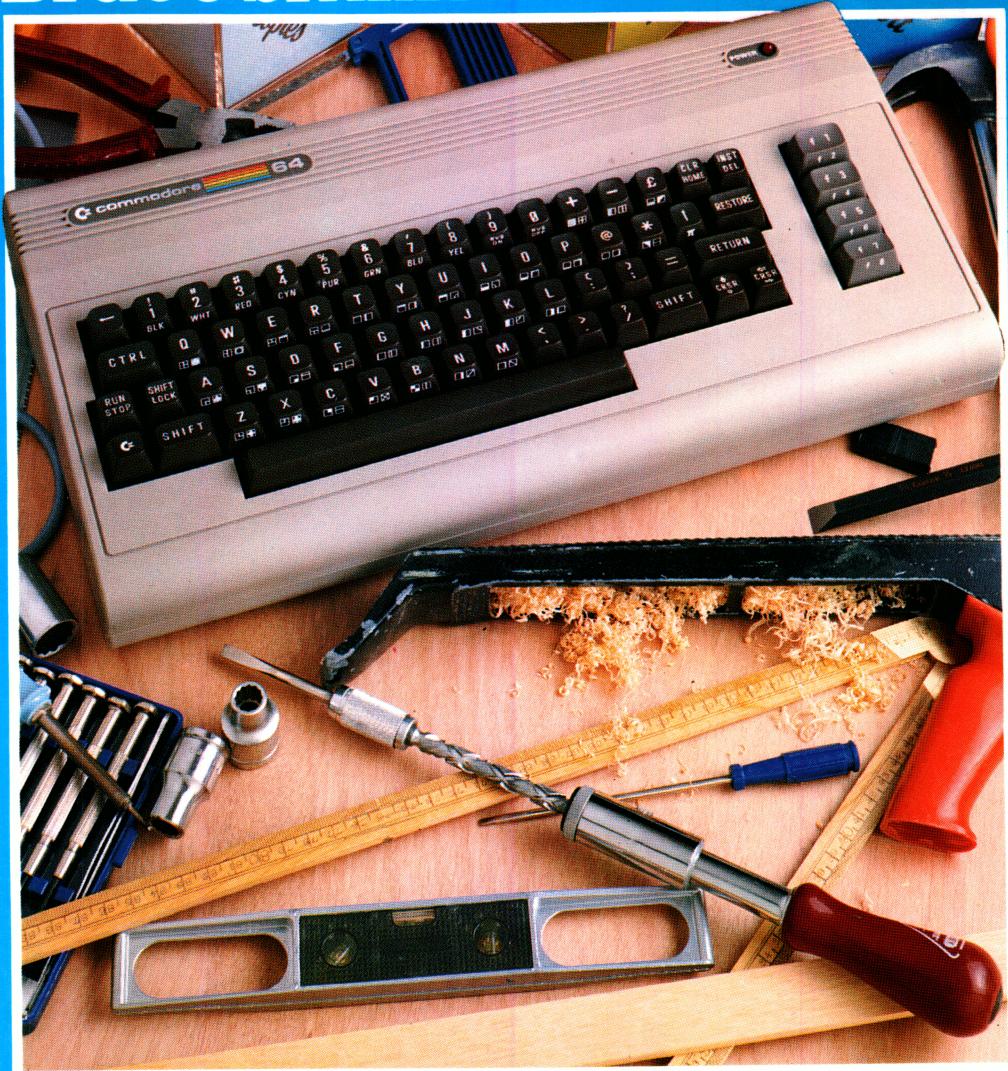


Commodore 64 Assembler Workshop

SHIVA'S
friendly
micro
series

Bruce Smith



Commodore 64 Assembler Workshop

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Bruce Smith



Shiva Publishing Limited

SHIVA PUBLISHING LIMITED
64 Welsh Row, Nantwich, Cheshire CW5 5ES, England

© Bruce Smith, 1984

ISBN 1 85014 004 9

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An interface was used to produce this book from a microcomputer disc, which ensures direct reproduction of error-free program listings.

Typeset and printed by Devon Print Group, Exeter

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Introduction

The *Commodore 64 Assembler Workshop* is aimed at those of you who have been delving into the delights of programming at machine code level. It is a natural progression from *Commodore 64 Assembly Language*, but will be invaluable even if you learned assembler and machine code using any of the other relevant books available. It provides a bench full of useful assembly language routines and utilities programs and examines the techniques involved.

Extensive use of vectored addresses is made throughout the Commodore's operation, allowing modifications to be made to the manner in which the micro operates. Chapter 2 demonstrates how the CHRGET subroutine can be used to allow new RAM-based commands to be added to the already extensive facilities provided by the machine. A short 'wedge' interpreter is provided and the techniques for adding your own commands examined, and to get you going, three commands come supplied with the wedge interpreter: @CLS, @UP and @LOW.

Conversion between ASCII based numerical character strings and their two-byte binary equivalents and vice versa is not straightforward. Such conversions are fully described in Chapters 3 and 4, and working routines are listed.

Any program which handles strings of data must be able to manipulate the strings, whether it is an adventure game or the latest stock control reports. Routines for comparing, copying, deleting and inserting strings are included, and Chapter 6 goes on to show the various ways in which text can be printed to the screen.

Sorting data lists into order is a task which it is often necessary to perform within a program, so the technique of bubble sorting is investigated.

Many other processors provide operations that would be useful to have available when using the 6510. A software stack implementation similar to that found on the 6809 precessor is produced in Chapter 8, allowing up to eight selected registers to be pushed on to a memory-based stack.

Routines to move, fill and produce a hex and ASCII dump of memory are then examined and the final chapter provides a few hi-resolution graphics utilities to speed you along the way.

Many of the chapters suggest projects for you to undertake at your leisure, while every program has a detailed line-by-line description of its operation. Program listings are provided using BASIC loaders so that they can be used directly as they are. Included in each line is a REM statement giving the mnemonic representation of the instruction should you be using an assembler.

In fact, all the tools for using the *Assembler Workshop* are supplied—assuming of course you have the workbench!

Highbury, November 1984

Bruce Smith

1 Opening the Tool Box

The routines included in this book are designed to make your life that much easier when writing machine code. Quite often, after mastering the delights of the Commodore 64's microprocessor, programmers become frustrated because the techniques involved in, say, converting between ASCII characters and their equivalent binary values are not known. Nor are they readily available in a published form, so the painful process of sitting down armed with pencil and paper and working out the conversion through trial and error begins.

This is just one example of the type of assembler program you will find within these pages. Wherever possible, they are supplied in a form that will make them relocatable, the only addresses requiring alteration being those specified by JSR or JMP.

Each listing is in the form of a BASIC loader program, using a loop to READ and POKE decimal machine code data into memory. This will allow those of you who have not yet splashed out your hard earned cash on a suitable assembler program to get underway. For those lucky ones among you who do have an assembler, each data statement has been followed by a REM line containing the standard mnemonic representation of the instruction (see Appendix 1 for a summary). This can be entered directly and assembled as required.

Although the programs are typeset they have been spooled direct as ASCII files and loaded into my word processor so all should run as they are.

BASIC is used freely to demonstrate the machine code's operation—rather than repeating sections of assembler code, BASIC is often used to shorten the overall listing, and it is left to you to add further sections of assembler from other programs within the book or from your own resources. For example, many programs require you to input a decimal address. In the demonstrations, this is indicated by means of a one-line INPUT statement. In Chapter 3, however, there is a routine for inputting a string of five ASCII

decimal characters and converting it into a two-byte binary number. This can be inserted into the assembler text of the program, to go some way to making it a full machine code program available for use as a completely self-contained section of machine code.

WRITING MACHINE CODE

You have an idea that you wish to convert into machine code—so what's the best way to go about it? Firstly, make some brief notes about its operation. Will it use the screen? If so, what mode? Will it require the user to input values from the keyboard? If so, what keys do you use? What will the screen presentation look like? Will you want to use sound?...and so on. Once you have decided on the effects you want, put them down in flowchart form. This need not be the normal flowchart convention of boxes and diamonds—I find it just as easy to write each operation I want the program to perform in a list and then join the flow of these up afterwards.

Quite often, the next step is to write the program in BASIC! This may sound crazy, but it allows you to examine various aspects of the program's operation in more detail. An obvious example of this is obtaining the correct screen layout—you might find after running the routine that the layout does not look particularly good. Finding this out at an early stage will save you a lot of time later, avoiding the need to rewrite the screen layout portion of your machine code—rewriting BASIC is much easier! If you write the BASIC tester as a series of subroutines, it will greatly simplify the process of conversion to machine code. Consider the main loop of such a BASIC tester, which takes the form:

```
10 GOSUB 200 : REM SET UP VARIABLES
20 GOSUB 300 : REM SET UP SCREEN
30             REM LOOP
40 GOSUB 400 : REM INPUT VALUES
50 GOSUB 500 : REM CONVERT AS NEEDED
60 GOSUB 600 : REM DISPLAY VALUES
70 GOSUB 700 : REM DO UPDATE
80 IF TEST=NOTDONE THEN GOTO 30
90 END
```

Each module can be taken in turn, converted into assembler and tested. Once performing correctly the next procedure can be examined. Debugging is made easier because the results of each module are known having used the BASIC tester. The final main loop of the assembler might then look something like this:

```
JSR $C200 : REM SET UP VARIABLES
JSR $C300 : REM SET UP SCREEN
    REM LOOP
JSR $C400 : REM INPUT VALUES
JSR $C500 : REM CONVERT AS NEEDED
JSR $C600 : REM DISPLAY VALUES
JSR $C700 : REM DO UPDATE
BNE LOOP
```

You might be surprised to learn that this technique of testing machine code programs by first using BASIC is employed by many software houses the world over.

DEBUGGING

A word or two about debugging machine code programs that will not perform as you had hoped: if this happens to you, before pulling your hair out and throwing the latest copy of *Machine Code Nuclear Astrophysics Weekly* in the rubbish bin, a check of the following points may reveal the bug!

1. If you are using a commercial assembler, check that your labels have all been declared and correctly assigned. If you are assembling 'by hand', double-check all your branch displacements and JMP and JSR destination addresses. You can normally ascertain exactly where the problem is by examining how much of the program works before the error occurs, rather than checking it all.
2. If your program uses immediate addressing, ensure you have prefixed the mnemonic with a hash (#) to inform the assembler or, if compiling by hand, check that you have used the correct opcode. It is all too easy to assemble the coding for LDA \$41 when you really want the coding for LDA #\$41.
3. Check that you have set or cleared the Carry flag before subtraction or addition.
4. My favourite now—ensure that you save the result of a subtraction or an addition. The sequence:

```
CLC
LDA $FB
ADC #1
BCC OVER
INC $FC
OVER
RTS
```

is not much good if you don't save the result of the addition with:

```
STA $FB
```

before the RTS!

5. Does the screen clear to the READY prompt whenever you perform a SYS call, seemingly without executing any of the machine code? The bug that often causes this is due to an extra comma being inserted into a series of DATA statements. For example the DATA line:

```
DATA 169,0,,162,255
```

with an extra comma between the 0 and 162, would assemble the following:

```
LDA #$00  
BRK  
LDX #$FF
```

as the machine has interpreted ‘,,’ as ‘,0,’ and assembled the command which has zero as its opcode—BRK!

6. Does the program ‘hang up’ every time you run it, when you are quite certain that the data statements are correct? This is often caused by a full stop instead of a comma being used between DATA statements, e.g.

```
DATA 169,6,162.5,96
```

Here, if a full stop has been used instead of a comma between the 162 and the 5, the READ command interprets this as a single number, 162.5, rounds it down to 162, and assembles this ignoring the 5 and using the 96 (RTS) as the operand, as follows:

```
LDA #$06  
LDX #$60  
XXX
```

When executed, the garbage after the last executable instruction results in the system hanging up. This error should not occur if you calculate your loop count correctly, so always double-check this value before running your program.

If none of these errors is the cause of the problem, then I'm afraid you must put your thinking cap on. Well-commented assembler will make debugging very much easier.

2 Commodore Command

One of the disadvantages of using random access memory-based machine code routines as utilities within a BASIC program is that it is left to you, the programmer, to remember just where they are stored, and to use the appropriate SYS call to implement them. This doesn't usually pose any problems if only one or two machine code utilities are present; the problems occur when several are being used. Normally you would need to keep a written list of these next to you, looking up the address of each routine as you need it. Great care must be taken to ensure that the SYS call is made to the correct address, as a mis-typed or wrongly called address can send the machine into an infinite internal loop, for which the only cure is a hard reset, which would destroy all your hard work.

The program offered here provides a useful and exciting solution to the problem, enabling you to add new commands to your Commodore 64's vocabulary. Each of your routines can be given a command name, and the machine code comprising any command will be executed by simply entering its command name. The routine is written so that these new commands can be used either directly from the keyboard or from within programs.

The trick in 'teaching' the Commodore 64 new commands is to get the machine to recognize them. If an unrecognized command is entered at the keyboard, the almost immediate response from the 64 is '?SYNTAX ERROR'. If you have any expansion cartridges you'll know that it is possible to expand the command set, and the *Programmer's Reference Guide* gives a few hints on how to do this, on pages 307 and 308—the method pursued here is by resetting the system CHRGET subroutine.

CHRGET

The CHRGET routine is, in fact, a subroutine which is called by the main BASIC Interpreter. You can think of it as a loop of code, protruding from the machine, into which we can wedge our own

bits of code, thereby allowing fundamental changes to be made to the manner in which the Commodore operates. Let's have a look at how the normal CHRGET subroutine (which is located in zero page from \$73) operates:

Table 2.1

Address	Machine code	Assembler
\$0073	E6 7A	INC \$7A
\$0075	D0 02	BNE \$0079
\$0077	E6 7B	INC \$7B
\$0079	AD xx xx	LDA \$xxxx
\$007C	C9 3A	CMP #\$3A
\$007E	B0 0A	BCS \$008A
\$0080	C9 20	CMP #\$20
\$0082	F0 EF	BEQ \$0073
\$0084	38	SEC
\$0085	E9 30	SBC #\$30
\$0087	38	SEC
\$0088	E9 D0	SBC #\$D0
\$008A	60	RTS

The subroutine begins by incrementing the byte located at \$7A. This address forms a vector which holds the address of the interpreter within the BASIC program that is currently being run. If there is no carry over into the high byte, which must therefore itself be incremented, a branch occurs to location \$0079. You will notice that the bytes which have just been incremented lie within the subroutine itself. These are signified in the above listing by 'xx xx', because they are being updated continually by the routine. The reason for this should be fairly self-evident: looking at the opcode, we can see that it is LDA, therefore each byte is, in turn, being extracted from the program.

The next two bytes at \$007C perform a compare, CMP #\$3A. The operand here, \$3A, is the ASCII code for a colon, so CHRGET is checking for a command delimiter. The BCS \$008A will occur if the accumulator contents are greater than \$3A, effectively returning control back to within the BASIC Interpreter ROM. The next line, CMP #\$20, checks whether a space has been encountered within the program. If it has, the branch is executed back to \$0073 and the code rerun.

The rest of the coding is checking that the byte is a legitimate

one—it should be an ASCII character code in the range \$30 to \$39, that is, a numeric code. If it is, the coding will return to the main interpreter with the Carry flag clear. If the accumulator contains less than \$30 (it could, of course, have ASCII \$20 in it, as we have already checked for this) then the Carry flag is set.

It is important to understand what is happening here, as we will need to overwrite part of this code to point it in the direction of our own ‘wedge’ interpreter. This has to perform the ‘deleted’ tasks before returning to the main interpreter to ensure the smooth and correct running of the Commodore 64.

THE WEDGE OPERATING SYSTEM

To distinguish the Wedge Operating System (WOS) commands from normal commands (and illegal ones!), we must prefix them with a special character—one which is not used by the Commodore 64. The *Programmer’s Reference Guide* suggests the use of the the ‘@’ sign, so that’s what we will use.

Program 1a lists the coding for the WOS. I have chosen to place it well out of the way, in the free RAM area from 49666 (3C202) onwards. As we shall see the memory below (bis to 49152 (\$C000)) is also used by the WOS.

Program 1a

```
10 REM *** WEDGE OPERATING SYSTEM - WOS ***
20 REM *** WOS INTERPRETER FOR COMMODORE 64 ***
30 :
40 CODE=49666
50 FOR LOOP=0 TO 188
60 READ BYTE
70 POKE CODE+LOOP,BYTE
80 NEXT LOOP
90 :
100 REM ** M/C DATA **
110 DATA 169,0      : REM LDA #$00
120 DATA 160,192    : REM LDY #$C0
130 DATA 32,30,171  : REM JSR $AB1E
140 DATA 169,76     : REM LDA #$4C
150 DATA 133,124    : REM STA $7C
160 DATA 169,24     : REM LDA #$18
170 DATA 133,125    : REM STA $7D
180 DATA 169,194    : REM LDA #$C2
190 DATA 133,126    : REM STA $7E
```

200	DATA	108,2,3	:	REM	JMP (\$0302)
205	::			REM	WOS STARTS HERE
210	DATA	201,64	:	REM	CMP #\$40
220	DATA	208,68	:	REM	BNE \$44
230	DATA	165,157	:	REM	LDA \$9D
240	DATA	240,40	:	REM	BEQ \$28
250	DATA	173,0,2	:	REM	LDA \$0200
260	DATA	201,64	:	REM	CMP #\$40
270	DATA	208,28	:	REM	BNE \$1C
280	DATA	32,114,194	:	REM	JSR \$C272
290	DATA	160,0	:	REM	LDY #\$00
300	DATA	177,122	:	REM	LDA (\$7A),Y
310	DATA	201,32	:	REM	CMP #\$20
320	DATA	240,9	:	REM	BEQ \$09
330	DATA	230,122	:	REM	INC \$7A
340	DATA	208,246	:	REM	BNE \$F6
350	DATA	230,123	:	REM	INC \$7B
360	DATA	56	:	REM	SEC
370	DATA	176,241	:	REM	BCS \$F1
380	DATA	32,116,164	:	REM	JSR \$A474
390	DATA	169,0	:	REM	LDA #\$00
400	DATA	56	:	REM	SEC
410	DATA	176,29	:	REM	BCS \$1D
420	DATA	169,64	:	REM	LDA #\$40
430	DATA	56	:	REM	SEC
440	DATA	176,24	:	REM	BCS \$18
445	::			REM	PROGRAM-MODE
450	DATA	32,114,194	:	REM	JSR \$C272
460	DATA	160,0	:	REM	LDY #\$00
470	DATA	177,122	:	REM	LDA (\$7A),Y
480	DATA	201,0	:	REM	CMP #\$00
490	DATA	240,13	:	REM	BEQ \$0D
500	DATA	201,58	:	REM	CMP #\$3A
510	DATA	240,9	:	REM	BEQ \$09
520	DATA	230,122	:	REM	INC \$7A
530	DATA	208,242	:	REM	BNE \$F2
540	DATA	230,123	:	REM	INC \$7B
550	DATA	56	:	REM	SEC
560	DATA	176,237	:	REM	BCS \$ED

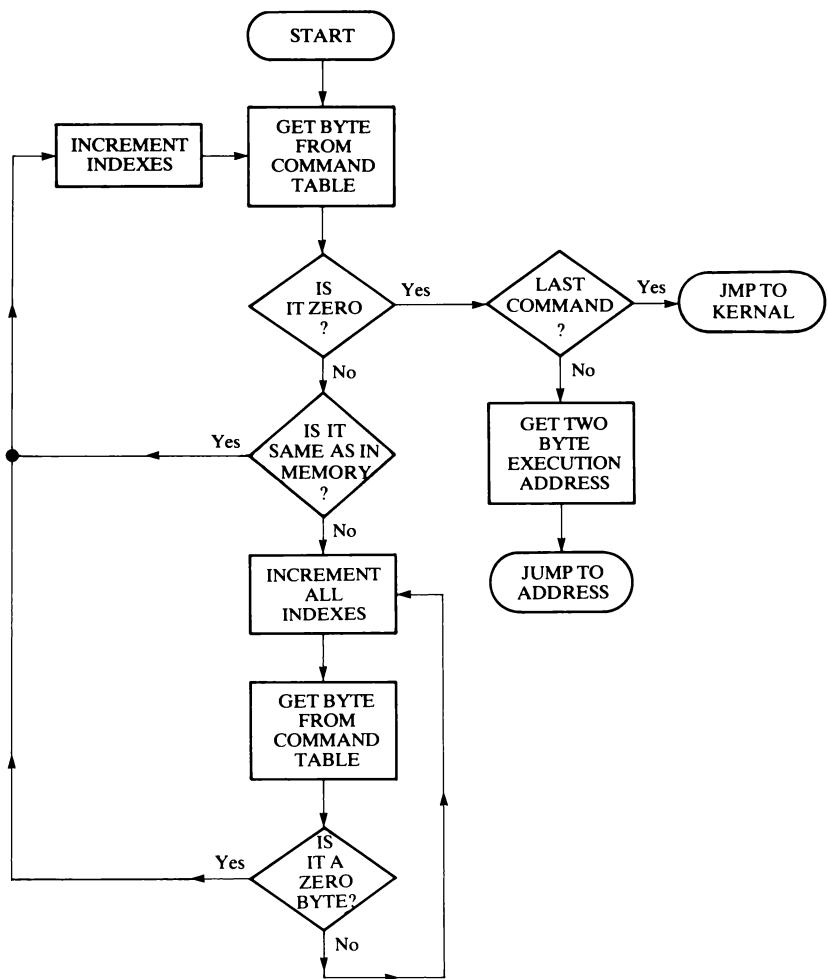


Figure 2.1 The wedge operating system flowchart

570 DATA 201,58	:	REM	CMP #\$3A
580 DATA 176,10	:	REM	BCS \$0A
590 DATA 201,32	:	REM	CMP #\$20
600 DATA 240,7	:	REM	BEQ \$07
610 DATA 56	:	REM	SEC
620 DATA 233,48	:	REM	SBC #\$30
630 DATA 56	:	REM	SEC
640 DATA 233,208	:	REM	SBC #\$D0
650 DATA 96	:	REM	RTS
660 DATA 76,115,0	:	REM	JMP \$0073
665 ::			REM FIND-EXECUTE

670 DATA 169,0	:	REM LDA #\$00
680 DATA 133,127	:	REM STA \$7F
690 DATA 169,193	:	REM LDA #\$C1
700 DATA 133,128	:	REM STA \$80
710 DATA 230,122	:	REM INC \$7A
720 DATA 208,2	:	REM BNE \$02
730 DATA 230,133	:	REM INC \$7B
740 DATA 160,0	:	REM LDY #\$00
750 DATA 162,0	:	REM LDX #\$00
760 DATA 177,127	:	REM LDA (\$7F),Y
770 DATA 240,36	:	REM BEQ \$24
780 DATA 209,122	:	REM CMP (\$7A),Y
790 DATA 208,4	:	REM BNE \$04
800 DATA 200	:	REM INY
810 DATA 56	:	REM SEC
820 DATA 176,244	:	REM BCS \$F4
830 DATA 177,127	:	REM LDA (\$7F),Y
840 DATA 240,4	:	REM BEQ \$04
850 DATA 200	:	REM INY
860 DATA 56	:	REM SEC
870 DATA 176,248	:	REM BCS \$F8
880 DATA 200	:	REM INY
890 DATA 152	:	REM TYA
900 DATA 24	:	REM CLC
910 DATA 101,127	:	REM ADC \$7F
920 DATA 133,127	:	REM STA \$7F
930 DATA 169,0	:	REM LDA #\$00
940 DATA 101,128	:	REM ADC \$80
950 DATA 133,128	:	REM STA \$80
960 DATA 160,0	:	REM LDY #\$00
970 DATA 232	:	REM INX
980 DATA 232	:	REM INX
990 DATA 56	:	REM SEC
1000 DATA 176,216	:	REM BCS \$D8
1010 DATA 189,80,192	:	REM LDA \$C050,X
1020 DATA 133,128	:	REM STA \$80
1030 DATA 232	:	REM INX
1040 DATA 189,80,192	:	REM LDA \$C050,X
1050 DATA 133,129	:	REM STA \$81
1060 DATA 108,128,0	:	REM JMP (\$0080)

```

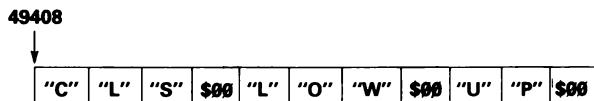
1065 :: REM ILLEGAL
1070 DATA 162,11 : REM LDX #$0B
1080 DATA 108,0,3 : REM JMP ($300)
1090 :
1100 REM ** SET UP COMMAND TABLE **
1110 TABLE=49408
1120 FOR LOOP=0 TO 10
1130 READ BYTE
1140 POKE TABLE+LOOP,BYTE
1150 NEXT LOOP
1160 :
1170 REM ** ASCII COMMAND DATA **
1180 DATA 67,76,83,0 : REM CLS
1190 DATA 76,79,87,0 : REM LOW
1200 DATA 85,80,0 : REM UP

```

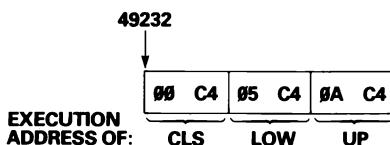
To enable the WOS to identify a wedge command, it needs a complete list to which it can compare the one it is interpreting in the program—this is done with the aid of a command table, which is formed by the program lines from 1100 to 1200. This ASCII table is based at 49408 (\$C100) and, as you can see from the listing, three commands are provided: @CLS, @LOW and @UP. Note that the @ is omitted from the front of each command in the table—it is unnecessary at the comparison stage, as by this time it has already been established that it is a WOS command—and that each command is terminated by a zero. A table listing the execution address of each command must also be constructed, but more of this later.

The main program consists of two parts, an initialization routine and the interpreter proper.

The initialization routine is embodied in lines 110 to 200. Its function is to reset the CHRGET subroutine investigated earlier. Lines 110 to 130 issue a heading on the screen indicating that the



The Command Table



The Address Table

Figure 2.2 The Command and Address Tables.

WOS has been initialized. The subroutine at \$AB1E, called by line 130, prints out an ASCII string located at the address given by the index registers. In this instance it is located at \$C000 (49152), and is assembled into memory by the second part of the listing. Lines 140 to 190 poke three bytes into the CHRGET subroutine which effectively assembles the code:

```
JMP $C218
```

The address \$C218 is the address of the start of the WOS interpreter at line 210. Finally, line 200 does an indirect jump through the IMAN vector at \$0302 to perform a warm BASIC start. The CHRGET subroutine, complete with wedge jump, now looks like this:

Table 2.2

Address	Machine code	Assembler
\$0073	E6 7A	INC \$7A
\$0075	D0 02	BNE \$0079
\$0077	E6 7B	INC \$7B
\$0079	AD xx xx	LDA \$xxxx
\$007C	4C 18 C2	JMP \$C218

When the WOS is entered, the byte in the accumulator is checked to see if it is an @ (line 210), signifying a wedge command. If it is not, then a branch to line 570 is performed. As you can see, the code from line 570 to 650 performs the normal function of the CHRGET routine, with control returning to the BASIC Interpreter.

If the byte is an @, the interpretation continues. The byte at \$9D is located, to detect whether the command is within a program or has been issued in direct mode. A zero indicates that the command has been called from within a program and the branch of line 240 to line 450 is performed. In both instances the interpretation follows similar lines—for descriptive purposes, we will assume program mode and resume the commentary from line 450.

The subroutine at \$C272 is the interpreter proper. Starting at line 665 it locates the command and executes it. The first eight bytes (lines 670 to 700) set up a zero page vector to point to the command table at \$C100. Lines 710 to 730 update the zero page bytes at \$7A and \$7B, which hold the address of the current point within the program. After initializing both index registers, the first

byte within the command table is located (lines 740 to 760), and compared to the byte within the program, immediately after the @ (line 780). If the comparison fails, the branch to line 830 is performed, locating the zero and therefore the next command in the command table. When a comparison is successful (the command is identified) and the terminating zero located by line 770, the branch to line 1010 is performed. Lines 1010 to 1060 locate the execution address of the command from the address table located at \$C050. The X register is used as an offset into this, being incremented by two each time a command table comparison fails (lines 970 and 980). The two address bytes are loaded to form a zero page vector and the machine code is executed via an indirect jump.

On completion of the routine, its terminating RTS returns control to line 460, and the next byte after the command is sought out. When a zero is found, the branch of line 490 is performed and the CHRGET routine is completed, control being returned to the BASIC Interpreter.

THE NEW COMMANDS

Program 1b provides the assembly routines to construct the initialization prompts, the machine code for the new commands and the address table:

Program 1b

```
1210 REM ** TITLE MESSAGE DISPLAYED ON SYS
        49666 **
1220 HEAD=49152
1230 FOR LOOP=0 TO 40
1240 READ BYTE
1250 POKE HEAD+LOOP,BYTE
1260 NEXT LOOP
1270 :
1280 REM ** ASCII CHARACTER DATA **
1290 DATA 147,13,32,32,42,42,32,67,54,52,32
1300 DATA 69,88,84,69,78,68,69,68,32,83,85
1200 DATA 80,69,82,32,66,65,83,73,67,32,86,49
1310 DATA 46,48,32,42,42,13,0
1320 :::
1360 REM ** SET UP M/C FOR COMMANDS **
1370 MC=50176
1380 FOR LOOP=0 TO 14
1390 READ BYTE
```

```

1400 POKE MC+LOOP, BYTE
1410 NEXT LOOP
1420 :
1430 REM ** COMMAND M/C **
1440 :: REM CLS
1450 DATA 169,147 : REM LDA #$93
1460 DATA 76,210,255 : REM JMP $FFD2
1470 :: REM LOW
1480 DATA 169,14 : REM LDA #$0E
1490 DATA 76,210,255 : REM JMF $FFD2
1500 :: REM UP
1510 DATA 169,142 : REM LDA #$8D
1520 DATA 76,210,255 : REM JMP $FFD2
1530 :: 
1540 REM ** SET UP ADDRESS TABLE **
1550 ADDR=49232
1560 FOR LOOP=0 TO 5
1570 READ BYTE
1580 POKE ADDR+LOOP, BYTE
1590 NEXT LOOP
1600 :
1610 REM ** ADDRESS DATA **
1620 DATA 0,196 : REM CLS $C400
1630 DATA 5,196 : REM LOW $C405
1640 DATA 10,196 : REM UP $C40A

```

Each command's machine code is located from 50176 (\$C400). The three new commands and their functions are:

CLS	:	clear screen and home cursor
LOW	:	select lower case character set
UP	:	select upper case character set

Nothing to set the house alight, admittedly, but the techniques involved are more important at present. These are simple to implement and, once understood, enable more useful and complex commands to be added. The code associated with each command is responsible simply for printing its ASCII code. The final section of listing (lines 1540 to 1650) pokes the execution address of each command into memory. The final address points to the code at line 170, and the program jumps to this position if the command is not found within the command table. This code performs an indirect jump to the BASIC Interpreter's error handler.

