

# CS 97SI: INTRODUCTION TO PROGRAMMING CONTESTS

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# Today's Lecture

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- Algebra
- Number Theory
- Combinatorics
- (non-computational) Geometry
  
- Emphasis on “how to compute”

# Sum of Powers

- $\sum_{k=1}^n k^2 = \frac{1}{6}n(n+1)(2n+1)$
- $\sum k^3 = (\sum k)^2 = \left(\frac{1}{2}n(n+1)\right)^2$
- Pretty useful in many random situations
- Memorize above!

# Fast Exponentiation

- $a^n =$ 
  - 1, if  $n = 0$
  - $a$ , if  $n = 1$
  - $(a^{n/2})^2$ , if  $n$  is even
  - $(a^{(n-1)/2})^2 \cdot a$ , if  $n$  is odd
  
- Can be computed recursively

# Implementation (recursive)

```
double pow(double a, int n) {  
    if(n == 0) return 1;  
    if(n == 1) return a;  
    double t = pow(a, n/2);  
    return t * t * pow(a, n%2);  
}
```

□ Running time:  $O(\log n)$

# Implementation (non-recursive)

```
double pow(double a, int n) {  
    double ret = 1;  
    while(n) {  
        if(n%2 == 1) ret *= a;  
        a *= a; n /= 2;  
    }  
    return ret;  
}
```

□ You should understand how it works

# Linear Algebra

- Gaussian Elimination
  - ▣ Solve a system of linear equations
  - ▣ Invert a matrix
  - ▣ Find the rank of a matrix
  - ▣ Compute the determinant of a matrix
- Cramer's Rule

# Greatest Common Divisor (GCD)

- $\gcd(a, b)$ : greatest integer divides both  $a$  and  $b$
- Used very frequently in number theoretical problems
- Some facts:
  - ▣  $\gcd(a, b) = \gcd(a, b - a)$
  - ▣  $\gcd(a, 0) = a$
  - ▣  $\gcd(a, b)$  is the smallest positive number in  $\{ax + by \mid x, y \in \mathbb{Z}\}$



# Euclidean Algorithm

- Repeated use of  $\gcd(a, b) = \gcd(a, b - a)$
- $\gcd(1989, 867)$ 
  - $= \gcd(1989 - 2 \times 867, 867)$
  - $= \gcd(255, 867)$
  - $= \gcd(255, 867 - 3 \times 255)$
  - $= \gcd(255, 102)$
  - $= \gcd(255 - 2 \times 102, 102)$
  - $= \gcd(51, 102)$
  - $= \gcd(51, 102 - 2 \times 51)$
  - $= \gcd(51, 0)$
  - $= 51$

# Implementation

```
int gcd(int a, int b) {  
    while(b) {int r = a % b; a = b; b = r;}  
    return a;  
}
```

- Running time:  $O(\log(a + b))$
- Be careful:  $a \% b$  follows the sign of  $a$ 
  - ▣  $5 \% 3 == 2$
  - ▣  $-5 \% 3 == -2$

# Congruence & Modulo Operation

- $x \equiv y \pmod{n}$  means  $x$  and  $y$  have the same remainder when divided by  $n$
  
- Multiplicative inverse
  - ▣  $x^{-1}$  is the inverse of  $x$  modulo  $n$  if  $xx^{-1} \equiv 1 \pmod{n}$
  - ▣  $5^{-1} \equiv 3 \pmod{7}$  because  $5 \cdot 3 \equiv 15 \equiv 1 \pmod{7}$
  - ▣ May not exist (e.g. Inverse of 2 mod 4)
  - ▣ Exists iff  $\gcd(x, n) = 1$

# Multiplicative Inverse

- All intermediate numbers computed by Euclidean algorithm are integer combinations of  $a$  and  $b$ 
  - ▣ Therefore,  $\gcd(a, b) = ax + by$  for some integers  $x, y$
  - ▣ If  $\gcd(a, n) = 1$ , then  $ax + ny = 1$  for some  $x, y$
  - ▣ Taking modulo  $n$  gives  $ax \equiv 1 \pmod{n}$
  
- We will be done if we can find such  $x$  and  $y$

# Extended Euclidean Algorithm

- Main idea: keep the original algorithm, but write all intermediate numbers as integer combinations of  $a$  and  $b$
- Exercise: implementation!

