

Module 8

Topics:

- Iterative structure in Python
- Dictionaries
- Classes

Readings: ThinkP 7

Python can be like Scheme ...

```
## count_down: int[>=0] ->
##                (listof int[>=0])
## Produces the list
## [x, x-1, x-2, ..., 1,0]
def count_down(x):
    if x == 0:
        return [0]
    else:
        return [x] + count_down(x-1)
```

But it can be different → Iteration

```
def count_down(x):          #L1
    answer = []             #L2
    while x >= 0:            #L3
        answer.append(x)    #L4
        x = x - 1           #L5
    return answer            #L6
```

What happens when we call `count_down(3)`?

Calling `count_down(3)`

- L1, L2: $x \leftarrow 3$, $\text{answer} \leftarrow []$
- L3: Since $x \geq 0$, execute L4, L5:
 - $\text{answer} \leftarrow [3]$, $x \leftarrow 2$
- Now, return to L3: since $x \geq 0$, execute L4, L5:
 - $\text{answer} \leftarrow [3, 2]$, $x \leftarrow 1$
- Now, return to L3: since $x \geq 0$, execute L4, L5:
 - $\text{answer} \leftarrow [3, 2, 1]$, $x \leftarrow 0$
- Now, return to L3: since $x \geq 0$, execute L4, L5:
 - $\text{answer} \leftarrow [3, 2, 1, 0]$, $x \leftarrow -1$
- Now, return to L3: since $x < 0$, do not execute L4, L5
- L6: return $[3, 2, 1, 0]$

`while` loop basics

- If the continuation test is **True**,
 - Execute the loop body
- If the continuation test is **False**,
 - Do not execute the loop body
- After completing the loop body:
 - Evaluate the continuation test again
- The body usually includes a mutation of variables used in the continuation test

`while` loop template

```
## initialize loop variables
while test:
    ## body, including statements to:
    ## - update variables used in test
    ## - update value being calculated
## additional processing
```

Steps for writing a **while** loop

You must determine

- how to initialize variables outside the loop
- when the loop body should be executed, or, when it should stop
- what variables must be updated in the loop body so the loop will eventually stop
- what other actions are needed within the loop body

Note: these can be determined in any order – just fill in the template!

Beware of “infinite loops”

```
while 0 == 0:
    print 'runs forever'

x = -5
total = 0
while x < 0:
    total = 2.0 ** x
    x = x-1
print total
```

Notes:

- *it is impossible to write a program that identifies if a loop will run indefinitely (more in CS360)*
- *The code will eventually be terminated in WingIDE with an error – it isn't really “infinite”*

Example: Checking Primality

A number $n \geq 2$ is prime if it has no factors other than 1 and itself.

To test if a number n is prime:

- Check every number from 2 to $n-1$
- If you find a factor of n , stop and return **False**
- If none of them are, stop and return **True**

Implementation of **prime**

```
## is_prime: int[>=2] -> bool
def is_prime (n):
    test_factor = 2
    while test_factor < n:
        if n % test_factor == 0:
            return False
        else:
            test_factor = test_factor + 1
    ## tried all the numbers from 2 to n-1
    return True
```

Testing a **while** loop

Include tests, when possible, for which the body executes

- zero times
- exactly one time
- a “typical” number of times
- the maximum number of times

Also, if the continuation test involves multiple conditions, test each way that the loop may terminate

Testing **is_prime**

Consider the following test cases:

- **n=2** (loop body does not execute)
- **n=3** (loop body executes once, terminates because **test_factor** equals **n**)
- **n=4** (loop terminates because 2 is a factor)
- **n=5** (maximum iterations, no factors found)
- **n=117** (larger composite number)
- **n=127** (larger prime number)

Exercise: **factorial**

Write a Python function to calculate **n**!

- Use a **while** loop that counts from 1 to **n**
- Use a **while** loop that counts down from **n** to 1

Why use loops instead of recursion?

- Iteration, like accumulative recursion, may allow for a more “natural” solution
- Python won’t let us recurse thousands of times
- Iteration is memory efficient
 - for each recursive call, we need memory for parameters
 - for an iterative call, we may just need to update an existing variable
- Iteration will generally run faster

Another type of loop: **for**

- While loops are called *guarded* iteration:
 - If the test evaluates to **True**, execute the body
- Another approach:
 - Iterate over all members in a collection
 - Called *bounded* iteration

```
for item in collection:  
    loop_body
```

for loop examples

```
for food in ['avocado', 'banana',  
            'cabbage']:  
    print food.upper()
```

```
for base in 'ACGGGTCG':  
    print base
```

```
for i in range(2,5):  
    print i*i
```

for and while

while

- Loop counter should be initialized outside loop
- Includes continuation test before body
- Should update loop variables in body of loop
- Body contains steps to repeat

for

- Loop counter initialized automatically
- Continues while more elements in collection
- Loop variable updated automatically – do not update in loop
- Body contains steps to repeat