USING THE WOS

Using the Wedge Operating System is easy: enter the program as shown, run it to assemble the code into memory, and if all goes well, save the program. To initialize the WOS enter:

```
SYS 49666
```

The screen will clear, and the following message be printed across the top of the screen:

```
** C64 EXTENDED SUPER BASIC V1.Ø **
```

The wedge commands are now available for immediate use. Remember that pressing RUN/STOP and RESTORE together will reset the CHRGET routine to its default value making the WOS invisible. To relink it, simply execute the SYS 49666 call again.

Line-by-line

A line-by-line description of the WOS now follows, to enable you to examine its operation in more detail:

```
line 11Ø : load accumulator with low byte message address
line 12Ø : load accumulator with high byte message address
line 13Ø : print start up message
line 14Ø : reset CHRGET subroutine
line 2ØØ : do a BASIC warm start
line 2Ø5 : main entry for WOS
line 21Ø : is it an '@' and therefore a WOS command?
line 22Ø : no, so branch to line 57Ø to update
line 23Ø : yes, check for direct or program mode
line 24Ø : if zero, then WOS command is within program,
           so branch to line 45Ø
line 25Ø : else direct mode so get byte from buffer
line 26Ø : recheck that it is a WOS command
line 27Ø : if error, branch to line 41Ø
line 28Ø : find and execute the command else issue
           appropriate error message
line 29Ø : initialize index
line 3ØØ : get byte from buffer
line 31Ø : is it a space?
line 32Ø : yes, so branch to line 38Ø
```

line 330 : increment low byte of address
line 340 : branch back to line 300 if high byte does not need
 to be updated
line 350 : else increment high byte of address
line 360 : set Carry flag and do a forced branch back to
 line 300
line 380 : print 'READY' prompt
line 390 : clear accumulator
line 400 : set Carry flag and force a branch to line 500 to
 update and return
line 420 : get '@' into accumulator
line 430 : set Carry flag and force a branch to line 570
line 445 : entry point for PROGRAM-MODE
line 450 : locate and execute command or print appropriate
 error message
line 460 : clear indexing register
line 470 : get byte from program
line 480 : is it a 0 and therefore end of line?
line 490 : yes, branch to line 500
line 500 : no, is it the command delimiter ':'?
line 510 : yes, branch to line 570
line 520 : no, increment low byte of address
line 530 : if not zero, branch back to line 470 to redo loop
line 540 : increment high byte of address
line 550 : set Carry flag and force a branch back to line 470.
line 570 : is it a command delimiter ':'?
line 580 : if greater than or equal to ':' then branch to line 650
line 590 : is it a space?
line 600 : yes, so branch to line 650
line 610 : set Carry flag
line 620 : subtract ASCII base code
line 630 : set Carry flag
line 640 : subtract token and ASCII set bits
line 650 : return to BASIC Interpreter
line 660 : jump to CHRGET
line 665 : entry for FIND-EXECUTE subroutine
line 670 : seed address of command table (\$C100) into vector
 at \$7F
line 710 : increment low byte of command address
line 720 : branch over if no carry into high byte

line 730 : else increment high byte of address
line 740 : back together, initialize Y register
line 750 : and X register
line 760 : get byte from the command table
line 770 : if zero byte, then command is identified, branch to
line 1010
line 780 : is it the same as the byte pointed to in the command
table?
line 790 : no, branch to line 830
line 800 : increment index
line 820 : set Carry flag and force a branch back to line 760
line 830 : command not identified—seek out zero byte. Get
byte from command table
line 840 : if zero, branch to line 880
line 850 : increment index
line 860 : set Carry flag and force a branch to line 830
line 880 : increment index
line 890 : transfer into accumulator
line 900 : clear Carry flag
line 910 : add to low byte of vector address
line 920 : save result
line 930 : clear accumulator
line 940 : add carry to high byte of vectored address
line 950 : and save the result
line 960 : initialize index
line 970 : add two to X to move onto next address in the
line 980 : command address table
line 990 : set Carry flag and force a branch to line 760
line 1010 : get low byte of command execution address
line 1020 : save it in a vector
line 1030 : increment index
line 1040 : get high byte of command execution address
line 1050 : save it in vector
line 1060 : jump to vectored address to execute machine code
of identified command
line 1065 : entry for ILLEGAL—unrecognized WOS command
line 1070 : get error code into X register
line 1080 : and jump to error handling routine

3 ASCII to Binary Conversions

An important aspect of interactive machine code is the ability to convert strings of ASCII characters into their hexadecimal equivalents, so that they may be manipulated by the processor. In this chapter we shall examine, with program examples, how this is performed. The routines provide the following conversions:

1. Single ASCII hex characters into binary.
2. Four ASCII hex digits into two hex bytes.
3. Signed ASCII decimal string into two signed hex bytes.

ASCII HEX TO BINARY CONVERSION

This routine converts a hexadecimal ASCII character in the accumulator into its four-bit binary equivalent. For example, if the accumulator contains \$37 (that is, ASC“7”), the routine will result in the accumulator holding \$7, or ~~0000~~0111 binary. Similarly, if the accumulator holds \$46 (ASC“F”) the routine will return \$F, or ~~0000~~1111, in the accumulator.

Conversion is quite simple, and Table 3.1 gives some indication of what is required.

Table 3.1

Hex	Binary value	ASCII value	ASCII binary
0	00000000	\$30	00110000
1	00000001	\$31	00110001
2	00000010	\$32	00110010
3	00000011	\$33	00110011
4	00000100	\$34	00110100

Table 3.1 (contd.)

5	00000101	\$35	00110101
6	00000110	\$36	00110110
7	00000111	\$37	00110111
8	000001000	\$38	00111000
9	000001001	\$39	00111001
A	000001010	\$41	01000001
B	000001011	\$42	01000010
C	000001100	\$43	01000011
D	000001101	\$44	01000100
E	000001110	\$45	01000101
F	000001111	\$46	01000110

The conversion of ASCII characters 0 to 9 is straightforward. All we need to do is mask off the high nibble of the character's ASCII code. For example ASC “1” is \$31 or 00110001 binary—masking the high nibble with AND \$OF results in 00000001. Converting ASCII characters A and F is a little less obvious, however. If the high nibble of the code is masked off, then the remaining bits are 9 less than the hex required. For example, the ASCII for the letter ‘D’ is \$44 or 01000100. Masking the high nibble with AND \$OF gives 4, or 00000100, and adding 9 to this gives:

$$\begin{array}{r}
 00000100 \\
 + 00001001 \\
 \hline
 00001101
 \end{array}$$

the binary value for \$D.

Program 2

```

10 REM ** CONVERT ASCII CHARACTER IN **
20 REM ** ACCUMULATOR TO BINARY **
30 REM ** REQUIRES 20 BYTES OF MEMORY **
40 :
50 CODE=49152
60 FOR LOOP=0 TO 20
70 READ BYTE
80 POKE CODE+LOOP,BYTE
90 NEXT LOOP
100 :

```

```
110 REM ** M/C DATA **
120 DATA 201,48      REM CMP #$30
130 DATA 144,15      REM BCC $0F
140 DATA 201,58      REM CMP #$3A
150 DATA 144,8       REM BCC $08
160 DATA 233,7       REM SBC #$07
170 DATA 144,7       REM BCC $07
180 DATA 201,64      REM CMP #$40
190 DATA 176,2       REM BCS $02
200 ::               REM ZERO-NINE
210 DATA 41,15       REM AND #$0F
220 ::               REM RETURN
230 DATA 96          REM RTS
240 ::               REM ILLEGAL
250 DATA 56          REM SEC
260 DATA 96          REM RTS
270 :
280 :
290 REM ** TESTING ROUTINE **
300 TEST=49184
310 FOR LOOP=0 TO 14
320 READ BYTE
330 POKE TEST+LOOP,BYTE
340 NEXT LOOP
350 :
360 REM ** M/C TEST DATA **
370 ::               REM TEST
380 DATA 32,228,255  REM JSR $FFE4
390 DATA 240,251     REM BEQ $FB
400 DATA 32,0,192    REM JSR $C000
410 DATA 144,2       REM BCC $02
420 DATA 169,255    REM LDA #$FF
430 ::               REM OVER
440 DATA 133,251    REM STA $FB
450 DATA 96          REM RTS
460 :
470 PRINT CHR$(147)
480 PRINT" HIT A HEX CHARACTER KEY, AND ITS
           BINARY"
490 PRINT" EQUIVALENT VALUE WILL BE PRINTED"
```

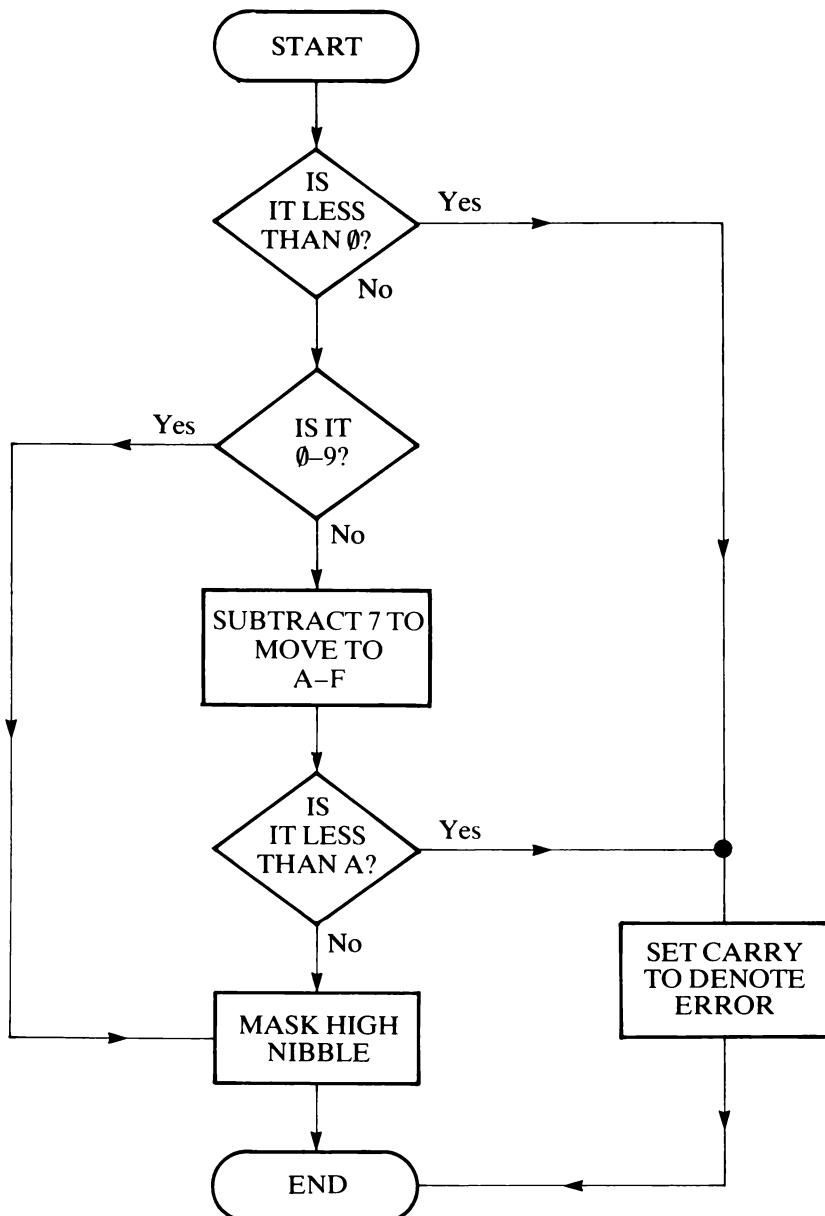


Figure 3.1 Conversion flowchart

5ØØ :

51Ø SYS TEST

52Ø :

53Ø PRINT "RESULT = "PEEK(251)

Program 2 contains a short demonstration, prompting for a hexadecim al value key to be pressed (i.e. Ø to F) and returning its hexadecimal code. Thus, pressing the 'A' key will produce a result of 41.

The ASCII-BINARY routine begins by checking for the legality of the character, by comparing it with 48 (\$30). If the value in the accumulator is less than ASC“Ø”, the Carry flag will be cleared, signalling an error. If the character is legal, the contents are then compared with 58 (\$3A), which is one greater than the ASCII code for 9. This part of the routine ascertains whether the accumulator’s contents are in the range \$30 to \$39. If they are, the Carry flag will be cleared and the branch to ZERO-NINE (lines 15Ø and 12Ø) performed. The high nibble is then masked off to complete the conversion.

If the branch of line 15Ø fails, a legality check for the hex characters A to F is performed. This is done by subtracting 7 from the accumulator’s contents, which should bring the value it holds down below 64 (\$40), or one less than the ASCII code for the letter ‘A’. At this point the Carry flag is set (it was previously set as the branch of the previous line was not performed), and the CMP #\$4Ø of line 18Ø clears it if the contents are higher than 64. The routine then masks off the high nibble, leaving the correct binary.

The following example shows how the conversion of ASC“F” to \$F works:

Mnemonic	Accumulator	Carry flag
	\$46 (ASC" F ")	
CMP #\$3Ø	\$46	1
BCC ILLEGAL		
CMP #\$3A	\$46	1
BCC ZERO-NINE		
SBC #7	\$3F	1
BCC ILLEGAL		
CMP #\$4Ø	\$3F	Ø
BCS RETURN		
AND \$ØF	\$ØF	Ø
RTS		

Note that this routine indicates an error by returning with the Carry flag set, so any calls to the conversion routine should always check for this on return. The short test routine does this, and loads the accumulator with \$FF to signal the fact.

Using two calls to this routine would allow two-byte hex values to be input and converted into a full eight-byte value. On completion of the first call, the accumulator’s contents would need to be shifted into the high nibble.

The coding might look like this:

```
: REM WAIT
JSR GETIN      : REM GET FIRST CHARACTER
BEQ WAIT1
JSR ASCII-BINARY : REM CONVERT TO BINARY
BCS REPORT-ERROR : REM NON-HEX IF C=1
ASL A          : REM MOVE INTO HIGHER
                  NIBBLE
ASL A
ASL A
ASL A
STA HIGH-NIBBLE : REM SAVE RESULT
                  : REM WAIT2
JSR GETIN      : REM GET SECOND CHARACTER
BEQ WAIT2
JSR ASCII-BINARY : REM CONVERT TO BINARY
BCS REPORT-ERROR : REM NON-HEX IF C=1
ORA HIGH-NIBBLE : REM ADD HIGH NIBBLE
                  : REM ALL BINARY NOW IN
                  ACCUMULATOR
```

Using this routine and entering, say, \$FE will return 1111110 in the accumulator.

Line-by-line

A line-by-line description of Program 2 follows:

```
line 120 : is it >= than ASC“0”?
line 130 : no, branch to ILLEGAL
line 140 : is it in range 0-9?
line 150 : yes, branch to ZERO-NINE to skip A-F
            translation.
line 160 : move onto ASCII codes for A-F
line 170 : branch to ILLEGAL if Carry flag clear
line 180 : is it higher than ASC“@”?
line 190 : no, branch to ILLEGAL
line 200 : entry for ZERO-NINE
line 210 : clear high nibble
line 220 : entry for RETURN
```

line 230 : return with binary in accumulator
line 240 : entry for ILLEGAL
line 250 : set Carry flag to denote an error
line 260 : return to BASIC
line 370 : entry for TEST
line 380 : read keyboard
line 390 : if null string, branch to TEST
line 400 : call conversion at \$C000
line 410 : if no errors, branch OVER
line 420 : else error, place 255 in accumulator
line 430 : entry for OVER
line 440 : save accumulator in \$FB
line 450 : and return to BASIC

FOUR ASCII DIGITS TO HEX

We can use the ASCII-BINARY routine as the main subroutine in a piece of coding which will convert four ASCII digits into a two-byte hexadecimal number, making the routine most useful for inputting two-byte hexadecimal addresses. For example, the routine would convert the ASCII string “CAFE” into a two-byte binary number 11001010 11111110 or \$CAFE. Program 3 lists the entire coding:

Program 3

```
10 REM ** CONVERT FOUR ASCII DIGITS INTO **
20 REM ** A TWO-BYTE HEXADECIMAL NUMBER **
30 CODE=49152
40 FOR LOOP=0 TO 62
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 DATA 160,0 : REM LDY #0
110 DATA 162,251 : REM LDX #$FB
120 DATA 148,0 : REM STY $00,X
130 DATA 148,1 : REM STY $01,X
140 DATA 148,2 : REM STY $02,X
150 :: REM NEXT-CHARACTER
```

160 DATA 185,60,3 : REM LDA \$33C,Y
170 DATA 32,42,192 : REM JSR \$C02A
180 DATA 176,21 : REM BCS \$15
190 DATA 10,10 : REM ASL A : ASL A
200 DATA 10,10 : REM ASL A : ASL A
210 DATA 148,2 : REM STY \$02,X
220 DATA 160,4 : REM LDY #\$04
225 :: REM AGAIN
230 DATA 10 : REM ASL A
240 DATA 54,0 : REM ROL \$00,X
250 DATA 54,1 : REM ROL \$01,X
260 DATA 136 : REM DEY
270 DATA 208,248 : REM BNE \$F8
280 DATA 180,2 : REM LDY \$02,Y
290 DATA 200 : REM INY
300 DATA 208,227 : REM BNE \$E3
310 :: REM ERROR
320 DATA 181,2 : REM LDA \$02,X
330 DATA 96 : REM RTS
340 ::
350 REM *** ASCII-BINARY CONVERSION ***
360 DATA 201,48 : REM CMP #\$30
370 DATA 144,15 : REM BCC \$0F
380 DATA 201,58 : REM CMP #\$3A
390 DATA 144,8 : REM BCC \$08
400 DATA 233,7 : REM SBC \$07
410 DATA 144,7 : REM BCC \$07
420 DATA 201,64 : REM CMP #\$40
430 DATA 176,2 : REM BCS \$02
440 :: REM ZERO-NINE
450 DATA 41,15 : REM AND \$0F
460 :: REM RETURN
470 DATA 96 : REM RTS
480 :: REM ILLEGAL
490 DATA 56 : REM SEC
500 DATA 96 : REM RTS
510 ::
520 REM *** SET UP A TEST PROCEDURE ***
530 TEST=49232
540 FOR LOOP=0 TO 34

```

550 READ BYTE
560 POKE TEST+LOOP, BYTE
570 NEXT LOOP
580 :
590 REM ** TEST M/C DATA **
600 DATA 160,0 : REM LDY #$00
610 DATA 162,4 : REM LDX #$04
620 :: REM OVER
630 DATA 142,52,3 : REM STX $334
640 DATA 140,53,3 : REM STY $335
650 :: REM INNER
660 DATA 32,228,255 : REM JSR $FFE4
670 DATA 240,251 : REM BEQ $FB
680 DATA 174,52,3 : REM LDX $334
690 DATA 172,53,3 : REM LDY $335
700 DATA 153,60,3 : REM STA $33C,Y
710 DATA 32,210,255 : REM JSR $FFD2
720 DATA 200 : REM INY
730 DATA 202 : REM DEX
740 DATA 208,229 : REM BNE $E5
750 DATA 32,0,192 : REM JSR $C000
760 DATA 96 : REM RTS
770 :
780 PRINT CHR$(147)
790 PRINT "INPUT A FOUR DIGIT HEX NUMBER : $";
800 SYS TEST
810 PRINT
820 PRINT "THE FIRST BYTE WAS :"; PEEK(251)
830 PRINT "THE SECOND BYTE WAS :"; PEEK(252)

```

The machine code begins by clearing three bytes of zero page RAM pointed to by the contents of the X register (lines 100 to 140). The ASCII characters are accessed one by one from a buffer which may be resident anywhere in memory (line 160), though in this case it is the four bytes at the start of the cassette buffer. Conversion and error-detection are performed (lines 170 and 180) and the four returned bits shifted into the high four bits of the accumulator. The buffer index, which keeps track of the character position in the buffer, is saved in the third of the three bytes cleared.

The loop between lines 250 and 300 is responsible for moving the four bits through the two zero page bytes which hold the final result. In fact, with the accumulator, the whole process of the loop is to perform the operation of a 24-bit shift register. Figure 3.2

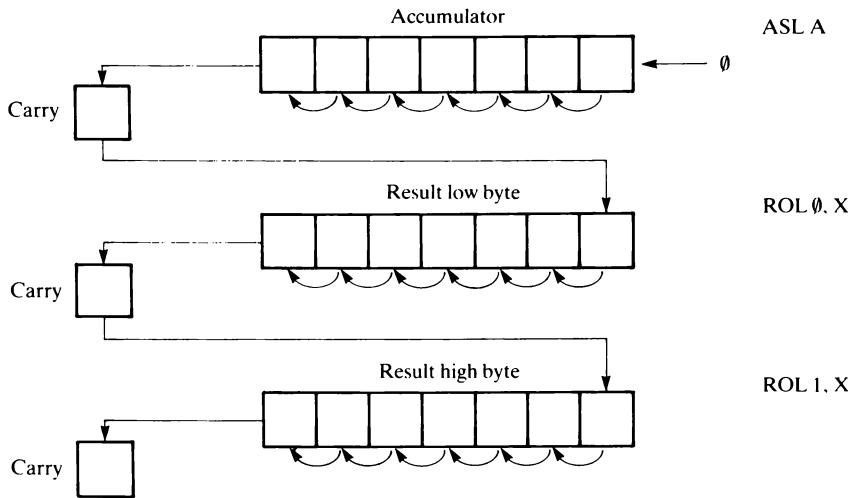


Figure 3.2 Movement of bits through a 3-byte shift register

illustrates the procedure.

The ASL A instruction shuffles the bits in the accumulator one bit to the left, with the dislodged bit 7 moving across into the Carry flag bit. This carry bit is then rotated into bit 0 of the result address low byte, which in turn rotates its bit 7 into the Carry flag. The next ROL instruction repeats this movement on the high byte. The net effect of all this is that as the process is executed four times, the returned conversions are shifted through the result address to reside in the correct place, as Figure 3.3 illustrates.

	1, X	0, X	Accumulator
Entry	00000000	00000000	11110000
1st pass	00000000	00001111	00000000
2nd pass	00000000	11110000	00000000
3rd pass	00001111	00000000	00000000
4th pass	11110000	00000000	00000000

Figure 3.3 A 24-bit shift register, showing passage of the bits in the number \$F000

Error-checking is provided for, the routine aborting when it encounters an illegal hex character, leaving the accumulator containing the index into the buffer, pointing to the illicit value. In fact, this method is used to complete the execution of the conversion-rotate loop, using a RETURN character placed at the end of the

ASCII hex string.

The test routine (lines 590 to 800) prompts for four hex-based characters to be input. These are placed in the buffer (line 610) and printed to the VDU. On completion of the input, the address-binary routine is called, and the result placed in the first two bytes of the user area, for printing or manipulation purposes.

Line-by-line

A line-by-line decription of Program 3 follows:

```
line 100 : clear indexing register
line 110 : get byte destination
line 120 : clear three bytes
line 150 : entry for NEXT-CHARACTER
line 160 : get character from buffer
line 170 : call ASCII-BINARY to convert
line 180 : branch to ERROR if Carry flag is set
line 190 : move low nibble into high nibble
line 210 : save index into buffer
line 220 : moving four bits
line 225 : entry for AGAIN
line 230 : move bit 7 into Carry flag
line 240 : move carry into bit 0 and bit 7 into Carry flag
line 250 : move carry into bit 0 and bit 7 into Carry flag
line 260 : decrement bit count
line 270 : and do until four bits done
line 280 : restore index into buffer
line 290 : increment it to point to next character
line 300 : do branch to NEXT-CHARACTER
line 310 : entry for ERROR
line 320 : get illegal character
line 330 : return to calling routine
```

CONVERT DECIMAL ASCII STRING TO BINARY

This routine takes a signed decimal string of ASCII characters and transforms it into a two-byte hexadecimal number. For example, entering -32,678 will return the value \$8000, where \$8000 is its signed binary equivalent. Entry requirements to the conversion routine are obtained by the BASIC text in lines 880 to 940. Note

that in addition to obtaining the characters for insertion into the string buffer, the number of characters for conversion is required, this being placed in the first byte of the buffer.

Program 4

```
10 REM ** DECIMAL ASCII TO BINARY **
20 REM ** READ & POKE M/C DATA **
30 CODE=49152
40 FOR LOOP=0 TO 155
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 :
110 DATA 174,60,3      REM LDX $33C
120 DATA 208,3          REM BEQ $03
125 DATA 76,154,192 :  REM JMP $C09A
130 DATA 160,0          REM LDY #0
140 DATA 140,55,3       REM STY $337
150 DATA 140,53,3       REM STY $335
160 DATA 150,54,3       REM STY $336
170 DATA 200              REM INY
180 DATA 140,52,3       REM STY $334
190 DATA 185,60,3          REM LDA $33C,Y
200 DATA 201,45            REM CMP #$2D
210 DATA 208,14            REM BNE $0E
220 DATA 169,255           REM LDA #&FF
230 DATA 141,55,3          REM STA $337
240 DATA 238,52,3          REM INC $334
250 DATA 202              REM DEX
260 DATA 240,113           REM BEQ $71
270 DATA 76,54,192 :      REM JMP $C036
280 ::                   REM POSITIVE
290 DATA 201,43             REM CMP #$2B
300 DATA 208,12             REM BNE $06
310 DATA 238,52,3           REM INC $334
320 DATA 202              REM DEX
330 DATA 240,100           REM BEQ $64
```

340 ::	REM CONVERT-CHARACTER
350 DATA 172,52,3	REM LDY \$334
360 DATA 185,60,3	REM LDA \$33C,Y
370 ::	REM CHECK-LEGALITY
380 DATA 201,58	REM CMP #\$3A
390 DATA 16,90	REM BPL \$5A
400 DATA 201,48	REM CMP #\$30
410 DATA 48,86	REM BMI \$56
420 DATA 72	REM PHA
430 DATA 14,53,3	REM ASL \$335
440 DATA 46,54,3	REM ROL \$336
450 DATA 173,53,3	REM LDA \$335
460 DATA 172,53,3	REM LDY \$336
470 DATA 14,53,3	REM ASL \$335
480 DATA 46,54,3	REM ROL \$336
490 DATA 14,53,3	REM ASL \$335
500 DATA 46,54,3	REM ROL \$336
510 DATA 24	REM CLC
520 DATA 109,53,3	REM ADC \$335
530 DATA 141,53,3	REM STA \$335
540 DATA 152	REM TYA
550 DATA 109,54,3	REM ADC \$336
560 DATA 141,54,3	REM STA \$336
570 DATA 56	REM SEC
580 DATA 104	REM PLA
590 DATA 233,48	REM SBC #\$30
600 DATA 24	REM CLC
610 DATA 109,53,3	REM ADC \$335
620 DATA 141,53,3	REM STA \$335
630 DATA 144,3	REM BCC \$03
640 DATA 238,54,3	REM INC \$336
650 ::	REM NO-CARRY
660 DATA 238,52,3	REM INC \$334
670 DATA 202	REM DEX
680 DATA 208,181	REM BNE \$B5
690 DATA 173,55,3	REM LDA \$337
700 DATA 16,17	REM BPL \$11
710 DATA 56	REM SEC
720 DATA 169,0	REM LDA #0
730 DATA 237,53,3	REM SBC \$335

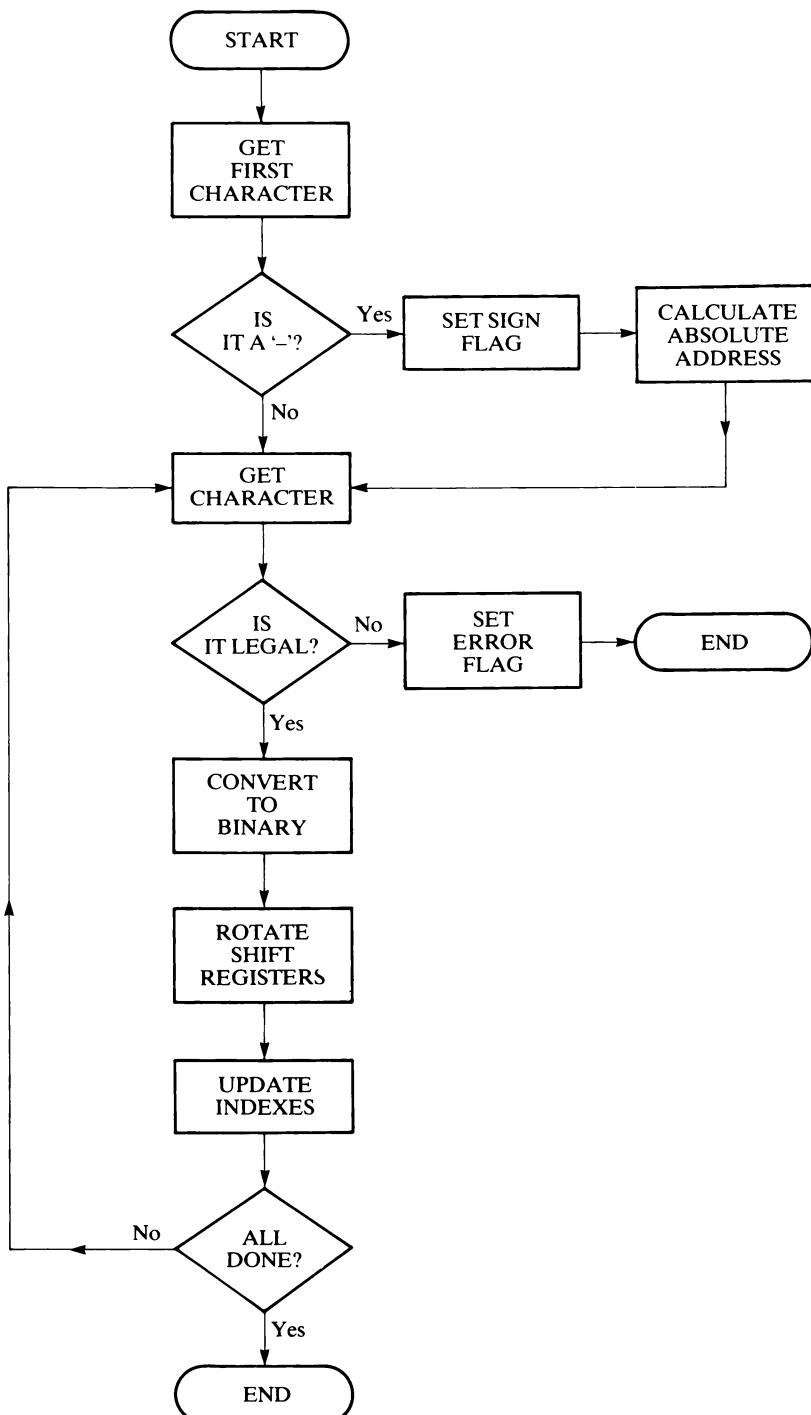


Figure 3.4 ASCII string to binary conversion flowchart

```

740 DATA 141,53,3      REM STA $335
750 DATA 169,0          REM LDA #0
760 DATA 237,54,3      REM SBC $336
770 DATA 141,54,3      REM STA $336
780 ::                  REM NO-COMPLEMENT
790 DATA 24             REM CLC
800 DATA 144,1          REM BCC $1
810 ::                  REM ERROR
820 DATA 56             REM SEC
830 ::                  REM FINISH
840 DATA 96             REM RTS
850 :
860 REM ** SET UP SCREEN AND GET NUMBER **
870 PRINT CHR$(147)
880 INPUT"NUMBER FOR CONVERSION";A$
890 FOR LOOP=1 TO LEN(A$)
900 TEMP$=MID$(A$,LOOP,1)
910 B=ASC(TEMP$)
920 POKE 828+LOOP,B
930 NEXT LOOP
940 POKE 828,LEN(A$)
950 :
960 SYS CODE
970 :
980 PRINT"THE TWO BYTES ARE AS FOLLOWS"
990 PRINT"LOW BYTE ";PEEK(821)
1000 PRINT"HIGH BYTE ";PEEK(822)

```

Bytes are designated as follows:

820 (\$334)	:	string index
821 (\$335)		current count
823 (\$336)	:	sign flag
828 (\$33C)		length of string
829 (\$33D)	:	start of character string

The machine code begins by obtaining the character count from the X register. An error is signalled if this count is zero, otherwise the

program progresses, clearing the sign flag (used to signal positive or negative values) and result destination bytes at ‘current’ (lines 130 to 160). Location \$70 is used to hold the string index, pointing to the next character for conversion. This byte is initially loaded with 1 so that it skips over the count byte in the buffer.