# Chinese Remainder Theorem

- Given  $a, b, n, m$  such that  $n$  and  $m$  are coprime
- Find  $x$  such that  $x \equiv a \pmod{m}$ ,  $x \equiv b \pmod{n}$
  
- Solution:
  - ▣ Let  $n^{-1}$  be the inverse of  $n$  modulo  $m$
  - ▣ Let  $m^{-1}$  be the inverse of  $m$  modulo  $n$
  - ▣ Set  $x = ann^{-1} + bmm^{-1}$  (check this yourself)
- Extension: solving for more simultaneous equations

# Binomial Coefficients

- $\binom{n}{k}$  is the number of ways to choose  $k$  objects out of  $n$  distinguishable objects
  - ▣ or, the coefficient of  $x^k y^{n-k}$  in the expansion of  $(x + y)^n$
- Appears everywhere in combinatorics

# Computing $\binom{n}{k}$

- Solution 1: Compute using the following formula

$$\binom{n}{k} = \frac{n(n-1)\cdots(n-k+1)}{k!}$$

- Solution 2: Use Pascal's triangle

- Case 1: Both  $n$  and  $k$  are small

- ▣ Use either solution

- Case 2:  $n$  is big, but  $k$  or  $n - k$  is small

- ▣ Use Solution 1 (carefully)



# Fibonacci sequence

- Definition:
  - $F_0 = 0, F_1 = 1$
  - $F_n = F_{n-1} + F_{n-2}$ , where  $n \geq 2$
- Appears in many different contexts

# Closed Form

- $F_n = \frac{\varphi^n - \bar{\varphi}^n}{\sqrt{5}}$ , where  $\varphi = \frac{1+\sqrt{5}}{2}$
- Bad because  $\varphi$  and  $\sqrt{5}$  are irrational
- Cannot compute the exact value of  $F_n$  for large  $n$
- There is a more stable way to compute  $F_n$ 
  - ▣ ... and any other recurrence of a similar form

# Better Closed Form

- $\begin{bmatrix} F_{n+1} \\ F_n \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} F_n \\ F_{n-1} \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^n \begin{bmatrix} F_1 \\ F_0 \end{bmatrix}$
- Use fast exponentiation to compute the matrix power
- Can be extended to support any linear recurrence with constant coefficients

# Geometry



- In theory: not that hard
- In programming contests: avoid if you can...
- Will cover basic stuff today
  - ▣ Computational geometry in week 9

# When Solving Geometry Problems

- Precision, precision, precision!
  - ▣ If possible, don't use floating-point numbers
  - ▣ If you have to, always use `double` and never use `float`
  - ▣ Avoid division whenever possible
  - ▣ Introduce small constant  $\epsilon$  in (in)equality tests
    - e.g. Instead of `if (x == 0)`, write `if (abs(x) < EPS)`
- No hacks!
  - ▣ In most cases, randomization, probabilistic methods, small perturbations won't help

# 2D Vector Operations

- Have a vector  $(x, y)$
- Norm (distance from the origin):  $\sqrt{x^2 + y^2}$
- Counterclockwise rotation by  $\theta$ :  
Left-multiply by  $\begin{bmatrix} \cos \theta & -\sin \theta \\ \sin \theta & \cos \theta \end{bmatrix}$ 
  - Make sure to use correct units (degrees, radians)
- Normal vectors:  $(y, -x), (-y, x)$
- Memorize all of them!

# Line-Line Intersection

- Have two lines:  $ax + by + c = 0$ ,  $dx + ey + f = 0$
- Write in matrix form:
  - ▣  $\begin{bmatrix} a & b \\ d & e \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = - \begin{bmatrix} c \\ f \end{bmatrix}$
- Invert the matrix on the left using the following:
  - ▣  $\begin{bmatrix} a & b \\ d & e \end{bmatrix}^{-1} = \frac{1}{ae-bd} \begin{bmatrix} e & -b \\ -d & a \end{bmatrix}$
  - ▣ Memorize this!
- Edge case:  $ae = bd$

# Circumcircle of a Triangle

- Have three points  $A, B, C$
- Want to compute  $P$  that is equidistance from  $A, B, C$
- Don't try to solve the system of quadratic equations!
- Instead, do the following:
  - ▣ Find the bisectors of  $AB$  and  $BC$
  - ▣ Compute their intersection



# Area of a Triangle

- Have three points  $A, B, C$
- Use cross product:  $2 \cdot \text{area} = |(B - A) \times (C - A)|$
- Cross product:
  - ▣  $(x_1, y_1) \times (x_2, y_2) = \begin{vmatrix} x_1 & x_2 \\ y_1 & y_2 \end{vmatrix} = x_1 y_2 - x_2 y_1$
  - ▣ Very important in computational geometry. Memorize!

# Area of a Simple Polygon

- Given  $P_1, P_2, \dots, P_n$  around perimeter of polygon  $P$
- If  $P$  is convex, we can decompose  $P$  into triangles:
  - ▣  $2 \cdot \text{area} = \left| \sum_{i=2}^{n-1} (P_{i+1} - P_1) \times (P_i - P_1) \right|$
- It turns out that the formula above works for non-convex polygons too
  - ▣ Area is the absolute value of the sum of “signed area”
- Alternative formula:
  - ▣  $2 \cdot \text{area} = \left| \sum_{i=1}^n (x_i y_{i+1} - x_{i+1} y_i) \right|$

# Conclusion

- No need to look for one-line closed form solutions
- Multi-step algorithms are good enough
  - ▣ This is a CS class, after all
- Have fun with the exercise problems
  - ▣ ... and come to the practice contest if you can!