Multidimensional Lists and Nested Loops

```
M = [[1,2,3], [[4],[5, 6]], [7,8,9],  
      [10, []]]  
print M[0]  
print M[0][2]  
print M[1][1][1]  
print M[3][0]  
print len(M)  
print len(M[0])  
print len(M[1])  
print len(M[1][1])  
print len(M[3][1])
```

```

## nested_max:
##(listof (listof int) [nonempty]) [nonempty]
##      -> int
## find the largest value in a list of lists
## Examples:
## nested_max([[1,5,3], [3], [35,1,2]]) => 35
def nested_max(alol):
    ## set the initial value
    cur_max = alol[0][0]
    for L in alol:      # each list in alol
        for elem in L:  # each value in L
            if elem > cur_max:
                cur_max = elem
    return cur_max

```

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Determining Run-time of code involving loops

- Determine the number steps performed outside the loop
- Determine the number of iterations of the loop
- Determine the number of steps performed in each iteration
- Add everything together to get the overall running time

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What is the run-time of this function
for a list of length n ?

```

def mystery(L):
    M = []
    index_range = range(len(L)):
    for index in index_range:
        rest = L[index:]
        new_val = max(rest)
        M = M + [new_val]
    return M

```

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Application: Searching a list

Suppose you have a list L . How could you determine if a particular value is in that list, if L is in no particular order?

Algorithm (called Linear Search)

- Check the first element in L : is it the one?
 - If Yes, return True
 - Else, check the next value
- The value is not in the list if you don't find it

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Implementing Linear Search

```
## linear_search: (listof X) X
##                  -> bool
## produces True if target is in L,
## False otherwise
## Note: equivalent to: target in L
def linear_search (L, target):
    for val in L:
        if val == target:
            return True
    return False
```

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Running Time of `linear_search`

- Let $n = \text{len}(L)$
- Best Case:
 - If target is in first position, we find it right away $\rightarrow O(1)$
- Worst Case:
 - If target is not in L , we have to check all n elements $\rightarrow O(n)$
 - What is the other worst case?

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Alternatives to Linear Search

- If L is unsorted, we can't do any better than Linear Search.
- How could we improve Linear Search if L was sorted into increasing order?
 - Are there situations in which we could stop earlier?
 - Is this any faster in the worst case?

A better approach: Binary Search

- Suppose L is a listing of the taxpayers in Canada, sorted by Social Insurance numbers.
- Approximately 22,000,000 entries
- Look at $L[11000000]$
- Is it the **target** taxpayer?
 - If yes, stop.
 - If not, is **target** $< L[11000000]$?
 - If yes, then **target** is in the first half of L
 - If not, then **target** is in the second half of L
 - Repeat this process for the half containing **target**

Developing **binary_search**

- We need to determine how to keep track of the section of the list still being searched
 - Variables **beginning**, **end**
 - Determine their initial values
- Determine the **middle** position
- If $L[\text{middle}]$ is target, return **True**
- Otherwise, update **beginning** and **end**
- Determine when we to continue (or stop) searching

Starting the implementation

```
def binary_search(L, target):
    beginning = ...
    end = ...
    while ...:
        middle = ...
        if L[middle] == target:
            return True
        elif L[middle] > target :
            ...
        else:
            ...
    return False
```

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`binary_search` tests should include

- empty list
- list of length 1: target in list and not in list
- small list, both even and odd lengths
- larger list
 - **target** “outside” list, i.e. **target** < **L[0]** or **target** > **L[len(L)-1]**
 - **target** in the list, various positions (first, last, middle)
 - **target** not in the list, value between two list consecutive values

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Worst Case running time of `binary_search`

Suppose $n = 2^k$:

- First comparison reduces list size to 2^{k-1}
- Second comparison reduces list size to 2^{k-2}
- Third comparison reduces list size to 2^{k-3}
- ...
- m^{th} comparison reduces list size to 2^{k-m}
- How many comparisons until we reduce list size to 1?

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Comments on running time for **binary_search**

- Worst case running time is $O(\log n)$
 - For $n \sim 1000$, will consider at most 11 elements ($2^{10} = 1024$)
 - For $n \sim 100,000$, will consider at most 17 elements ($2^{17} = 131072$)
 - For $n \sim 22,000,000$, will consider at most 25 elements ($2^{25} = 33,554,432$)
 - Doubling the size of list requires 1 more comparison worst-case!!!!

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Comments and Questions on running time for **binary_search**

- Binary search isn't faster than linear search all the time (*Q1: in what situation is linear faster?*)
- Could modify it to produce something other than a boolean (*Q2: what would be a good value?*)
- Could be written recursively instead in Python and still have worst case run-time of $O(\log n)$ (*Q3: would be the worst cast for a recursive implementation in Scheme still be $O(\log n)$?*)

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Didn't we do something like this in CS115?

- In CS115, you studied searching in binary search trees (bst)
- You looked at the root value, and either searched the left or right
- The technique is the similar, but
 - In binary search tree search, there is no guarantee that the left and right subtrees were about the same size, so its worst case is not $O(\log n)$

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Goals of Module 8

- Understand that iteration is central to Python
- Understand the difference between `while` and `for` loops
- Be able to write a loop to solve a problem
- Understand how binary search works and why it is much faster than linear search
- Understand the basics of determining run-time of code involving loops