The first byte of the string is tested for a ‘+’ or ‘-’ sign, the former being an optional item in the string, and the sign flag is set accordingly (lines 190 to 230). The CONVERT-CHARACTER loop starts by testing the character about to be manipulated to ensure it is a decimal value, i.e. 0 to 9 inclusive. Converting the byte into binary form is achieved by multiplying the byte by 10. This multiplication is readily available using four arithmetic shifts and an addition: $2 * 2 * 2 + 2 = 10$.

Because we are dealing with a two-byte result, the arithmetic shift must be performed on the two bytes, allowing bits to be transferred from one byte to the other. This is performed by using an ASL followed by a ROL. As figure 3.5 illustrates, this acts exactly like a 16-bit ASL. The first pass through this character-conversion loop has little effect, as it is operating on characters already converted, of which there are none first time round!

Lines 570 to 620 carry out the conversion of ASCII to binary and store the result. This is performed, as we know from earlier examples, by masking off the high nibble. Another technique for doing this is simply to subtract the ASCII code for ‘0’: \$30.

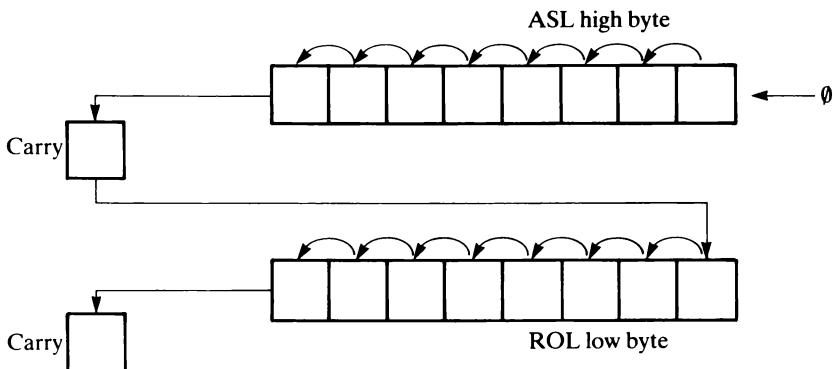


Figure 3.5 A 16-bit arithmetic shift

Once all the characters have been processed, the sign flag at \$334 (820) is checked for a negative value. If this is indicated (lines 690 and 700), the value of current is subtracted from zero, thereby converting the absolute value into a signed negative byte (lines 710 to 770). The Carry flag is used to indicate any error conditions—if it is set an error occurred, and the string index at \$334 points to the illegal character.

Line-by-line

A line-by-line description of Program 4 now follows:

line 11 \emptyset : get length of string
line 12 \emptyset : branch if not zero
line 125 : else jump to ERROR
line 13 \emptyset : clear Y register
line 14 \emptyset : sign flag
line 15 \emptyset : and store bytes
line 17 \emptyset : increment Y
line 18 \emptyset : set index to first ASCII character
line 19 \emptyset : get first character
line 20 \emptyset : is it a minus sign?
line 21 \emptyset : no, branch to POSITIVE
line 22 \emptyset : yes, get negative byte
line 23 \emptyset : and set the sign flag
line 24 \emptyset : move to next character
line 25 \emptyset : decrement length counter
line 26 \emptyset : branch to ERROR if zero
line 27 \emptyset : else jump to CONVERT-CHARACTER
line 28 \emptyset : entry for POSITIVE
line 29 \emptyset : is first character a +?
line 30 \emptyset : no, branch to CHECK-LEGALITY
line 31 \emptyset : yes, move to next character
line 32 \emptyset : decrement length counter
line 33 \emptyset : branch to ERROR if zero
line 34 \emptyset : entry for CONVERT-CHARACTER
line 35 \emptyset : restore index
line 36 \emptyset : get character from buffer
line 37 \emptyset : entry for CHECK-LEGALITY
line 38 \emptyset : is it \leq ASC“9”?
line 39 \emptyset : no, it’s bigger, branch to ERROR
line 40 \emptyset : is it \geq ASC“0”?
line 41 \emptyset : no, branch to ERROR
line 42 \emptyset : save code on stack
line 43 \emptyset : multiply both bytes by two
line 45 \emptyset : save low byte

line 46 \emptyset : save high byte
line 47 \emptyset : multiply by two again (now *4)
line 49 \emptyset : and again (now *8)
line 51 \emptyset : clear Carry flag
line 52 \emptyset : add low byte *2
line 53 \emptyset : and save result
line 54 \emptyset : transfer high byte *2
line 55 \emptyset : and add to *8 high byte
line 56 \emptyset : save it. Now *10
line 57 \emptyset : set Carry flag
line 58 \emptyset : restore ASCII code from stack
line 59 \emptyset : convert ASCII to binary
line 60 \emptyset : clear Carry flag
line 61 \emptyset : add it to low byte current
line 62 \emptyset : save result
line 63 \emptyset : branch if NO-CARRY
line 64 \emptyset : else increment high byte
line 65 \emptyset : entry for NO-CARRY
line 66 \emptyset : move index on to next byte
line 67 \emptyset : decrement length counter
line 68 \emptyset : branch to CONVERT-CHARACTER if not finished
line 69 \emptyset : completed so get sign flag
line 70 \emptyset : if clear branch to NO-COMPLEMENT
line 71 \emptyset : else set Carry flag
line 72 \emptyset : clear accumulator
line 73 \emptyset : and obtain two's complement
line 74 \emptyset : save low byte result
line 75 \emptyset : clear accumulator
line 76 \emptyset : subtract high byte from \emptyset
line 77 \emptyset : and save result
line 78 \emptyset : entry for NO-COMPLEMENT
line 79 \emptyset : clear Carry flag
line 80 \emptyset : and force branch to FINISH
line 81 \emptyset : entry for ERROR
line 82 \emptyset : set Carry flag to denote error
line 83 \emptyset : entry for FINISH
line 84 \emptyset : return to BASIC

4 Binary to Hex ASCII

This chapter complements the previous one and illustrates how memory-based hex values can be converted into their ASCII representation. The routines provide the following conversions:

1. Print accumulator as two ASCII hex characters.
2. Print two hex bytes as four ASCII hex characters.
3. Print two-byte signed binary number as signed decimal number.

PRINT ACCUMULATOR

To convert an eight-bit binary number into its ASCII hex equivalent characters, the procedure described in Chapter 3 must be reversed. However, because text is printed on the screen from left to right, we must deal with the high nibble of the byte first. Program 5 uses the hexprint routine to print the hexadecimal value of any key pressed at the keyboard.

Program 5

```
10 REM ** PRINT ACCUMULATOR AS A HEX NUMBER **
20 :
30 CODE=49152
40 FOR LOOP=0 TO 21
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 :
```

```

110 DATA 72           REM PHA
120 DATA 74, 74       REM ASL A : ASL A
130 DATA 74, 74       REM ASL A : ASL A
140 DATA 32,9,192     REM JSR $C009
150 DATA 104          REM PLA
160 :                REM FIRST $C009
170 DATA 41,15         REM AND #$0F
180 DATA 201,10         REM CMP #$0A
190 DATA 144,02         REM BCC $02
200 DATA 105,6          REM ADC #$06
210 :                REM OVER
220 DATA 105,48         REM ADC #$30
230 DATA 76,210,255     REM JMP $FFD2
240 :
250 REM ** SET UP DEMO AT 828 **
260 REM LDA $FB : JMP $C000
270 POKE 828,165 : POKE 829,251
280 POKE 830,76 : POKE 831,0 : POKE 832,192
290 PRINT CHR$(147)
300 PRINT "HIT ANY KEY AND ITS HEX VALUE IN"
310 PRINT "ASCII WILL BE DISPLAYED"
320 GET A$: IF A$="" THEN GOTO 320
330 A=ASC(A$)
340 POKE 251,A
350 :
360 SYS 828
370 REM CALL 'SYS CODE' TO USE DIRECTLY

```

The hexprint routine is embedded between lines 110 and 230. The accumulator's contents are first pushed on to the hardware stack. This procedure is necessary as it will have to be restored before the second pass, which calculates the ASCII code for the second character. The first pass through the routine sets about moving the upper nibble of the accumulator byte into the lower nibble (lines 120 and 130). The FIRST subroutine ensures that the high nibble is cleared by logically ANDing it with \$0F. This is, of course, surplus to requirement on the first pass, but is needed on the second pass to isolate the low nibble. Comparing the accumulator's contents with 10 will ascertain whether the value is in the range 0 to 9 or A to F. If the Carry flag is clear, it falls in the lower range (0 to 9) and simply setting bits 4 and 5, by adding \$30, will give the required ASCII code. A further 7 must be added to skip non-hex ASCII codes to arrive at the ASCII codes for A to F (\$41 to \$46). You may have

noticed that line 200 does not add 7 but in fact adds one less, 6. This is because, for this section of coding to be executed, the carry must have been set, and the 6510 addition opcode references the Carry flag in addition. Therefore, the addition performed is: accumulator + 6 + 1.

The JMP of line 230 will return the program back to line 150. Remember, FIRST was called with a JSR, so the RTS from completion of the CHROUT call returns control here. The accumulator is restored and the process repeated for the second ASCII digit.

A short test routine is established in lines 250 to 340. This requests you to hit a key, the value of which is placed in a free zero page byte. The ‘hand-POKEd’ routine at 828 is called by line 360, and puts the key’s value into the accumulator before performing a jump to the main routine.

The following example illustrates the program’s operation, assuming the accumulator holds the value 01001111, \$4F:

Mnemonic	Accumulator	Carry flag
	\$4F	
LSR A	\$27	1
LSR A	\$13	1
LSR A	\$09	1
LSR A	\$04	1
JSR FIRST		
AND #\$0F	\$04	1
CMP #\$0A	\$04	Ø
BCC OVER		
OVER		
ADC #\$30	\$34 (ASC"4")	Ø
JMP CHROUT		
PLA	\$4F	Ø
AND #\$0F	\$0F	Ø
CMP #\$0A		

Line-by-line

A line-by-line description of Program 5 follows:

line 110 : save accumulator on stack

line 120 : move high nibble into low nibble

line 14 \emptyset : call FIRST subroutine
line 15 \emptyset : restore accumulator
line 16 \emptyset : entry for FIRST
line 17 \emptyset : ensure only low nibble set
line 18 \emptyset : is it < 10?
line 19 \emptyset : yes, branch to OVER
line 20 \emptyset : no, add 7, value \$A to \$F
line 21 \emptyset : entry for OVER
line 22 \emptyset : add 48 to convert to ASCII code
line 23 \emptyset : and print, returning to line 14 \emptyset or BASIC

PRINT A HEXADECIMAL ADDRESS

The hexprint routine can be extended to enable two zero page bytes to be printed out in hexadecimal form. This is an especially important procedure when writing machine based utilities, such as a hex dump or disassembler. The revamped program is listed below:

Program 6

```
1 $\emptyset$  REM ** PRINT TWO HEX BYTES AS **  
2 $\emptyset$  REM ** A TWO-BYTE ADDRESS **  
3 $\emptyset$  CODE=49152  
4 $\emptyset$  FOR LOOP= $\emptyset$  TO 34  
5 $\emptyset$  READ BYTE  
6 $\emptyset$  POKE CODE+LOOP,BYTE  
7 $\emptyset$  NEXT LOOP  
8 $\emptyset$  :  
9 $\emptyset$  REM ** M/C DATA **  
10 $\emptyset$  REM ** CALL WITH $FB,$FC HOLDING BYTES **  
11 $\emptyset$  :: REM ADDRESS-PRINT  
12 $\emptyset$  DATA 162,251 : REM LDX #$FB  
13 $\emptyset$  DATA 181,1 : REM LDA $01,X  
14 $\emptyset$  DATA 32,13,192 : REM JSR $C00D  
15 $\emptyset$  DATA 181, $\emptyset$  : REM LDA $00,X  
16 $\emptyset$  DATA 32,13,192 : REM JSR $C00D  
17 $\emptyset$  DATA 96 : REM RTS  
18 $\emptyset$  :: REM HEXPRINT  
19 $\emptyset$  DATA 72 : REM PHA  
20 $\emptyset$  DATA 74,74 : REM LSR A : LSR A
```

```

210 DATA 74,74      : REM   LSR A : LSR A
220 DATA 32,22,192  : REM   JSR   $C016
230 DATA 104        : REM   PLA
240 ::              REM FIRST
250 DATA 41,15      : REM   AND # $0F
260 DATA 201,10     : REM   CMP # $0A
270 DATA 144,2       : REM   BCC $02
280 DATA 105,6       : REM   ADC # $06
290 ::              REM OVER
300 DATA 105,48      : REM   ADC # $30
310 DATA 76,210,255 : REM   JMP  $FFD2

```

Zero paged indexed addressing is used to access the two bytes, the crucial location being given in the X register, which acts as the index for the high byte, LDA \$01,X (line 130), and the low byte, LDA \$00,X (line 150). The all-important address in this instance is \$FB (line 130), so the bytes accessed by ADDRESS-PRINT are \$FB (\$FB+0) and \$FC (\$FB+1). Using this method, various addresses can be housed within zero page and any one reached simply by seeding the X register with the location value.

Project

Adapt Program 6 to accept a five character decimal number from the keyboard, printing its hexadecimal value on the screen. Remember—no BASIC, and the input routine must be able to accept numbers in the range 0 to 65!

BINARY SIGNED NUMBER TO SIGNED ASCII DECIMAL STRING

This conversion utility takes a two-byte hexadecimal number and converts it into its equivalent decimal based ASCII character string. For example, if the two-byte value is \$7FFF, the decimal string is 32,767, \$7FFF being 32,767 in decimal. The coding uses signed binary values so that if the most significant bit is set, a negative value is interpreted. This is relayed in the string with a minus sign. This means that the routine can handle values in the range 32,767 to -32,768. When using the routine, remember that the two's complement representation is used, so that a hex value of \$FFFF is converted to the string -1, and \$8000 returns the character string -32,767.

The two address bytes are located at \$334 and \$335 and the string buffer from \$FB onwards. The length of the string buffer will vary, but its maximum length will not exceed six digits, so this number of bytes should be reserved.

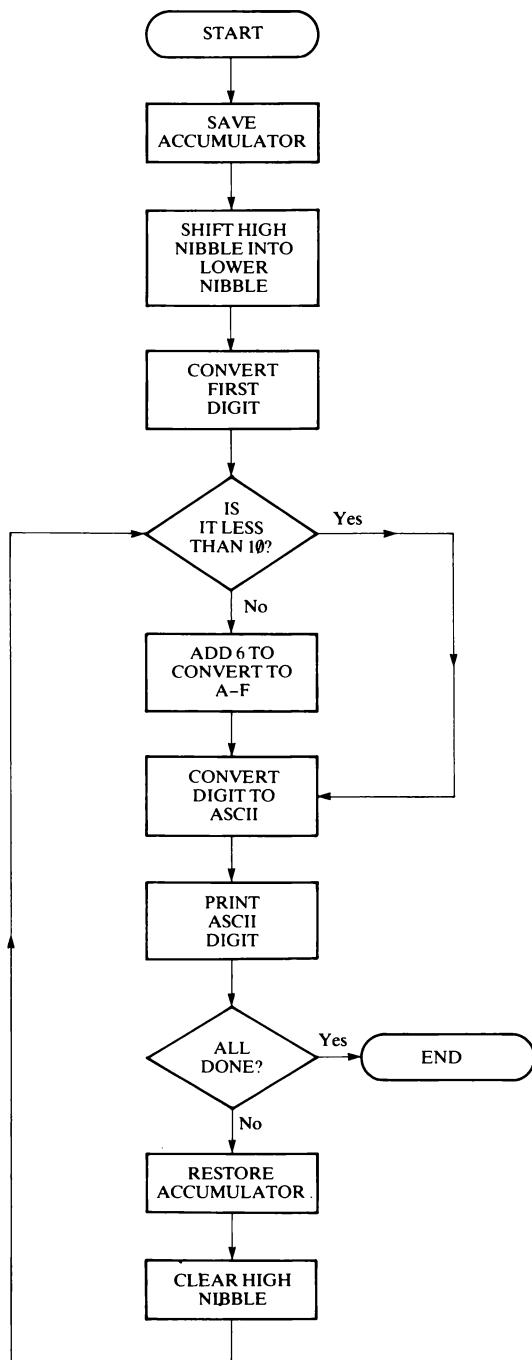


Figure 4.1 Hex to ASCII conversion flowchart

Program 7

```
10 REM ** BINARY SIGNED NUMBER CONVERSION **
20 REM ** INTO SIGNED DECIMAL ASCII STRING **
30 CODE=49152 : OUTPUT=49301
40 FOR LOOP=0 TO 163
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 DATA 160,0 : REM LDY #$00
110 DATA 152 : REM TYA
120 DATA 133,251 : REM STA $FB
130 DATA 133,252 : REM STA $FC
140 DATA 133,253 : REM STA $FD
150 DATA 133,254 : REM STA $FE
160 DATA 133,255 : REM STA $FF
170 DATA 173,53,3 : REM LDA $335
180 DATA 141,56,3 : REM STA $338
190 DATA 16,15 : REM BPL $0F
200 DATA 56 : REM SEC
210 DATA 152 : REM TYA
220 DATA 237,52,3 : REM SBC $334
230 DATA 141,52,3 : REM STA $334
240 DATA 152 : REM TYA
250 DATA 237,53,3 : REM SBC $335
260 DATA 141,53,3 : REM STA $335
270 :: REM CONVERSION
280 DATA 169,0 : REM LDA #$00
290 DATA 141,54,3 : REM STA $336
300 DATA 141,55,3 : REM STA $337
310 DATA 24 : REM CLC
320 DATA 162,16 : REM LDX #$10
330 :: REM LOOP
340 DATA 46,52,3 : REM ROL $334
350 DATA 46,53,3 : REM ROL $335
360 DATA 46,54,3 : REM ROL $336
370 DATA 46,55,3 : REM ROL $337
380 DATA 56 : REM SEC
```

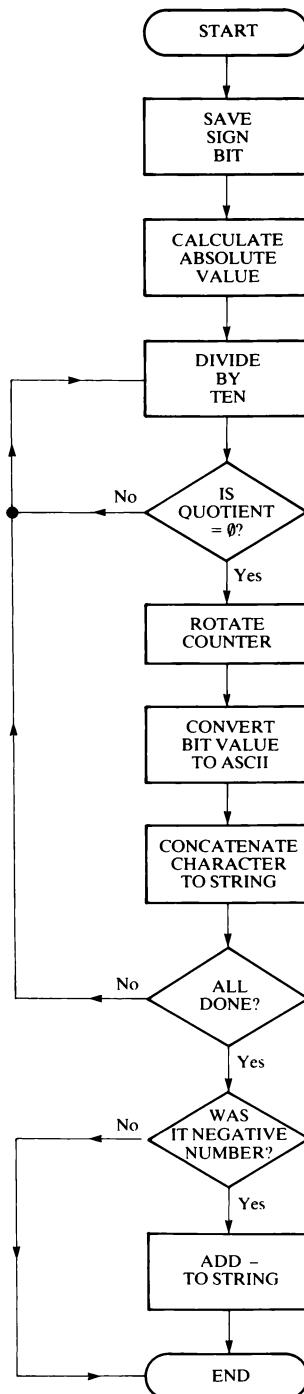


Figure 4.2 Binary to ASCII string conversion flowchart

```

390 DATA 173,54,3 : REM LDA $336
400 DATA 233,10 : REM SBC #$0A
410 DATA 168 : REM TAY
420 DATA 173,55,3 : REM LDA $337
430 DATA 233,0 : REM SBC #$00
440 DATA 144,6 : REM BCC $06
450 DATA 140,54,3 : REM STY $336
460 DATA 141,55,3 : REM STA $337
470 :: REM LESS-THAN
480 DATA 202 : REM DEX
490 DATA 208,221 : REM BNE $DD
500 DATA 46,52,3 : REM ROL $334
510 DATA 46,53,3 : REM ROL $335
520 :: REM ADD-ASCII
530 DATA 24 : REM CLC
540 DATA 173,54,3 : REM LDA $336
550 DATA 105,48 : REM ADC #$30
560 DATA 32,116,192 : REM JSR $C074
570 DATA 173,52,3 : REM LDA $334
580 DATA 13,53,3 : REM ORA $335
590 DATA 208,187 : REM BNE $BB
600 :: REM FINISHED
610 DATA 173,56,3 : REM LDA $338
620 DATA 16,5 : REM BPL $05
630 DATA 169,45 : REM LDA #$2D
640 DATA 32,116,192 : REM JSR $C074
650 :: REM POSITIVE
660 DATA 96 : REM RTS
670 REM SUBROUTINE TO FORM ASCII CHARACTER
      STRING IN $FB
680 :: REM CONCATENATE
690 DATA 72 : REM PHA
700 DATA 160,0 : REM LDY #$00
710 DATA 185,251,0 : REM LDA $00FB,Y
720 DATA 168 : REM TAY
730 DATA 240,11 : REM BEQ $0B
740 :: REM SHUFFLE-ALONG
750 DATA 185,251,0 : REM LDA $00FB,Y
760 DATA 200 : REM INY
770 DATA 153,251,0 : REM STA $00FB,Y

```

```

780 DATA 136,136      : REM DEY : DEY
790 DATA 208,245      : REM BNE $F5
800 ::                  REM ZERO-FINISH
810 DATA 104           : REM PLA
820 DATA 160,1          : REM LDY #$01
830 DATA 153,251,0      : REM STA $00FB,Y
840 DATA 136           : REM DEY
850 DATA 182,251       : REM LDX $FB,Y
860 DATA 232           : REM INX
870 DATA 150,251       : REM STX $FB,Y
880 DATA 96            : REM RTS
890 REM STRING PRINTING ROUTINE
900 ::                  REM STRING-PRINT
910 DATA 166,251       : REM LDX $FB
920 DATA 160,1          : REM LDY #$01
930 ::                  REM PRINT-LOOP
940 DATA 185,251,0      : REM LDA $FB,Y
950 DATA 32,210,255     : REM JSR $FFD2
960 DATA 200             : REM INY
970 DATA 202             : REM DEX
980 DATA 208,246        : REM BNE $F6
990 DATA 96            : REM RTS
1000 :
1010 REM ** GET IN A HEX NUMBER **
1020 PRINT CHR$(147) : PRINT
1030 PRINT"INPUT A HEX NUMBER :$$";
1040 GOSUB 2000
1050 POKE 820,LOW        : REM LOW BYTE HEX
    NUMBER
1060 GOSUB 2000
1070 POKE 821,HIGH       : REM HIGH BYTE HEX
    NUMBER
1080 :
1090 SYS CODE            : REM CALL CONVERSION
1100 :
1110 PRINT"ITS DECIMAL EQUIVALENT IS :";
1120 SYS OUTPUT
1130 END
1140 :
1999 REM ** HEX INPUT CONTROL **

```

```

2000 GOSUB 2500
2010 F=NUM : PRINT Z$;
2020 GOSUB 2500
2030 S=NUM : PRINT Z$;
2040 HIGH=F*16+S
2050 GOSUB 2500
2060 F=NUM : PRINT Z$;
2070 GOSUB 2500
2080 S=NUM : PRINT Z$;
2090 LOW=F*16+S
2100 RETURN
2200 :
2499 REM ** GET HEX ROUTINE **
2500 GET Z$
2510 IF Z$="" THEN GOTO 2500
2520 IF Z$>"F" THEN GOTO 2500
2530 IF Z$="A" THEN NUM=10: RETURN
2540 IF Z$="B" THEN NUM=11: RETURN
2550 IF Z$="C" THEN NUM=12: RETURN
2560 IF Z$="D" THEN NUM=13: RETURN
2570 IF Z$="E" THEN NUM=14: RETURN
2580 IF Z$="F" THEN NUM=15: RETURN
2590 NUM=VAL(Z$): RETURN

```

Functional bytes:

251–255	(\$FB–\$FF)	: ASCII string buffer
820–821	(\$334–\$335)	: binary address for conversion
822–823	(\$336–\$337)	: temporary storage
824	(\$338)	: sign flag

To demonstrate the routine's workings, the program first prompts for a hexadecimal number using the BASIC hex loader subroutine at line 2000. This is evaluated and placed at BINARY-ADDRESS by lines 1050 and 1070.

The program proper begins by clearing the string buffer area (lines 100 to 160), an important procedure which ensures no illicit characters find their way into the ASCII string. The sign of the number is tested by loading the high byte of the address byte into the accumulator and saving its value in the sign flag byte. This process will condition the Negative flag. If it is set, a negative number is interpreted and the plus branch to CONVERSION (line

190) fails. The next seven operations obtain the absolute value of the two-byte number by subtracting it from itself and the set carry bit. Thus \$FFFF will result in an absolute value of 1 and \$8000 an absolute value of 32,678.

The two flows of the program rejoin at line 280, where the two temporary bytes are cleared. These bytes are used in conjunction with the binary address bytes to form a 32-bit shift register, allowing bits to flow from the low byte address to the high byte of temporary.

The loop of lines 340 to 510 performs the conversion, by successively dividing through by ten until the quotient has a value of zero. By this time the binary equivalent of this ASCII character being processed will have been placed in the temporary byte. To produce this, the loop needs sixteen iterations so the X register is used to count these out. Converting the binary to hex involves simply adding \$30 or ASC“0” to it (lines 530 to 550).

Because it may not be immediately clear what is happening, Table 4.1 shows the values of the accumulator and four associated bytes after each of the 16 passes of the loop, when converting \$FFFF into its absolute ASCII value of 1. It should be clear from this how the bits shuffle their way through the four byte ‘register’.

Table 4.1

Iteration	Accumulator	\$334	\$335	\$336	\$337
1	00	01	00	00	00
2	FF	02	00	00	00
3	FF	04	00	00	00
4	FF	08	00	00	00
5	FF	10	00	00	00
6	FF	20	00	00	00
7	FF	40	00	00	00
8	FF	80	00	00	00
9	FF	00	00	01	00
10	FF	00	00	01	00
11	FF	00	00	01	00
12	FF	00	00	01	00
13	FF	00	00	01	00
14	FF	00	00	01	00
15	FF	00	00	01	00
16	FF	00	00	01	00

All that is now required is for this character to be added to the string buffer. This concatenation is completed by the code of lines 690 to 880. This began by obtaining the buffer index, which contains the current number of characters already concatenated. This is stored in the first byte of the buffer, \$FB in this instance. It is then moved across into the accumulator. Next, lines 750 to 790 move any characters present in the buffer up memory one byte, thereby opening up a gap of one byte into which the newly formed character can be placed (lines 810 to 870). The buffer index is also incremented and restored at this point, before an RTS is made back to the main body of the program.

End of program operation is tested for by logically ORing the contents of the high and low bytes of the address. If the result is zero, all bits have been rotated and dealt with, in which case the sign flag byte is tested to ascertain whether a minus sign need be placed at the start of the ASCII string (lines 600 to 660).

Line-by-line

A line-by-line description of Program 7 follows:

```
line 100 : clear Y register
line 110 : and accumulator
line 120 : and then the five buffer bytes
line 170 : get high byte for conversion
line 180 : save in sign flag
line 190 : if positive branch to CONVERSION
line 200 : else set Carry flag
line 210 : clear accumulator
line 220 : obtain absolute value of low byte
line 230 : and save
line 240 : clear accumulator
line 250 : obtain absolute value of high byte
line 260 : and save
line 270 : entry for CONVERSION
line 280 : clear accumulator
line 290 : clear temporary storage bytes
line 310 : clear Carry flag
line 320 : sixteen bits to process
line 330 : entry for LOOP
line 340 : move bit 7 into Carry flag
line 350 : and on into bit 0
line 360 : move bit 7 into Carry flag
```

line 370 : and on into bit 0
line 380 : set Carry flag
line 390 : get low byte of temp
line 400 : subtract 10
line 410 : save result in Y
line 420 : get high byte of temporary
line 430 : subtract carry bit
line 440 : branch to LESS-THAN if divisor>dividend
line 450 : else save result of operation in temporary
line 470 : entry for LESS-THAN
line 480 : decrement bit count
line 490 : branch to LOOP until 16 bits done
line 500 : rotate bit 7 into Carry flag
line 510 : and on into bit 0
line 520 : entry for ADD-ASCII
line 530 : clear Carry flag
line 540 : get low byte from temporary
line 550 : convert into ASCII character
line 560 : concatenate on to string in buffer
line 570 : get low byte of binary number
line 580 : OR with high byte. If 0 then all done
line 590 : if not finished branch to CONVERSION
line 600 : entry for FINISHED
line 610 : get sign
line 620 : if N = 0 branch to POSITIVE
line 630 : otherwise get ASC“-”
line 640 : and add it to final string
line 650 : entry for POSITIVE
line 660 : back to BASIC
line 680 : entry for CONCATENATE, \$C074
line 690 : save accumulator
line 700 : initialize index
line 710 : and get buffer length
line 720 : move it into Y for indexing
line 730 : if 0 branch to ZERO-LENGTH
line 740 : entry for SHUFFLE-ALONG
line 750 : get character from buffer
line 760 : increment index
line 770 : save character one byte along
line 780 : restore original address minus one

line 790 : branch to SHUFFLE-ALONG until completed
line 800 : entry for ZERO-FINISH
line 810 : restore accumulator
line 820 : index past length byte
line 830 : add character to buffer
line 840 : decrement index
line 850 : get length byte
line 860 : increment it
line 870 : save it
line 880 : back to calling routine
line 900 : entry for OUTPUT
line 910 : get length of string as counter
line 920 : set index to first character
line 930 : entry for PRINT-LOOP
line 940 : get character
line 950 : print it
line 960 : increment index
line 970 : decrement count
line 980 : branch to PRINT-LOOP until all done
line 990 : back to BASIC

```
110 INPUT"WHICH DIRECTION
```

5 String Manipulation

In this chapter we will look at how ASCII character strings can be manipulated using machine code routines to perform the following operations:

1. Compare two strings.
2. Concatenate one string onto another.
3. Copy a substring from within a main string.
4. Insert a substring into a main string.

These types of routines are essential if you intend to write any programs that manipulate data and information. Adventure games are a typical example of this kind of program.

COMPARING STRINGS

String comparison is normally performed after the computer user has input some information from the keyboard. In BASIC this might be written as:

```
100 A$="MOVE LEFT"  
110 INPUT"WHICH DIRECTION ?"; B$  
120 IF A$=B$ THEN PRINT "CORRECT!"
```

We do not always wish to test for equality, however. In BASIC, we are able to test for unlike items using the NOT operators '<>'. Thus, line 120 could have been written as:

```
120 IF A$ <> B$ PRINT "WRONG!"
```

At other times, we may wish to test which of two strings has a greater length, and this is possible in BASIC using the LEN statement:

```
210 IF LEN(A$) > LEN(B$) THEN PRINT "FIRST"
```

Program 8 gives the assembler and BASIC listing for the string comparison routine, which puts all the functions described above at your disposal whenever the program is used. The Status register holds these answers in the Zero and Carry flags. The Zero flag is used to signal equality: if it is set ($Z=1$), the two strings compared were identical; if it is cleared ($Z=0$) they were dissimilar.

The Carry flag returns information as to which of the two strings was the longer: if it is set ($C=1$), they were identical in length or the first string was the larger. The actual indication required here is evaluated in conjunction with the Zero flag. If $Z=0$ and $C=1$, then a longer string rather than an equal-length string is indicated, but if the Carry flag is returned clear ($C=0$), then the second string was longer than the first.

Program 8

```
10 REM ** STRING COMPARISON ROUTINE **
20 CODE=49152
30 TEST=49184
40 FOR LOOP=0 TO 41
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 DATA 173,52,3 : REM LDA $334
110 DATA 205,53,3 : REM CMP $335
120 DATA 144,3 : REM BCC $03
130 DATA 174,53,3 : REM LDX $335
140 :: : REM COMPARE-STRING
150 DATA 240,12 : REM BEQ $0C
160 DATA 160,0 : REM LDY #$00
170 :: : REM COMPARE-BYTES
180 DATA 177,251 : REM LDA ($FB),Y
190 DATA 209,253 : REM CMP ($FD),Y
200 DATA 208,10 : REM BNE $0A
210 DATA 200 : REM INY
220 DATA 202 : REM DEX
230 DATA 208,246 : REM BNE $F6
240 :: : REM CONDITION-FLAGS
250 DATA 173,52,3 : REM LDA $334
```

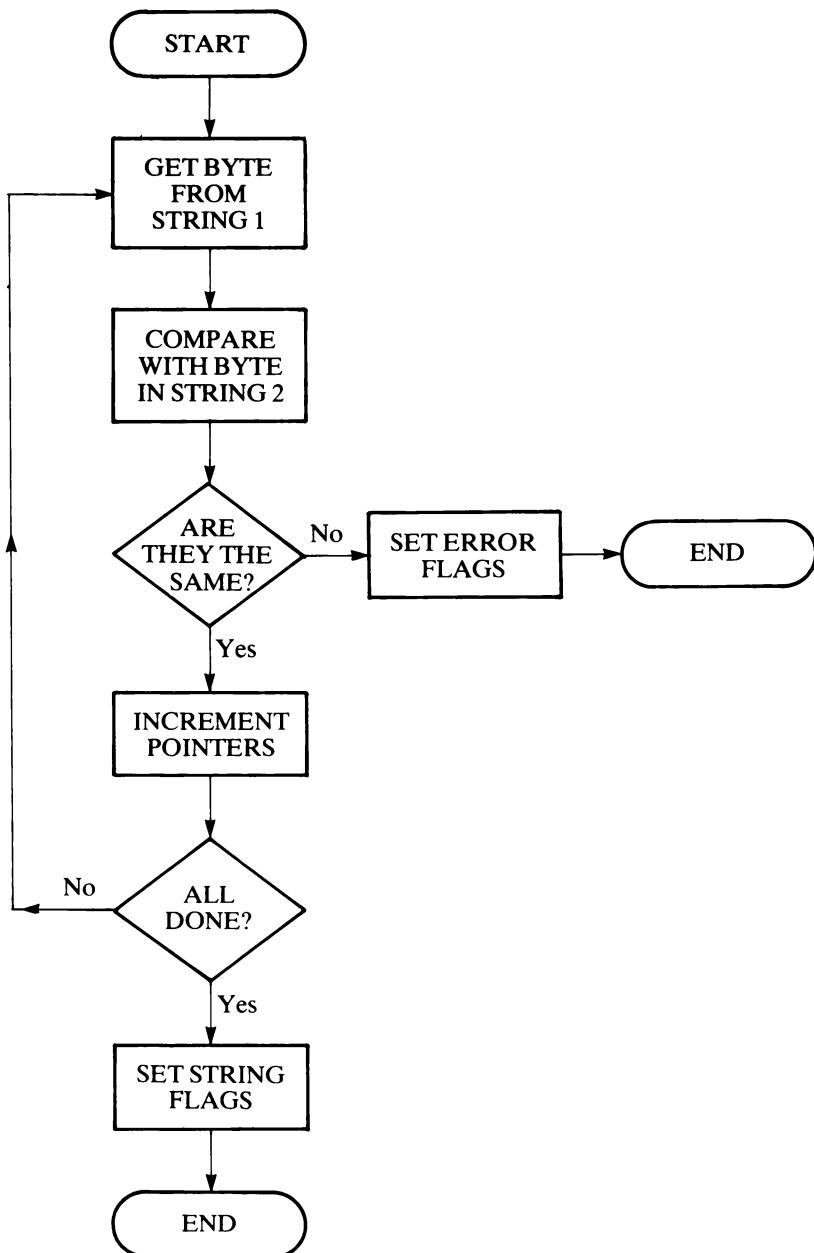


Figure 5.1 Compare strings flowchart

```

260 DATA 205,53,3 : REM CMP $335
270 :: : REM FINISH
280 DATA 96 : REM RTS
290 :
300 :: : REM TEST ROUTINE
310 DATA 32,0,192 : REM JSR $C000

```

```

320 DATA 8 : REM PHP
330 DATA 104 : REM PLA
340 DATA 41,3 : REM AND #$03
350 DATA 133,251 : REM STA $FB
360 DATA 96 : REM RTS
370 :
380 REM ** SET UP STRINGS FOR COMPARISON **
390 PRINT CHR$(147)
400 INPUT "FIRST STRING :";A$
410 FOR LOOP=1 TO LEN(A$)
420 TEMP$=MID$(A$,LOOP,1)
430 A=ASC(TEMP$)
440 POKE 50432+LOOP-1,A
450 NEXT LOOP
460 :
470 INPUT "SECOND STRING :";B$
480 FOR LOOP=1 TO LEN(B$)
490 TEMP$=MID$(B$,LOOP,1)
500 B=ASC(TEMP$)
510 POKE 50688+LOOP-1,B
520 NEXT LOOP
530 :
540 POKE 251,0 : POKE 252,197
550 POKE 253,0 : POKE 254,198
560 POKE 820,LEN(A$) : POKE 821,LEN(B$)
570 :
580 SYS TEST
590 :
600 PRINT "RESULT IS : ";PEEK(251)

```

Bytes reserved:

251-252	(\$FB-\$FC)	: address of first string
253-254	(\$FD-\$FE)	: address of second string
820	(\$334)	: length of first string
821	(\$335)	: length of second string

Once run, the BASIC text of lines 380 to 520 calls for two strings to be input. These are stored in memory from \$C500 and \$C600. Note that the routine cannot handle strings greater than 256 characters in length (though it could of course be expanded to do so). The length

of each string is also required by the routine, so this is ascertained and stored in the appropriate zero page bytes at \$334 and \$335 (line 560).

To allow the string buffers to be fully relocatable, the string addresses are held in two zero page vectors (lines 540 and 550).

String comparison proper starts by evaluating the length bytes to find out if they are the same length. If they are not equal, then the strings cannot be identical. However, as the routine returns information about the lengths of the strings it is still completed—in this case the program compares bytes through the length of the smaller of the two strings.

Byte comparison is performed by lines 170 to 190, using post-indexed indirect addressing. On the first non-equal characters the main loop is exited to FINISH. Assuming the entire comparison works, and the X register, which holds the working string length, has been decremented to zero, the length bytes (lines 250 and 260) are compared to condition the Zero and Carry flags before the routine completes.

The short test routine returns the Zero and Carry flag values and prints them out, indicating the following results:

Returned	Z	C	Result
Ø	Ø	Ø	Strings <> and string 1 larger
1	Ø	1	Strings <> and string 2 larger
3	1	1	Strings =

Line-by-line

A line-by-line description of Program 8 follows:

```
line 100 : get length of first string
line 110 : is it the same length as the second string?
line 120 : no, it's longer, so branch to COMPARE-STRING
line 130 : yes, so get length of second string
line 140 : entry for COMPARE-STRING
line 150 : if zero, branch to CONDITION-FLAGS
line 160 : initialize indexing register
line 170 : entry for COMPARE-BYTES
line 180 : get character from first string
line 190 : compare to same character in second string
line 200 : if dissimilar, branch to FINISH
line 210 : increment index
```

```
line 220 : decrement string counter
line 230 : branch back to COMPARE-BYTES until zero
line 240 : entry for CONDITION-FLAG
line 250 : get length of first string
line 260 : compare with length of the second string
line 270 : entry for finish
line 280 : back to calling routine
line 300 : entry for TEST routine
line 310 : push status onto stack
line 320 : pull into accumulator
line 330 : save Z and C
line 340 : save at location $FB
line 350 : back to BASIC
```

STRINGS UNITE

Strings may be joined together by a process called 'concatenation'. In BASIC the addition operator '+' performs this function. Thus the program:

```
100 A$="REM"
110 B$="ARK"
120 C$=A$+B$
```

assigns the string 'REMARK' to the string C\$. If line 120 were rewritten as:

```
120 C$=B$+A$
```

the resultant value assigned to C\$ would be 'ARKREM'. We can see from this that one string is simply tagged on to the end of the other, overwriting the former's RETURN character, but preserving the latter's.

This process of concatenation can be performed quite readily as Program 9 illustrates. However, the actual BASIC equivalent of the operation we are performing here is:

```
A$=A$+B$
```

In other words, we are adding the second string on to the first string, rather than summing the two to give a separate final string, although this is possible with slight modifications to the assembler text.

Program 9

```
10 REM ** STRING CONCATENATION **
20 CODE=49152
30 FOR LOOP=0 TO 96
40 READ BYTE
50 POKE CODE+LOOP,BYTE
60 NEXT LOOP
70 :
80 REM ** M/C DATA **
90 :: REM STRING-CONCATENATION
100 DATA 173,52,3 : REM LDA $334
110 DATA 141,54,3 : REM STA $336
120 DATA 169,0 : REM LDA #$00
130 DATA 141,55,3 : REM STA $337
140 DATA 24 : REM CLC
150 DATA 173,53,3 : REM LDA $335
160 DATA 109,52,3 : REM ADC $334
170 DATA 176,3 : REM BCS $03
180 DATA 76,45,192 : REM JMP $C02D
190 :: REM TOO-LONG
200 DATA 169,255 : REM LDA #$FF
210 DATA 141,57,3 : REM STA $339
220 DATA 56 : REM SEC
230 DATA 237,52,3 : REM SBC $334
240 DATA 144,51 : REM BCC $33
250 DATA 141,56,3 : REM STA $338
260 DATA 169,255 : REM LDA #$FF
270 DATA 141,52,3 : REM STA $334
280 DATA 76,59,192 : REM JMP $C03B
290 :: REM GOOD-LENGTH
300 DATA 141,52,3 : REM STA $334
310 DATA 169,0 : REM LDA #$00
320 DATA 141,57,3 : REM STA $339
330 DATA 173,53,3 : REM LDA $335
340 DATA 141,56,3 : REM STA $338
350 :: REM CONCATENATION
360 DATA 173,56,3 : REM LDA $338
370 DATA 240,21 : REM BEQ $15
380 :: REM LOOP
```

```
390 DATA 172,55,3 : REM LDY $337
400 DATA 177,253 : REM LDA ($FD),Y
410 DATA 172,54,3 : REM LDY $336
420 DATA 145,251 : REM STA ($FB),Y
430 DATA 238,54,3 : REM INC $336
440 DATA 238,55,3 : REM INC $337
450 DATA 206,56,3 : REM DEC $338
460 DATA 208,235 : REM BNE $EB
470 :: REM FINISHED
480 DATA 172,52,3 : REM LDY $334
490 DATA 169,13 : REM LDA #$0D
500 DATA 145,251 : REM STA ($FB),Y
510 DATA 173,57,3 : REM LDA $339
520 DATA 106 : REM ROR A
530 DATA 96 : REM RTS
540 :
600 PRINT CHR$(147)
610 INPUT "FIRST STRING ";A$
620 INPUT "SECOND STRING ";B$
630 :
640 F=49664 : REM $C200
650 S=49920 : REM $C300
660 :
670 FOR LOOP=1 TO LEN(A$)
680 TEMP$=MID$(A$,LOOP,1)
690 A=ASC(TEMP$)
700 POKE F+LOOP-1,A
710 NEXT LOOP
720 :
730 FOR LOOP=1 TO LEN(B$)
740 TEMP$=MID$(B$,LOOP,1)
750 B=ASC(TEMP$)
760 POKE S+LOOP-1,B
770 NEXT LOOP
780 :
790 POKE 251,0      POKE 252,194
800 POKE 253,0 : POKE 254,195
810 POKE 820,LEN(A$)
```

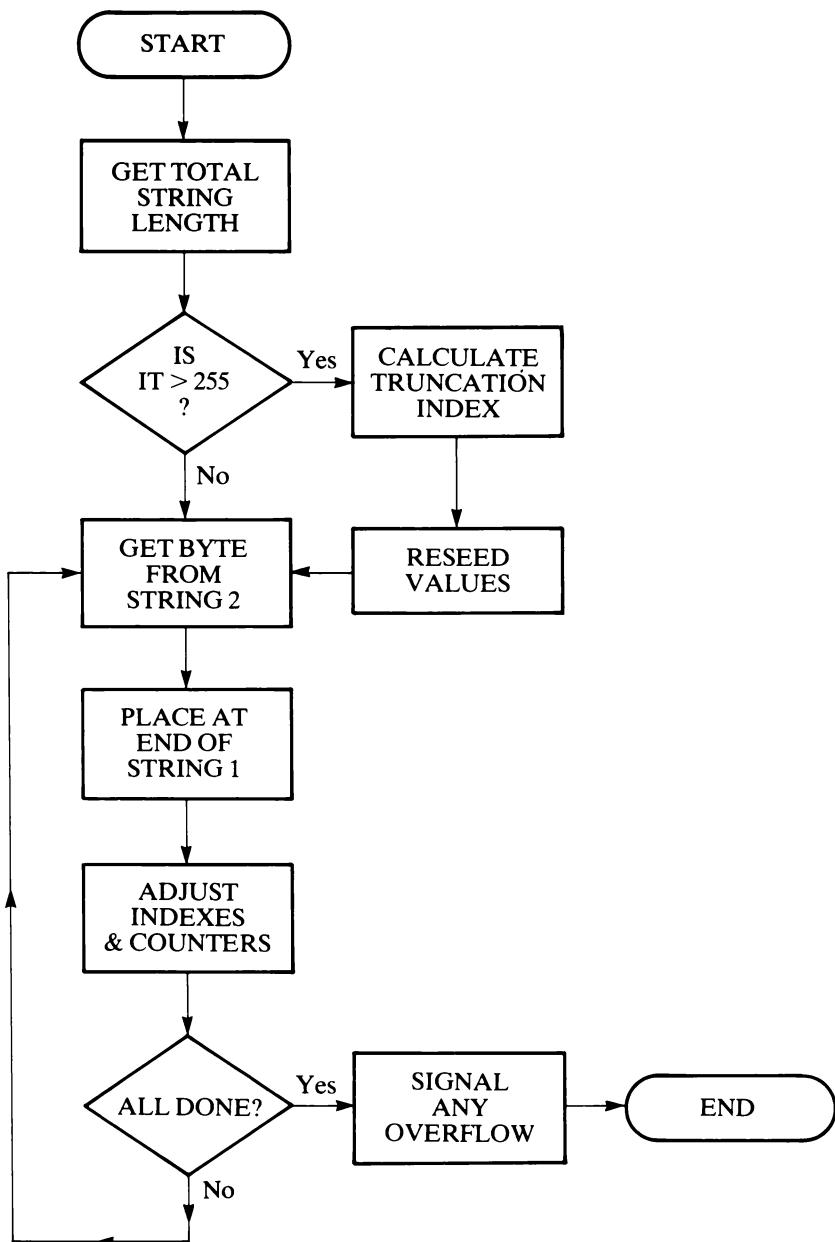


Figure 5.2 Concatenate strings flowchart

82 \emptyset POKE 821, LEN(B\$)

83 \emptyset :

84 \emptyset SYS CODE

85 \emptyset :

```

860 REM *** PRINT OUT FINAL STRING ***
870 PRINT "FINAL STRING IS :";
880 LOOP=0
890 REM ** REPEAT **
900 BYTE=PEEK(F+LOOP)
910 PRINT CHR$(BYTE);
920 LOOP=LOOP+1
930 IF BYTE=13 THEN END
940 GOTO 900

```

This program allows a final string of 256 characters in length to be manipulated. Therefore, as the program stands, the combined lengths of the two strings should not exceed this length. If they do, then only as many characters as space allows will be concatenated on to the first string, leaving the second string truncated. The Carry flag is used to signal whether any truncation has taken place, being set if it has and cleared otherwise. As with the string comparison routine, the string buffers are accessed via two zero page vectors (lines 790 and 800) and two bytes are reserved to hold the length of each string. A further two bytes are used to save index values.

The first nine machine code operations (lines 100 to 180) determine the final length of the string, by adding the length of the first string to that of the second string. A sum greater than 256 is signalled in the Carry flag and the branch of line 170 is performed, in which case the number of characters which can be inserted into the first string buffer is ascertained. The overflow indicator is loaded with \$FF if a truncation occurs; otherwise it is cleared with \$00.

The concatenating loop is held between lines 350 and 460. This simply moves a byte from the vectored address plus the index of the second string and places it at the end of the first string, as pointed to by the first string index byte. This process is reiterated until the value of 'count' has reached zero. Lines 480 and 500 place a RETURN character at the end of the string to facilitate printing from BASIC or machine code. The Overflow flag is loaded into the accumulator and bit 7 rotated across into the Carry flag, thereby signalling whether truncation has occurred. Lines 610 to 770 hold the BASIC test routine that reads in and then pokes the character strings into memory at \$C200 and \$C300. After the SYS call (line 840), the final BASIC routine prints the concatenated string from memory.

Project

Adapt the program to perform the BASIC equivalent of C\$=A\$+B\$ or C\$=B\$+A\$ on request.

Line-by-line

A line-by-line description of Program 9 now follows:

```
line 100 : get first string's length
line 110 : string one's index
line 120 : clear accumulator
line 130 : set string two's index to zero
line 140 : clear Carry flag
line 150 : get second string's length
line 160 : and add to length of first string
line 170 : branch to TOO-LONG if total greater than 256 bytes
line 180 : otherwise jump to GOOD-LENGTH
line 190 : entry for TOO-LONG
line 200 : load accumulator with 255
line 210 : and store to indicate overflow
line 220 : set Carry flag and subtract
line 230 : string one's length from maximum length
line 240 : branch to FINISH if first string is greater than
           256 bytes in length
line 250 : save current count
line 260 : restore maximum length
line 270 : store in string one's length
line 280 : jump to concatenation routine
line 290 : entry for GOOD-LENGTH
line 300 : save accumulator in string one's length
line 310 : load with 0 to clear
line 320 : overflow indicator
line 330 : get string two's length
line 340 : save in count
line 350 : entry for CONCATENATION
line 360 : get count value
line 370 : if zero, then finish
line 380 : entry for LOOP
line 390 : get index for string two
line 400 : and get character from second string
line 410 : get string one's index
line 420 : and place character into first string
line 430 : increment first string's index
line 440 : increment second string's index
line 450 : decrement count
```

```
line 46Ø : branch to LOOP until count=Ø
line 47Ø : entry for FINISHED
line 48Ø : get final length of first string
line 49Ø : load accumulator with ASCII return
line 50Ø : place at end of string
line 51Ø : get overflow indicator
line 52Ø : and move it into Carry flag
line 53Ø : back to calling routine
```

COPY·CAT

String manipulation routines must include a method of copying substrings of characters from anywhere within a string of characters. In BASIC, three such commands are provided. They are MID\$, LEFT\$ and RIGHT\$, although with the first of these, any point in a string can be accessed. The following shows the sort of thing possible in BASIC:

```
1ØØ A$="CONCATENATE"
11Ø B$=MID$(A$,Ø,3)
12Ø PRINT B$
```

Running this will output the string 'CON'. What the code has done is to take the three characters from the first character in the Main\$. Program 10 produces the same type of operation from machine code.

Program 10

```
1Ø REM ** COPY A SUBSTRING FROM WITHIN **
2Ø REM ** A MAIN ASCII STRING **
3Ø CODE=49152
4Ø MAIN=5Ø432      : REM $C5ØØ
5Ø SUB=5Ø688       : REM $C6ØØ
6Ø REM ** READ AND POKE M/C DATA **
7Ø FOR LOOP=Ø TO 123
8Ø READ BYTE
9Ø POKE CODE+LOOP,BYTE
1ØØ NEXT LOOP
11Ø :
12Ø REM ** M/C DATA **
13Ø DATA 16Ø,Ø      : REM LDY #$ØØ
14Ø DATA 14Ø,52,3   : REM STY $334
```

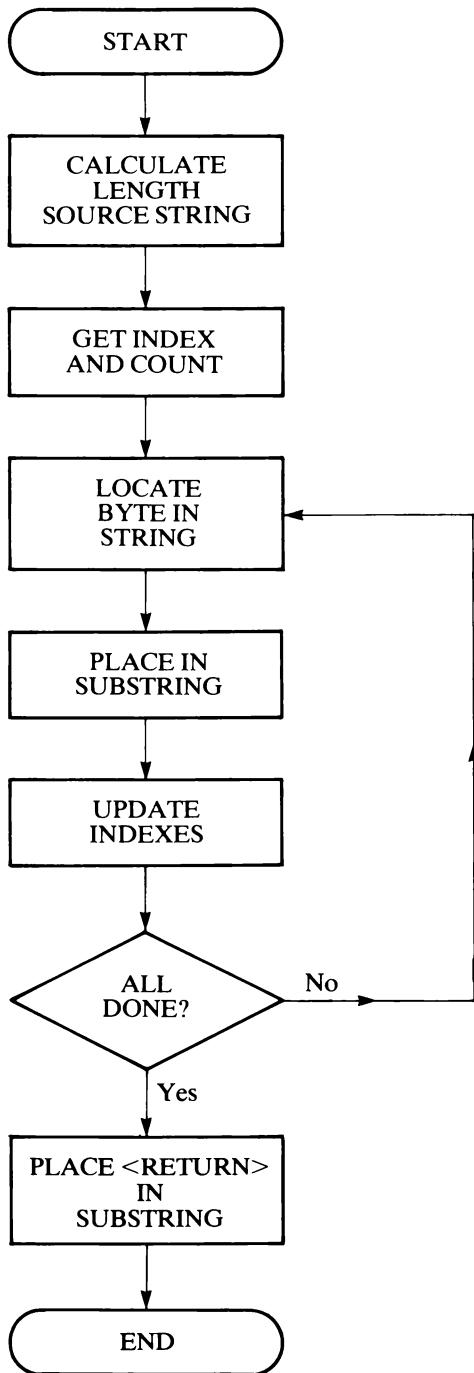


Figure 5.3 Copy string flowchart

```

150 DATA 140,56,3 : REM STY $338
160 DATA 173,54,3 : REM LDA $336
170 DATA 240,98     : REM BEQ $62

```

180 DATA 173,53,3 : REM LDA \$335
190 DATA 205,55,3 : REM CMP \$337
200 DATA 144,93 : REM BCC \$5D
210 DATA 24 : REM CLC
220 DATA 173,55,3 : REM LDA \$337
230 DATA 109,54,3 : REM ADC \$336
240 DATA 176,9 : REM BCS \$09
250 DATA 170 : REM TAX
260 DATA 202 : REM DEX
270 DATA 236,53,3 : REM CPX \$335
280 DATA 144,20 : REM BCC \$14
290 DATA 240,18 : REM BEQ \$12
300 :: REM TRUNCATION
310 DATA 56 : REM SEC
320 DATA 173,53,3 : REM LDA \$335
330 DATA 237,55,3 : REM SBC \$337
340 DATA 141,54,3 : REM STA \$336
350 DATA 238,54,3 : REM INC \$336
360 DATA 169,255 : REM LDA #\$FF
370 DATA 141,56,3 : REM STA \$338
380 :: REM GREATER-EQUAL
390 DATA 173,54,3 : REM LDA \$336
400 DATA 201,255 : REM CMP #\$FF
410 DATA 144,10 : REM BCC \$0A
420 DATA 240,8 : REM BEQ \$08
430 DATA 169,255 : REM LDA #\$FF
440 DATA 141,54,3 : REM STA \$336
450 DATA 141,56,3 : REM STA \$338
460 :: REM COPY-SUBSTRING
470 DATA 174,54,3 : REM LDX \$336
480 DATA 240,35 : REM BEQ \$23
490 DATA 169,0 : REM LDA #\$00
500 DATA 141,52,3 : REM STA \$334
510 :: REM LOOP
520 DATA 172,55,3 : REM LDY \$337
530 DATA 177,251 : REM LDA (\$FB),Y
540 DATA 172,52,3 : REM LDY \$334
550 DATA 145,253 : REM STA (\$FD),Y
560 DATA 238,55,3 : REM INC \$337
570 DATA 238,52,3 : REM INC \$334

```
580 DATA 202 : REM DEX
590 DATA 208,237 : REM BNE $ED
600 DATA 206,52,3 : REM DEC $334
610 DATA 173,56,3 : REM LDA $338
620 DATA 208,3 : REM BNE $03
630 :: REM FINISH
640 DATA 24 : REM CLC
650 DATA 144,1 : REM BCC $01
655 :: REM ERROR
660 DATA 56 : REM SEC
670 :: REM OUT
680 DATA 169,13 : REM LDA #$0D
690 DATA 172,52,3 : REM LDY $334
700 DATA 200 : REM INY
710 DATA 145,253 : REM STA ($FD),Y
720 DATA 96 : REM RTS
730 :
740 REM ** SET UP MAIN STRING **
750 PRINT CHR$(147)
760 :: REM ERROR
770 INPUT "MAIN STRING ";B$
780 FOR LOOP=1 TO LEN(B$)
790 TEMP$=MID$(B$,LOOP,1)
800 B=ASC(TEMP$)
810 POKE MAIN+LOOP-1,B
820 NEXT LOOP
830 :
840 INPUT"INDEX INTO STRING ";X
850 INPUT"NUMBER OF BYTES TO COPY ";Y
860 :
870 REM ** SET UP BYTES FOR M/C **
880 POKE 251,0 : POKE 252,197
: REM $C500 VECTOR
890 POKE 253,0 : POKE 254,198
: REM $C600 VECTOR
900 POKE 821,LEN(B$)
910 POKE 822,Y
920 POKE 823,X
930 :
940 SYS CODE
```

```

950 :
960 REM ** READ COPIED SUBSTRING **
970 FOR LOOP=1 TO Y
980 Z=PEEK(SUB+LOOP-1)
990 PRINT CHR$(Z);
1000 NEXT LOOP

```

Bytes are designated as follows:

251-252	(\$FB-\$FC)	: main string vector
253-254	(\$FD-\$FE)	: substring vector
820	(\$334)	: length of substring
821	(\$335)	: length of main string
822	(\$336)	: number of bytes to be copied
823	(\$337)	: index into main string
824	(\$338)	: error flag

Once again, a few lines of BASIC demonstrate the operation of the routine, requesting the source string, starting index and length of substring, or rather the number of bytes to be copied into the substring from the starting index. The main string is in a buffer located at \$C500 and the substring is copied into its own buffer at \$C600. As always, these addresses may be changed to suit user needs, as they are vectored through zero page (lines 880 and 890).

Error-checking is allowed, as the Carry flag is set on exit if an error has occurred. Normally, an error will occur only if the starting index is beyond the length of the source string, or the number of bytes to be copied from the main string is zero. If the number of bytes requested in the length exceeds the number left from the indexed position to the end of the main string, then only the bytes available will be copied to the substring buffer.

On entry to the routine, error-checking is performed (lines 160 to 240) and if any are found, the program exits. Lines 300 to 370 perform a truncation if the number of bytes to be copied exceeds those available. The COPY-SUBSTRING loop (lines 460 to 590) copies each string byte from the vectored address in the main string to the substring buffer. Each time a character is copied, the substring length byte is incremented. On completion of this loop, controlled by the X register, the error flag is restored and the Carry flag conditioned accordingly (lines 610 to 660). Finally (lines 690 to 730), an ASCII RETURN character is placed at the end of the substring.

The following example shows the resultant substrings produced from the main string 'CONCATENATE' for different indexes. Figure 5.4 illustrates the index value for each of the main string's characters.

Index	Length	Substring
Ø	3	CON
3	3	CAT
4	3	ATE

String	C	O	N	C	A	T	E	N	A	T	E
Index	Ø	1	2	3	4	5	6	7	8	9	1Ø

Figure 5.4 String Index

Line-by-line

A line-by-line description of Program 10 follows:

```

line 13Ø : initialize Y register
line 14Ø : clear substring length
line 15Ø : and error flag
line 16Ø : get substring length
line 17Ø : if a null string, branch to FINISH
line 18Ø : get main string's length
line 19Ø : compare it with index byte
line 20Ø : branch to ERROR if index is greater
line 21Ø : clear the Carry flag
line 22Ø : get index
line 23Ø : add it to substring length
line 24Ø : branch to TRUNCATION if result is greater than 255
line 25Ø : move index across into X register
line 26Ø : decrement it by one
line 27Ø : compare result with string length
line 28Ø : branch to GREATER-EQUAL if result is
line 29Ø : greater than or equal to string length
line 30Ø : entry for TRUNCATION
line 31Ø : set the Carry flag
line 32Ø : get string length
line 33Ø : subtract the index from it
line 34Ø : save the new length
line 35Ø : and increment it by one

```

line 360 : denote an error by
line 370 : setting the error flag
line 380 : entry for GREATER-EQUAL
line 390 : get length into accumultor
line 400 : compare with maximum length
line 410 : branch if count is
line 420 : greater or equal to maximum length
line 430 : put maximum length in accumulator
line 440 : store in bytes to copy
line 450 : and also in error flag
line 460 : entry for COPY-SUBSTRING
line 470 : get the index position
line 480 : branch to ERROR if zero
line 490 : clear accumulator
line 500 : and substring length
line 510 : entry for LOOP
line 520 : get main string index into Y register
line 530 : get character from main string
line 540 : get substring index
line 550 : copy character into substring
line 560 : increment main string index
line 570 : increment substring index
line 580 : decrement bytes to move counter
line 590 : branch to LOOP if still bytes to be copied
line 600 : decrement final substring count
line 610 : get error flag into accumulator
line 620 : branch to ERROR if not zero
line 630 : FINISH entry
line 640 : clear Carry flag as no error
line 650 : branch to OUT
line 655 : entry for ERROR
line 660 : set Carry flag to indicate error
line 670 : entry for OUT
line 680 : place RETURN in accumlator
line 690 : get substring index into Y
line 700 : increment Y
line 710 : place RETURN at end of substring
line 720 : return to BASIC.

INSERTION

This final routine provides the facility for inserting a string within the body of another string, allowing textual material—for example, in word processing applications—to be manipulated. If the main string held 'ELIZABETH OKAY', this routine could be called to insert the string 'RULES', so that the final string would read 'ELIZABETH RULES OKAY'. As with the COPY routine, the position of the insertion is pointed to by an index byte, and the Carry flag is set if an error is detected—that is, if an index of 0 or a null substring is specified.

The maximum length of the final string is 256 characters. If the insertion of the substring would cause this length to be exceeded, the substring is truncated to the length given by (256 minus length of main string) and only these characters are inserted.

As always, a BASIC primer demonstrates the routine's use. The string buffers are held at \$C500 and \$C600 and in this instance they are accessed directly, although there is no reason why vectored addresses could not be used.

Program 11

```
10 REM ** INSERT ONE ASCII STRING **
20 REM ** INTO ANOTHER ASCII STRING **
30 MAIN=50432           : REM $C500
40 SUB=50688            : REM $C600
50 CODE=49152
60 REM ** READ AND POKE DATA **
70 FOR LOOP=0 TO 141
80 READ BYTE
90 POKE LOOP+CODE,BYTE
100 NEXT LOOP
110 :
120 REM ** M/C DATA**
130 DATA 160,0           : REM LDY #0
140 DATA 140,53,3        : REM STY $335
150 DATA 165,252          : REM LDA $FC
160 DATA 208,3            : REM BNE $03
170 DATA 76,137,192       : REM JMP $C089
180 ::                   REM ZERO-LENGTH
190 DATA 165,253          : REM LDA $FD
200 DATA 240,124          : REM BEQ $7C
210 DATA ::                REM CHECK
```

220	DATA	24	:	REM	CLC
230	DATA	165,252	:	REM	LDA \$FC
240	DATA	101,251	:	REM	ADC \$FB
250	DATA	176,6	:	REM	BCS \$06
260	DATA	201,255	:	REM	CMP #\$FF
270	DATA	240,18	:	REM	BEQ \$12
280	DATA	144,16	:	REM	BCC \$10
290	:	:		REM	CUT-OFF
300	DATA	169,255	:	REM	LDA #\$FF
310	DATA	56	:	REM	SEC
320	DATA	229,251	:	REM	SBC \$FB
330	DATA	240,104	:	REM	BEQ \$68
340	DATA	144,102	:	REM	BCC \$66
350	DATA	133,252	:	REM	STA \$FC
360	DATA	169,255	:	REM	LDA #\$FF
370	DATA	141,53,3	:	REM	STA \$335
380	:	:		REM	CALC-LENGTH
390	DATA	165,251	:	REM	LDA \$FB
400	DATA	197,253	:	REM	CMP \$FD
410	DATA	176,20	:	REM	BCS \$14
420	DATA	166,251	:	REM	LDX \$FB
430	DATA	232	:	REM	INX
440	DATA	134,253	:	REM	STX \$FD
450	DATA	169,255	:	REM	LDA #\$FF
460	DATA	141,53,3	:	REM	STA \$335
470	DATA	24	:	REM	CLC
480	DATA	165,251	:	REM	LDA \$FB
490	DATA	101,252	:	REM	ADC \$FC
500	DATA	133,251	:	REM	STA \$FB
510	DATA	76,109,192	:	REM	JMP \$C06D
520	:	:		REM	NO-PROBLEMS
530	DATA	56	:	REM	SEC
540	DATA	165,251	:	REM	LDA \$FB
550	DATA	229,253	:	REM	SBC \$FD
560	DATA	170	:	REM	TAX
570	DATA	232	:	REM	INX
580	DATA	165,251	:	REM	LDA \$FB
590	DATA	133,254	:	REM	STA \$FE
600	DATA	24	:	REM	CLC
610	DATA	101,252	:	REM	ADC \$FB

620 DATA 133,251 : REM STA \$FB
630 DATA 141,52,3 : REM STA \$334
640 :: REM MAKE-SPACE
650 DATA 164,254 : REM LDY \$FE
660 DATA 185,0,197 : REM LDA \$C500,Y
670 DATA 172,52,3 : REM LDY \$334
680 DATA 153,0,197 : REM STA \$C500,Y
690 DATA 206,52,3 : REM DEC \$334
700 DATA 198,254 : REM DEC \$FE
710 DATA 202 : REM DEX
720 DATA 208,237 : REM BNE \$ED
730 :: REM INSERT-SUBSTRING
740 DATA 169,0 : REM LDA #\$00
750 DATA 133,254 : REM STA \$FE
760 DATA 166,252 : REM LDX \$FC
770 :: REM TRANSFER
780 DATA 164,254 : REM LDY \$FE
790 DATA 185,0,198 : REM LDA \$C600,Y
800 DATA 164,253 : REM LDY \$FE
810 DATA 153,0,197 : REM STA \$C500,Y
820 DATA 230,253 : REM INC \$FD
830 DATA 230,254 : REM INC \$FE
840 DATA 202 : REM DEX
850 DATA 208,239 : REM BNE \$EF
860 DATA 173,53,3 : REM LDA \$335
870 DATA 208,3 : REM BNE \$03
880 :: REM GOOD
890 DATA 24 : REM CLC
900 DATA 144,1 : REM BCC \$01
910 :: REM ERROR
920 DATA 56 : REM SEC
930 :: REM FINISH
940 DATA 96 : REM RTS
950 :
960 REM ** GET MAIN STRING AND STORE AT
 \$C500 **
970 PRINT CHR\$(147)
980 INPUT"MAIN STRING";B\$
990 FOR LOOP=1 TO LEN(B\$)
1000 TEMP\$=MID\$(B\$,LOOP,1)

```

1010  B=ASC(TEMP$)
1020  POKE MAIN+LOOP-1,B
1030  NEXT LOOP
1040  :
1050  REM ** GET SUBSTRING AND STORE AT $C600 **
1060  INPUT"SUB STRING";C$
1070  FOR LOOP=1 TO LEN(C$)
1080  TEMP$=MID$(C$,LOOP,1)
1090  B=ASC(TEMP$)
1100  POKE SUB+LOOP-1,B
1110  NEXT LOOP
1120  :
1130  REM ** GET INSERTION INDEX **
1140  INPUT"INSERTION INDEX"; X
1150  :
1160  REM ** POKE VALUES INTO ZERO PAGE **
1170  POKE 251,LEN(B$)
1180  POKE 252,LEN(C$)
1190  POKE 253,X
1200  :
1210  SYS CODE
1220  :
1230  REM ** READ FINAL STRING **
1240  COUNT=LEN(B$)+LEN(C$)-1
1250  FOR LOOP=0 TO COUNT
1260  Z=PEEK(MAIN+LOOP)
1270  PRINT CHR$(Z);
1280  NEXT LOOP

```

The program begins by checking the length bytes to ensure that no null strings are present (lines 150 to 200) and then sums the two lengths to obtain the final length. If the addition results in the Carry flag being set (line 250), the total length will exceed 256 bytes and, as a result, the inserted substring will be truncated (lines 310 to 390).

If the insertion index is greater than the length of the string, the substring is actually concatenated on to the end of the main string. This evaluation is performed through lines 400 to 530. Before inserting the substring, all characters to the left of the index must be shuffled up through memory to make space for it. These calculations are carried out in lines 550 to 650, ready for the shuffling process (lines 660 to 740). Inserting the substring now involves simply copying it from its buffer into the space opened up for it

(lines 750 to 870), the X register being used as the characters-moved counter.

Finally, the error flag is restored and the Carry flag conditioned to signal any errors.

Line-by-line

A line-by-line description of Program 11 follows:

```
line 130 : clear indexing register
line 140 : clear error flag
line 150 : get substring length
line 160 : branch to ZERO-LENGTH if Z=0
line 170 : otherwise carry on
line 180 : entry for ZERO-LENGTH
line 190 : get offset
line 200 : branch to ERROR if Z=1
line 210 : entry for CHECK
line 220 : clear Carry flag
line 230 : get substring length
line 240 : add it to main string length
line 250 : branch to CUT-OFF if greater than 256
line 260 : is it maximum length?
line 270 : branch to CALC-LENGTH if
line 280 : it is equal to or greater than
line 290 : entry for CUT-OFF
line 300 : get the maximum length allowed
line 310 : set Carry flag
line 320 : subtract length of string
line 330 : branch to ERROR if
line 340 : length is equal to or greater than string
line 350 : save characters free
line 360 : set error flag
line 380 : entry for CALC-LENGTH
line 390 : get main string length
line 400 : is offset within string?
line 410 : branch to NO-PROBLEMS if it is
line 420 : else place substring
line 430 : at end of main string
line 440 : save X in offset
line 450 : and flag the error
```

line 460 : in error flag byte
line 470 : clear Carry flag
line 480 : get length of string
line 490 : calculate total length
line 500 : and save result
line 510 : jump to INSERT-SUBSTRING
line 520 : entry for NO-PROBLEMS
line 530 : set Carry flag
line 540 : get length of substring
line 550 : subtract offset
line 560 : move index into X
line 570 : increment index
line 580 : get length
line 590 : save in source
line 600 : clear Carry flag
line 610 : find total length
line 620 : save result
line 630 : and for index
line 640 : entry for MAKE-SPACE
line 650 : get source index
line 660 : get byte from main
line 670 : get offset into string
line 680 : move byte along
line 690 : decrement both indexes
line 710 : decrement counter
line 720 : branch to MAKE-SPACE until done
line 730 : entry for INSERT-SUBSTRING
line 740 : clear accumulator
line 750 : and source
line 760 : get counter
line 770 : entry for TRANSFER
line 780 : get index
line 790 : get byte from substring
line 800 : get offset into main string
line 810 : and place byte in main
line 820 : increment both indexes
line 840 : do until substring inserted
line 850 : branch to TRANSFER
line 860 : get error flag
line 870 : branch to ERROR

line 88Ø : entry for GOOD
line 89Ø : signal no error
line 9ØØ : branch to FINISH
line 91Ø : entry for ERROR
line 92Ø : denote error
line 93Ø : entry for FINISH
line 94Ø : return to calling routine

6 Printing Print!

Every machine code program sooner or later requires text to be printed on to the screen. In most instances, this is a fairly simple process and often involves merely indexing into an ASCII string table and printing the characters, using one of the Operating System calls, until either a RETURN character or zero byte is encountered. Program 12 uses this method.

Program 12

```
1Ø REM ** PRINT STRING FROM MEMORY **
2Ø CODE=49152
3Ø FOR LOOP=Ø TO 13
4Ø READ BYTE
5Ø POKE CODE+LOOP,BYTE
6Ø NEXT LOOP
7Ø :
8Ø REM ** M/C DATA **
9Ø :: : REM STRING-PRINT
10Ø DATA 162,Ø : REM LDX #$ØØ
11Ø :: : REM NEXT-CHARACTER
12Ø DATA 189,Ø,197 : REM LDA $C5ØØ,X
13Ø DATA 32,21Ø,255 : REM JSR $FFD2
14Ø DATA 232 : REM INX
15Ø DATA 2Ø1,13 : REM CMP #$ØD
16Ø DATA 2Ø8,245 : REM BNE $F5
17Ø DATA 96 : REM RTS
18Ø :
19Ø REM ** GET STRING TO BE PRINTED **
2Ø STRING=5Ø432
```

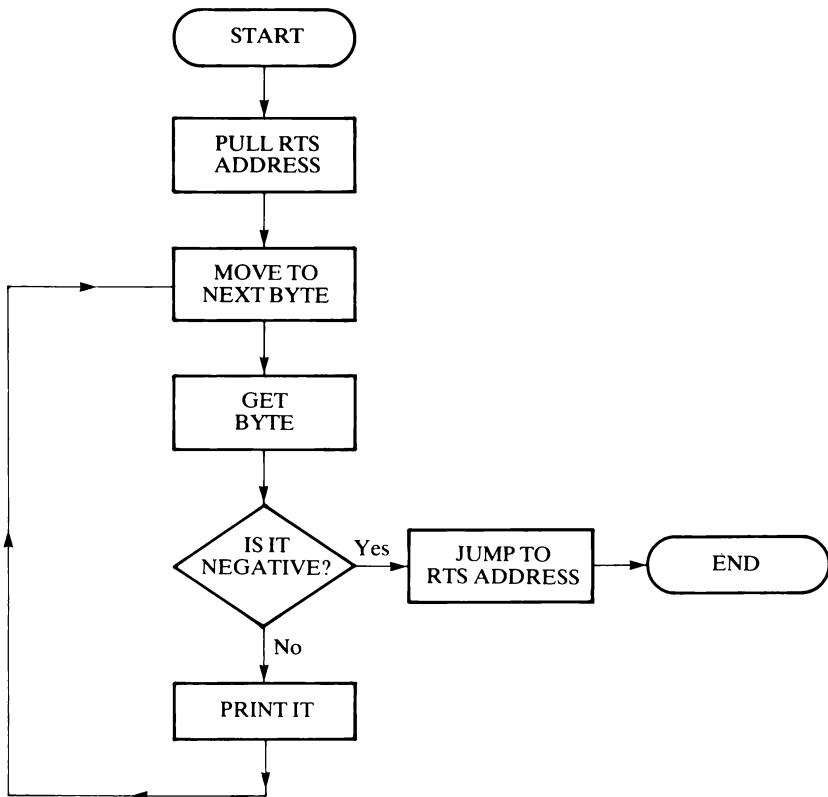


Figure 6.1 Printing embedded code flowchart

```

210 PRINT CHR$(147)
220 INPUT "INPUT STRING :";A$
230 FOR LOOP=1 TO LEN(A$)
240 TEMP$=MID$(A$,LOOP,1)
250 B=ASC(TEMP$)
260 POKE STRING+LOOP-1,B
270 NEXT LOOP
280 PRINT:PRINT
290 PRINT"YOUR STRING WAS AS FOLLOWS :";
300 SYS CODE
  
```

Here, a string buffer is located at \$C500 (50432) and the requirement for printing the string is that it must be terminated with an ASCII RETURN character, \$0D. The program begins by initializing an index, the X register (line 100), and loading the byte at \$C500+X into the accumulator. This is printed using the Kernal's CHROUT routine, the index is incremented and then the accumulator's contents are compared to see whether the character just output was a RETURN (line 150). If not, the loop branches back and the next character is sought.

Program 13 shows how several strings may be printed to the screen using a loop similar to that described above. The number of strings for printing may be variable, the desired number being passed into the routine via the Y register. The string data has been entered using the DATA statement. If a large amount of string data is to be stored, and the amount to be printed at any one time varied, a vectored address should be used to access the table. Positioning of the text on the screen can be performed by embedding the relative number of RETURNS and spaces into the DATA, or more neatly by using the Kernal's PLOT routine to set the X and Y tab co-ordinates.

Program 13

```
1Ø REM ** PRINT Y NUMBER OF STRINGS **
2Ø CODE=49152
3Ø FOR LOOP=Ø TO 18
4Ø READ BYTE
5Ø POKE CODE+LOOP,BYTE
6Ø NEXT LOOP
7Ø :
8Ø REM ** M/C DATA **
9Ø DATA 162,Ø : REM LDX #$ØØ
10Ø DATA 16Ø,4 : REM LDY #$Ø4
11Ø :: REM NEXT-CHARACTER
12Ø DATA 189,Ø,197 : REM LDA $C5ØØ,X
13Ø DATA 32,21Ø,255 : REM JSR $FFD2
14Ø DATA 232 : REM INX
15Ø DATA 2Ø1,13 : REM CMP #$ØD
16Ø DATA 2Ø8,245 : REM BNE $F5
17Ø DATA 136 : REM DEY
18Ø DATA 2Ø8,242 : REM BNE $F2
19Ø DATA 96 : REM RTS
2ØØ :
21Ø REM ** SET UP FOUR SIMPLE STRINGS **
22Ø STRING=5Ø432
23Ø FOR LOOP=Ø TO 31
24Ø READ BYTE
25Ø POKE STRING+LOOP,BYTE
26Ø NEXT
27Ø :
28Ø REM ** ASCII DATA **
```

```
290 DATA 32,65,65,65,65,65,65,13  
300 DATA 32,32,66,66,66,66,66,13  
310 DATA 32,32,32,67,67,67,67,13  
320 DATA 32,32,32,32,68,68,68,13
```

The final program in this chapter shows the way I find easiest to store and print character strings, stowing them directly within the machine code. The two main advantages of this method are that the string is inserted directly at the point it is needed, avoiding the need to calculate indexes into look-up tables, and that because it manipulates its own address it is fully relocatable.

Program 14

```
10 REM ** ASCII STRING OUTPUT ROUTINE **  
20 CODE=49152  
30 FOR LOOP=0 TO 26  
40 READ BYTE  
50 POKE CODE+LOOP,BYTE  
60 NEXT LOOP  
70 :  
80 REM ** M/C DATA **  
90 DATA 104 : REM PLA  
100 DATA 133,251 : REM STA $FB  
110 DATA 104 : REM PLA  
120 DATA 133,252 : REM STA $FC  
130 :: REM REPEAT  
140 DATA 160,0 : REM LDY #$0  
150 DATA 230,251 : REM INC $FB  
160 DATA 208,2 : REM BNE $02  
170 DATA 230,252 : REM INC $FC  
180 :: REM OVER  
190 DATA 177,251 : REM LDA ($FB),Y  
200 DATA 48,6 : REM BMI $06  
210 DATA 32,210,255 : REM JSR $FFD2  
220 DATA 76,6,192 : REM JMP $C006  
230 :: REM FINISH  
240 DATA 108,251,0 : REM JMP ($FB)  
250 :  
260 REM ** DEMO ROUTINE LOCATED AT $C200 **  
270 DEMO=49664
```

```

280 FOR LOOP=0 TO 38
290 READ BYTE
300 POKE DEMO+LOOP, BYTE
310 NEXT LOOP
320 :
330 REM ** DEMO M/C DATA **
340 DATA 169,147 : REM LDA #$93
350 DATA 32,210,255 : REM JSR $FFD2
360 DATA 32,0,192 : REM JSR $C000
370 REM ** NOW STORE ASCII CODES FOR PRINTING **
380 DATA 13 : REM CARRIAGE-RETURN
390 DATA 83,84,82,73,78,71,83,32
: REM STRINGS<SPACE>
400 DATA 87,73,84,72,73,78,32
: REM WITHIN<SPACE>
410 DATA 77,65,67,72,73,78,69,32
: REM MACHINE<SPACE>
420 DATA 67,79,68,69,33
: REM CODE!
430 DATA 234 : REM NOP
440 DATA 96 : REM RTS
450 :
460 SYS DEMO

```

The ASCII character string is placed in memory by leaving the machine code assembly (line 360) and POKEing the ASCII codes of the string directly into successive memory locations (lines 380 to 420).

For this routine to work, it is imperative that the first byte following the string is a negative byte—that is, one with bit 7 set. The opcode for NOP, \$EA, is ideal for this purpose as it has its most significant bit set (\$EA=11101010) and its only effect is to cause a very short delay.

The ASCII print routine is just 27 bytes in length and it should be called as a subroutine immediately before the string is encountered (line 360). On entry into the subroutine, the first four operations pull the return address from the stack and save it in a zero page vector at \$FB and \$FC. These bytes are then incremented by one to point at the byte following the subroutine call.

Because the string data follows on immediately after the ASCII print subroutine call, post-indexed indirect addressing can be used to load the first string character into the accumulator (line 190). The string terminating negative byte is tested for (line 200), and if not found the byte is printed with a CHROUT call. A JMP to

REPEAT is then performed and the loop reiterated. When the negative byte is encountered, and the branch of line 200 succeeds, an indirect jump (line 240) via the current vectored address is executed, returning control back to the calling machine code at the end of the ASCII string.

Line-by-line

A line-by-line description of Program 14 follows:

```
line  90 : set low byte RTS address
line 100 : save in $FB
line 110 : get high byte RTS address
line 120 : save in $FC
line 130 : entry for REPEAT
line 140 : initialize index to zero
line 150 : increment low byte of vectored address
line 160 : branch to OVER if not zero
line 170 : else increment page value
line 180 : entry for OVER
line 190 : get byte from within program
line 200 : if negative, branch to FINISH
line 210 : else print it
line 220 : jump to REPEAT
line 230 : entry for FINISH
line 240 : jump back into main program
line 340 : load accumulator with clear screen code
line 350 : and print it
line 360 : call string printing routine at $C000
line 380 : ASCII code for RETURN
line 390 : ASCII string 'STRINGS '
line 400 : ASCII string 'WITHIN '
line 410 : ASCII string 'MACHINE '
line 420 : ASCII string 'CODE!'
line 430 : negative byte
line 440 : back to BASIC
```

7 A Bubble of Sorts

Any program written to handle quantities of data will, at some time, require the data in a data table to be sorted into ascending or descending order. Several algorithms are available to facilitate this manipulation of data, of which the bubble sort is perhaps the simplest to implement in BASIC or machine code.

The technique involves moving through the data list and comparing pairs of bytes. If the first byte is smaller than the next byte in the list, the next pair of bytes is sought. If, on the other hand, the second byte is less than the first, the two bytes are swapped. This procedure is repeated until a pass is executed in which no elements are exchanged, so all are in ascending order. Program 15 is the BASIC version of such a bubble sort.

Program 15

```
10 REM ** BASIC BUBBLE SORT **
20 TABLE=828
30 FOR LOOP=0 TO 19
40 READ BYTE
50 POKE TABLE+LOOP,BYTE
60 NEXT LOOP
70 :
80 REM ** BUBBLE-UP ROUTINE **
90 FOR BUBBLE=0 TO 19
100 TEMP=BUBBLE
110 :
120 IF PEEK(TABLE+TEMP)>PEEK(TABLE+(TEMP-1))
    THEN GOTO 180
130 HOLD=PEEK(TABLE+TEMP)
140 POKE TABLE+TEMP,PEEK(TABLE+(TEMP-1))
```

```

150 POKE TABLE+(TEMP-1), HOLD
160 TEMP=TEMP-1
170 IF TEMP<>0 THEN GOTO 120
180 NEXT
190 :
200 REM ** DATA FOR SORTING **
210 DATA 1,255,67,89,120
220 DATA 6,200,85,45,199
230 DATA 0,123,77,98,231
240 DATA 9,234,99,98,100
250 :
260 REM ** PRINT SORTED DATA **
270 FOR LOOP=0 TO 19
280 PRINT PEEK(TABLE+LOOP)
290 NEXT LOOP

```

The data bytes for sorting are held within the four data lines from 210 to 240 and these are read into a memory array called TABLE. The sorting procedure is performed through lines 90 to 180, line 120 checking to see if a swap is required. If a swap is unnecessary, GOTO 180 is executed and the swap routine bypassed. If it is required, however, the GOTO statement is not encountered, and the swap is performed in lines 130 to 160. The byte currently being pointed to is PEEKed into the variable HOLD (line 130) and the next byte is PEEKed and then POKEed into the location immediately before it (line 140). The swap is completed by POKEing the value of HOLD into the now 'vacant' location. The variable TEMP is used to keep track of the number of passes through the loop.

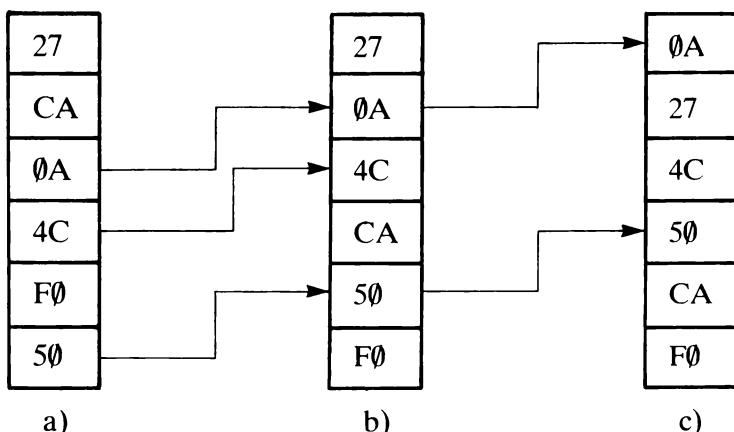


Figure 7.1 Numbers bubbling up

Figure 7.1 illustrates how small numbers bubble up through a data list using this sorting method. In this example, the data list consists of six numbers 27, CA, 0A, 4C, F0 and 50 (Figure 7.1a). After the first pass of the bubble sort three swaps have occurred (Figure 7.1b), thus:

1. $27 < CA$ therefore no change.
2. $CA > 0A$ therefore swap items.
3. $CA > 4C$ therefore swap items.
4. $CA < F0$ therefore no change.
5. $F0 > 50$ therefore swap items.

The next pass through the data list produces the ordered list of Figure 7.1c in which just two swaps occurred, as follows:

1. $27 > 0A$ therefore swap items.
2. $27 < 4C$ therefore no change.
3. $4C < 50$ therefore no change.
4. $CA > 50$ therefore swap items.
5. $CA < F0$ therefore no change.

All the data elements are now in their final order, so the next pass through the list will have no effect. We can signal this by using an exchange flag to indicate whether the last pass produced any swaps, the sort routine exiting when the flag is cleared. This detail is included in the BASIC loader listed below as Program 16.

Program 16

```
10 REM *** BUBBLE SORT ***
20 CODE=49152
30 TABLE=50432
40 FOR LOOP=0 TO 44
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 DATA 206,52,3      : REM DEC $334
110 ::                   : REM BUBBLE-LOOP
120 DATA 160,0          : REM LDY #$00
130 DATA 140,53,3       : REM STY $335
140 DATA 174,52,3       : REM LDX $334
150 ::                   : REM LOOP
```

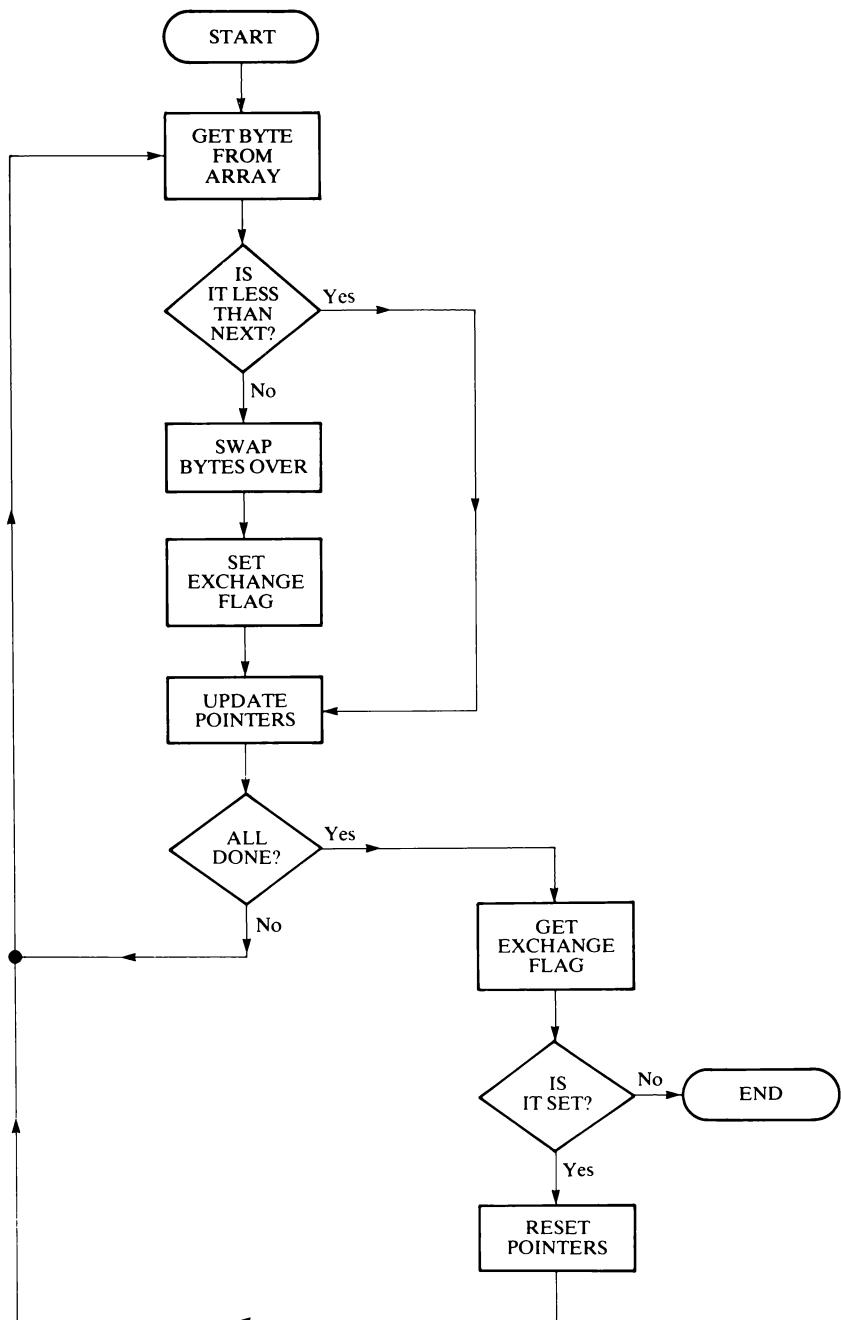


Figure 7.2 Bubble sort flowchart

```

160 DATA 177,253      : REM LDA ($FD),Y
170 DATA 209,251      : REM CMP ($FB),Y
180 DATA 176,13       : REM BCS $0D
190 DATA 72            : REM PHA
200 DATA 177,251      : REM LDA ($FB),Y
210 DATA 145,253      : REM STA ($FD),Y
220 DATA 104           : REM PLA
230 DATA 145,251      : REM STA ($FB),Y
240 DATA 169,1         : REM LDA #$01
250 DATA 141,53,3     : REM STA $335
260 ::                 REM SECOND-FIRST
270 DATA 200           : REM INY
280 DATA 202           : REM DEX
290 DATA 208,233       : REM BNE $E9
300 DATA 173,53,3     : REM LDA $335
310 DATA 240,5         : REM BEQ $05
320 DATA 206,52,3     : REM DEC $334
330 DATA 208,215       : REM BNE $D7
335 ::                 REM FINISH
340 DATA 96            : REM RTS
350 :
360 REM ** SET UP VECTORS **
370 REM $FB=$C500, $FD=$C501
380 POKE 251,0 : POKE 252,197
390 POKE 253,1 : POKE 253,197
400 :
410 REM ** SET UP SCREEN AND ARRAY **
420 PRINT CHR$(147)
430 PRINT "**** MACHINE CODE BUBBLE SORT ****"
440 PRINT:PRINT
450 INPUT"NUMBER OF ELEMENTS IN ARRAY ";N
460 POKE 820,N          : REM LENGTH OF ARRAY
                           AT $334
470 FOR LOOP=0 TO N-1
480 PRINT"INPUT ELEMENT ";LOOP+1;
490 INPUT A
500 POKE TABLE+LOOP,A
510 NEXT LOOP
520 :
530 REM ** CALL CODE THEN PRINT SORTED TABLE **

```

```

540 SYS CODE
550 PRINT" SORTED VALUES ARE AS FOLLOWS"
560 FOR LOOP=0 TO N-1
570 PRINT PEEK(TABLE+LOOP)
580 NEXT LOOP

```

After POKEing the machine code data into memory at \$C000, two zero page vectors are created to hold the address of the TABLE and TABLE+1 (lines 370 to 390). The program then requests (in BASIC!) the number of elements in the array, which should be a series of integer values less than 256. These are then POKEd into memory (lines 450 to 510). The machine code begins by decrementing the length of array byte by one.(line 100), because the last element in the array will have no element beyond it to swap with. The swap flag is then cleared (line 130) and the main loop entered using the X register to count the iterations.

The LOOP begins by loading the data byte into the accumulator (line 160) and comparing it with the one immediately preceding it. If the byte+1 is greater than the byte, the Carry flag will be set and no swap required, in which case the branch to SECOND-FIRST is executed (line 180).

If a swap is required, the second byte is saved, pushing it on to the hardware stack. The first byte is then transferred to the second byte's position (lines 200 and 210) and the accumulator is restored from the stack and transferred to the position of the first byte (lines 220 to 230). To denote that a swap has occurred, the swap flag is set (lines 240 and 250). The index and counters are then adjusted (lines 270 and 280) and the loop continues until all the array elements have been compared. Upon completion of a full pass through the array, the swap flag is checked. If it is clear, no exchanges took place during the last pass, so the data list is now ordered and the sort finished (line 300 and 310). If the flag is set, the length of array byte is decremented and the procedure repeated once more (lines 320 and 330). On return from the SYS call, the now ordered list is printed out to the screen.

Line-by-line

A line-by-line description of Program 16 now follows:

- line 100 : subtract one from the length of the array
- line 110 : entry for BUBBLE-LOOP
- line 120 : initialize indexing register
- line 130 : clear the swap flag
- line 140 : get the array size into the X register to act as a loop counter

line 150 : entry for LOOP
line 160 : get the byte at the byte+1 position
line 170 : compare it with the previous byte
line 180 : branch to SECOND-FIRST if the second byte
 (byte+1) is larger than the first (byte)
line 190 : save accumulator on hardware stack
line 200 : get first byte at 'byte' position
line 210 : place in current location (byte+1)
line 220 : restore accumulator
line 230 : and complete swap of bytes
line 240 : load accumulator with 1
line 250 : and set the swap flag to denote that a swap has been
 performed
line 260 : entry for SECOND-FIRST
line 270 : move index on to next byte
line 280 : decrement loop counter
line 290 : branch to LOOP until done
line 300 : get the swap flag into the accumulator
line 310 : if clear, branch to FINISH
line 320 : decrement outer counter
line 330 : branch to BUBBLE-LOOP until all done
line 335 : entry to FINISH
line 340 : back to calling routine

Projects

Rewrite the BASIC sections of the program to make it a complete machine code routine.

Adapt the sorting routine to handle 16-bit numbers.

8 Software Stack

One of the criticisms of the 6510 processor is that it has a very limited set of operation instructions—only 56, though addressing modes extend this to 152 functions. With some thought, however, it is possible to implement operations present on other processors, such as the Z80 or 6809, and build up a set of very useful subroutines which can ultimately be strung together to perform quite sophisticated operations, as well as making the conversion of programs written for other processors much easier.

The routine described below mimics an instruction in the 6809 instruction set which allows the contents of up to eight registers to be pushed on to a stack in memory. This stack is often known as the user stack. I said ‘up to eight registers’, because the ones to be pushed can be selected, this being determined by the bit pattern of the byte after the user stack subroutine call. But more of that in a moment. First, which registers are we going to push? Obviously all the processor registers: the Program Counter, Status register, accumulator, and Index registers. The three remaining ones, we will implement as three two-byte ‘pseudo-registers’ from the user area of zero page. These are:

```
PR1 : $80 and $81  
PR2 : $82 and $83  
PR3 : $84 and $84
```

This now enables us to save the contents of these locations when required.

As already stated, the byte after the user stack subroutine call determines by its bit pattern which registers are to be pushed, as follows:

```
bit 0 : pseudo-register 1  
bit 1 : pseudo-register 2  
bit 2 : pseudo-register 3
```

```
bit 3 : Y register  
bit 4 : X register  
bit 5 : accumulator  
bit 6 : Status register  
bit 7 : Program Counter
```

The rule here is that if the bit is set, the related register is pushed. Thus the instructions:

```
JSR USER-STACK  
.BYTE $FF
```

would push all registers on to the user stack, the embedded byte being \$FF or 11111111. Alternatively, the coding:

```
JSR USER-STACK  
.BYTE $1E
```

where \$1E = 00011110 would push only the accumulator, Status and Index registers. Perhaps at this point a question is running through your mind: ‘won’t the embedded byte cause my program to crash?’. That’s true on face value, but what we do is get the user stack coding to move the Program Counter on one byte, to pass over it, as Program 17 shows:

Program 17

```
1Ø REM ** USER STACK **  
2Ø CODE=49152  
3Ø FOR LOOP=Ø TO 116  
4Ø READ BYTE  
5Ø POKE CODE+LOOP,BYTE  
6Ø NEXT LOOP  
7Ø :  
8Ø REM ** M/C DATA **  
9Ø DATA 8 : REM PHP  
10Ø DATA 72 : REM PHA  
11Ø DATA 138,72 : REM TXA : PHA  
12Ø DATA 152,72 : REM TYA : PHA  
13Ø DATA 186 : REM TSX  
14Ø DATA 16Ø,6 : REM LDY #$Ø6  
15Ø :: REM PUSH-ZERO-PAGE  
16Ø DATA 185,138,Ø : REM LDA $ØØ8A,Y
```

170	DATA	72	:	REM	PHA
180	DATA	136	:	REM	DEY
190	DATA	208,249	:	REM	BNE \$F9
200	DATA	254,5,1	:	REM	INC \$105,X
210	DATA	189,5,1	:	REM	LDA \$105,X
220	DATA	133,139	:	REM	STA \$8B
230	DATA	208,3	:	REM	BNE \$03
240	DATA	254,6,1	:	REM	INC \$106,X
250	:			REM	PC-LOW
260	DATA	189,6,1	:	REM	LDA \$106,X
270	DATA	133,140	:	REM	STA \$8C
280	DATA	169,135	:	REM	LDA #\$87
290	DATA	133,141	:	REM	STA \$8D
300	DATA	177,139	:	REM	LDA (\$8B),Y
310	DATA	133,142	:	REM	STA \$8E
320	DATA	169,8	:	REM	LDA #\$08
330	DATA	133,143	:	REM	STA \$8F
340	DATA	136	:	REM	DEY
350	DATA	198,252	:	REM	DEC \$FC
360	:			REM	ROTATE-BYTE
370	DATA	38,142	:	REM	ROL \$8E
380	DATA	144,16	:	REM	BCC \$10
390	DATA	189,6,1	:	REM	LDA \$106,X
400	DATA	145,251	:	REM	STA (\$FB),Y
410	DATA	136	:	REM	DEY
420	DATA	36,141	:	REM	BIT \$8D
430	DATA	16,6	:	REM	BPL \$06
440	DATA	189,5,1	:	REM	LDA \$105,X
450	DATA	145,251	:	REM	STA (\$FB),Y
460	DATA	136	:	REM	DEY
470	:			REM	BIT-CLEAR
480	DATA	202	:	REM	DEX
490	DATA	38,141	:	REM	ROL \$8D
500	DATA	144,1	:	REM	BCC \$01
510	DATA	202	:	REM	DEX
520	:			REM	OVER
530	DATA	198,143	:	REM	DEC \$8F
540	DATA	208,226	:	REM	BNE \$E2
550	DATA	56	:	REM	SEC
560	DATA	152	:	REM	TYA

```

570 DATA 101,251      : REM ADC $FB
580 DATA 133,251      : REM STA $FB
590 DATA 144,2         : REM BCC $02
600 DATA 230,252      : REM INC $FC
610 ::                 REM CLEAR-STACK
620 DATA 162,0         : REM LDX #0
630 ::                 REM REPEAT
640 DATA 104           : REM PLA
650 DATA 149,139       : REM STA $8B,X
660 DATA 232           : REM INX
670 DATA 224,6         : REM CPX #$06
680 DATA 208,248       : REM BNE $F8
690 DATA 104,168       : REM PLA : TAY
700 DATA 104,170       : REM PLA : TAX
710 DATA 104           : REM PLA
720 DATA 40            : REM PLP
730 DATA 96            : REM RTS
740 ::                 REM TEST-ROUTINE
750 DATA 169,240       : REM LDA #$F0
760 DATA 162,15         : REM LDX #$0F
770 DATA 160,255       : REM LDY #$FF
780 DATA 32,0,192      : REM JSR $C000
790 DATA 255           : REM EMBEDDED-BYTE
800 DATA 96            : REM RTS
810 :
820 REM ** SET UP ZERO PAGE AND FREE RAM **
830 PRINT CHR$(147)
840 POKE 251,12    POKE 252,197
850 FOR N=139 TO 144    : POKE N,N : NEXT
860 FOR N=50432 TO 50440 : POKE N,0 : NEXT
870 :
880 SYS 49258          : REM SYS TEST-ROUTINE
890 :
900 REM ** READ RESULTS **
910 FOR LOOP=50432 TO 50443
920 READ NAME$;
930 PRINT NAME$;
940 PRINT PEEK(LOOP)
950 NEXT LOOP
960 :

```

```

970 DATA "ZERO PAGE  ", "ZERO PAGE+1"
980 DATA "ZERO PAGE+2", "ZERO PAGE+3"
990 DATA "ZERO PAGE+4", "ZERO PAGE+5"
1000 DATA "Y REGISTER ", "X REGISTER "
1010 DATA "ACCUMULATOR", "STATUS      "
1020 DATA "PC LOW      ", "PC HIGH     "

```

The problem to solve next is that of where to place the user stack. This will depend on your own requirements, so to make the whole thing flexible, a vectored address in the bytes at \$FB and \$FC contains the stack address. In the program listed above, this is \$C512 (line 840). The vectored address is, in fact, the address + 12. This is because the stack is pushed in reverse (decreasing) order.

When executed, the coding first pushes all the processor registers on to the hardware stack and moves the stack pointer across into the X register (lines 90 to 140). Next, the six zero page pseudo-registers are pushed there (lines 150 to 190). The return address from the subroutine call is then incremented on the stack, using the contents of the X register (stack pointer) to access it (lines 200 to 240). The two bytes that form the RTS address are copied into pseudo-register 1 (now safely on the hardware stack) to form a vector through which the embedded data byte can be loaded into the accumulator and then saved for use in zero page (lines 250 to 310).

In line 280, a pre-defined byte was loaded into the accumulator and saved in zero page. This byte holds a bit code that will inform the program as to whether the register being pulled from the hardware stack for transfer to the software stack is one or two bytes long. The byte value, \$87, is 10000111 in binary and the set bits correspond to the two-byte registers, the Program Counter and the three pseudo-registers. By rotating this byte left after each pull operation and using the BIT operation, the Negative flag can be tested to see if a further pull is needed. All this and the copy hardware stack/push software stack is handled by lines 320 to 550.

Finally, the registers and pseudo-registers are restored to their original values (lines 620 to 730). The test routine between lines 750 and 800 shows the way the program is used. When run, the test procedure produces the following output on the screen:

ZERO PAGE	139
ZERO PAGE+1	140
ZERO PAGE+2	141
ZERO PAGE+3	142
ZERO PAGE+4	143
ZERO PAGE+5	144
Y REGISTER	255
X REGISTER	15

ACCUMULATOR	240
STATUS	176
PC LOW	115
PC HIGH	192

As can be seen, the zero page bytes contain the values POKEd into them by the FOR...NEXT loop of line 830 while the accumulator and Index registers display their seeded values (lines 750 to 770). The Program Counter holds $192 * 256 + 115$, or \$C073, which was the point in the program where its contents where pushed at line 780.

This program could be extended to provide a routine to perform a pull user stack, to copy the contents of a software stack into the processor and pseudo-registers.

Line-by-line

A line-by-line description of Program 17 follows:

```

line  90 : save all processor registers on hardware stack
line 140 : move stack pointer into X for index
line 150 : entry for PUSH-ZERO-PAGE
line 160 : get zero page byte
line 170 : push on to hardware stack
line 180 : decrement index
line 190 : branch to PUSH-ZERO-PAGE until done
line 200 : increment low byte of RTS address
line 210 : get it from stack
line 220 : and save in zero page
line 230 : if not equal branch to PC-LOW
line 240 : else increment page byte of RTS address
line 250 : entry for PC-LOW
line 260 : get high byte of RTS address
line 270 : and save it to form vector
line 280 : get bit code to indicate register size
line 290 : and save it
line 300 : get embedded code after subroutine call
line 310 : and save it
line 320 : eight bits in embedded byte to test
line 330 : save bit count
line 340 : decrement index to $FF
line 350 : decrement high byte of vectored address at $FB

```

line 360 : entry for ROTATE-BYTE
line 370 : move next coded bit into Carry flag
line 380 : if bit clear skip it, branch to BIT-CLEAR
line 390 : otherwise get byte from stack
line 400 : save it on user stack
line 410 : decrement index
line 420 : is it a two byte register?
line 430 : no, so branch to BIT-CLEAR
line 440 : yes, so get the second byte from the stack
line 450 : and save it on the user stack
line 460 : decrement index
line 470 : entry for BIT-CLEAR
line 480 : decrement hardware stack index
line 490 : move bit of register code into Carry flag
line 500 : if clear, branch to OVER
line 510 : else decrement hardware stack index
line 520 : entry for OVER
line 530 : decrement bit counter
line 540 : and repeat until all done
line 550 : set Carry flag
line 560 : move user stack pointer into accumulator
line 570 : add to low byte of address
line 580 : and save
line 590 : branch to CLEAR-STACK if carry is clear
line 600 : else increment high byte of address
line 610 : entry for CLEAR-STACK
line 620 : initialize X register
line 630 : entry for REPEAT
line 640 : pull byte from stack
line 650 : and restore zero page
line 660 : increment index
line 670 : all bytes restored?
line 680 : no, branch to REPEAT
line 690 : yes, restore all registers
line 730 : back to calling routine
line 740 : entry for TEST-ROUTINE
line 750 : seed registers
line 780 : call user stack routine
line 790 : embedded byte
line 800 : back to BASIC

BINARY INS AND OUTS

Sometimes when printing the values of registers, it is necessary to have their binary representation—for example, in the case of the Status register, because we are concerned with the state of the particular bits within it, rather than the overall value of the contents. Program 18 provides a short routine which produces such a binary output from a decimal input. This could easily be adapted for use within a program such as the software stack given above.

Program 18

```
10 REM ** PRINT ACCUMULATOR AS A **
20 REM ** BINARY NUMBER **
30 CODE=49152
40 FOR LOOP=0 TO 17
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 DATA 162,0 : REM LDX #$08
110 DATA 72 : REM PHA
120 :: REM NEXT-BIT
130 DATA 104 : REM PLA
140 DATA 10 : REM ASL A
150 DATA 72 : REM PHA
160 DATA 169,48 : REM LDA #$30
170 DATA 105,0 : REM ADC #$00
180 DATA 32,210,255 : REM JSR $FFD2
190 DATA 202 : REM DEX
200 DATA 208,243 : REM BNE $F3
210 DATA 104 : REM PLA
220 DATA 96 : REM RTS
230 :
240 REM ** SET UP DEMO RUN **
250 REM LDA $FB : JSR $C000 : RTS
260 POKE 820,165 :POKE 821,251
270 POKE 822,32 : POKE 823,0
280 POKE 824,192 : POKE 825,96
290 PRINT CHR$(147) PRINT
300 INPUT "INPUT A NUMBER ";A$
```

```

310 A=VAL(A$)
320 POKE 251,A
330 PRINT"BINAY VALUE IS :";
340 SYS 820

```

Line-by-line

The following line-by-line description should make the program's operation clear. It is simply moving each bit of the accumulator in turn into the Carry flag position, using the arithmetic shift left operation (see Figure 8.1) and adding its value to the ASCII code for \emptyset , i.e.

`accumulator=48+carry`

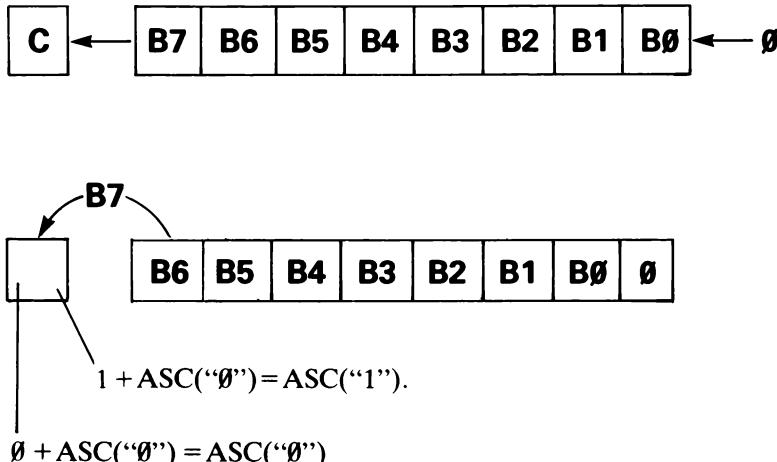


Figure 8.1 Arithmetic shift left

If the Carry flag is clear, the result will be $48+\emptyset=48$, so the CHROUT routine will print a \emptyset . On the other hand, if the Carry flag is set, the result of the addition will be $48+1=49$, so a 1 will be printed by CHROUT.

```

line 100 : eight bits in a byte
line 110 : push accumulator on to stack
line 120 : entry for NEXT-BIT

```

```

line 13Ø : restore accumulator
line 14Ø : shift bit 7 into carry
line 15Ø : save shifted accumulator on stack
line 16Ø : get ASCII code for Ø
line 17Ø : add carry
line 18Ø : print either Ø or 1
line 19Ø : decrement bit counter
line 2ØØ : do NEXT-BIT until complete
line 21Ø : pull stack to balance push
line 22Ø : back to BASIC

```

COME IN

By reversing this process, it is possible to input a number directly into the accumulator in binary form as Program 19 shows. The program scans the keyboard for a pressed 1 or Ø key and the Carry flag is set or cleared respectively. A copy of the accumulator, initially cleared, is kept on the hardware stack and restored each time round to rotate the carry bit into it using the rotate left operation (see Figure 8.2). The loop is executed eight times, once for each bit, and on completion, the accumulator holds the hexadecimal value of the binary number.

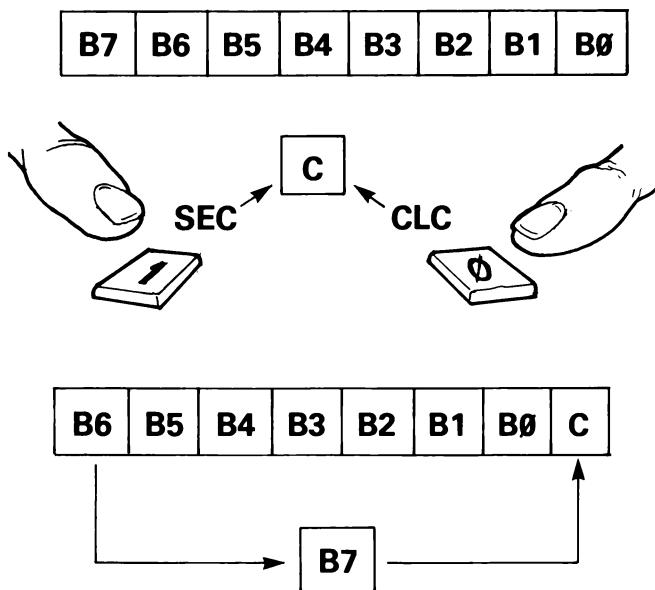


Figure 8.2 Input a number directly into the accumulator

Program 19

```
10 REM ** INPUT A HEX NUMBER IN BINARY FORM **
20 CODE=49152
30 FOR LOOP=0 TO 41
40 READ BYTE
50 POKE CODE+LOOP,BYTE
60 NEXT LOOP
70 :
80 REM ** M/C DATA **
90 DATA 162,8 : REM LDX #$08
100 DATA 169,0 : REM LDA #$00
110 DATA 72 : REM PHA
120 DATA 24 : REM CLC
130 :: : REM MAINLOOP
140 DATA 134,243 : REM STX $FD
150 :: : REM LOOP
160 DATA 32,228,255 : REM JSR $FFE4
170 DATA 240,251 : REM BEQ $FB
180 DATA 201,49 : REM CMP #$31
190 DATA 240,7 : REM BEQ $07
200 DATA 201,48 : REM CMP #$30
210 DATA 208,243 : REM BNE $F3
220 DATA 24 : REM CLC
230 DATA 144,1 : REM BCC $01
240 :: : REM SET
250 DATA 56 : REM SEC
260 :: : REM OVER
270 DATA 8 : REM PHP
280 DATA 32,210,255 : REM JSR $FFD2
290 DATA 40 : REM PLP
300 DATA 104 : REM PLA
310 DATA 42 : REM ROL A
320 DATA 72 : REM PHA
330 DATA 166,253 : REM LDX $FD
340 DATA 202 : REM DEX
350 DATA 208,224 : REM BNE $E0
360 DATA 104 : REM PLA
370 DATA 133,251 : REM STA $FB
380 DATA 96 : REM RTS
```

```
390 :
400 PRINT CHR$(147)
410 PRINT
420 PRINT"INPUT YOUR BINARY NUMBER :";
430 SYS CODE
440 PRINT PEEK(251)
```

Line-by-line

A line-by-line explanation of Program 19 now follows:

line 90 : eight bits to read
line 100 : clear accumulator—shift register
line 110 : push it on to stack
line 120 : clear the Carry flag
line 130 : entry for MAINLOOP
line 140 : save X register
line 150 : entry for LOOP
line 160 : jump to GETIN
line 170 : if null, branch to LOOP
line 180 : is it ASC“1”?
line 190 : yes, branch to SET
line 200 : is it ASC“0”?
line 210 : no, branch to LOOP
line 220 : yes, clear Carry flag
line 230 : and force branch to OVER
line 240 : entry for SET
line 250 : set Carry flag
line 260 : entry for OVER
line 270 : save Carry flag on stack
line 280 : print 0 or 1
line 290 : restore Carry flag
line 300 : restore accumulator
line 310 : move Carry flag into bit 0
line 320 : save accumulator
line 330 : restore bit count
line 340 : decrement it by one
line 350 : branch to MAINLOOP until all done
line 360 : restore accumulator
line 370 : save in zero page
line 380 : back to BASIC

Project

Convert the software stack program to print the binary values of each register upon completion.

Modify it further to allow register values to be seeded into the software stack test routine, using the binary input routine. Note that you should only attempt seeding the accumulator and Index registers. Why?

9 Move, Fill and Dump

MOVE IT!

The ability to move blocks of memory around within the bounds of the memory map is a necessity. When manipulating hi-resolution graphics, for example, large blocks of memory need to be moved around quickly and smoothly. The program could also be used to relocate sections of machine code rather than rewriting the assembler that created them—assuming, of course, that your code has been designed to make it portable.

At first sight, it may seem that the simplest method of moving a block of memory is to take the first byte to be moved and store it at the destination address, take the second byte and place it at the destination address + 1, and so forth. There would be no problem here if the destination address was outside the source address, but consider what would happen if the destination address was within the bounds to be searched by the source address—that is, the two regions overlapped. Figure 9.1 illustrates the problem using this straightforward method to move a block of five bytes forward by just a single byte, relocating the five bytes from \$C500 to \$C501.

Using the obvious method, the first character, ‘S’, is moved from \$C500 to \$C501 thereby overwriting the ‘A’. The program then takes the next character at location START+1 (\$C501), the ‘S’ that has just been written there, and places it at START+2 (\$C502)

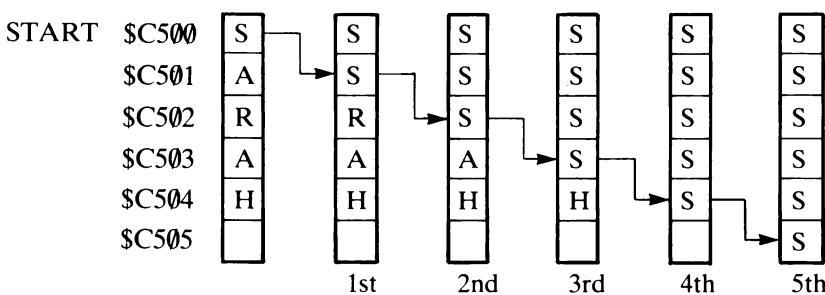


Figure 9.1 The overwriting move sequence

overwriting the 'R'. As you can see, the end result is SSSSS—the whole block is full of 'S's—not the required effect!

To avoid this problem, the MOVE routine acts 'intelligently' and if it calculates that an overwrite would occur, performs the movement of bytes in the reverse order, starting at the highest address and moving down the memory map as Figure 9.2 shows.

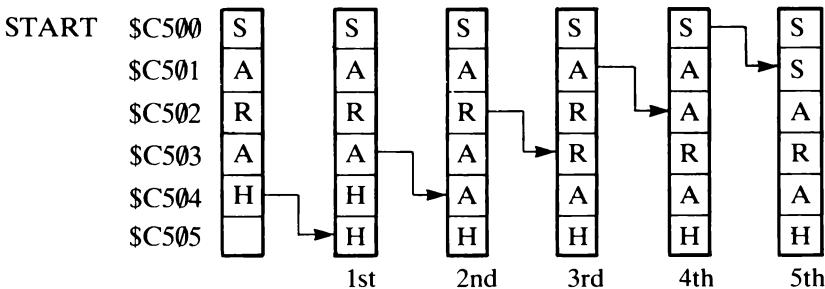


Figure 9.2 The correct move sequence

Program 20

```

10 REM ** MEMORY BLOCK MOVE RQUTINE **
20 REM ** 109 BYTES LONG WHEN ASSEMBLED **
30 REM ** PLUS 5 DATA BYTES IN ZERO PAGE **
40 CODE=49152
50 FOR LOOP=0 TO 108
60 READ BYTE
70 POKE CODE+LOOP,BYTE
80 NEXT LOOP
90 :
100 REM ** M/C DATA **
110 DATA 56 : REM SEC
120 DATA 165,251 : REM LDA $FB
130 DATA 229,253 : REM SBC $FD
140 DATA 170 : REM TAX
150 DATA 165,252 : REM LDA $FC
160 DATA 229,254 : REM SBC $FE
170 DATA 168 : REM TAY
180 DATA 138 : REM TXA
190 DATA 205,52,3 : REM CMP $334
200 DATA 152 : REM TYA
210 DATA 237,53,3 : REM SBC $335

```

220 DATA 176,2	:	REM BCS \$02
230 DATA 144,35	:	REM BCC \$23
240 ::		REM MOVE-LEFT
250 DATA 160,0	:	REM LDY #\$00
260 DATA 174,53,3	:	REM LDX \$335
270 DATA 240,14	:	REM BEQ \$0E
280 ::		REM LEFT-COMPLETE-PAGES
290 DATA 177,253	:	REM LDA (\$FD),Y
300 DATA 145,251	:	REM STA (\$FB),Y
310 DATA 200	:	REM INY
320 DATA 208,249	:	REM BNE \$F9
330 DATA 230,254	:	REM INC \$FE
340 DATA 230,252	:	REM INC \$FC
350 DATA 202	:	REM DEX
360 DATA 208,242	:	REM BNE \$F2
370 ::		REM LEFT-PARTIAL-PAGE
380 DATA 174,52,3	:	REM LDX \$334
390 DATA 240,8	:	REM BEQ \$08
400 ::		REM LAST-LEFT
410 DATA 177,253	:	REM LDA (\$FD),Y
420 DATA 145,251	:	REM STA (\$FB),Y
430 DATA 200	:	REM INY
440 DATA 202	:	REM DEX
450 DATA 208,248	:	REM BNE \$F8
460 ::		REM EXIT
470 DATA 96	:	REM RTS
480 ::		
490 ::		REM MOVE-RIGHT
500 DATA 24	:	REM CLC
510 DATA 173,53,3	:	REM LDA \$335
520 DATA 72	:	REM PHA
530 DATA 101,254	:	REM ADC \$FE
540 DATA 133,254	:	REM STA \$FE
550 DATA 24	:	REM CLC
560 DATA 104	:	REM PLA
570 DATA 101,252	:	REM ADC \$FC
580 DATA 133,252	:	REM STA \$FC
590 DATA 172,52,3	:	REM LDY \$334
600 DATA 240,9	:	REM BEQ \$09

```

610 :: REM TRANSFER
620 DATA 136 : REM DEY
630 DATA 177,253 : REM LDA ($FD),Y
640 DATA 145,251 : REM STA ($FB),Y
650 DATA 192,0 : REM CPY #$00
660 DATA 208,247 : REM BNE $F7
670 :: REM RIGHT-COMPLETE-PAGES
680 DATA 174,53,3 : REM LDX $335
690 DATA 240,221 : REM BEQ $DD
700 :: REM UPDATE
710 DATA 198,254 : REM DEC $FE
720 DATA 198,252 : REM DEC $FC
730 :: REM PAGE
740 DATA 136 : REM DEY
750 DATA 177,253 : REM LDA ($FD),Y
760 DATA 145,251 : REM STA ($FB),Y
770 DATA 192,0 : REM CPY #$00
780 DATA 208,247 : REM BNE $F7
790 DATA 202 : REM DEX
800 DATA 208,240 : REM BNE $F0
810 DATA 96 : REM RTS
820 :
830 REM ** SET UP VARIABLES **
840 PRINT CHR$(147)
850 PRINT" *** MEMORY MOVER V1.1 ***"
860 INPUT"START ADDRESS ";S
870 INPUT"DESTINATION ";D
880 INPUT"LENGTH IN BYTES ";L
890 :
900 S1=INT(S/256) : S2=S-(S1*256)
910 D1=INT(D/256) : D2=D-(D1*256)
920 L1=INT(L/256) : L2=L-(L1*256)
930 :
940 POKE 251,D2 : POKE 252,D1
950 POKE 253,S2 : POKE 254,S1
960 POKE 820,L2 : POKE 821,L1
970 :
980 REM ** SET UP DEMO **
990 FOR N=0 TO 15
1000 POKE 828+N,N

```

```

1010 POKE 900+N,0
1020 NEXT N
1030 :
1040 SYS CODE
1050 :
1060 REM ** PRINT THE RESULTS! **
1070 FOR N=0 TO 15
1080 PRINT PEEK(828+N);"
1090 PRINT PEEK(900+N)
1100 NEXT N

```

Bytes reserved:

251-252	(\$FB-\$FC)	: Destination vector
253-254	(\$FD-\$FE)	: Source vector
820-821	(\$334-\$335)	: Length of block to be moved

When run, the BASIC test requests three inputs: the START address of the memory block to be moved, its DESTINATION address and its LENGTH in bytes. All values should be entered as decimal values. Thus, to move a 1K block of memory from 49152 to 56000, the values to input are:

START ADDRESS	:	49152
DESTINATION	:	56000
LENGTH IN BYTES		1024

For reasons already explained, the coding begins by ascertaining whether a left-move or a right-move operation is required. It calculates this (lines 110 to 210) by subtracting the source address from the destination address. If the result is less than the number of bytes to be moved, overwriting would occur using the MOVE-LEFT routine, so the MOVE-RIGHT coding is called (line 230). If the memory locations do not overlap, the quicker MOVE-LEFT routine is selected (line 220). For further description purposes we will examine the MOVE-LEFT routine (lines 240 to 470).

Memory movement is performed in two phases: complete memory pages are first relocated, and then any remaining bytes in the final partial page are moved. These details are held in the length of block bytes \$334 and \$335.

The routine begins by loading the number of pages to be moved into the X register (line 260), branching to LEFT-PARTIAL-PAGE if it is zero (line 280). Transfer of data bytes is completed using post-indexed indirect addressing through the zero page vectors. When all the whole pages have been transferred, any

remaining bytes are transferred by the LEFT-PARTIAL-PAGE loop (lines 370 to 450).

The MOVE-RIGHT routine is similar in operation, except that it starts at the highest memory location referenced and moves down through memory, the highest address of the source and destination being calculated in lines 500 to 650.

Line-by-line

A line-by-line description of Program 20 now follows:

```
line 110 : set Carry flag
line 120 : get low byte destination address
line 130 : subtract low byte source address
line 140 : transfer result into X register
line 150 : get high byte destination address
line 160 : subtract high byte source address
line 170 : save result in X register
line 180 : restore result of low byte subtraction
line 190 : compare it with low byte of length
line 200 : restore result of high byte subtraction
line 210 : subtract high byte of length from it
line 220 : if Carry flag set, branch to MOVE-LEFT
line 230 : else branch to MOVE-RIGHT
line 240 : entry for MOVE-LEFT
line 250 : initialize index
line 260 : get number of pages to be moved
line 270 : if zero, branch to LEFT-PARTIAL-PAGE
line 280 : entry for LEFT-COMPLETE-PAGES
line 290 : get source byte
line 300 : store at destination
line 310 : increment index
line 320 : branch to LEFT-COMPLETE-PAGES until page
           done
line 330 : increment source page
line 340 : increment destination page
line 350 : decrement page counter
line 360 : branch to LEFT-COMPLETE-PAGES until all moved
line 370 : entry for LEFT-PARTIAL-PAGE
line 380 : get number of bytes on page to be moved
line 390 : if zero, branch to EXIT
```

line 400 : entry for LAST-LEFT
line 410 : get source byte
line 420 : store at destination
line 430 : increment index
line 440 : decrement byte count
line 450 : branch to LAST-LEFT until done
line 460 : entry for EXIT
line 470 : back to BASIC
line 490 : entry for MOVE-RIGHT
line 500 : clear Carry flag
line 510 : get number of pages to be moved
line 520 : save on stack
line 530 : add it to source high byte
line 540 : and save result
line 550 : reclear Carry flag
line 560 : get length high byte off stack
line 570 : add it to destination high byte
line 580 : and save the result
line 590 : get low byte of length into Y register
line 600 : branch to RIGHT-COMPLETE-PAGES if zero
line 610 : entry for TRANSFER
line 620 : decrement index
line 630 : get source byte
line 640 : and copy to destination
line 650 : is Y = 0?
line 660 : no, branch to TRANSFER
line 670 : entry for RIGHT-COMPLETE-PAGES
line 680 : get number of pages to be moved
line 690 : if zero, branch to EXIT
line 700 : entry for UPDATE
line 710 : decrement number of pages to do
line 720 : and also destination
line 730 : entry for PAGE
line 740 : decrement index
line 750 : get source byte
line 760 : copy to destination
line 770 : is Y = 0?
line 780 : no, branch to PAGE

line 790 : decrement page counter
line 800 : if not zero, branch to UPDATE
line 810 : return to BASIC

FILL

Program 21 provides the BASIC loader listing to implement a memory FILL routine, which is particularly useful for clearing sections of RAM with a pre-determined value.

Program 21

```
10 REM ** MEMORY FILL ROUTINE **
20 REM ** 30 BYTES LONG WHEN ASSEMBLED **
30 REM ** PLUS 5 DATA BYTES IN ZERO PAGE **
40 CODE=49152
50 FOR LOOP=0 TO 30
60 READ BYTE
70 POKE CODE+LOOP,BYTE
80 NEXT LOOP
90 :
100 REM ** M/C DATA **
110 DATA 165,255 : REM LDA $FF
120 DATA 166,252 : REM LDX $FC
130 DATA 240,12   : REM BEQ $0C
140 DATA 160,0    : REM LDY #$00
150 ::           REM COMPLETE-PAGE
160 DATA 145,253 : REM STA ($FD),Y
170 DATA 200       : REM INY
180 DATA 208,251  : REM BNE $FB
190 DATA 230,254  : REM INC $FE
200 DATA 202       : REM DEX
210 DATA 208,246  : REM BNE $F6
220 ::           REM PARTIAL-PAGE
230 DATA 166,251  : REM LDX $FB
240 DATA 240,8    : REM BEQ $08
250 DATA 160,0    : REM LDY #$00
260 ::           REM AGAIN
270 DATA 145,253  : REM STA ($FD),Y
```

```

280 DATA 200      : REM INY
290 DATA 202      : REM DEX
300 DATA 208,250  : REM BNE $FA
310 ::           REM FINISH
320 DATA 96       : REM RTS
330 :
340 REM ** GET DETAILS **
350 PRINT CHR$(147)
360 INPUT" FILL DATA      :" ;F
370 INPUT" START ADDRESS   :" ;S
380 INPUT" NUMBER OF BYTES :" ;L
390 :
400 S1=INT(S/256)    : S2=S-(S1*256)
410 L1=INT(L/256)    : L2=L-(L1*256)
420 :
430 POKE 251,L2     : POKE 252,L1
440 POKE 253,S2     : POKE 254,S1
450 POKE 255,F
460 :
470 SYS CODE

```

Bytes reserved:

251–252	(\$FB–\$FC)	: number of bytes to be filled
253–254	(\$FD–\$FE)	: start of address of bytes to be filled
255	(\$FF)	: value to fill with

When executed, the machine code expects to find the fill value, the start address and the amount of memory to be filled, in five zero page bytes of memory from \$FB. Input of each of these is handled by a few lines of BASIC from line 360. To clear a 1K block of RAM from \$C500 with zero, the following information should be entered in response to the 64's prompt:

FILL DATA	:	Ø
START ADDRESS	:	49152
NUMBER OF BYTES	:	1024

The FILL routine works in a similar manner to the MOVE routine described above, dealing with whole and partial pages separately. The main fill loop is embodied in lines 150 to 300.

Line-by-line

A line-by-line description of the program now follows:

```
line 110 : get data with which to fill
line 120 : get number of complete pages to be filled
line 130 : if zero, branch to PARTIAL-PAGE
line 140 : initialize index
line 150 : entry for COMPLETE-PAGE
line 160 : fill byte
line 170 : increment index
line 180 : branch to COMPLETE-PAGE until all of page is
           done
line 190 : increment page
line 200 : decrement page counter
line 210 : branch to COMPLETE-PAGE until all pages are
           filled
line 220 : entry for PARTIAL-PAGE
line 230 : get number of bytes left to be filled
line 240 : if zero, branch to FINISH
line 250 : else clear index
line 260 : entry for AGAIN
line 270 : fill byte
line 280 : increment index
line 290 : decrement bytes left to do count
line 300 : branch to AGAIN until all filled
line 310 : entry for FINISH
line 320 : back to BASIC
```

A MEMORY DUMP

A hex and ASCII dump of memory can be extremely useful, not only within machine code programs, but also when used from a BASIC program. Most often it provides information about the way a program is manipulating numeric and string variables and tables. Figure 9.3 shows the type of dump produced by the routine: twenty-four lines of eight bytes each. The example shows some text stored in memory. Each line starts with the current address, followed by the eight bytes stored in memory from that point. The far right of the listing provides the ASCII equivalents of each byte. Any non-ASCII character (that is, one greater than \$7F) or control code (those less than \$20) is represented by a full stop.

C108 : 54 68 69 73 20 69 73 20 This is
C110 : 61 20 73 69 6D 70 6C 65 a simple
C118 : 20 65 78 61 6D 70 6C 65 example
C120 : 20 6F 66 20 68 6F 77 20 of how
C128 : 74 68 65 20 8D 64 75 6D the .dum
C130 : 70 20 72 6F 75 74 69 6E p routin
C138 : 65 20 66 6F 72 20 74 68 e for th
C140 : 65 20 43 6F 6D 6D 6F 64 e Commod
C148 : 6F 72 65 20 36 34 20 8D ore 64 .
C150 : 77 6F 72 6B 73 2E 0D 54 works..T
C158 : 68 65 20 64 75 6D 70 20 he dump
C160 : 63 61 6E 20 62 65 20 64 can be d
C168 : 69 76 69 64 65 64 20 69 ivided i
C170 : 6E 74 6F 20 74 68 72 65 nto thre
C178 : 65 20 8D 73 65 63 74 69 e .secti
C180 : 6F 6E 73 2E 20 54 68 65 ons. The
C188 : 20 66 69 72 73 74 20 63 first c
C190 : 6F 6C 75 6D 6E 20 6C 69 olumn li
C198 : 73 74 73 20 74 68 65 20 sts the
C1A0 : 8D 73 74 61 72 74 20 61 .start a
C1A8 : 64 64 72 65 73 73 20 6F ddress o
C1B0 : 66 20 74 68 65 20 62 6C f the bl
C1B8 : 6F 63 6B 2E 20 54 68 65 ock. The
C1C0 : 20 73 65 63 6F 6E 64 20 second
C1C8 : 8D 63 6F 6C 75 6D 6E 20 .column
C1D0 : 69 73 20 69 6E 20 66 61 is in fa
C1D8 : 63 74 20 74 68 65 20 68 ct the h
C1E0 : 65 78 61 64 65 63 69 6D exadecim
C1E8 : 61 6C 20 8D 76 61 6C 75 al .valu
C1F0 : 65 73 20 6F 66 20 65 69 es of ei
C1F8 : 67 68 74 20 62 79 74 65 ght byte
C200 : 73 20 66 72 6F 6D 20 74 s from t
C208 : 68 69 73 20 8D 61 64 64 his .add
C210 : 72 65 73 73 2E 20 46 69 ress. Fi
C218 : 6E 61 6C 6C 79 20 74 68 nally th
C220 : 65 20 6C 61 73 74 20 63 e last c
C228 : 6F 6C 75 6D 6E 20 8D 64 olumn .d
C230 : 65 70 69 63 74 73 20 74 epicts t
C238 : 68 65 20 41 53 43 49 49 he ASCII
C240 : 20 76 61 6C 75 65 73 20 values
C248 : 6F 66 20 74 68 65 73 65 of these
C250 : 20 8D 62 79 74 65 73 2E .bytes.
C258 : 20 75 6E 6C 65 73 73 20 unless
C260 : 74 68 65 20 62 79 74 65 the byte
C268 : 20 69 73 20 6E 6F 6E 2D is non-
C270 : 41 53 43 49 49 20 8D 77 ASCII .w
C278 : 68 69 63 68 20 69 73 20 hich is
C280 : 74 68 65 6E 20 64 69 73 then dis
C288 : 70 6C 61 79 65 64 20 61 played a
C290 : 73 20 61 20 66 75 6C 6C s a full
C298 : 20 73 74 6F 70 21 0D 00 stop!..
C2A0 : 00 4C 00 C9 A9 FF 85 22 .L....."
C2A8 : 08 60 00 00 00 00 00 00 ..'.....

Figure 9.3 Memory dump

As it stands, the routine requires three zero page data bytes, two for the start address and one for the number of eight byte lines to be dumped. The routine also employs the ADDRESS-PRINT and HEXPRINT routines discussed earlier.

Program 22

```
10 REM ** DUMP LINES OF 8 BYTES OF **
20 REM ** MEMORY IN HEX AND ASCII  **
30 CODE=49152
40 FOR LOOP=0 TO 111
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 DATA 32,71,192 : REM JSR $C047
110 :: : REM HEX-BYTES
120 DATA 162,7 : REM LDX #$07
130 DATA 160,0 : REM LDY #$00
140 : REM HEX-LOOP
150 DATA 177,251 : REM LDA ($FB),Y
160 DATA 32,90,192 : REM JSR $C05A
170 DATA 32,66,192 : REM JSR $C042
180 DATA 200 : REM INY
190 DATA 202 : REM DEX
200 DATA 16,244 : REM BPL $F4
210 DATA 32,66,192 : REM JSR $C042
220 :: : REM ASCII-BYTES
230 DATA 162,7 : REM LDX #$07
240 DATA 160,0 : REM LDY #$00
250 :: : REM ASCII-LOOP
260 DATA 177,251 : REM LDA ($FB),Y
270 DATA 201,32 : REM CMP #$20
280 DATA 48,4 : REM BMI $04
290 DATA 201,128 : REM CMP #$80
300 DATA 144,2 : REM BCC $02
310 :: : REM FULL-STOP
320 DATA 169,46 : REM LDA #$2E
330 :: : REM LEAP-FROG
```

340 DATA 32,210,255 : REM JSR \$FFD2
350 DATA 200 : REM INY
360 DATA 202 : REM DEX
370 DATA 16,237 : REM BPL \$ED
380 DATA 169,13 : REM LDA #\$0D
390 DATA 32,210,255 : REM JSR \$FFD2
400 DATA 24 : REM CLC
410 DATA 165,251 : REM LDA \$FB
420 DATA 105,8 : REM ADC #\$08
430 DATA 133,251 : REM STA \$FB
440 DATA 144,2 : REM BCC \$02
450 DATA 230,252 : REM INC \$FC
460 :: REM NO-CARRY
470 DATA 198,254 : REM DEC \$FE
480 DATA 208,191 : REM BNE \$BF
490 DATA 96 : REM RTS
500 :: REM SPACE
510 DATA 169,32 : REM LDA #\$20
520 DATA 76,210,255 : REM JMP \$FFD2
530 :: REM ADDRESS-PRINT
540 DATA 162,251 : REM LDX #\$FB
550 DATA 181,1 : REM LDA 1,X
560 DATA 32,90,192 : REM JSR \$C05A
570 DATA 181,0 : REM LDA 0,X
580 DATA 32,90,192 : REM JSR \$C05A
590 DATA 32,66,192 : REM JSR \$C042
600 DATA 32,66,192 : REM JSR \$C042
610 DATA 96 : REM RTS
620 :: REM HEXPRINT
630 DATA 72 : REM PHA
640 DATA 74,74 : REM LSR A : LSR A
650 DATA 74,74 : REM LSR A : LSR A
660 DATA 32,99,192 : REM JSR \$C063
670 DATA 104 : REM PLA
680 :: REM FIRST
690 DATA 41,15 : REM AND #\$0F
700 DATA 201,10 : REM CMP #\$0A
710 DATA 144,2 : REM BCC \$02
720 DATA 105,6 : REM ADC #\$06
730 :: REM OVER

```

740 DATA 105,48      : REM ADC #$30
750 DATA 76,210,255  : REM JMP $FFD2
760 :
770 REM ** INPUT DETAILS FOR DUMP **
780 PRINT CHR$(147)
790 INPUT"Dump Start Address ";A
800 HIGH=INT(A/256)
810 LOW=A-(HIGH*256)
820 POKE 251,LOW : POKE 252,HIGH
830 INPUT"Number of Lines (20/Screen) ";B
840 POKE 254,B
850 SYS CODE

```

The program's operation is quite simple, using the X register to count the bytes as they are printed across the screen using HEXPRINT (lines 120 to 210). The second section of code (lines 220 to 370) is responsible for printing either the ASCII character contained in the byte, or a full stop if an unprintable character or a control code is encountered. The final section of code moves the cursor down one line and increments the address counter. The whole loop is repeated until the line count reaches zero.

Line-by-line

A line-by-line description of the Program 22 now follows:

```

line 100 : print start address of current line
line 110 : entry for HEX-BYTES
line 120 : eight bytes to do (0-7)
line 130 : clear index
line 140 : entry for HEX-LOOP
line 150 : get byte through vectored address
line 160 : print it as two hex digits
line 170 : print a space
line 180 : increment index
line 190 : decrement bit count
line 200 : branch to HEX-LOOP until all done
line 210 : print a space
line 220 : entry for ASCII-BYTES
line 230 : eight bytes to redo
line 240 : set index
line 250 : entry for ASCII-LOOP

```

line 260 : get byte through vectored address
line 270 : is it less than ASC“ ”?
line 280 : yes, branch to FULL-STOP
line 290 : is it greater than 128?
line 300 : no, branch to LEAP-FROG
line 310 : entry for FULL-STOP
line 320 : get ASC“.” into accumulator
line 330 : entry for LEAP-FROG
line 340 : print accumulator's contents
line 350 : increment index
line 360 : decrement bit count
line 370 : branch to ASCII-LOOP until all done
line 380 : get ASCII code for RETURN
line 390 : print new line
line 400 : clear Carry flag
line 410 : get low byte of address
line 420 : add 8 to it
line 430 : save result
line 440 : if no carry, branch to NO-CARRY
line 450 : else increment high byte of address
line 460 : entry for NO-CARRY
line 470 : decrement line counter
line 480 : branch to start at \$C100 until all lines done
line 490 : return to BASIC
line 500 : entry to SPACE
line 510 : get ASCII code for space
line 520 : print it and return through jump
line 530 : entry to ADDRESS-PRINT
line 540 : load index into X register
line 550 : get high byte of address
line 560 : print it as two hex digits
line 570 : get low byte of address
line 580 : print it as two hex digits
line 590 : print a space
line 600 : print a second space
line 610 : return to main program
line 620 : entry to HEXPRINT
line 630 : save accumulator on stack
line 640 : move high nibble into low nibble position
line 660 : call FIRST subroutine

line 670 : restore accumulator to do low byte
line 680 : entry for FIRST
line 690 : mask off high nibble
line 700 : is it less than 10?
line 710 : yes, so jump OVER
line 720 : add 7 to convert to A-F
line 730 : entry to OVER
line 740 : add 48 to convert to ASCII code
line 750 : print it and return

10 Hi-res Graphics

The Commodore 64 can support hi-resolution graphics. However, as you are no doubt aware, setting up the hi-res screen prior to using it can be a rather long-winded process, requiring several lines of BASIC text. In fact, four routines are normally required:

1. Move start of BASIC user area and set position for hi-res screen.
2. Clear screen memory.
3. Select screen colour and clear to that colour.
4. Reselect normal character mode.

All of these can be performed quite simply at machine level, and the routines for each follow. They can be compiled as DATA at the end of a graphics program, poked into memory at RUN time and executed via a SYS call. This does have one of the original disadvantages, in that a large chunk of program is required. However, the main advantage is speed, particularly in clearing the screen. Alternatively, any of these routines would make an admirable addition to the Wedge Operating System, allowing it to be called by name from within your programs. Suitable command names might be:

- | | |
|-----------------|---|
| @MOVEBAS | : move BASIC program area to make room for
hi-res screen |
| @HIRES | : select hi-res screen |
| @CLEAR | : clear hi-res screen |
| @GCOL | : clear to graphics colour specified in a dedicated
byte |
| @MODE | : select normal character mode |

Let us now examine each command in turn.

A BASIC MOVE

You may be wondering why we should bother to move the BASIC program area at all—why not just position the hi-res screen midway in memory? The reason for the careful positioning of the routine is as a matter of safety—placing the hi-res screen above the BASIC program area could lead to it being corrupted, especially if it is being used in conjunction with the program, because adding a line or two to the program could cause it to extend into the hi-res screen. Making sure the BASIC program fits in is no real safeguard either, as variables, strings and arrays all eat up memory at an incredible rate, and these could find their way into the screen memory. All these problems can be avoided by moving the start of BASIC up enough bytes to allow the hi-res screen to be tucked in underneath.

To do this requires a machine code program. The *Programmer's Reference Guide* lists five vectors associated with BASBAS (that's my mnemonic for BASIC's base!), as follows:

\$2B-\$2C	TXTTAB	:	start of BASIC text
\$2D-\$2E	VARTAB	:	start of BASIC variables
\$2F-\$30	ARYTAB	:	start of BASIC arrays
\$31-\$32	STREND	:	end of BASIC arrays+1
\$281-\$282	MEMSTR	:	bottom of memory

To move BASIC, each of these vectors must be reset to point to the new start area and the first three bytes of the new start area must be cleared to keep the Kernal happy.

Program 23 performs each of these functions. The address of the new BASIC area is \$4000, which allows room for the hi-res screen plus 32 sprites.

Program 23

```
10 REM ** MOVE BASIC PROGRAM AREA START **
20 REM ** UP TO 16348 TO FREE HI-RES SCREEN **
30 :
40 CODE=49152
50 FOR LOOP=0 TO 39
60 READ BYTE
70 POKE CODE+LOOP,BYTE
80 NEXT LOOP
90 :
100 REM ** M/C DATA **
110 DATA 169,0 : REM LDA #$00
```

```

120 DATA 141,2,64 : REM STA $4002
130 DATA 141,1,64 : REM STA $4001
140 DATA 141,0,64 : REM STA $4000
150 DATA 141,129,2 : REM STA $0281
160 DATA 169,64 : REM LDA #$40
170 DATA 133,44 : REM STA $2C
180 DATA 133,46 : REM STA $2E
190 DATA 133,48 : REM STA $30
200 DATA 133,50 : REM STA $32
210 DATA 141,130,2 : REM STA $0282
220 DATA 169,1 : REM LDA #$01
230 DATA 133,43 : REM STA $2B
240 DATA 169,3 : REM LDA #$03
250 DATA 133,45 : REM STA $2D
260 DATA 133,47 : REM STA $2F
270 DATA 133,49 : REM STA $31
280 DATA 96 : REM RTS

```

Line-by-line

A line-by-line description of Program 23 follows:

```

line 110 : initialize accumulator
line 120 : and clear first four bytes of new program area
line 150 : set low byte of MEMSTR (bottom of memory pointer)
line 160 : load high byte of new program area address into
           accumulator
line 170 : set high byte of TXTTAB
line 180 : set high byte of VARTAB
line 190 : set high byte of ARYTAB
line 200 : set high byte of STREND
line 210 : set high byte of MEMSTR
line 220 : load accumulator with 1
line 230 : store in low byte of TXTTAB
line 240 : load accumulator with 3
line 250 : set low bytes of all vectored addresses

```

SELECTING HI-RES

Before selecting the hi-resolution screen mode, it is necessary to point the VIC chip to the start of screen memory. This is done by writing to the VIC Memory Control register located at \$D018 (57272). The actual location is controlled by the condition of bits 3, 2 and 1. Table 10.1 details their settings for various addresses.

Table 10.1

Bit code	Value	Address selected	
xxxx000x	0	0-2047	(\$0000-\$07FF)
xxxx001x	2	2048-4095	(\$0800-\$0FFF)
xxxx010x	4	4096-6143	(\$1000-\$17FF)
xxxx011x	6	6144-8191	(\$1800-\$1FFF)
xxxx100x	8	8192-10239	(\$2000-\$27FF)
xxxx101x	10	10240-12287	(\$2800-\$2FFF)
xxxx110x	12	12288-14335	(\$3000-\$37FF)
xxxx111x	14	14336-16383	(\$3800-\$3FFF)

You can see from the table that the screen memory may be moved around in 2K block steps. An 'x' in each of the other bits denotes that these bits may be in either state. However, remember that these bits are controlling other aspects of the VIC's function, so that any reprogramming of bits 3, 2 and 1 must preserve the other bits. This is best done with the logical OR function. Looking at Table 10.1 we can see that bit 3 must be set to point the Memory Control register at location 8192. In BASIC this would simplify to:

```
100 A=PEEK(53727) : REM GET VALUE  
110 A=A OR 8       : REM SET BIT 3  
120 POKE 53727,A   : REM REPROGRAM
```

which translates to assembler as:

```
LDA #$08  
ORA $D018  
STA $D018
```

Now that the hi-res screen has been defined, it can be switched in by setting bit 5 of the VIC Control register at \$D011 (53265).

Again, the other bits in the register must be preserved, so the byte must be ORed with 32 (**00100000** binary). In BASIC this is:

```
13Ø A=PEEK(53265) : REM GET VALUE  
14Ø A=A OR 32       : REM SET BIT 5  
15Ø POKE 53265,A   : REM REPROGRAM
```

and in assembler:

```
LDA #$2Ø  
ORA $DØ11  
STA $DØ11
```

A CLEAR VIEW

Once hi-res mode has been selected, it will be filled with junk (often referred to as garbage). To clear this, each location must in turn be POKEd with zero. A BASIC program to do this would take the form:

```
2ØØ SB=8192  
21Ø FOR L=SB TO SB+7999  
22Ø POKE L,Ø  
23Ø NEXT L
```

Previously, in normal character mode, locations 1Ø24 to 2Ø23 were used to control which character was displayed—for example, POKEing a 1 into location 1Ø24 would make a letter A appear at the top left hand corner of the screen. When in hi-res mode, this area of memory is used to hold the colour information of that byte. Note that the colour information does not now come from the colour memory—colour details are taken directly from the hi-res screen itself. The high nibble of the byte (that is, bits 4 to 7) holds the colour code of any bit that is set in that 8 by 8 bit matrix, while the lower nibble (bits 3 to Ø) holds the colour of any bits that are clear in the same area.

To clear the hi-res screen to black ink on green paper in BASIC we would use:

```
24Ø FOR C=1Ø24 TO 2Ø23  
25Ø POKE C,13  
26Ø NEXT C
```

If all the above BASIC program lines were to be combined and RUN, the resulting hi-res screen would take around 20 seconds to construct—a bit slow, you'll agree! Program 24 provides the

machine code equivalent. Note that the value assigned to CODE is 49408 and NOT 49152 as we have been using previously. This is to allow the program to be used in conjunction with the MOVEBAS program described earlier. After you have entered and RUN MOVEBAS, try this one for an instant hi-res screen!

Program 24

```
10 REM ** HI-RES GRAPHICS SCREEN SET AND
    CLEAR **
20 CODE=49408
30 FOR LOOP=0 TO 105
40 READ BYTE
50 POKE CODE+LOOP,BYTE
60 NEXT LOOP
70 :
80 REM ** M/C DATA **
85 ::                      REM SELECT-HI-RES
90 DATA 169,8      : REM LDA #$08
100 DATA 13,24,208 : REM ORA $D018
110 DATA 141,24,208 : REM STA $D018
120 DATA 169,32     : REM LDA #$20
130 DATA 13,17,208 : REM ORA $D011
140 DATA 141,17,208 : REM STA $D011
150 ::                  REM CLEAR-SCREEN-MEMORY
160 DATA 169,0      : REM LDA #$00
170 DATA 133,251   : REM STA $FC
180 DATA 169,32     : REM LDA #$20
190 DATA 133,252   : REM STA $FC
200 DATA 169,64     : REM LDA #$40
210 DATA 133,253   : REM STA $FD
220 DATA 169,63     : REM LDA #$3F
230 DATA 133,254   : REM STA $FE
240 ::                  REM IN
250 DATA 165,252   : REM LDA $FC
260 DATA 197,254   : REM CMP $FE
270 DATA 208,9      : REM BNE $09
280 DATA 165,251   : REM LDA $FB
290 DATA 197,253   : REM CMP $FD
300 DATA 208,3      : REM BNE $03
```

```

310 DATA 76,62,192 : REM JMP $C03E
320 :: REM CLEAR
330 DATA 160,0 : REM LDY #$00
340 DATA 169,0 : REM LDA #$00
350 DATA 145,251 : REM STA ($FB),Y
360 DATA 230,251 : REM INC $FB
370 DATA 208,231 : REM BNE $E7
380 DATA 230,252 : REM INC $FC
390 DATA 56 : REM SEC
400 DATA 176,226 : REM BCS $E2
410 :
420 :: REM COLOUR
430 DATA 169,0 : REM LDA #$00
440 DATA 133,251 : REM STA $FB
450 DATA 169,4 : REM LDA #$04
460 DATA 133,252 : REM STA $FC
470 DATA 169,231 : REM LDA #$E7
480 DATA 133,253 : REM STA $FD
490 DATA 169,7 : REM LDA #$07
500 DATA 133,254 : REM STA $FE
510 :: REM CIN
520 DATA 165,252 : REM LDA $FC
530 DATA 197,254 : REM CMP $FE
540 DATA 208,7 : REM BNE $07
550 DATA 165,251 : REM LDA $FB
560 DATA 197,253 : REM CMP $FD
570 DATA 208,1 : REM BNE $01
580 DATA 96 : REM RTS
590 :: REM GREEN
600 DATA 160,0 : REM LDY #$00
610 DATA 169,13 : REM LDA #$0D
620 DATA 145,251 : REM STA ($FB),Y
630 DATA 230,251 : REM INC $FB
640 DATA 208,233 : REM BNE $E9
650 DATA 230,252 : REM INC $FC
660 DATA 56 : REM SEC
670 DATA 176,228 : REM BCS $E4

```

Line-by-line

A line-by-line description of Program 24 follows:

line 90 : load accumulator with mask 00001000
line 100 : force bit 3 to select 8196 as bit map start address
line 110 : and program VIC Memory Control register
line 120 : load accumulator with mask 00100000
line 130 : force bit 5 to select bit map mode
line 140 : and program CIC Control register
line 150 : entry for bit map CLEAR-SCREEN-MEMORY routine
line 160 : set up vector to point to screen start address \$2000
line 200 : set up vector to point to screen end address \$403F
line 240 : entry for IN
line 250 : get high byte current address
line 260 : is it same as high byte end address?
line 270 : no, so branch to CLEAR
line 280 : yes, get low byte current address
line 290 : is it same as low byte end address
line 300 : no, so branch to CLEAR
line 310 : yes, all done jump to COLOUR
line 320 : entry for CLEAR
line 330 : initialize index
line 340 : clear accumulator
line 350 : clear byte of screen memory
line 360 : increment low byte of current screen address
line 370 : branch to IN if no carry over
line 380 : increment high byte
line 390 : set Carry flag
line 400 : force branch to IN
line 420 : entry for COLOUR
line 430 : set up vector to point to start of colour memory
line 470 : set up vector to point to end of colour memory
line 510 : entry for CIN
line 520 : get high byte of current address
line 530 : is it the same as high byte end address?
line 540 : no, branch to GREEN

line 550 : get low byte of current address
line 560 : is it the same as the low byte end address?
line 570 : no, branch to GREEN
line 580 : back to calling routine
line 590 : entry for GREEN
line 600 : clear indexing register
line 610 : get code for green into accumulator
line 620 : POKE it into colour memory
line 630 : increment low byte of current address
line 640 : branch to CIN if no carry over
line 650 : increment high byte
line 660 : set Carry flag
line 670 : and force branch to CIN

Appendix 1: 6510 Complete Instruction Set

ADC	Add with carry	NZCV	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$69	2	2
Zero page	\$65	2	3
Zero page,X	\$75	2	4
Absolute	\$6D	3	4
Absolute,X	\$7D	3	4 or 5
Absolute,Y	\$79	3	4 or 5
(Indirect,X)	\$61	2	6
(Indirect),Y	\$71	2	5

AND	AND with accumulator	NZ	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$29	2	2
Zero page	\$25	2	3
Zero page,X	\$35	2	4
Absolute	\$2D	3	4
Absolute,X	\$3D	3	4 or 5
Absolute,Y	\$39	3	4 or 5
(Indirect,X)	\$21	2	6
(Indirect),Y	\$31	2	5

ASL	Shift left	NZC	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Accumulator	\$0A	1	2
Zero page	\$06	2	5
Zero page,X	\$16	2	6
Absolute	\$0E	3	6
Absolute,X	\$1E	3	7

BCC	Branch if C = 0	Flags unaltered	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Relative	\$90	2	3 or 2

BCS	Branch if C = 1	Flags unaltered	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Relative	\$B0	2	3 or 2

BEQ	Branch if Z = 1	Flags unaltered	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Relative	\$F0	2	3 or 2

BIT		Z,N,V
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>
Zero page	\$24	2
Absolute	\$2C	3
BMI	Branch if N = 1	Flags unaltered
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>
Relative	\$30	2
BNE	Branch if Z = 0	Flags unaltered
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>
Relative	\$D0	2
BPL	Branch if N = 0	Flags unaltered
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>
Relative	\$10	2

BRK	Break	B flag = 1
<i>Address mode</i> Implied	<i>Op-code</i> \$00	<i>Bytes</i> 1
BVC	Branch if V = 0	Flags unaltered
<i>Address mode</i> Relative	<i>Op-code</i> \$50	<i>Bytes</i> 2
BVS	Branch if V = 1	Flags unaltered
<i>Address mode</i> Relative	<i>Op-code</i> \$70	<i>Bytes</i> 2
CLC	Clear Carry flag	C flag = 0
<i>Address mode</i> Implied	<i>Op-code</i> \$18	<i>Bytes</i> 1

CLD	Clear Decimal flag	D flag = Ø
------------	--------------------	------------

<i>Address mode</i> Implied	<i>Op-code</i> \$D8	<i>Bytes</i> 1	<i>Cycles</i> 2
--------------------------------	------------------------	-------------------	--------------------

CLI	Clear Interrupt flag	I flag = Ø
------------	----------------------	------------

<i>Address mode</i> Implied	<i>Op-code</i> \$58	<i>Bytes</i> 1	<i>Cycles</i> 2
--------------------------------	------------------------	-------------------	--------------------

CLV	Clear Overflow flag	V flag = Ø
------------	---------------------	------------

<i>Address mode</i> Implied	<i>Op-code</i> \$B8	<i>Bytes</i> 1	<i>Cycles</i> 2
--------------------------------	------------------------	-------------------	--------------------

CMP	Compare accumulator	NZC
------------	---------------------	-----

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$C9	2	2
Zero page	\$C5	2	3
Zero page,X	\$D5	2	4
Absolute	\$CD	3	4
Absolute,X	\$DD	3	4 or 5
Absolute,Y	\$D9	3	4 or 5
(Indirect,X)	\$C1	2	6
(Indirect),Y	\$D1	2	5 or 6

CPX	Compare X register	NZC	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$ EØ	2	2
Zero page	\$ E4	2	3
Absolute	\$ EC	3	4

CPY	Compare Y register	NZC	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$ CØ	2	2
Zero page	\$ C4	2	3
Absolute	\$ CC	3	4

DEC	Decrement memory	NZ	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Zero page	\$ C6	2	5
Zero page,X	\$ D6	2	6
Absolute	\$ CE	3	6
Absolute,X	\$ DE	3	7

DEX	Decrement X register	NZ	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$ CA	1	2

DEY	Decrement Y register		NZ
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$88	1	2
EOR	Exclusive-OR		NZ
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$49	2	2
Zero page	\$45	2	3
Zero page,X	\$55	2	4
Absolute	\$4D	3	4
Absolute,X	\$5D	3	4 or 5
Absolute,Y	\$59	3	4 or 5
(Indirect,X)	\$41	2	6
(Indirect),Y	\$51	2	5
INC	Increment memory		NZ
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Zero page	\$E6	2	5
Zero page,X	\$F6	2	6
Absolute	\$EE	3	6
Absolute,X	\$FE	3	7

INX	Increment X register	NZ	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$ E8	1	2
<hr/>			
INY	Increment Y register	NZ	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$ C8	1	2
<hr/>			
JMP	Jump	Flags unaltered	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Absolute	\$ 4C	3	3
Indirect	\$ 6C	3	5
<hr/>			
JSR	Jump to subroutine	Flags unaltered	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Absolute	\$ 20	3	6
<hr/>			

LDA	Load accumulator		NZ
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$A9	2	2
Zero page	\$A5	2	3
Zero page,X	\$B5	2	4
Absolute	\$AD	3	4
Absolute,X	\$BD	3	4 or 5
Absolute,Y	\$B9	3	4 or 5
(Indirect,X)	\$A1	2	6
(Indirect),Y	\$B1	2	5 or 6

LDX	Load X register		NZ
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$A2	2	2
Zero page	\$A6	2	3
Zero page,Y	\$B6	2	4
Absolute	\$AE	3	4
Absolute,Y	\$BE	3	4 or 5

LDY	Load Y register		NZ
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$A0	2	2
Zero page	\$A4	2	3
Zero page,X	\$B4	2	4
Absolute	\$AC	3	4
Absolute,X	\$BC	3	4 or 5

LSR	Logical shift right	N = Ø,ZC	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Accumulator	\$4A	1	2
Zero page	\$46	2	5
Zero page,X	\$56	2	6
Absolute	\$4E	3	6
Absolute,X	\$5E	3	7

NOP	No operation	Flags unaltered	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$EA	1	2

ORA	Inclusive OR	NZ	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$09	2	2
Zero page	\$05	2	3
Zero page,X	\$15	2	4
Absolute	\$0D	3	4
Absolute,X	\$1D	3	4 or 5
Absolute,Y	\$19	3	4 or 5
(Indirect,X)	\$01	2	6
(Indirect),Y	\$11	2	5

PHA	Push accumulator	Flags unaltered	
<i>Address mode</i> Implied	<i>Op-code</i> \$48	<i>Bytes</i> 1	<i>Cycles</i> 3
<hr/>		<hr/>	
PHP	Push Status register	Flags unaltered	
<i>Address mode</i> Implied	<i>Op-code</i> \$08	<i>Bytes</i> 1	<i>Cycles</i> 3
<hr/>		<hr/>	
<hr/>		<hr/>	
PLA	Pull accumulator	NZ	
<i>Address mode</i> Implied	<i>Op-code</i> \$68	<i>Bytes</i> 1	<i>Cycles</i> 4
<hr/>		<hr/>	
<hr/>		<hr/>	
PLP	Pull Status register	Flags as status	
<i>Address mode</i> Implied	<i>Op-code</i> \$28	<i>Bytes</i> 1	<i>Cycles</i> 4
<hr/>		<hr/>	

ROL	Rotate left	NZC	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Accumulator	\$2A	1	2
Zero page	\$26	2	5
Zero page,X	\$36	2	6
Absolute	\$2E	3	6
Absolute,X	\$3E	3	7

ROR	Rotate right	NZC	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Accumulator	\$6A	1	2
Zero page	\$66	2	5
Zero page,X	\$76	2	6
Absolute	\$6E	3	6
Absolute,X	\$7E	3	7

RTI	Return from interrupt	Flags as pulled	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$40	1	6

RTS	Return from subroutine	Flags unaltered	
<i>Address mode</i> Implied	<i>Op-code</i> \$60	<i>Bytes</i> 1	<i>Cycles</i> 6
SBC	Subtract from accumulator		NZCV
<i>Address mode</i> Immediate	<i>Op-code</i> \$E9	<i>Bytes</i> 2	<i>Cycles</i> 2
Zero page	\$E5	2	3
Zero page,X	\$F5	2	4
Absolute	\$ED	3	4
Absolute,X	\$FD	3	4 or 5
Absolute,Y	\$F9	3	4 or 5
(Indirect,X)	\$E1	2	6
(Indirect),Y	\$F1	2	5 or 6
SEC	Set Carry flag	C = 1	
<i>Address mode</i> Implied	<i>Op-code</i> \$38	<i>Bytes</i> 1	<i>Cycles</i> 2
SED	Set Decimal flag	D = 1	
<i>Address mode</i> Implied	<i>Op-code</i> \$F8	<i>Bytes</i> 1	<i>Cycles</i> 2

SEI	Set Interrupt flag	I = 1
------------	--------------------	--------------

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$ 78	1	2

STA	Store accumulator	Flags unaltered
------------	-------------------	-----------------

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Zero page	\$ 85	2	3
Zero page,X	\$ 95	2	4
Absolute	\$ 8D	3	4
Absolute,X	\$ 9D	3	5
Absolute,Y	\$ 99	3	5
(Indirect,X)	\$ 81	2	6
(Indirect),Y	\$ 91	2	6

STX	Store X register	Flags unaltered
------------	------------------	-----------------

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Zero page	\$ 86	2	3
Zero page,Y	\$ 96	2	4
Absolute	\$ 8E	3	4

STY	Store Y register	Flags unaltered	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Zero page	\$84	2	3
Zero page,X	\$94	2	4
Absolute	\$8C	3	4
<hr/>			
TAX	Transfer accumulator to X	NZ	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$AA	1	2
<hr/>			
TAY	Transfer accumulator to Y	NZ	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$A8	1	2
<hr/>			
TSX	Transfer Stack Pointer to X	NZ	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$BA	1	2

TXA	Transfer X to accumulator	NZ
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>
Implied	\$8A	1
		<i>Cycles</i>
		2
<hr/>		
TXS	Transfer X to Stack Pointer	Flags unaltered
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>
Implied	\$9A	1
		<i>Cycles</i>
		2
<hr/>		
TYA	Transfer Y to accumulator	NZ
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>
Implied	\$98	1
		<i>Cycles</i>
		2

Appendix 2: 6510 Opcodes

All numbers are hexadecimal.

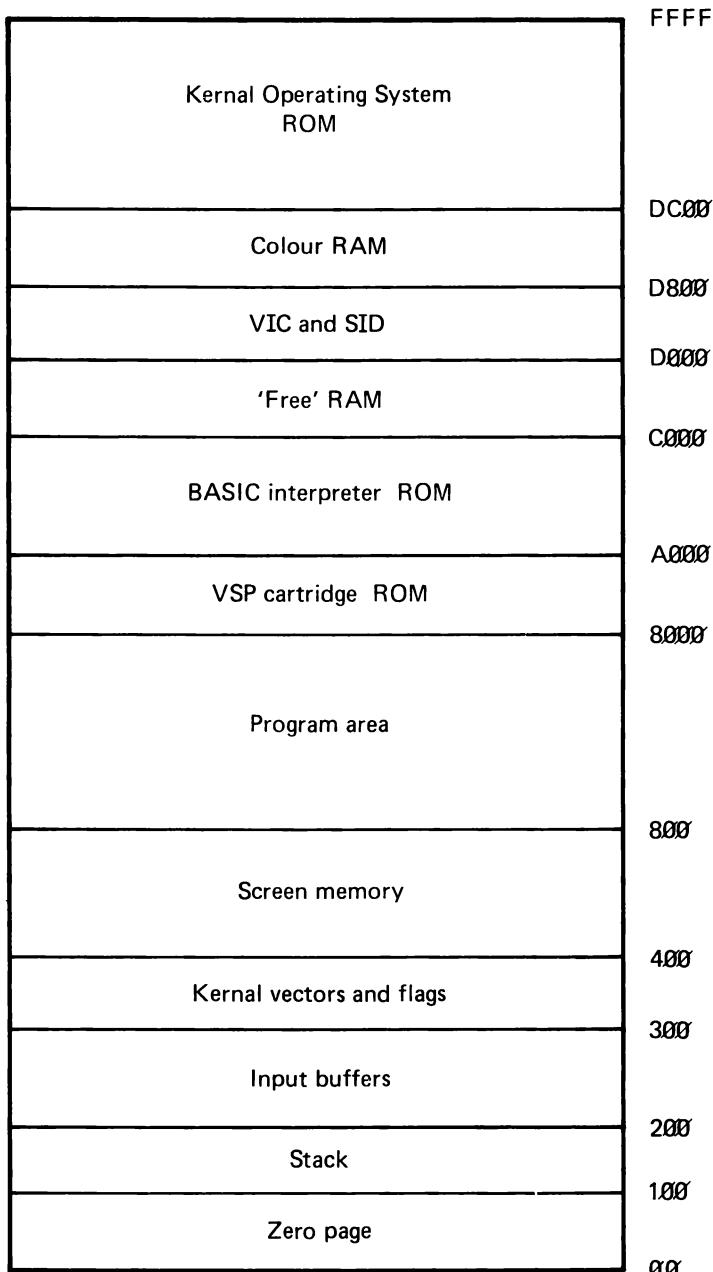
00	BRK implied	1C	Future expansion
01	ORA (zero page, X)	1D	ORA absolute, X
02	Future expansion	1E	ASL absolute, X
03	Future expansion	1F	Future expansion
04	Future expansion	20	JSR absolute
05	ORA zero page	21	AND (zero page, X)
06	ASL zero page	22	Future expansion
07	Future expansion	23	Future expansion
08	PHP implied	24	BIT zero page
09	ORA #immediate	25	AND zero page
0A	ASL accumulator	26	ROL zero page
0B	Future expansion	27	Future expansion
0C	Future expansion	28	PLP implied
0D	ORA absolute	29	AND #immediate
0E	ASL absolute	2A	ROL accumulator
0F	Future expansion	2B	Future expansion
10	BPL relative	2C	BIT absolute
11	ORA (zero page), Y	2D	AND absolute
12	Future expansion	2E	ROL absolute
13	Future expansion	2F	Future expansion
14	Future expansion	30	BMI relative
15	ORA zero page, X	31	AND (zero page), Y
16	ASL zero page, X	32	Future expansion
17	Future expansion	33	Future expansion
18	CLC implied	34	Future expansion
19	ORA absolute, Y	35	AND zero page, X
1A	Future expansion	36	ROL zero page, X
1B	Future expansion	37	Future expansion

38	SEC implied	5D	EOR absolute, X
39	AND absolute, Y	5E	LSR absolute, X
3A	Future expansion	5F	Future expansion
3B	Future expansion	60	RTS implied
3C	Future expansion	61	ADC (zero page, X)
3D	AND absolute, X	62	Future expansion
3E	ROL absolute, X	63	Future expansion
3F	Future expansion	64	Future expansion
40	RTI implied	65	ADC zero page
41	EOR (zero page, X)	66	ROR zero page
42	Future expansion	67	Future expansion
43	Future expansion	68	PLA implied
44	Future expansion	69	ADC #immediate
45	EOR zero page	6A	ROR accumulator
46	LSR zero page	6B	Future expansion
47	Future expansion	6C	JMP (indirect)
48	PHA implied	6D	ADC absolute
49	EOR #immediate	6E	ROR absolute
4A	LSR accumulator	6F	Future expansion
4B	Future expansion	70	BVS relative
4C	JMP absolute	71	ADC (zero page), Y
4D	EOR absolute	72	Future expansion
4E	LSR absolute	73	Future expansion
4F	Future expansion	74	Future expansion
50	BVC relative	75	ADC zero page, X
51	EOR (zero page), Y	76	ROR zero page, X
52	Future expansion	77	Future expansion
53	Future expansion	78	SEI implied
54	Future expansion	79	ADC absolute, Y
55	EOR zero page, X	7A	Future expansion
56	LSR zero page, X	7B	Future expansion
57	Future expansion	7C	Future expansion
58	CLI implied	7D	ADC absolute, X
59	EOR absolute, Y	7E	ROR absolute, X
5A	Future expansion	7F	Future expansion
5B	Future expansion	80	Future expansion
5C	Future expansion	81	STA (zero page, X)

82	Future expansion	A7	Future expansion
83	Future expansion	A8	TAY implied
84	STY zero page	A9	LDA #immediate
85	STA zero page	AA	TAX implied
86	STX zero page	AB	Future expansion
87	Future expansion	AC	LDY absolute
88	DEY implied	AD	LDA absolute
89	Future expansion	AE	LDX absolute
8A	TXA implied	AF	Future expansion
8B	Future expansion	B0	BCS relative
8C	STY absolute	B1	LDA (zero page), Y
8D	STA absolute	B2	Future expansion
8E	STX absolute	B3	Future expansion
8F	Future expansion	B4	LDY zero page, X
90	BCC relative	B5	LDA zero page, X
91	STA (zero page), Y	B6	LDX zero page, Y
92	Future expansion	B7	Future expansion
93	Future expansion	B8	CLV implied
94	STY zero page, X	B9	LDA absolute, Y
95	STA zero page, X	BA	TSX implied
96	STX zero page, Y	BB	Future expansion
97	Future expansion	BC	LDY absolute, X
98	TYA implied	BD	LDA absolute, X
99	STA absolute, Y	BE	LDX absolute, Y
9A	TXS implied	BF	Future expansion
9B	Future expansion	C0	CPY #immediate
9C	Future expansion	C1	CMP (zero page, X)
9D	STA absolute, X	C2	Future expansion
9E	Future expansion	C3	Future expansion
9F	Future expansion	C4	CPY zero page
A0	LDY #immediate	C5	CMP zero page
A1	LDA (zero page, X)	C6	DEC zero page
A2	LDX #immediate	C7	Future expansion
A3	Future expansion	C8	INY implied
A4	LDY zero page	C9	CMP #immediate
A5	LDA zero page	CA	DEX implied
A6	LDX zero page	CB	Future expansion

CC	CPY absolute	E6	INC zero page
CD	CMP absolute	E7	Future expansion
CE	DEC absolute	E8	INX implied
CF	Future expansion	E9	SBC #immediate
D0	BNE relative	EA	NOP implied
D1	CMP (zero page), Y	EB	Future expansion
D2	Future expansion	EC	CPX absolute
D3	Future expansion	ED	SBC absolute
D4	Future expansion	EE	INC absolute
D5	CMP zero page, X	EF	Future expansion
D6	DEC zero page, X	F0	BEQ relative
D7	Future expansion	F1	SBC (zero page), Y
D8	CLD implied	F2	Future expansion
D9	CMP absolute, Y	F3	Future expansion
DA	Future expansion	F4	Future expansion
DB	Future expansion	F5	SBC zero page, X
DC	Future expansion	F6	INC zero page, X
DD	CMP absolute, X	F7	Future expansion
DE	DEC absolute, X	F8	SED implied
DF	Future expansion	F9	SBC absolute, Y
E0	CPX #immediate	FA	Future expansion
E1	SBC (zero page, X)	FB	Future expansion
E2	Future expansion	FC	Future expansion
E3	Future expansion	FD	SBC absolute, X
E4	CPX zero page	FE	INC absolute, X
E5	SBC zero page	FF	Future expansion

Appendix 3: Commodore 64 Memory Map



Appendix 4: Branch Calculators

The branch calculators are used to give branch values in hex. First, count the number of bytes you need to branch. Then locate this number in the centre of the appropriate table, and finally, read off the high and low hex nibbles from the side column and top row respectively.

Example For a backward branch of 16 bytes:

Locate 16 in the centre of Table A4.1 (bottom row), then read off high nibble (#F) and low nibble (#0) to give displacement value (#F0).

Table A4.1 Backward branch calculator

LSD MSD	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
8	128	127	126	125	124	123	122	121	120	119	118	117	116	115	114	113
9	112	111	110	109	108	107	106	105	104	103	102	101	100	99	98	97
A	96	95	94	93	92	91	90	89	88	87	86	85	84	83	82	81
B	80	79	78	77	76	75	74	73	72	71	70	69	68	67	66	65
C	64	63	62	61	60	59	58	57	56	55	54	53	52	51	50	49
D	48	47	46	45	44	43	42	41	40	39	38	37	36	35	34	33
E	32	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17
F	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1

Table A4.2 Forward branch calculator

LSD MSD	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
1	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
2	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47
3	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63
4	64	65	66	67	68	69	70	71	72	73	74	75	76	77	78	79
5	80	81	82	83	84	85	86	87	88	89	90	91	92	93	94	95
6	96	97	98	99	100	101	102	103	104	105	106	107	108	109	110	111
7	112	113	114	115	116	117	118	119	120	121	122	123	124	125	126	127

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